

Echo's notebook

```

1  FLAGS=-Wall -Wextra -Wshadow -Wno-unused-result -D_GLIBCXX_DEBUG -fsanitize=address -fsanitize=undefined
   ↪ -fno-sanitize-recover
2
3  @g++ A.cpp $(FLAGS) -DJUNCO_DEBUG && ./a.out < z.in

```

```

1  // Iterate over all submasks of a mask. CONSIDER SUBMASK = 0 APART.
2  for(submask = mask; submask > 0; submask = (submask-1)&mask) {}

```

LIS

```

1  vll v_LIS(vll &v) {
2      int i, j, n = v.size();
3      vll lis, lis_time(n), ans;
4      if(!n) return ans;
5      lis.pb(v[0]); lis_time[0] = 1;
6      for(i = 1; i < n; i++) {
7          if(v[i] > lis.back()) {lis.pb(v[i]); lis_time[i] = lis.size(); continue;}
8          int pos = upper_bound(lis.begin(), lis.end(), v[i]) - lis.begin();
9          // if(pos > 0 && lis[pos-1] == v[i]) continue; // USE IF YOU WANT STRICTLY INCREASING.
10         lis[pos] = v[i];
11         lis_time[i] = pos+1;
12     }
13     j = lis.size();
14     for(i = n-1; i >= 0; i--) {
15         if(lis_time[i] == j && (ans.empty() || v[i] <= ans.back())) {ans.pb(v[i]); j--;} // <= or <.
16     }
17     reverse(ans.begin(), ans.end());
18     return ans;
19 }

```

IO

```

1  ios::sync_with_stdio(false); cin.tie(nullptr); cout.tie(nullptr);
2
3  stringstream ss;
4  ss << "Hello world";
5  ss.str("Hello world");
6  while(ss >> s) cout << s << endl;
7  ss.clear();

```

Dates

<pre> 1 // Change here and date_to_num. 2 ll is_leap_year(ll y) { 3 // if(y%4 (y%100==0 && y%400)) return 0; // ↪ Complete leap year. 4 if(y%4 != 0) return 0; // Restricted leap year. 5 return 1; 6 } 7 ll days_month[12] = {31, 28, 31, 30, 31, 30, 31, ↪ 31, 30, 31, 30, 31}; 8 ll days_month_accumulate[12] = {31, 59, 90, 120, ↪ 151, 181, 212, 243, 273, 304, 334, 365}; 9 10 // d 1-index, m 1-index. 11 ll date_to_num(ll d, ll m, ll y) { 12 ll sum = d; 13 m -= 2; 14 if(m >= 1) sum += is_leap_year(y); 15 y--; 16 if(m >= 0) sum += days_month_accumulate[m]; 17 if(y >= 0) { </pre>	<pre> 18 sum += 365*y; 19 // sum += y/4 -y/100 + y/400; // Complete ↪ leap year. 20 sum += y/4; // Restricted leap year. 21 } 22 return sum; 23 } 24 25 // Tiny optimization, binary search the year, month ↪ and day. 26 void num_to_date(ll num, ll &d, ll &m, ll &y) { 27 d = 1; m = 1; y = 0; // The date searched is >= ↪ this date. 28 while(date_to_num(d, m, y) <= num) y++; 29 y--; 30 while(date_to_num(d, m, y) <= num) m++; 31 m--; 32 while(date_to_num(d, m, y) <= num) d++; 33 d--; 34 } </pre>
--	---

Geometry

```

1  template<typename T>
2  class Point {
3      public:
4          static const int LEFT_TURN = 1;
5          static const int RIGHT_TURN = -1;
6          T x = 0, y = 0;
7          Point() = default;
8          Point(T _x, T _y) {
9              x = _x;
10             y = _y;
11         }
12         friend ostream &operator << (ostream &os,
13             ↪ Point<T> &p) {
14             os << "(" << p.x << " " << p.y << ")";
15             return os;
16         }
17         bool operator == (const Point<T> other) const {
18             return x == other.x && y == other.y;
19         }
20         // Get the (1º) bottom (2º) left point.
21         bool operator < (const Point<T> other) const {
22             if(y != other.y) return y < other.y;
23             return x < other.x;
24         }
25         T euclidean_distance(Point<T> other) {
26             T dx = x - other.x;
27             T dy = y - other.y;
28             return sqrt(dx*dx + dy*dy);
29         }
30         T euclidean_distance_squared(Point<T> other) {
31             T dx = x - other.x;
32             T dy = y - other.y;
33             return dx*dx + dy*dy;
34         }
35     };
36
37     T manhatan_distance(Point<T> other) {
38         return abs(other.x - x) + abs(other.y - y);
39     }
40
41     // Get the height of the triangle with base b1,
42     ↪ b2.
43     T height_triangle(Point<T> b1, Point<T> b2) {
44         if(b1 == b2 || *this == b1 || *this == b2)
45             ↪ return 0; // It's not a triangle.
46         T a = euclidean_distance(b1);
47         T b = b1.euclidean_distance(b2);
48         T c = euclidean_distance(b2);
49         T d = (c*c-b*b-a*a)/(2*b);
50         return sqrt(a*a - d*d);
51     }
52
53     int get_quadrant() {
54         if(x > 0 && y >= 0) return 1;
55         if(x <= 0 && y > 0) return 2;
56         if(x < 0 && y <= 0) return 3;
57         if(x >= 0 && y < 0) return 4;
58         return 0; // Point (0, 0).
59     }
60
61     // Relative quadrant respect the point other,
62     ↪ not the origin.
63     int get_relative_quadrant(Point<T> other) {
64         Point<T> p(other.x - x, other.y - y);
65         return p.get_quadrant();
66     }
67
68     // Orientation of points *this -> a -> b.
69     int get_orientation(Point<T> a, Point<T> b) {
70         T prod = (a.x - x)*(b.y - a.y) - (a.y -
71             ↪ y)*(b.x - a.x);
72         if(prod == 0) return 0;
73         return prod > 0? LEFT_TURN : RIGHT_TURN;
74     }
75
76     // True if a have less angle than b, if *this->a->b is a left turn.
77     bool angle_cmp(Point<T> a, Point<T> b) {
78         if(get_relative_quadrant(a) != get_relative_quadrant(b))
79             return get_relative_quadrant(a) < get_relative_quadrant(b);
80         int ori = get_orientation(a, b);
81         if(ori == 0) return euclidean_distance_squared(a) < euclidean_distance_squared(b);
82         return ori == LEFT_TURN;
83     }
84
85     // Anticlockwise sort starting at 1º quadrant, respect to *this point.
86     void polar_sort(vector<Point<T>> &v) {
87         sort(v.begin(), v.end(), [&](Point<T> a, Point<T> b) {return angle_cmp(a, b);});
88     }
89
90     // Convert v to its convex hull, Do a Graham Scan. O(n log n).
91     void convert_convex_hull(vector<Point<T>> &v) {
92         if(v.size() < 3) return;
93         Point<T> bottom_left = v[0], p2;
94         for(auto p : v) bottom_left = min(bottom_left, p);
95         bottom_left.polar_sort(v);
96         vector<Point<T>> v_input = v; v.clear();
97         for(auto p : v_input) {
98             while(v.size() >= 2) {
99                 p2 = v.back(); v.pop_back();
100                 if(v.back().get_orientation(p2, p) == LEFT_TURN) {
101                     v.pb(p2);
102                     break;
103                 }
104             }
105             v.pb(p);
106         }
107     }
108 };

