

## CITED REFERENCES AND FURTHER READING:

- Wilkinson, J.H., and Reinsch, C. 1971, *Linear Algebra*, vol. II of *Handbook for Automatic Computation* (New York: Springer-Verlag). [1]
- Smith, B.T., et al. 1976, *Matrix Eigensystem Routines — EISPACK Guide*, 2nd ed., vol. 6 of *Lecture Notes in Computer Science* (New York: Springer-Verlag). [2]
- Stoer, J., and Bulirsch, R. 1980, *Introduction to Numerical Analysis* (New York: Springer-Verlag), §6.5.4. [3]

## 11.6 The QR Algorithm for Real Hessenberg Matrices

Recall the following relations for the  $QR$  algorithm with shifts:

$$\mathbf{Q}_s \cdot (\mathbf{A}_s - k_s \mathbf{1}) = \mathbf{R}_s \quad (11.6.1)$$

where  $\mathbf{Q}$  is orthogonal and  $\mathbf{R}$  is upper triangular, and

$$\begin{aligned} \mathbf{A}_{s+1} &= \mathbf{R}_s \cdot \mathbf{Q}_s^T + k_s \mathbf{1} \\ &= \mathbf{Q}_s \cdot \mathbf{A}_s \cdot \mathbf{Q}_s^T \end{aligned} \quad (11.6.2)$$

The  $QR$  transformation preserves the upper Hessenberg form of the original matrix  $\mathbf{A} \equiv \mathbf{A}_1$ , and the workload on such a matrix is  $O(n^2)$  per iteration as opposed to  $O(n^3)$  on a general matrix. As  $s \rightarrow \infty$ ,  $\mathbf{A}_s$  converges to a form where the eigenvalues are either isolated on the diagonal or are eigenvalues of a  $2 \times 2$  submatrix on the diagonal.

As we pointed out in §11.3, shifting is essential for rapid convergence. A key difference here is that a nonsymmetric real matrix can have complex eigenvalues. This means that good choices for the shifts  $k_s$  may be complex, apparently necessitating complex arithmetic.

Complex arithmetic can be avoided, however, by a clever trick. The trick depends on a result analogous to the lemma we used for implicit shifts in §11.3. The lemma we need here states that if  $\mathbf{B}$  is a nonsingular matrix such that

$$\mathbf{B} \cdot \mathbf{Q} = \mathbf{Q} \cdot \mathbf{H} \quad (11.6.3)$$

where  $\mathbf{Q}$  is orthogonal and  $\mathbf{H}$  is upper Hessenberg, then  $\mathbf{Q}$  and  $\mathbf{H}$  are fully determined by the first column of  $\mathbf{Q}$ . (The determination is unique if  $\mathbf{H}$  has positive subdiagonal elements.) The lemma can be proved by induction analogously to the proof given for tridiagonal matrices in §11.3.

The lemma is used in practice by taking two steps of the  $QR$  algorithm, either with two real shifts  $k_s$  and  $k_{s+1}$ , or with complex conjugate values  $k_s$  and  $k_{s+1} = k_s^*$ . This gives a real matrix  $\mathbf{A}_{s+2}$ , where

$$\mathbf{A}_{s+2} = \mathbf{Q}_{s+1} \cdot \mathbf{Q}_s \cdot \mathbf{A}_s \cdot \mathbf{Q}_s^T \cdot \mathbf{Q}_{s+1}^T \quad (11.6.4)$$

The  $\mathbf{Q}$ 's are determined by

$$\mathbf{A}_s - k_s \mathbf{1} = \mathbf{Q}_s^T \cdot \mathbf{R}_s \quad (11.6.5)$$

$$\mathbf{A}_{s+1} = \mathbf{Q}_s \cdot \mathbf{A}_s \cdot \mathbf{Q}_s^T \quad (11.6.6)$$

$$\mathbf{A}_{s+1} - k_{s+1} \mathbf{1} = \mathbf{Q}_{s+1}^T \cdot \mathbf{R}_{s+1} \quad (11.6.7)$$

