

Homework 1 CS2233

1. Propositions

- $p := \text{"You passed CS2233"}$
- $q := \text{"You passed CS3333"}$
- $r := \text{"You can register for CS3343"}$
- $s := \text{"You understand Boolean Algebra"}$

a: You have not passed CS2233 but you understand boolean algebra.

Rewrite : You understand boolean algebra and you have not passed CS2233.

$$s \wedge \neg p$$

b: You can not register for CS3343 only if you have not passed both CS2233 and CS3333

Rewrite : If you have not passed CS2233 and CS3333 then you cannot register for CS3343

$$\neg p \wedge \neg q \rightarrow \neg r$$

c: If you can register for CS3343 then you have passed CS2233 and you understand boolean algebra if you passed CS2233.

Rewrite: If you passed CS2233 then you can register for CS3343 and you understand boolean algebra.

$$p \rightarrow (r \wedge s)$$

2: Equivalences

1: Show that $(\neg q \wedge (p \vee \neg p)) \rightarrow \neg q$ is a tautology (i.e. $(\neg q \wedge (p \vee \neg p)) \rightarrow \neg q \equiv T$)

1a: Show using a truth table

p	q	$\neg q$	$(\neg q \wedge p)$	$(\neg q \wedge p) \rightarrow \neg q$	T	$(\neg q \wedge p) \rightarrow \neg q$	T
T	T	F	F	T	T	T	T
T	F	T	T	T	T	T	T
F	T	F	F	T	T	T	T
F	F	T	T	T	T	T	T

1b: Show using a sequence of equivalences

$$\begin{aligned}
 & (\neg q \wedge p) \rightarrow \neg q \quad T \\
 & \neg(\neg q \wedge p) \rightarrow \neg q \quad T \\
 \text{LHS} &= \neg(\neg q \wedge p) \rightarrow \neg q, \text{ RHS} = T \\
 \text{LHS} &= \neg(\neg q) \rightarrow \neg p \rightarrow \neg q \\
 \text{LHS} &= q \rightarrow \neg p \rightarrow \neg q \\
 \text{LHS} &= T \rightarrow \neg p \\
 \text{LHS} &= T \\
 \text{LHS} &= \text{RHS} \quad \text{Tautology}
 \end{aligned}$$

Idempotent Law : $p \vee p = p$
Table 7 Line 1
Distribute \neg
Double Negation Law : $\neg(\neg q) = q$
Negation Law : $p \rightarrow \neg p = T$
Domination law : $p \vee T = T$

1b alternate sequence:

Idempotent Law	$p \vee p = p$	$\neg q \vee p \vee \neg q$
Table 7, Line 1	$u \vee \neg w = \neg u \vee \neg w, u = (\neg q \vee p), w = (\neg q)$	$\neg(\neg q \vee p) \vee \neg q$
De Morgan's Law	$\neg(\neg q \vee p) = q \vee \neg p$	$q \vee \neg p \vee \neg q$
Negation Law	$q \vee \neg q = T$	$\neg p \vee T$
Domination Laws	$u \vee T = T, u = (\neg p)$	T

2: Show that $\neg q \rightarrow (p \vee r) \equiv (\neg q \rightarrow r) \wedge (q \vee p)$

2a: Show using Truth tables

p	q	r	$\neg q$	$p \vee r$	$\neg q \rightarrow (p \vee r)$	$\neg q \rightarrow r$	$(q \vee p) \wedge (\neg q \rightarrow r)$
T	T	T	F	T	T	T	T
T	T	F	F	T	T	F	F
T	F	T	T	T	T	T	T
T	F	F	T	F	F	F	F
F	T	T	F	T	T	T	T
F	T	F	F	T	T	F	F
F	F	T	T	T	T	T	T
F	F	F	T	F	F	F	F

2b: Show using a sequence of equivalences

$$\neg q \rightarrow (p \vee r) \equiv (\neg q \rightarrow r) \wedge (q \vee p)$$

$q \vee (p \vee r)$ Table 7, line 1
 $(q \vee p) \wedge (q \vee r)$ Distributive Law
 LHS = $(q \vee p) \wedge (q \vee r)$

RHS = $(\neg q \rightarrow r) \wedge (p \vee q)$ Commutative law
 $(q \vee r) \wedge (p \vee q)$ Table 7 Line 1
 RHS = $(q \vee p) \wedge (q \vee r)$
 LHS = RHS
 Tautology

2b alternate sequence :

Table 7, Line 1	$u \rightarrow w = \neg u \vee w$	$\neg q \rightarrow (p \vee r) \equiv (\neg q \rightarrow r) \wedge (q \vee p)$
Double Negation Law	$\neg(\neg q) = q$	$\neg(\neg q) \vee (p \vee r) \equiv (\neg(\neg q) \vee r) \wedge (q \vee p)$
Distributive Law	$(q \vee r) \wedge (q \vee p) = q \vee (r \wedge p)$	$q \vee (p \vee r) \equiv (q \vee p) \wedge (q \vee r)$

3. Additional Operators

1: Using truth tables, show that \wedge is associative. i.e., $p \wedge (q \wedge r) \equiv (p \wedge q) \wedge r$

p	q	r	$q \wedge r$	$(p \wedge (q \wedge r))$	p	q	$((p \wedge q) \wedge r)$	$((p \wedge q) \wedge r)$	$(p \wedge (q \wedge r))$
T	T	T	T	T	T	T	T	T	T
T	T	F	F	F	T	F	F	F	F
T	F	T	F	F	F	T	F	F	F
T	F	F	F	F	F	F	F	F	F
F	T	T	T	F	F	T	F	F	F
F	T	F	F	F	F	F	F	F	F
F	F	T	F	F	F	T	F	F	F
F	F	F	F	F	F	F	F	F	F

2: Using truth tables show that \uparrow is not associative,
i.e., $p \uparrow (q \uparrow r) \neq (p \uparrow q) \uparrow r$

p	q	r	$q \uparrow r$	$p \uparrow (q \uparrow r)$	p	q	$(p \uparrow q) \uparrow r$	$((p \uparrow q) \uparrow r)$	$(p \uparrow (q \uparrow r))$
T	T	T	F	T	T	T	T	T	T
T	T	F	T	F	T	F	T	F	F
T	F	T	T	F	T	T	F	T	T
T	F	F	T	F	T	F	T	F	F
F	T	T	F	T	F	T	F	F	T
F	T	F	T	T	F	F	T	T	T
F	F	T	T	T	F	T	F	F	T
F	F	F	T	T	F	F	T	T	T