```

Graphs

Articulation points and bridges

```

1 vector<vi> adyList; // Graph
2 vi num, low; // num and low for DFS
3 int cnt; // Counter for DFS
4 int root, rchild; // Root and number of (DFS)
  ↳ children
5 vi artic; // Contains the articulation
  ↳ points at the end
6 set<pii> bridges; // Contains the bridges at the
  ↳ end
7
8 void dfs(int nparent, int nnode) {
9     num[nnode] = low[nnode] = cnt++;
10    rchild += (nparent == root);
11
12    for (auto a : adyList[nnode]) {
13        if (num[a] == -1) { // Tree edge
14            dfs(nnode, a);
15            low[nnode] = min(low[nnode], low[a]);
16            if (low[a] >= num[nnode]) {
17                artic[nnode] = true;
18            }
19            if (low[a] > num[nnode]) {
20                bridges.insert((nnode < a) ?
  ↳ mp(nnode, a) : mp(a, nnode));
21
22        } else if (a != nparent) { // Back edge
23            low[nnode] = min(low[nnode], num[a]);
24        }
25    }
26 }
27 void findArticulations(int n) {
28     cnt = 0;
29     low = num = vi(n, -1);
30     artic = vi(n, 0);
31     bridges.clear();
32
33     for (int i = 0; i < n; ++i) {
34         if (num[i] != -1) {
35             continue;
36         }
37         root = i;
38         rchild = 0;
39         dfs(-1, i);
40         artic[root] = rchild > 1; //Special case
41     }
42 }

```

Max Flow: Edmond Karp's $\mathcal{O}(VE^2)$

```

1 vector<vector<ll>> adjList;
2 vector<vector<ll>> adjMat;
3
4 void initialize(int n) {
5     adjList = decltype(adjList)(n);
6     adjMat = decltype(adjMat)(n, vector<ll>(n, 0));
7 }
8
9 map<int, int> p;
10 bool bfs(int source, int sink) {
11     queue<int> q;
12     vi visited(adjList.size(), 0);
13     q.push(source);
14     visited[source] = 1;
15     while (!q.empty()) {
16         int u = q.front();
17         q.pop();
18         if (u == sink)
19             return true;
20         for (auto v : adjList[u]) {
21             if (adjMat[u][v] > 0 && !visited[v]) {
22                 visited[v] = true;
23                 q.push(v);
24                 p[v] = u;
25             }
26         }
27     }
28     return false;
29 }
30 int max_flow(int source, int sink) {
31     ll max_flow = 0;
32     while (bfs(source, sink)) {
33         ll flow = inf;
34         for (int v = sink; v != source; v = p[v]) {
35             flow = min(flow, adjMat[p[v]][v]);
36         }
37         for (int v = sink; v != source; v = p[v]) {
38             adjMat[p[v]][v] -= flow; // Decrease
  ↳ capacity forward edge
39             adjMat[v][p[v]] += flow; // Increase
  ↳ capacity backward edge
40         }
41         max_flow += flow;
42     }
43     return max_flow;
44 }
45 void addEdgeUni(int orig, int dest, ll flow) {
46     adjList[orig].pb(dest);
47     adjMat[orig][dest] = flow;
48     adjList[dest].pb(orig); //Add edge for residual
  ↳ flow
49 }
50 void addEdgeBi(int orig, int dest, ll flow) {
51     adjList[orig].pb(dest);
52     adjList[dest].pb(orig);
53     adjMat[orig][dest] = flow;
54     adjMat[dest][orig] = flow;
55 }

```

Bellman Ford's

```

1 for(i = 0; i < n - 1; i++) { // Iterate n - 1 times.
2     for(auto e : edge) {
3         if(dist[e.fi.fi] != inf)
4             dist[e.fi.se] = min(dist[e.fi.se], dist[e.fi.fi] + e.se);
5     }
6 }

```

Floyd cycle detection

```

1 ll f(ll x) {return (x + 1) % 4;} // Example.
2 // mu is the first index of the node in the cycle. lambda is the length of the cycle.
3 pll floyd_cycle_detection(ll x0) {
4     ll tortoise = f(x0), hare = f(f(x0)), mu = 0, lambda = 1;
5     while(tortoise != hare) tortoise = f(tortoise), hare = f(f(hare));
6     tortoise = x0;
7     while(tortoise != hare) tortoise = f(tortoise), hare = f(hare), mu++;
8     hare = f(hare);
9     while(tortoise != hare) hare = f(hare), lambda++;
10    return mp(mu, lambda);
11 }

```

Max Flow: Dinic's $\mathcal{O}(V^2E)$

```

1 //  $\mathcal{O}(V^2E)$  max flow algorithm. For bipartite
  ↳ matching  $\mathcal{O}(\sqrt{V} * E)$ , always faster than
  ↳ Edmond-Karp.
2 // Creates layer's graph with a BFS and then it
  ↳ tries all possibles DFS, branching while the
  ↳ path doesn't reach the sink
3 struct EdgeFlow {
4     ll u, v;
5     ll cap, flow = 0; //capacity and current flow
6     EdgeFlow(ll _u, ll _v, ll _cap) : u(_u), v(_v),
  ↳ cap(_cap) {}
7 };
8
9 struct Dinic {
10     vector<EdgeFlow> edge; //keep the edges
11     vector<vll> graph; //graph[u] is the list of
  ↳ their edges
12     ll n, n_edges = 0;
13     ll source, sink, inf_flow = inf;
14     vll lvl; //lvl of the node to the source
15     vll ptr; //ptr[u] is the next edge you have to
  ↳ take in order to branch the DFS
16     queue<ll> q;
17
18     Dinic(ll _n, ll _source, ll _sink) : n(_n),
  ↳ source(_source), sink(_sink) { //n nodes
19         graph.assign(_n, vll());
20     }
21
22     void add_edge(ll u, ll v, ll flow) { //u->v
  ↳ with cost x
23         EdgeFlow uv(u, v, flow), vu(v, u, 0);
24         edge.pb(uv);
25         edge.pb(vu);
26         graph[u].pb(n_edges);
27         graph[v].pb(n_edges+1);
28         n_edges += 2;
29     }
30
31     bool BFS() {
32         ll u;
33         while(q.empty() == false) {
34             u = q.front(); q.pop();
35             for(auto e1 : graph[u]) {
36                 if(lvl[edge[e1].v] != -1) {
37                     continue;
38                 }
39                 if(edge[e1].cap - edge[e1].flow <=
  ↳ 0) {
40                     continue;
41                 }
42                 lvl[edge[e1].v] = lvl[edge[e1].u] +
  ↳ 1;
43                 q.push(edge[e1].v);
44             }
45         }
46         return lvl[sink] != -1;
47     }
48
49     ll dfs(ll u, ll min_flow) {
50         if(u == sink) return min_flow;
51         ll pushed, el;
52         for(; ptr[u] < (int)graph[u].size();
  ↳ ptr[u]++) { //if you can pick ok, else
  ↳ you crop that edge for the current bfs
  ↳ layer
53             el = graph[u][ptr[u]];
54             if(lvl[edge[el].v] != lvl[edge[el].u] +
  ↳ 1 || edge[el].cap - edge[el].flow
  ↳ <= 0) {
55                 continue;
56             }
57             pushed = dfs(edge[el].v, min(min_flow,
  ↳ edge[el].cap - edge[el].flow));
58             if(pushed > 0) {
59                 edge[el].flow += pushed;
60                 edge[el^1].flow -= pushed;
61                 return pushed;
62             }
63         }
64         return 0;
65     }
66
67     ll max_flow() {
68         ll flow = 0, pushed;
69         while(true) {
70             lvl.assign(n, -1);
71             lvl[source] = 0;
72             q.push(source);
73             if(!BFS()) {
74                 break;
75             }
76             ptr.assign(n, 0);
77             while(true) {
78                 pushed = dfs(source, inf_flow);
79                 if(!pushed) break;
80                 flow += pushed;
81             }
82         }
83         return flow;
84     }
85 };