Using (11.6.6), equation (11.6.7) can be rewritten

$$\mathbf{A}_s - k_{s+1} \mathbf{1} = \mathbf{Q}_s^T \cdot \mathbf{Q}_{s+1}^T \cdot \mathbf{R}_{s+1} \cdot \mathbf{Q}_s \quad (11.6.8)$$

Hence, if we define

$$\mathbf{M} = (\mathbf{A}_s - k_{s+1} \mathbf{1}) \cdot (\mathbf{A}_s - k_s \mathbf{1}) \quad (11.6.9)$$

equations (11.6.5) and (11.6.8) give

$$\mathbf{R} = \mathbf{Q} \cdot \mathbf{M} \quad (11.6.10)$$

where

$$\mathbf{Q} = \mathbf{Q}_{s+1} \cdot \mathbf{Q}_s \quad (11.6.11)$$

$$\mathbf{R} = \mathbf{R}_{s+1} \cdot \mathbf{R}_s \quad (11.6.12)$$

Equation (11.6.4) can be rewritten

$$\mathbf{A}_s \cdot \mathbf{Q}^T = \mathbf{Q}^T \cdot \mathbf{A}_{s+2} \quad (11.6.13)$$

Thus suppose we can somehow find an upper Hessenberg matrix  $\mathbf{H}$  such that

$$\mathbf{A}_s \cdot \overline{\mathbf{Q}}^T = \overline{\mathbf{Q}}^T \cdot \mathbf{H} \quad (11.6.14)$$

where  $\overline{\mathbf{Q}}$  is orthogonal. If  $\overline{\mathbf{Q}}^T$  has the same first column as  $\mathbf{Q}^T$  (i.e.,  $\overline{\mathbf{Q}}$  has the same first row as  $\mathbf{Q}$ ), then  $\overline{\mathbf{Q}} = \mathbf{Q}$  and  $\mathbf{A}_{s+2} = \mathbf{H}$ .

The first row of  $\mathbf{Q}$  is found as follows. Equation (11.6.10) shows that  $\mathbf{Q}$  is the orthogonal matrix that triangularizes the real matrix  $\mathbf{M}$ . Any real matrix can be triangularized by premultiplying it by a sequence of Householder matrices  $\mathbf{P}_1$  (acting on the first column),  $\mathbf{P}_2$  (acting on the second column), ...,  $\mathbf{P}_{n-1}$ . Thus  $\mathbf{Q} = \mathbf{P}_{n-1} \cdots \mathbf{P}_2 \cdot \mathbf{P}_1$ , and the first row of  $\mathbf{Q}$  is the first row of  $\mathbf{P}_1$  since  $\mathbf{P}_i$  is an  $(i-1) \times (i-1)$  identity matrix in the top left-hand corner. We now must find  $\overline{\mathbf{Q}}$  satisfying (11.6.14) whose first row is that of  $\mathbf{P}_1$ .

The Householder matrix  $\mathbf{P}_1$  is determined by the first column of  $\mathbf{M}$ . Since  $\mathbf{A}_s$  is upper Hessenberg, equation (11.6.9) shows that the first column of  $\mathbf{M}$  has the form  $[p_1, q_1, r_1, 0, \dots, 0]^T$ , where

$$\begin{aligned} p_1 &= a_{11}^2 - a_{11}(k_s + k_{s+1}) + k_s k_{s+1} + a_{12} a_{21} \\ q_1 &= a_{21}(a_{11} + a_{22} - k_s - k_{s+1}) \\ r_1 &= a_{21} a_{32} \end{aligned} \quad (11.6.15)$$

Hence

$$\mathbf{P}_1 = \mathbf{I} - 2\mathbf{w}_1 \cdot \mathbf{w}_1^T \quad (11.6.16)$$

where  $\mathbf{w}_1$  has only its first 3 elements nonzero (cf. equation 11.2.5). The matrix  $\mathbf{P}_1 \cdot \mathbf{A}_s \cdot \mathbf{P}_1^T$  is therefore upper Hessenberg with 3 extra elements:

$$\mathbf{P}_1 \cdot \mathbf{A}_1 \cdot \mathbf{P}_1^T = \begin{bmatrix} \times & \times & \times & \times & \times & \times & \times \\ \times & \times & \times & \times & \times & \times & \times \\ \mathbf{x} & \times & \times & \times & \times & \times & \times \\ \mathbf{x} & \mathbf{x} & \times & \times & \times & \times & \times \\ & & & \times & \times & \times & \times \\ & & & & \times & \times & \times \\ & & & & & \times & \times \end{bmatrix} \quad (11.6.17)$$

