

Final Assignment : Schedule Search

Student: Level:

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1 Introduction

This report presents the analysis and scheduling of a given task set using three different approaches: preemptive scheduling, non-preemptive scheduling and non-preemptive scheduling with permutations. The task set consists of 7 periodic tasks with varying computation times and periods.

Link to the git repository to access the project codes:

https://github.com/JulienMaeckg/Project_In422

2 Task Set and Schedulability Analysis

The given task set is:

Task	C	T_i
$ au_1$	2	10
$ au_2$	3	10
$ au_3$	2	20
$ au_4$	2	20
$ au_5$	2	40
$ au_6$	2	40
$ au_7$	3	80

Table 1 – Task Set Parameters

2.1 Schedulability Test

The schedulability is first checked using this condition:

$$\sum_{i=1}^{n} \frac{C_i}{T_i} \le 1$$

For our task set:

$$\frac{2}{10} + \frac{3}{10} + \frac{2}{20} + \frac{2}{20} + \frac{2}{40} + \frac{2}{40} + \frac{3}{80} = 0.2 + 0.3 + 0.1 + 0.1 + 0.05 + 0.05 + 0.0375 = 0.8375$$

Since $0.8375 \le 1$, the task set is schedulable.

3 Hyperperiod Calculation

The hyperperiod is the least common multiple (LCM) of all task periods, which represents the time after which the schedule repeats. For our task set:

$$LCM(10, 10, 20, 20, 40, 40, 80) = 80$$

All scheduling implementations use this hyperperiod to generate complete schedules.

4 Scheduling Approaches

4.1 Preemptive Scheduling

The C++ implementation (schedule_preemptive.cpp) uses a preemptive approach :

- Tasks can be interrupted at any time
- At each time unit, the scheduler selects the task with the smallest remaining time to deadline

Here is the schedule calculated with this algorithm:

Hyperperiod: 80
11222334451122256677112223344711222-1-1-1-1-1-1112223344511222566-1-1112223344-111222-1-1-1-1-1-1

FIGURE 1 – Preemptive schedule visualization $(-1 = idle \ time)$

4.2 Non-Preemptive Scheduling

The first implementation (schedule_non_preemptive.py) uses a non-preemptive scheduling algorithm that:

- Calculates the hyperperiod (80 time units)
- At each time step, checks which tasks are ready to execute
- Selects the task with the smallest remaining time until its deadline
- Executes the selected task to completion without interruption

Here is the schedule calculated with this algorithm:

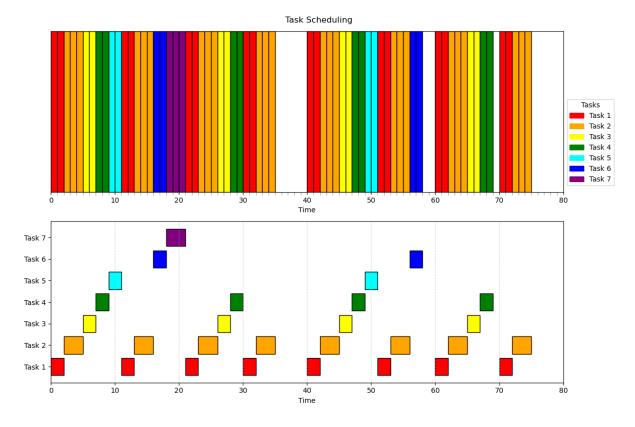


FIGURE 2 – Non-preemptive schedule visualization

0.8375 ==> Schedulable
Waiting time: 140
Responding time: 140

FIGURE 3 – Waiting time & Respond time of this schedule

4.3 Non-Preemptive with Permutations

The second implementation (schedule_non_preemptive_permutations.py) explores all possible task execution orders to find the schedule with minimum total response time :

- Generates all permutations of task execution orders
- For each permutation, simulates the schedule
- Tracks response times for all tasks
- Selects the permutation with the smallest total response time

The computational complexity is $O(n! \times \text{hyperperiod})$ (whith n the number of tasks), making it impractical for large task sets but suitable for our 7-task set.

Here is the schedule calculated with this algorithm:

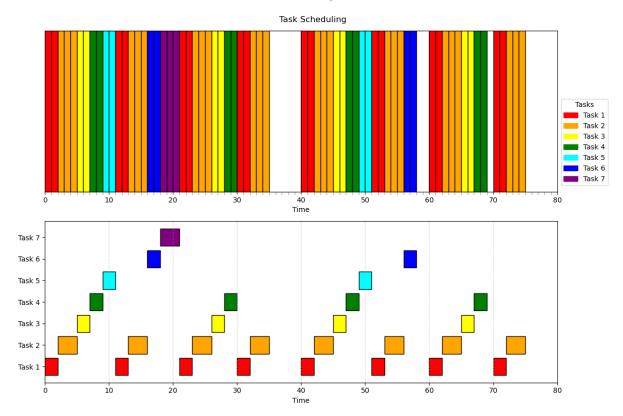


FIGURE 4 – Non-preemptive schedule visualization

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0.8375 ==> Schedulable
Minimum total response time: 140
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FIGURE 5 – Respond time of this schedule

We notice that the previous algorithm found the same scheduling.

5 Results and Analysis

5.1 Performance Metrics

All implementations were evaluated based on:

- Total waiting time
- Processor idle time
- Response times for individual jobs

The non-preemptive with and without permutations approach found the optimal schedule with minimum total response time of 140 units.

6 Computational Complexity

- Preemptive : $O(\text{hyperperiod} \times n)$
- Non-preemptive : $O(\text{hyperperiod} \times n)$
- Non-preemptive with permutations : $O(n! \times \text{hyperperiod})$

With n the number of tasks.

7 Conclusion

The task set was successfully scheduled using three different approaches. The preemptive implementation provided a more responsive schedule while the non-preemptive permutation method found the optimal schedule with minimum waiting time. All implementations confirmed the schedulability of the task set under normal conditions.