```

Hungarian Algorithm

```

1 // The rows are jobs, the columns are workers
2 pair<ll, vl> hungarian(vector<vl> &matrix) {
3     int n = matrix.size(), m = matrix[0].size();
4     vl jobP(n), workerP(m + 1), matched(m + 1, -1);
5
6     vl dist(m + 1, inf);
7     vi from(m + 1, -1), seen(m + 1, 0);
8
9     for (int i = 0; i < n; ++i) {
10         int cWorker = m;
11         matched[cWorker] = i;
12         std::fill(all(dist), inf);
13         std::fill(all(from), -1);
14         std::fill(all(seen), false);
15
16         while (matched[cWorker] != -1) {
17             seen[cWorker] = true;
18             int i0 = matched[cWorker];
19             int nextWorker = -1;
20             ll delta = inf;
21
22             for (int worker = 0; worker < m;
23                  ↪ ++worker) {
24                 if (seen[worker])
25                     continue;
26                 ll candidateDistance =
27                     ↪ matrix[i0][worker];
28                 candidateDistance += -jobP[i0] -
29                     ↪ workerP[worker];
30
31                 if (candidateDistance < dist[worker]) {
32                     dist[worker] =
33                         ↪ candidateDistance;
34                     from[worker] = cWorker;
35                 }
36                 if (dist[worker] < delta) {
37                     delta = dist[worker];
38                     nextWorker = worker;
39                 }
40             }
41             for (int j = 0; j <= m; ++j) {
42                 if (seen[j]) {
43                     jobP[matched[j]] += delta;
44                     workerP[j] -= delta;
45                 } else {
46                     dist[j] -= delta;
47                 }
48             }
49             cWorker = nextWorker;
50         }
51         while (cWorker != m) {
52             int prevWorker = from[cWorker];
53             matched[cWorker] = matched[prevWorker];
54             cWorker = prevWorker;
55         }
56     }
57     ll ans = -workerP[m];
58     vl rowMatchesWith(n);
59     for (int j = 0; j < m; ++j) {
60         if (matched[j] != -1) {
61             rowMatchesWith[matched[j]] = j;
62         }
63     }
64     return {ans, std::move(rowMatchesWith)};
65 }

```

Floyd - Warshall: k->i->j
Kosaraju

```

1 vector<vi> adyList; // Graph
2 vector<int> visited; // Visited for DFS
3 vector<vi> sccs; // Contains the SCCs at the
4 ↪ end
5 void dfs(int nnode, vector<int> &v, vector<vi>
6 ↪ &adyList) {
7     if (visited[nnode]) {
8         return;
9     }
10    visited[nnode] = true;
11    for (auto a : adyList[nnode]) {
12        dfs(a, v, adyList);
13    }
14    v.push_back(nnode);
15 }
16 void Kosaraju(int n) {
17     visited = vi(n, 0);
18     stack<int> s = stack<int>();
19     sccs = vector<vi>();
20
21     vector<int> postorder;
22
23     for (int i = 0; i < n; ++i) {
24         dfs(i, postorder, adyList);
25     }
26     reverse(all(postorder));
27
28     vector<vi> rAdyList = vector<vi>(n, vi());
29     for (int i = 0; i < n; ++i) {
30         for (auto v : adyList[i]) {
31             rAdyList[v].push_back(i);
32         }
33     }
34
35     visited = vi(n, 0);
36     vi data;
37     for (auto a : postorder) {
38         if (!visited[a]) {
39             data = vi();
40             dfs(a, data, rAdyList);
41             if (!data.empty())
42                 sccs.pb(data);
43         }
44     }
45 }

```

LCA tree

```

1  const int MAX_N = 1e5 + 5;
2  const int MAX_LOG_N = 18;
3  int n;
4  vector<vi> graph; // Directed graph, allways
   ↳ reserve memory for it.
5  vector<vi> bigraph; // Undirected graph, reserve
   ↳ memory only if needed.
6
7  int level[MAX_N]; // level of the node rooted.
8  int parent[MAX_N][MAX_LOG_N]; // parent[i][j] is
   ↳ the parent 2^j of the node i.
9
10 vector<bool> visited_bigraph;
11 // root_graph(u, -1) roots the bigraph at node u.
12 void root_graph(int u, int p) {
13     if(p == -1) visited_bigraph.assign(n, false);
14     for(auto v : bigraph[u]) {
15         if(v == p) continue;
16         graph[u].pb(v);
17         root_graph(v, u);
18     }
19 }
20
21 // Calculate the level and parent 1. Don't call.
22 void dfs_level(int u, int p) {
23     parent[u][0] = p;
24     level[u] = level[p] + 1;
25     for(auto v : graph[u]) {
26         if(v == p) continue;
27         dfs_level(v, u);
28     }
29 }
30 // Builds the LCA.
31
32 void build_lca(int root) {
33     int i, j;
34     level[root] = -1;
35     dfs_level(root, root); // The parent of the
   ↳ root is itself.
36     for(j = 1; j < MAX_LOG_N; j++) {
37         for(i = 0; i < MAX_N; i++) {
38             parent[i][j] = parent[parent[i][j -
   ↳ 1]][j - 1];
39         }
40     }
41     // Calculates the LCA(u, v) in O(log n).
42     int lca(int u, int v) {
43         if(level[u] > level[v]) swap(u, v);
44         int i, d = level[v] - level[u];
45         for(i = MAX_LOG_N - 1; i >= 0; i--) {
46             if(is_set(d, i)) v = parent[v][i];
47         }
48         if(u == v) return u;
49         for(i = MAX_LOG_N - 1; i >= 0; i--) {
50             if(parent[u][i] != parent[v][i])
51                 u = parent[u][i], v = parent[v][i];
52         }
53         return parent[u][0];
54     }
55     // Calculates the distance(u, v) in a tree in O(log
   ↳ n).
56     int dist(int u, int v) {
57         return level[u] + level[v] - 2 * level[lca(u,
   ↳ v)];
58     }

```

Mathematics

Binary operations

```

1  ll elevate(ll a, ll b) { // b >= 0.
2      ll ans = 1;
3      while(b) {
4          if(b & 1) ans = ans * a % mod;
5          b >>= 1;
6          a = a * a % mod;
7      }
8      return ans;
9  }
10 // a^(mod - 1) = 1, Euler.
11 ll inv(ll a) {
12     return elevate(((a%mod) + mod)%mod, mod - 2);
13 }
14
15 ll mul(ll a, ll b) {
16     ll ans = 0, neg = (a < 0) ^ (b < 0);
17     a = abs(a); b = abs(b);
18     while(b) {
19         if(b & 1) ans = (ans + a) % mod;
20         b >>= 1;
21         a = (a + a) % mod;
22     }
23     if(neg) return -ans;
24     return ans;
25 }

```

$$\text{Catalan numbers: } C_n = \frac{1}{n+1} \binom{2n}{n}$$

Combinatoric numbers

```

1  const int MAX_C = 1+66; // 66 is the for long
   ↳ long, C(66, x)
2  ll Comb[MAX_C][MAX_C];
3
4  void calc() {
5      int i, j;
6      for(i = 0; i < MAX_C; i++) {
7          Comb[i][0] = 1;
8          Comb[0][i] = 1;
9      }
10     for(i = 1; i < MAX_C; i++) {
11         for(j = 1; j < MAX_C; j++) {
12             if(i+j >= MAX_C) continue;
13             Comb[i][j] = Comb[i-1][j] +
   ↳ Comb[i][j-1];
14         }
15     }
16 }
17 ll C(ll i, ll j) {
18     return Comb[i-j][j];
19 }

