

# The Relational Model

Chapter #3  
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## Lecture 3 - Objectives

- Terminology of relational model.
- How tables are used to represent data.
- Connection between mathematical relations and relations in the relational model.
- Properties of database relations.
- How to identify candidate, primary, and foreign keys.
- Meaning of entity integrity and referential integrity.
- Purpose and advantages of views.

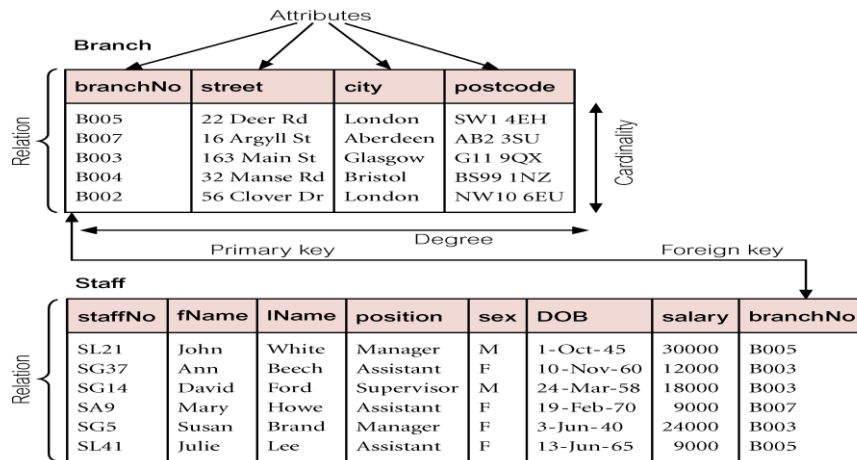
## Relational Model Terminology

- A **relation** is a table with columns and rows.
  - Only applies to logical structure of the database, not the physical structure.
- **Attribute** is a named column of a relation.
- **Domain** is the set of allowable values for one or more attributes.

## Relational Model Terminology

- **Tuple** is a row of a relation.
- **Degree** is the number of attributes in a relation.
- **Cardinality** is the number of tuples in a relation.
- **Relational Database** is a collection of normalized relations with distinct relation names.

## Instances of Branch and Staff (part) Relations



## Examples of Attribute Domains

Attribute	Domain Name	Meaning	Domain Definition
branchNo	BranchNumbers	The set of all possible branch numbers	character: size 4, range B001–B999
street	StreetNames	The set of all street names in Britain	character: size 25
city	CityNames	The set of all city names in Britain	character: size 15
postcode	Postcodes	The set of all postcodes in Britain	character: size 8
sex	Sex	The sex of a person	character: size 1, value M or F
DOB	DatesOfBirth	Possible values of staff birth dates	date, range from 1-Jan-20, format dd-mmm-yy
salary	Salaries	Possible values of staff salaries	monetary: 7 digits, range 6000.00–40000.00

## Alternative Terminology for Relational Model

**Table 3.1** Alternative terminology for relational model terms.

Formal terms	Alternative 1	Alternative 2
Relation	Table	File
Tuple	Row	Record
Attribute	Column	Field

## Database Relations

- **Relation schema**
  - Named relation defined by a set of attribute and domain name pairs.
- **Relational database schema**
  - Set of relation schemas, each with a distinct name.

## Properties of Relations

- Relation name is distinct from all other relation names in relational schema.
- Each cell of relation contains exactly one atomic (single) value.
- Each attribute has a distinct name.
- Values of an attribute are all from the same domain.

## Properties of Relations

- Each tuple is distinct; there are no duplicate tuples.
- Order of attributes has no significance.
- Order of tuples has no significance, theoretically.

## Relational Keys

- **Superkey**
  - An attribute, or a set of attributes, that uniquely identifies a tuple within a relation.
- **Candidate Key**
  - Superkey (K) such that no proper subset is a superkey within the relation.
  - In each tuple of R, values of K uniquely identify that tuple (uniqueness).
  - No proper subset of K has the uniqueness property (irreducibility).

## Relational Keys

- **Primary Key**
  - Candidate key selected to identify tuples uniquely within relation.
- **Alternate Keys**
  - Candidate keys that are not selected to be primary key.
- **Foreign Key**
  - Attribute, or set of attributes, within one relation that matches candidate key of some (possibly same) relation.

## Relational Integrity

- **Null**
  - Represents value for an attribute that is currently unknown or not applicable for tuple
  - Deals with incomplete or exceptional data.
  - Represents the absence of a value and is not the same as zero or spaces, which are values.

## Relational Integrity

- **Entity Integrity**
  - In a base relation, no attribute of a primary key can be null.
- **Referential Integrity**
  - If foreign key exists in a relation, either foreign key value must match a candidate key value of some tuple in its home relation or foreign key value must be wholly null.

## Relational Integrity

- **Enterprise Constraints**
  - Additional rules specified by users or database administrators.

## Views

- **Base Relation**
  - Named relation corresponding to an entity in conceptual schema, whose tuples are physically stored in database.
- **View**
  - Dynamic result of one or more relational operations operating on base relations to produce another relation.



## Views

- A virtual relation that does not necessarily actually exist in the database but is produced upon request, at time of request.
- Contents of a view are defined as a query on one or more base relations.
- Views are dynamic, meaning that changes made to base relations that affect view attributes are immediately reflected in the view.

## Purpose of Views

- Provides powerful and flexible security mechanism by hiding parts of database from certain users.
- Permits users to access data in a customized way, so that same data can be seen by different users in different ways, at same time.
- Can simplify complex operations on base relations.

## Updating Views

- All updates to a base relation should be immediately reflected in all views that reference that base relation.
- If view is updated, underlying base relation should reflect change.

## Updating Views

- There are restrictions on types of modifications that can be made through views:
  - Updates are allowed if query involves a single base relation and contains a candidate key of base relation.
  - Updates are not allowed involving multiple base relations.
  - Updates are not allowed involving aggregation or grouping operations.

## Updating Views

- Classes of views are defined as:
  - theoretically not updateable
  - theoretically updateable
  - partially updateable.