This matrix can be restored to upper Hessenberg form without affecting the first row by a sequence of Householder similarity transformations. The first such Householder matrix,  $\mathbf{P}_2$ , acts on elements 2, 3, and 4 in the first column, annihilating elements 3 and 4. This produces a matrix of the same form as (11.6.17), with the 3 extra elements appearing one column over:

$$\begin{bmatrix} \times & \times & \times & \times & \times & \times & \times \\ \times & \times & \times & \times & \times & \times & \times \\ & \times & \times & \times & \times & \times & \times \\ & & \mathbf{x} & \times & \times & \times & \times \\ & & \mathbf{x} & \mathbf{x} & \times & \times & \times \\ & & & & \times & \times & \times \\ & & & & & \times & \times \end{bmatrix} \quad (11.6.18)$$

Proceeding in this way up to  $\mathbf{P}_{n-1}$ , we see that at each stage the Householder matrix  $\mathbf{P}_r$  has a vector  $\mathbf{w}_r$  that is nonzero only in elements  $r$ ,  $r+1$ , and  $r+2$ . These elements are determined by the elements  $r$ ,  $r+1$ , and  $r+2$  in the  $(r-1)$ st column of the current matrix. Note that the preliminary matrix  $\mathbf{P}_1$  has the same structure as  $\mathbf{P}_2, \dots, \mathbf{P}_{n-1}$ .

The result is that

$$\mathbf{P}_{n-1} \cdots \mathbf{P}_2 \cdot \mathbf{P}_1 \cdot \mathbf{A}_s \cdot \mathbf{P}_1^T \cdot \mathbf{P}_2^T \cdots \mathbf{P}_{n-1}^T = \mathbf{H} \quad (11.6.19)$$

where  $\mathbf{H}$  is upper Hessenberg. Thus

$$\bar{\mathbf{Q}} = \mathbf{Q} = \mathbf{P}_{n-1} \cdots \mathbf{P}_2 \cdot \mathbf{P}_1 \quad (11.6.20)$$

and

$$\mathbf{A}_{s+2} = \mathbf{H} \quad (11.6.21)$$

The shifts of origin at each stage are taken to be the eigenvalues of the  $2 \times 2$  matrix in the bottom right-hand corner of the current  $\mathbf{A}_s$ . This gives

$$\begin{aligned} k_s + k_{s+2} &= a_{n-1,n-1} + a_{nn} \\ k_s k_{s+1} &= a_{n-1,n-1} a_{nn} - a_{n-1,n} a_{n,n-1} \end{aligned} \quad (11.6.22)$$

Substituting (11.6.22) in (11.6.15), we get

$$\begin{aligned} p_1 &= a_{21} \{[(a_{nn} - a_{11})(a_{n-1,n-1} - a_{11}) - a_{n-1,n}a_{n,n-1}]/a_{21} + a_{12}\} \\ q_1 &= a_{21}[a_{22} - a_{11} - (a_{nn} - a_{11}) - (a_{n-1,n-1} - a_{11})] \\ r_1 &= a_{21}a_{32} \end{aligned} \quad (11.6.23)$$

We have judiciously grouped terms to reduce possible roundoff when there are small off-diagonal elements. Since only the ratios of elements are relevant for a Householder transformation, we can omit the factor  $a_{21}$  from (11.6.23).