```

Chinese Remainder

```

1  const ll MAX = 10;
2  ll a[MAX], p[MAX], n;
3  // Given  $n$   $x \equiv a[i] \pmod{p[i]}$ , find  $x$ ,
4  // or -1 if it doesn't exist.
5  // Let  $q[i] = (\prod_{i=0}^{n-1} p[j]) / p[i]$ .
6  //  $x$  will be  $\sum_{i=0}^{n-1} a[i] * q[i]$ 
7  //  $*inv(q[i], \text{mod } p[i])$ 
8  ll chinese_remainder() {
9      ll i, j, g, ans = 0, inv1, inv2;
10     mod = 1;
11     for(i = 0; i < n; i++) {
12         // If the  $p[i]$  are not coprimes, do them
13         //  $\hookrightarrow$  coprimes.
14         a[i] %= p[i]; a[i] += p[i]; a[i] %= p[i];
15         for(j = 0; j < i; j++) {
16             g = __gcd(p[i], p[j]);
17             if((a[i]%g + g)%g != (a[j]%g + g)%g)
18                 return -1;
19         }
20         // Delete the repeated factor at the
21         //  $\hookrightarrow$  correct side.
22         if (__gcd(p[i]/g, p[j]) == 1) {p[i] /=
23             g; a[i] %= p[i];}
24         else {p[j] /= g; a[j] %= p[j];}
25     }
26     // If you have a supermod, take  $P = \min(P,$ 
27     //  $\hookrightarrow$  supermod);
28     for(i = 0; i < n; i++) {
29         mod *= p[i];
30     }
31     for(i = 0; i < n; i++) {
32         gcdEx(mod/p[i], p[i], &inv1, &inv2);
33         ans += mul(a[i], mul(mod/p[i], inv1));
34         ans %= mod;
35     }
36     return (ans%mod + mod) % mod;
37 }

```

Euclides

```

1  ll gcdEx(ll a, ll b, ll *x1, ll *y1) {
2      if(a == 0) {
3          *x1 = 0;
4          *y1 = 1;
5          return b;
6      }
7      ll x0, y0, g;
8      g = gcdEx(b%a, a, &x0, &y0);
9      *x1 = y0 - (b/a)*x0;
10     *y1 = x0;
11     return g;
12 }

```

Hash Set

```

1  const int MAX = 2*1e5+5;
2  ll val[MAX]; // For random numbers and not index
3  //  $\hookrightarrow$  use  $f$  with random xor.
4  void ini() { // CALL ME ONCE.
5      srand(time(0));
6      for(int i = 0; i < MAX; i++) val[i] = rand();
7  }
8  // Hash_set contains a set of indices  $[0..MAX-1]$ 
9  //  $\hookrightarrow$  with duplicates.
10 //  $a[i] = \text{sum\_x}\{val\_x\} \% \text{mod } p[i]$ .
11 class Hash_set {
12 public:
13     vll p = {1237273, 1806803, 3279209}; // Prime
14     //  $\hookrightarrow$  numbers.
15     vll a = {0, 0, 0};
16     int n = 3; //  $n = p.size()$ ;
17 }
18 void insert(int x) { // Insert index  $x$ .
19     for(int i = 0; i < n; i++) a[i] = (a[i] +
20         //  $\hookrightarrow$  val[x]) % p[i];
21     }
22 // Insert all the elements of  $hs$ .
23 void insert (Hash_set hs) {
24     for(int i = 0; i < n; i++) a[i] = (a[i] +
25         //  $\hookrightarrow$  hs.a[i]) % p[i];
26     }
27 bool operator == (Hash_set hs) {
28     for(int i = 0; i < n; i++) if(a[i] !=
29         //  $\hookrightarrow$  hs.a[i]) return false;
30     return true;
31 }
32 };

```

Hash of pairs

```

1  // Use unordered_set<pii, pair_hash> us or
2  //  $\hookrightarrow$  unordered_map<pii, int, pair_hash> um;
3  struct pair_hash
4  {
5      template <class T1, class T2>
6      size_t operator () (pair<T1, T2> const &pair)
7          //  $\hookrightarrow$  const
8      {
9          size_t h1 = hash<T1>()(pair.first);
10         size_t h2 = hash<T2>()(pair.second);
11         return (h1 ^ 0b11001001011001101) +
12             //  $\hookrightarrow$  (0b011001010011100111 ^ h2);
13     }
14 };

```

Linear Sieve

```

1  const int MAX_PRIME = 1e6+5;
2  bool num[MAX_PRIME]; // If num[i] = false  $\Rightarrow$   $i$  is prime.
3  int num_div[MAX_PRIME]; // Number of divisors of  $i$ .
4  int min_div[MAX_PRIME]; // The smallest prime that divide  $i$ .
5  vector<int> prime;
6

```



```

7 void linear_sieve(){
8     int i, j, prime_size = 0;
9     min_div[1] = 1;
10    for(i = 2; i < MAX_PRIME; ++i){
11        if(num[i] == false) {prime.push_back(i); ++prime_size; num_div[i] = 1; min_div[i] = i;}
12
13        for(j = 0; j < prime_size && i * prime[j] < MAX_PRIME; ++j){
14            num[i * prime[j]] = true;
15            num_div[i * prime[j]] = num_div[i] + 1;
16            min_div[i * prime[j]] = min(min_div[i], prime[j]);
17            if(i % prime[j] == 0) break;
18        }
19    }
20 }
21 bool is_prime(ll n) {
22     for(auto el : prime) {
23         if(n == el) return true;
24         if(n%el == 0) return false;
25     }
26     return true;
27 }
28 vll fact, nfact; // The factors of n and their
    ↪ exponent.
29 void factorize(int n) { // Up to
    ↪ MAX_PRIME*MAX_PRIME.
30     ll cont, prev_p;
31     fact.clear(); nfact.clear();
32     for(auto p : prime) {
33         if(n < MAX_PRIME) break;
34         if(n%p == 0) {
35             fact.pb(p);
36             cont = 0;
37             while(n%p == 0) n /= p, cont++;
38             nfact.pb(cont);
39         }
40     }
41     if(n >= MAX_PRIME) {
42         fact.pb(n);
43         nfact.pb(1);
44         return;
45     }
46     while(n != 1) { // When n < MAX_PRIME,
    ↪ factorization in almost O(1).
47         prev_p = min_div[n];
48         cont = 0;
49         while(n%prev_p == 0) n /= prev_p, cont++;
50         fact.pb(prev_p);
51         nfact.pb(cont);
52     }
53 }

