

DSK814: Algorithms and Data Structures

SDU - Kolding

Algorithm Analysis

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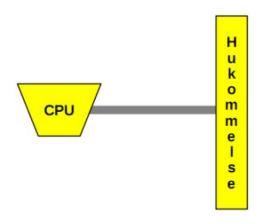
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In this course, we will describe existing algorithms, develop new algorithms, and analyze algorithms for correctness and quality.

We need for Algorithm Analysis:

- ▶ Model of problem. Individual for each problem.
- ▶ Model of machine. We are using the RAM model.
- ▶ Quality goals. We focus (mostly) on time consumption.
- ► Mathematical analysis tools: Loop invariants, induction, recursion equations, asymptotic notation.

The RAM model



- ► A CPU
- ► A memory (~infinite array of cells).
- ► A number of basic operations:add,sub,shift,compare,move data item,jump into program(loop,branching,method call). These are all assumed to take the same amount of time.
- ▶ Time for an algorithm: number of basic operations performed.
- ► Space for an algorithm: maximum number of occupied memory cells.

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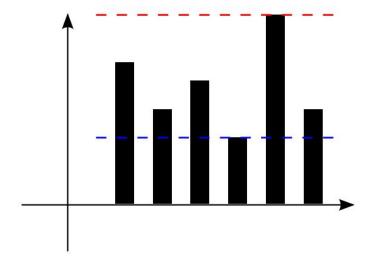
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8,7,6,5,4,3,2,1

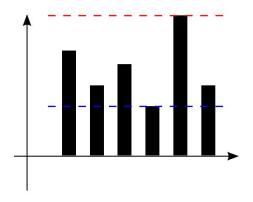
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What inputs should we use to evaluate the time consumption?

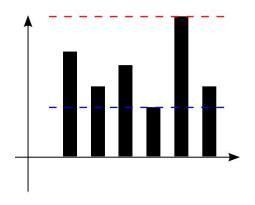
- Worst case(max over all inputs of size n)
- ► Average case (average over a distribution of inputs of size n)
- ► Best case(min over all inputs of size n)



Running time for the different input sizes *n*

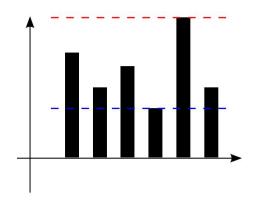


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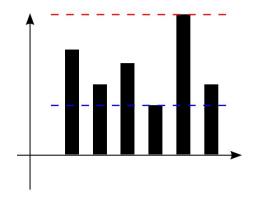
Average case: What is the distribution of input? Why is this distribution realistic/relevant? Analysis is often (mathematically) difficult to carry out.



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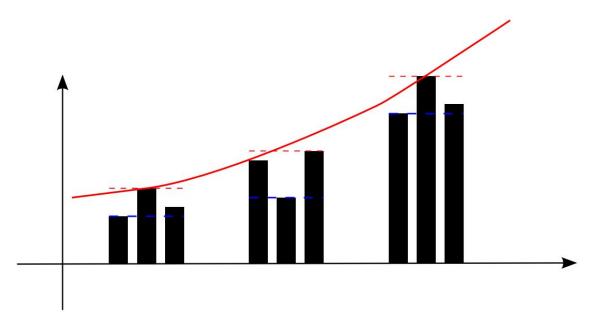
Best case: Often does not provide much relevant information.



Almost all analyses in this course are worst case.

Different input sizes

Worst-case running time is usually an increasing function of the input size n:

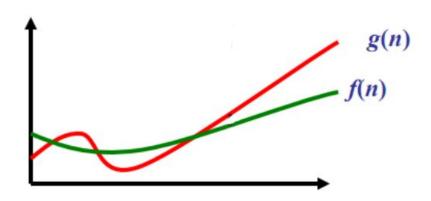


Running time for the various inputs of increasing size n

Growth rate

An analysis must therefore provide a function f(n) of the input size n.

We therefore need to compare functions. The relevant target is the growth rate of f(n) when n increases - a faster growing function will always overtake a slower growing function when n becomes big enough. And too small n(almost) all algorithms are fast.



Growth rate

Examples (increasing growth rate):

1,
$$\log n$$
, \sqrt{n} , n , $n \log n$, $n\sqrt{n}$, n^2 , n^3 , n^{10} , 2^n

Later: definition of asymptotic growth rate for functions and comparisons of them.