In summary, to carry out a double  $QR$  step we construct the Householder matrices  $\mathbf{P}_r$ ,  $r = 1, \dots, n-1$ . For  $\mathbf{P}_1$  we use  $p_1$ ,  $q_1$ , and  $r_1$  given by (11.6.23). For the remaining matrices,  $p_r$ ,  $q_r$ , and  $r_r$  are determined by the  $(r, r-1)$ ,  $(r+1, r-1)$ , and  $(r+2, r-1)$  elements of the current matrix. The number of arithmetic operations can be reduced by writing the nonzero elements of the  $2\mathbf{w} \cdot \mathbf{w}^T$  part of the Householder matrix in the form

$$2\mathbf{w} \cdot \mathbf{w}^T = \begin{bmatrix} (p \pm s)/(\pm s) \\ q/(\pm s) \\ r/(\pm s) \end{bmatrix} \cdot [1 \quad q/(p \pm s) \quad r/(p \pm s)] \quad (11.6.24)$$

where

$$s^2 = p^2 + q^2 + r^2 \quad (11.6.25)$$

(We have simply divided each element by a piece of the normalizing factor; cf. the equations in §11.2.)

If we proceed in this way, convergence is usually very fast. There are two possible ways of terminating the iteration for an eigenvalue. First, if  $a_{n,n-1}$  becomes “negligible,” then  $a_{nn}$  is an eigenvalue. We can then delete the  $n$ th row and column of the matrix and look for the next eigenvalue. Alternatively,  $a_{n-1,n-2}$  may become negligible. In this case the eigenvalues of the  $2 \times 2$  matrix in the lower right-hand corner may be taken to be eigenvalues. We delete the  $n$ th and  $(n-1)$ st rows and columns of the matrix and continue.

The test for convergence to an eigenvalue is combined with a test for negligible subdiagonal elements that allows splitting of the matrix into submatrices. We find the largest  $i$  such that  $a_{i,i-1}$  is negligible. If  $i = n$ , we have found a single eigenvalue. If  $i = n-1$ , we have found two eigenvalues. Otherwise we continue the iteration on the submatrix in rows  $i$  to  $n$  ( $i$  being set to unity if there is no small subdiagonal element).

After determining  $i$ , the submatrix in rows  $i$  to  $n$  is examined to see if the *product* of any two consecutive subdiagonal elements is small enough that we can work with an even smaller submatrix, starting say in row  $m$ . We start with  $m = n-2$  and decrement it down to  $i+1$ , computing  $p$ ,  $q$ , and  $r$  according to equations (11.6.23) with 1 replaced by  $m$  and 2 by  $m+1$ . If these were indeed the elements of the special “first” Householder matrix in a double  $QR$  step, then applying the Householder matrix would lead to nonzero elements in positions  $(m+1, m-1)$ ,  $(m+2, m-1)$ , and  $(m+2, m)$ . We require that the first two of these elements be

small compared with the local diagonal elements  $a_{m-1,m-1}$ ,  $a_{mm}$  and  $a_{m+1,m+1}$ . A satisfactory approximate criterion is

$$|a_{m,m-1}|(|q| + |r|) \ll |p|(|a_{m+1,m+1}| + |a_{mm}| + |a_{m-1,m-1}|) \quad (11.6.26)$$

Very rarely, the procedure described so far will fail to converge. On such matrices, experience shows that if one double step is performed with any shifts that are of order the norm of the matrix, convergence is subsequently very rapid. Accordingly, if ten iterations occur without determining an eigenvalue, the usual shifts are replaced for the next iteration by shifts defined by

$$\begin{aligned} k_s + k_{s+1} &= 1.5 \times (|a_{n,n-1}| + |a_{n-1,n-2}|) \\ k_s k_{s+1} &= (|a_{n,n-1}| + |a_{n-1,n-2}|)^2 \end{aligned} \quad (11.6.27)$$

The factor 1.5 was arbitrarily chosen to lessen the likelihood of an “unfortunate” choice of shifts. This strategy is repeated after 20 unsuccessful iterations. After 30 unsuccessful iterations, the routine reports failure.

The operation count for the  $QR$  algorithm described here is  $\sim 5k^2$  per iteration, where  $k$  is the current size of the matrix. The typical average number of iterations per eigenvalue is  $\sim 1.8$ , so the total operation count for all the eigenvalues is  $\sim 3n^3$ . This estimate neglects any possible efficiency due to splitting or sparseness of the matrix.