```

Suffix Array

```

1 class SuffixArray {
2     public:
3     int n;
4     string s;
5     vi p; // p[i] is the position in the order
    ↪ array of the ith suffix (s[i..n-1]).
6     vi c; // c[i] is the equivalence class of the
    ↪ ith suffix. When build, c[p[i]] = i,
    ↪ inverse.
7     // dont use lcp[0] = 0.
8     vi lcp; // lcp[i] is the longest common prefix
    ↪ in s[p[i-1]..n-1] and s[p[i]..n-1].
9     // To get lcp(s[i..n-1], s[j..n-1]) is
    ↪ min(lcp[c[i]+1], lcp[c[j]]) (use SegTree).
10    void radix_sort(vector<pair<pii, int>> &v) { //
    ↪ O(n).
11        vector<pair<pii, int>> v2(n);
12        vi freq(n, 0); // first frequency and then
    ↪ the index of the next item.
13        int i, sum = 0, temp;
14        for(i = 0; i < n; i++) freq[v[i].fi.se]++;
    ↪ // Sort by second component.
15        for(i = 0; i < n; i++) {temp = freq[i];
    ↪ freq[i] = sum; sum += temp;}
16        for(i = 0; i < n; i++)
    ↪ {v2[freq[v[i].fi.se]] = v[i];
    ↪ freq[v[i].fi.se]++;}
17        freq.assign(n, 0); sum = 0;
18        for(i = 0; i < n; i++) freq[v2[i].fi.fi]++;
    ↪ // Sort by first component.
19        for(i = 0; i < n; i++) {temp = freq[i];
    ↪ freq[i] = sum; sum += temp;}
20
21        for(i = 0; i < n; i++)
    ↪ {v[freq[v2[i].fi.fi]] = v2[i];
    ↪ freq[v2[i].fi.fi]++;}
22    }
23    SuffixArray() = default;
24    SuffixArray(string &s) {
25        s = _s;
26        s += "$"; // smaller char to end the
    ↪ string.
27        n = s.size();
28        int i, k;
29        p.assign(n, 0);
30        c.assign(n, 0);
31        vector<pii> v1(n); // temporal vector to
    ↪ sort.
32        for(i = 0; i < n; i++) v1[i] = mp(s[i], i);
33        sort(v1.begin(), v1.end());
34        for(i = 0; i < n; i++) p[i] = v1[i].se;
35        c[p[0]] = 0;
36        for(i = 1; i < n; i++) {
37            if(v1[i].fi == v1[i-1].fi) c[p[i]] =
    ↪ c[p[i-1]];
38            else c[p[i]] = c[p[i-1]] + 1;
39        }
40        k = 0; // in k+1 iterations sort strings of
    ↪ length 2^(k+1).
41        while(c[p[n-1]] != n-1) { // At most
    ↪ ceil(log2(n)).
42            vector<pair<pii, int>> v2(n); //
    ↪ temporal vector to sort.
43            for(i = 0; i < n; i++) v2[i] =
    ↪ mp(mp(c[i], c[(i + (1 << k)) % n]),
    ↪ i);

```



```

43         radix_sort(v2);
44         for(i = 0; i < n; i++) p[i] = v2[i].se;
45         c[p[0]] = 0;
46         for(i = 1; i < n; i++) {
47             if(v2[i].fi == v2[i - 1].fi)
48                 c[p[i]] = c[p[i - 1]];
49             else c[p[i]] = c[p[i - 1]] + 1;
50         }
51         k++;
52     }
53     void show_suffixes() { // IMPORTANT use this to
54         debug.
55         for(int i = 0; i < n; i++) cout << i << " "
56         << p[i] << " " << s.substr(p[i]) <<
57         endl;
58         if(!lcp.empty()) cout << "LCP: " << lcp <<
59         endl;
60     }
61     // cmp s with t. return -1 if s < t, 1 if s >
62     // t, 0 if s == t.
63     int cmp_string(int pos, string &t) {
64         for(int i = p[pos], j = 0; j < (int)
65         t.size(); i++, j++) {
66             if(s[i] < t[j]) return -1; // i < n
67             // because s[n-1] = '$'.
68             if(s[i] > t[j]) return 1;
69         }
70         return 0;
71     }
72     // Count the number of times t appears in s.
73     int count_substring(string &t) {
74         int l = -1, r = n, mid, L, R;
75         while(l + 1 < r) { //
76             // -1,...,-1=L,0,...,0,1=R...1.
77             mid = (l + r) / 2;
78             if(cmp_string(mid, t) < 0) l = mid;
79             else r = mid;
80         }
81         L = l;
82         l = -1; r = n;
83         while(l + 1 < r) {
84             mid = (l + r) / 2;
85             if(cmp_string(mid, t) <= 0) l = mid;
86             else r = mid;
87         }
88         R = r;
89         return R - L - 1;
90     }
91     // 0(n) build. At most 2n lcp++ and n lcp--;
92     void build_lcp() {
93         lcp.assign(n, 0);
94         for(int i = 0; i < n - 1; i++) {
95             if(i > 0) lcp[c[i]] = max(lcp[c[i - 1]]
96             <- -1, 0);
97             while(s[i + lcp[c[i]]] == s[p[c[i] - 1]
98             <- + lcp[c[i]]) lcp[c[i]]++;
99         }
100     }
101     ll number_substrings() {
102         ll ans = 0, i;
103         for(i = 1; i < n; i++) {
104             ans += n - p[i-1] - lcp[i]; // Length
105             // of the suffix - lcp with the next
106             // suffix.
107         }
108         ans += n - p[n - 1]; // Plus the last
109         // suffix.
110         return ans - n; // Remove the '$' symbol on
111         // n substrings.
112     }
113 };
114 string LCS(string s, string &t) {
115     int mx = 0, mxi = 0, i, n2 = t.length();
116     string ans = "";
117     s += "@" + t; // Concatenate with a special
118     // char.
119     SuffixArray sa(s);
120     sa.build_lcp();
121     for(i = 1; i < sa.n; i++) {
122         // Suffix of s and before suffix of t.
123         if(sa.n - sa.p[i] > n2 + 2 && sa.n -
124         sa.p[i-1] <= n2 + 1) {
125             if(sa.lcp[i] > mx) mx = sa.lcp[i], mxi
126             = i;
127         }
128         // Suffix of t and before suffix of s.
129         if(sa.n - sa.p[i] <= n2 + 1 && sa.n -
130         sa.p[i-1] > n2 + 2) {
131             if(sa.lcp[i] > mx) mx = sa.lcp[i], mxi
132             = i;
133         }
134     }
135     return sa.s.substr(sa.p[mxi], mx);
136 }

```

BIT Fenwick tree

```

1  template<typename T>
2  class BIT{
3      vector<T> bit;
4      int n;
5  public:
6      BIT(int _n) {
7          n = _n;
8          bit.assign(n+1, 0);
9      }
10     BIT(vector<T> v) {
11         n = v.size();
12         bit.assign(n+1, 0);
13         for(int i = 0; i < n; i++) update(i, v[i]);
14     }
15     // Point update.
16     void update(int i, T dx) {
17         for(i++; i < n+1; i += LSB(i)) bit[i] +=
18         dx;
19     }
20     T query(int r) { // query [0, r].
21         T ans = 0;
22         for(r++; r > 0; r -= LSB(r)) ans += bit[r];
23         return ans;
24     }
25     T query(int l, int r) { // query [l, r].
26         return query(r) - query(l-1);
27     }
28     // k-th smallest element inserted.
29     int k_element(ll k) { // k > 0 (1-indexed).
30         int l = 0, r = n+1, mid;
31         if(query(0) >= k) return 0;
32         while(l + 1 < r) {
33             mid = (l + r)/2;
34             if(query(mid) >= k) r = mid;
35             else l = mid;
36         }
37         return r;
38     }
39 };

```

Strings: KMP

```

1  template <typename T>
2  vi prefixFun(const T &s, int n) {
3      vi res(n);
4      for (int i = 1; i < n; ++i) {
5          int j = res[i - 1];
6          while (j > 0 && s[i] != s[j]) {
7              j = res[j - 1];
8          }
9          res[i] = j + (s[i] == s[j]);
10     }
11     return res;
12 }
13
14 template <typename T>
15 int kmpSearch(const T &text, int n,
16               const T &pattern, int m,
17               const vi &patternPre) {
18
19     int count = 0;
20     int j = 0;
21     for (int i = 0; i < n; ++i) {
22         while (j > 0 && text[i] != pattern[j]) {
23             j = max(0, patternPre[j] - 1);
24         }
25         j += (text[i] == pattern[j]);
26         if (j == m) {
27             count++;
28             j = patternPre[j - 1];
29         }
30     }
31     return count;
32 }

```

Longest Palindromic Substring

```

1  // LPS Longest Palindromic Substring, O(n).
2  void Manacher(string &str) {
3      char ch = '#'; // '#' a char not contained in str.
4      string s(1, ch), ans;
5      for(auto c : str) {s += c; s += ch;}
6      int i, n = s.length(), c = 0, r = 0;
7      vi lps(n, 0);
8      for(i = 1; i < n; i++) {
9          // lps[i] >= it's mirror, but falling in the interval [L..R]. L = c - (R - c).
10         if(i < r) lps[i] = min(r - i, lps[c - (i - c)]);
11         // Try to increase.
12         while(i-lps[i]-1 >= 0 && i+lps[i]+1 < n && s[i-lps[i]-1] == s[i+lps[i]+1]) lps[i]++;
13         // Update the interval [L..R].
14         if(i + lps[i] > r) c = i, r = i + lps[i];
15     }
16     // Get the longest palindrome in ans.
17     int pos = max_element(lps.begin(), lps.end()) - lps.begin();
18     for(i = pos - lps[pos]; i <= pos + lps[pos]; i++) {
19         if(s[i] != ch) ans += s[i];
20     }
21     //cout << ans.size() << "\n";
22 }

```

Z-algorithm

```

1  // Search the occurrences of t (pattern to search)
2  // in s (the text).
3  // O(n + m). It increases R at most 2n times
4  // and decreases at most n times.
5  // z[i] is the longest string s[i..i+z[i]-1]
6  // that is a prefix = s[0..z[i]-1].
7  void z_algorithm(string &s, string &t) {
8      s = t + "$" + s;
9      // "$" is a char not present in s nor t.
10     int n = s.length(), m = t.length(), i;
11     int L = 0, R = 0;
12     vi z(n, 0);
13     // s[L..R] = s[0..R-L], [L, R]
14     // is the current window.
15     for(i = 1; i < n; i++) {
16         if(i > R) { // Old window, recalculate.
17             L = R = i;
18             while(R < n && s[R] == s[R-L]) R++;
19             R--;
20             z[i] = R - L + 1;
21         } else {
22             // z[i] will fall in the window.
23             if(z[i-L] < R - i) z[i] = z[i-L];
24             // z[i] can fall outside the window,
25             // try to increase the window.
26             else {
27                 L = i;
28                 while(R < n && s[R] == s[R-L]) R++;
29                 R--;
30                 z[i] = R - L + 1;
31             }
32         }
33         if(z[i] == m) { // Match found.
34             //echo("Pattern found at: ", i-m-1);
35         }
36     }
37 }

```

Hours

```

1 // One day has 60*60*24 = 86400 seconds.
2 // Converts the hour to number of seconds since
  ↳ 00:00:00.
3 ll hours_to_seconds(ll h, ll m, ll s) {
4     return 60*60*h + 60*m + s;
5 }
6
7 // From sec seconds, get the hour. Just's for one
  ↳ day.
8 void seconds_to_hours(ll &h, ll &m, ll &s, ll sec)
  ↳ {
9     sec %= 86400; sec += 86400; sec %= 86400;
10    h = sec / (60*60);
11    sec %= 60*60;
12    m = sec / 60;
13    sec %= 60;
14    s = sec;
15 }
16
17 // Convert grades of the clock hand to hours and
  ↳ minutes. gh is grades of hours and gm grades of
  ↳ minutes.
18 // return mp(-1, -1) if no solution exists.
19 pair<ll, ll> grades_to_hour(ld gh, ld gm) {
20     ll h = gh/30, m = gm/6;
21     if((ld)30*h + (ld)m/2 != gh || (ld)6*m != gm)
22         ↳ return mp(-1, -1);
23     return mp(h, m);
24 }
25
26 // Convert hours and minutes to grades of the clock
  ↳ hand, mp(grade of large hour hand, small minute
  ↳ hand).
27 pair<ld, ld> hour_to_grades(ll h, ll m) {
28     return mp((ld)30*h + (ld)m/2, (ld)6*m);
29 }
30
31 // Convert hours and minutes to grades of the clock
  ↳ hand, mp(grade of large hour hand, small minute
  ↳ hand).
32 // Not tested.
33 pair<ld, pair<ld, ld>> hour_to_grades(ll h, ll m,
  ↳ ll s) {
34     return mp((ld)30*h + (ld)m/2 + (ld)s/120,
35         ↳ mp((ld)6*m + (ld)s/10, (ld)6*s));
36 }

```

Persistant Segment Tree

```

1 const int MAX_VERSION = 3*1e4+4; // Maximum number
  ↳ of versions.
2 template<typename T>
3 struct node {
4     node *pl = NULL, *pr = NULL;
5     int l, r, mid;
6     T value = 0; // Sum query.
7     node(int _l, int _r) {l = _l; r = _r; mid =
  ↳ (l+r)>>1;}
8     node(int _l, int _r, T _value) {l = _l; r = _r;
  ↳ value = _value; mid = (l+r)>>1;}
9     void update() { // Sum query.
10        value = 0;
11        if(pl) value += pl->value;
12        if(pr) value += pr->value;
13    }
14 }; // Declare outside, else static memory gives seg
  ↳ fault.
15 node<ll> *root[MAX_VERSION]; //it stores the i
  ↳ versions after updates, start at 0.
16
17 template<typename T>
18 class PersistentSegmentTree {
19     vector<T> arr; // Copy of the array to build
  ↳ SegmentTree.
20
21     void build(node<T> *n) { // O(n).
22         if(n->l == n->r) {n->value = arr[n->l];
23             ↳ return;}
24         n->pl = new node<T>(n->l, n->mid);
25         n->pr = new node<T>(n->mid+1, n->r);
26         build(n->pl);
27         build(n->pr);
28         n->update();
29     }
30
31     node<T>* update(node<T> *n, int q, ll x) { //
  ↳ O(logn).
32         if(n->l == n->r) {
33             return new node<T>(n->l, n->r, x);
34         }
35         node<T> *nod = new node<T>(n->l, n->r);
36         if(q <= n->mid) {
37             nod->pl = update(n->pl, q, x);
38             nod->pr = n->pr;
39         } else {
40             nod->pl = n->pl;
41             nod->pr = update(n->pr, q, x);
42         }
43         nod->update();
44         return nod;
45     }
46
47     T query(node<T> *n, int l, int r) { // O(logn).
48         if(l <= n->l && n->r <= r) return n->value;
49         if(r < n->l || n->r < l) return 0; // Sum
50         ↳ query.
51         return query(n->pl, l, r) + query(n->pr, l,
52             ↳ r);
53     }
54
55     public:
56     PersistentSegmentTree() = default;
57     PersistentSegmentTree(int n) { // Build from
58         ↳ empty vector of size n.
59         arr.assign(n, 0);
60         root[0] = new node<T>(0, n-1);
61         build(root[0]);
62     }
63
64     PersistentSegmentTree(vector<T> &v) { // Build
65         ↳ from vector v.
66         arr = v;
67         root[0] = new node<T>(0, (int)arr.size() -
68             ↳ 1);
69         build(root[0]);
70     }
71
72     T query(int version, int l, int r) {return
73         ↳ query(root[version], l, r);} // O(logn).
74     // Set v[idx] = x. Set update.
75     void update(int version, int new_version, int
76         ↳ idx, T x) { // update the segTree version
77         ↳ into new_version root.