The following routine `hqr` is based algorithmically on the above description, in turn following the implementations in [1,2].

```

SUBROUTINE hqr(a,n,np,wr,wi)
INTEGER n,np
REAL a(np,np),wi(np),wr(np)
  Finds all eigenvalues of an n by n upper Hessenberg matrix a that is stored in an np by np
  array. On input a can be exactly as output from elmhes §11.5; on output it is destroyed.
  The real and imaginary parts of the eigenvalues are returned in wr and wi, respectively.
INTEGER i,its,j,k,l,m,nn
REAL anorm,p,q,r,s,t,u,v,w,x,y,z
anorm=0.
do 12 i=1,n
  do 11 j=max(i-1,1),n
    anorm=anorm+abs(a(i,j))
  enddo 11
enddo 12
nn=n
t=0.
1 if(nn.ge.1)then
  its=0
2 do 13 l=nn,2,-1
  s=abs(a(l-1,l-1))+abs(a(l,l))
  if(s.eq.0.)s=anorm
  if(abs(a(l,l-1))+s.eq.s)then
    abs(a(l,l-1))=0.
    goto 3
  endif
enddo 13
3 x=a(nn,nn)
if(l.eq.nn)then
  wr(nn)=x+t
  wi(nn)=0.
  nn=nn-1
  Compute matrix norm for possible use
  in locating single small subdiagonal
  element.
  Gets changed only by an exceptional shift.
  Begin search for next eigenvalue.
  Begin iteration: look for single small sub-
  diagonal element.
  One root found.

```

```

else
  y=a(nn-1,nn-1)
  w=a(nn,nn-1)*a(nn-1,nn)
  if(1.eq.nn-1)then
    p=0.5*(y-x)
    q=p**2+w
    z=sqrt(abs(q))
    x=x+t
    if(q.ge.0.)then
      z=p+sign(z,p)
      wr(nn)=x+z
      wr(nn-1)=wr(nn)
      if(z.ne.0.)wr(nn)=x-w/z
      wi(nn)=0.
      wi(nn-1)=0.
    else
      wr(nn)=x+p
      wr(nn-1)=wr(nn)
      wi(nn)=z
      wi(nn-1)=-z
    endif
    nn=nn-2
  else
    if(its.eq.30)pause 'too many iterations in hqr'
    if(its.eq.10.or.its.eq.20)then
      t=t+x
      do 14 i=1,nn
        a(i,i)=a(i,i)-x
      enddo 14
      s=abs(a(nn,nn-1))+abs(a(nn-1,nn-2))
      x=0.75*s
      y=x
      w=-0.4375*s**2
    endif
    its=its+1
    do 15 m=nn-2,1,-1
      z=a(m,m)
      r=x-z
      s=y-z
      p=(r*s-w)/a(m+1,m)+a(m,m+1)
      q=a(m+1,m+1)-z-r-s
      r=a(m+2,m+1)
      s=abs(p)+abs(q)+abs(r)
      p=p/s
      q=q/s
      r=r/s
      if(m.eq.1)goto 4
      u=abs(a(m,m-1))*(abs(q)+abs(r))
      v=abs(p)*(abs(a(m-1,m-1))+abs(z)+abs(a(m+1,m+1)))
      if(u+v.eq.v)goto 4
    enddo 15
    do 16 i=m+2,nn
      a(i,i-2)=0.
      if(i.ne.m+2) a(i,i-3)=0.
    enddo 16
    do 19 k=m,nn-1
      if(k.ne.m)then
        p=a(k,k-1)
        q=a(k+1,k-1)
        r=0.
        if(k.ne.nn-1)r=a(k+2,k-1)
        x=abs(p)+abs(q)+abs(r)
        if(x.ne.0.)then
          p=p/x

```

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```

        q=q/x
        r=r/x
    endif
endif
s=sign(sqrt(p**2+q**2+r**2),p)
if(s.ne.0.)then
    if(k.eq.m)then
        if(l.ne.m)a(k,k-1)=-a(k,k-1)
    else
        a(k,k-1)=-s*x
    endif
    p=p+s          Equations (11.6.24).
    x=p/s
    y=q/s
    z=r/s
    q=q/p
    r=r/p
    do 17 j=k,nn   Row modification.
        p=a(k,j)+q*a(k+1,j)
        if(k.ne.nn-1)then
            p=p+r*a(k+2,j)
            a(k+2,j)=a(k+2,j)-p*z
        endif
        a(k+1,j)=a(k+1,j)-p*y
        a(k,j)=a(k,j)-p*x
    enddo 17
    do 18 i=1,min(nn,k+3) Column modification.
        p=x*a(i,k)+y*a(i,k+1)
        if(k.ne.nn-1)then
            p=p+z*a(i,k+2)
            a(i,k+2)=a(i,k+2)-p*r
        endif
        a(i,k+1)=a(i,k+1)-p*q
        a(i,k)=a(i,k)-p
    enddo 18
endif
enddo 19
goto 2          ...for next iteration on current eigenvalue.
endif
endif
goto 1          ...for next eigenvalue.
endif
return
END

```

## CITED REFERENCES AND FURTHER READING:

- Wilkinson, J.H., and Reinsch, C. 1971, *Linear Algebra*, vol. II of *Handbook for Automatic Computation* (New York: Springer-Verlag). [1]
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## 11.7 Improving Eigenvalues and/or Finding Eigenvectors by Inverse Iteration

The basic idea behind inverse iteration is quite simple. Let  $\mathbf{y}$  be the solution of the linear system

$$(\mathbf{A} - \tau \mathbf{1}) \cdot \mathbf{y} = \mathbf{b} \quad (11.7.1)$$

where  $\mathbf{b}$  is a random vector and  $\tau$  is close to some eigenvalue  $\lambda$  of  $\mathbf{A}$ . Then the solution  $\mathbf{y}$  will be close to the eigenvector corresponding to  $\lambda$ . The procedure can be iterated: Replace  $\mathbf{b}$  by  $\mathbf{y}$  and solve for a new  $\mathbf{y}$ , which will be even closer to the true eigenvector.

We can see why this works by expanding both  $\mathbf{y}$  and  $\mathbf{b}$  as linear combinations of the eigenvectors  $\mathbf{x}_j$  of  $\mathbf{A}$ :

$$\mathbf{y} = \sum_j \alpha_j \mathbf{x}_j \quad \mathbf{b} = \sum_j \beta_j \mathbf{x}_j \quad (11.7.2)$$

Then (11.7.1) gives

$$\sum_j \alpha_j (\lambda_j - \tau) \mathbf{x}_j = \sum_j \beta_j \mathbf{x}_j \quad (11.7.3)$$

so that

$$\alpha_j = \frac{\beta_j}{\lambda_j - \tau} \quad (11.7.4)$$

and

$$\mathbf{y} = \sum_j \frac{\beta_j \mathbf{x}_j}{\lambda_j - \tau} \quad (11.7.5)$$

If  $\tau$  is close to  $\lambda_n$ , say, then provided  $\beta_n$  is not accidentally too small,  $\mathbf{y}$  will be approximately  $\mathbf{x}_n$ , up to a normalization. Moreover, each iteration of this procedure gives another power of  $\lambda_j - \tau$  in the denominator of (11.7.5). Thus the convergence is rapid for well-separated eigenvalues.

Suppose at the  $k$ th stage of iteration we are solving the equation

$$(\mathbf{A} - \tau_k \mathbf{1}) \cdot \mathbf{y} = \mathbf{b}_k \quad (11.7.6)$$

where  $\mathbf{b}_k$  and  $\tau_k$  are our current guesses for some eigenvector and eigenvalue of interest (let's say,  $\mathbf{x}_n$  and  $\lambda_n$ ). Normalize  $\mathbf{b}_k$  so that  $\mathbf{b}_k \cdot \mathbf{b}_k = 1$ . The exact eigenvector and eigenvalue satisfy

$$\mathbf{A} \cdot \mathbf{x}_n = \lambda_n \mathbf{x}_n \quad (11.7.7)$$

so

$$(\mathbf{A} - \tau_k \mathbf{1}) \cdot \mathbf{x}_n = (\lambda_n - \tau_k) \mathbf{x}_n \quad (11.7.8)$$

Since  $\mathbf{y}$  of (11.7.6) is an improved approximation to  $\mathbf{x}_n$ , we normalize it and set

$$\mathbf{b}_{k+1} = \frac{\mathbf{y}}{|\mathbf{y}|} \quad (11.7.9)$$