```

```

64         root[new_version] = update(root[version],
65                                     ↪ idx, x); // O(logn).
66     }
67 };
68 template<typename T>
69 class NumberDistinctNumbers { // Works for queries
70     ↪ online. For offline can check MO's.
71     PersistentSegmentTree<T> pst;
72     static const int MAX_ELEMENT_VALUE = 1e6+4; //
73     ↪ for querying last[ei].
74 public:
75     NumberDistinctNumbers(vector<T> &v) { //
76         ↪ O(nlogn).
77         T last[MAX_ELEMENT_VALUE]; // last
78         ↪ occurrence of the i-number in the array,
79         ↪ updating from left to right.
80     };
81     int n = v.size();
82     fill(last, last+MAX_ELEMENT_VALUE, -1);
83     pst = PersistentSegmentTree<T>(n);
84     for(int r = 0; r < n; r++) {
85         pst.update(max(0, r-1), r, r, 1); //
86         ↪ Actualize r.
87         if(last[v[r]] != -1) pst.update(r, r,
88         ↪ last[v[r]], 0); // Remove
89         ↪ last[v[r]].
90         last[v[r]] = r; // Actualize
91         ↪ last[v[r]].
92     }
93     T query(int l, int r) { // Return the number of
94         ↪ Distinct numbers in [l..r], O(logn).
95         return pst.query(r, l, r);
96     }
97 };

```

Euler Circuit

```

1 // Euler circuit, visit all edges once.
2 // Condition: for every node in-degree =
3     ↪ out-degree. All edges are in the same SCC
4     ↪ (connected).
5 // For a Euler path the condition is all nodes
6     ↪ in-degree = out-degree, one out-degree+1 =
7     ↪ in-degree (start with this node) and one
8     ↪ out-degree = in-degree+1.
9 vi euler_tour;
10 void hierholzer(int u) {
11     int v;
12     while(!graph[u].empty()) {
13         v = graph[u].back();
14         graph[u].pop_back(); // DESTROYS THE GRAPH.
15         hierholzer(v);
16     }
17     euler_tour.pb(u);
18 }
19 void f_euler_tour() {
20     hierholzer(0);
21     reverse(euler_tour.begin(), euler_tour.end());
22 }
23 // Topo sort the n_sz first values of graph.
24 void topo_sort(int n_sz) {
25     vSorted.clear();
26     visited.assign(n_sz, false);
27     for(int i = 0; i < n_sz; i++) topo_rec(i);
28     reverse(vSorted.begin(), vSorted.end());
29 }
30 vi vSorted;
31 vector<bool> visited;
32 void topo_rec(int u) {
33     if(visited[u]) return;
34     visited[u] = true;
35     for(auto _v : graph[u]) topo_rec(_v);
36     vSorted.pb(u);
37 }

```

Hash String

```

1 // https://www.browerling.com/tools/prime-numbers.
2 // s = a[i], hash = a[0] + b*a[1] + b^2*a[2] +
3     ↪ b^n*a[n].
4 class HashStr {
5     public:
6         string s;
7         int n, n_p;
8         vector<vll> v; // contain the hash for [0..i].
9         vll p = {16532849, 91638611, 83157709}; //
10         ↪ prime numbers. // 15635513 77781229
11         vll base = {37, 47, 53}; // base numbers:
12         ↪ primes that > alphabet size. // 49 83
13         vector<vll> b; // b[i][j] = (b_i^j) % p_i.
14         vector<vll> b_inv; // b_inv[i][j] = (b_i^-j)^-1
15         ↪ % p_i.
16     ll elevate(ll a, ll _b, ll mod){
17         ll ans = 1;
18         while(_b){
19             if(_b & 1) ans = ans * a % mod;
20             _b >>= 1;
21             a = a * a % mod;
22         }
23         return ans;
24     }
25     // a^(mod - 1) = 1, Euler.
26     ll inv(int i, int j){
27         if(b_inv[i][j] != -1) return b_inv[i][j];
28         return b_inv[i][j] = elevate(b[i][j], p[i]
29         ↪ - 2, p[i]);
30     }
31     HashStr() = default; // Initialize later.
32     HashStr(string &s) { // not empty strings.
33         s = _s;
34         n = _s.length();
35         n_p = (int)p.size();
36         v.assign(n_p, vll(n, 0));
37         b.assign(n_p, vll(n, 0));
38         b_inv.assign(n_p, vll(n, -1));
39         int i, j;
40         for(i = 0; i < n_p; i++) {
41             b[i][0] = 1;
42             for(j = 1; j < n; j++) {
43                 b[i][j] = (b[i][j-1]*base[i]) %
44                 ↪ p[i];
45             }
46         }
47     }

```

```

41     }
42     char initial = 'A'; // change initial for
    ↪ range. 'a', 'A'.
43     for(i = 0; i < n_p; i++) {
44         v[i][0] = s[0]-initial+1;
45         for(j = 1; j < n; j++) {
46             v[i][j] = (b[i][j]*(s[j]-initial+1)
    ↪ + v[i][j-1]) % p[i];
47         }
48     }
49 }
50 ll getHash(int l, int r, int imod) {
51     ll ans = v[imod][r];
52     if(l > 0) ans -= v[imod][l-1];
53     ans *= inv(imod, l);
54     ans = ((ans%p[imod])+p[imod])%p[imod];
55     return ans;
56 }
57 // O(1).
58 bool equals(HashStr other, int l, int r) {
59     for(int i = 0; i < n_p; i++) {
60         if(getHash(l, r, i) != other.getHash(l,
    ↪ r, i)) return false;
61     }
62     return true;
63 }
64 // O(1).
65 bool operator == (HashStr other) {
66     if(n != other.n) return false;
67     return equals(other, 0, n-1);
68 }
69 // return the index of the Longest Comon
    ↪ Prefix, -1 if no Common Prefix.
    // O(log n).
70 int LCP(HashStr other) {
71     int l = 0, r = min(n, other.n), mid;
72     if(s[0] != other.s[0]) return -1;
73     if(*this == other) return n-1;
74     while(l + 1 < r) {
75         mid = (l + r) >> 1;
76         if(equals(other, 0, mid)) l = mid;
77         else r = mid;
78     }
79     return l;
80 }
81 }
82 bool operator < (HashStr other) {
83     int id = LCP(other);
84     if(id == -1) return s[0] < other.s[0];
85     if(*this == other) return false;
86     if(id == n) return true; // "ho" < "hol"
87     if(id == other.n) return false;
88     return s[id+1] < other.s[id+1];
89 }
90 }
91 };
92

```

FFT

```

1  typedef complex<double> cd;
2  typedef vector<cd> vcd;
3
4  void show(vcd &e) { //for debug
5      int cont = 0; for(auto el : e) {cout << " + " <<
    ↪ (el.real() > eps ? el.real() : 0) << "x^"
    ↪ << cont++;} cout << endl;
6  }
7  void convolution(vcd &a) { //insert a_i and get y_i
    ↪ = sum_j(a_j*w_i^j)
8      int i, n = a.size(); //n power of 2
9      if(n == 1) return;
10     vcd a_even, a_odd;
11     for(i = 0; i < n; i++) { //divide part of FFT
12         if(i%2) a_odd.pb(a[i]);
13         else a_even.pb(a[i]);
14     }
15     convolution(a_even); //recursive part
16     convolution(a_odd);
17     cd wn = polar(1.0, 2*(double)PI/n), w = 1.0;
    ↪ //w^n are the n roots of n-unity
    //cd w;
18     for(i = 0; i < n/2; i++) {
19         //w = polar(1.0, i*2*(double)PI/n); //avoid
    ↪ precision error, but slower
20         a[i] = a_even[i] + w*a_odd[i]; //A(w^n)^k =
    ↪ Aeven(w^n/2^k) + wn^k*Aodd(w^n/2^k)
21         a[i + n/2] = a_even[i] - w*a_odd[i];
    ↪ //A(w^n)^k = Aeven(w^n/2^(k-n/2)) -
    ↪ wn^(k-n/2)*Aodd(w^n/2^(k-n/2))
22         w = w*wn;
23     }
24 }
25 }
26 void deconvolution(vcd &a) { //insert y_i and get
    ↪ a_i = sum_j(y_j*w_i^-j)/n
27     for(auto &el : a) el = conj(el); //you can
    ↪ conjugate wn and do a[i]/n o can
    ↪ conj(a[i])/n
28     convolution(a); // The coefficients of the
    ↪ polynomial have to be are real
29     for(auto &el : a) el /= (double)a.size();
30 }
31 // Calculate \sum_{i=0}^{n-1} a[i]*b[n-i].
32 vcd FFT(vcd &a, vcd &b) { //multiply polynomial a*b
33     //vcd a = {1.0, 2.0}, b = {3.0}, c; // a and b
    ↪ examples of polynomials to multiply, real
    ↪ coefficients
34     vcd c;
35     if(a.size() < b.size()) swap(a, b);
36     int i, n = a.size();
37     while(n - LSB(n)) n++, a.pb(0.0); //add 0.0's
    ↪ to the next power of two of the next power
    ↪ of two, 3->8
38     n++, a.pb(0.0);
39     while(n - LSB(n)) n++, a.pb(0.0);
40     while((int) b.size() < n) b.pb(0.0); //the
    ↪ grade of a and b equal.
41     convolution(a);
42     convolution(b); //if you want a*a then delete
    ↪ this 2° call
43     for(i = 0; i < n; i++) c.pb(a[i]*b[i]);
44     deconvolution(c);
45     return c;
46 }

```

```

1  const ll mod = 31;
2  ll inverse[mod];
3  // O(mod) calculate inverse[i] % const mod.
4  void calc() {
5      inverse[1] = 1;
6
7      for(ll i = 2; i < mod; i++) {
8          inverse[i] = -(mod/i)*inverse[mod%i];
9          inverse[i] = (inverse[i]%mod + mod) % mod;
10     }

```

Matrix

```

1 // Determinant: https://cp-
  ↪ algorithms.com/linear\_algebra/determinant-
  ↪ gauss.html
2
3 template<typename T>
4 class Matrix {
5     public:
6         int nrow = 0;
7         int ncol = 0;
8         vector<vector<T>> v;
9         Matrix() {}
10        // Empty Matrix.
11        Matrix(int _nrow, int _ncol) {
12            nrow = _nrow;
13            ncol = _ncol;
14            v.assign(nrow, vector<T>(ncol, 0));
15        }
16        // Example: Matrix<ll> a({{1, 2}, {3, 4}}); //
17        ↪ Can't use for one column vector.
18        Matrix(vector<vector<T>> _v) {
19            nrow = _v.size();
20            ncol = _v[0].size();
21            v = _v;
22        }
23        friend ostream& operator << (ostream &os,
24        ↪ Matrix<T> m) {
25            int i, j;
26            for(i = 0; i < m.nrow; i++) {
27                for(j = 0; j < m.ncol; j++) {
28                    if(j) cout << " ";
29                    cout << m.v[i][j]; // Becareful
30                    ↪ with "-0".
31                }
32                cout << "\n";
33            }
34            return os;
35        }
36        Matrix<T> operator + (const Matrix<T> other) {
37            int i, j;
38            Matrix<T> ans(nrow, ncol);
39            for(i = 0; i < nrow; i++) {
40                for(j = 0; j < ncol; j++) {
41                    ans.v[i][j] = v[i][j] +
42                    ↪ other.v[i][j];
43                }
44            }
45            return ans.delete_negative_cero();
46        }
47        // Use this for an empty square Matrix to
48        ↪ create an identity Matrix.
49        Matrix<T> convert_to_identity() {
50            for(int i = 0; i < nrow; i++) v[i][i] = 1;
51            return *this;
52        }
53        Matrix<T> operator * (const Matrix<T> other) {
54            int i, j, k;
55            Matrix<T> ans(nrow, other.ncol);
56            for(i = 0; i < nrow; i++) {
57                for(j = 0; j < other.ncol; j++) {
58                    for(k = 0; k < ncol; k++) {
59                        ans.v[i][j] +=
60                        ↪ v[i][k]*other.v[k][j];
61                    }
62                }
63            }
64            return ans.delete_negative_cero();
65        }
66        Matrix<T> operator ^ (ll ex) {
67            if(ex == 0) {
68                Matrix<T> ans(nrow, ncol);
69                return ans.convert_to_identity();
70            }
71            Matrix<T> half = (*this) ^ (ex/2);
72            if(ex%2) return half * half * (*this);
73            else return half * half;
74        }
75        bool operator == (const Matrix<T> other) {
76            int i, j;
77            if(nrow != other.nrow || ncol !=
78            ↪ other.ncol) return false;
79            for(i = 0; i < nrow; i++) {
80                for(j = 0; j < ncol; j++) {
81                    if(abs(v[i][j] - other.v[i][j]) >
82                    ↪ eps) return false;
83                }
84            }
85            return true;
86        }
87        bool is_null_matrix() {
88            return ncol == 0 || nrow == 0;
89        }
90        // Change "-0" by "0".
91        Matrix<T> delete_negative_cero() {
92            int i, j;
93            for(i = 0; i < nrow; i++) {
94                for(j = 0; j < ncol; j++) {
95                    if(abs(v[i][j]) < eps) v[i][j] = 0;
96                }
97            }
98            return *this;
99        }
100        static Matrix<T> gaussian_elimination(Matrix<T>
101        ↪ mat, Matrix<T> dato) {
102            int i, j, k, imx;
103            T mx, val;
104            for(i = 0; i < mat.ncol; i++) {
105                mx = mat.v[i][i];
106                imx = i;
107                for(j = i+1; j < mat.nrow; j++) {
108                    if(mat.v[j][i] > mx) {
109                        mx = mat.v[j][i];
110                        imx = j;
111                    }
112                }
113                // If no pivot found, the matrix is not
114                ↪ invertible. Its determinant is 0.
115                if(mat.v[imx][i] == 0) return
116                ↪ Matrix<T>(0, 0);
117                // Swap the line with the highest
118                ↪ value.
119                for(j = i; j < mat.ncol; j++) {
120                    swap(mat.v[i][j], mat.v[imx][j]);
121                }
122                for(j = 0; j < dato.ncol; j++) {
123                    swap(dato.v[i][j], dato.v[imx][j]);
124                }
125                for(j = i+1; j < mat.nrow; j++) {
126                    T factor = - mat.v[j][i] /
127                    ↪ mat.v[i][i]; // Change if using
128                    ↪ modulus.
129                    for(k = i; k < mat.ncol; k++) {
130                        mat.v[j][k] += factor *
131                        ↪ mat.v[i][k];
132                    }
133                }
134            }
135        }

```



```

117         }
118         for(k = 0; k < dato.ncol; k++) {
119             dato.v[j][k] +=
120                 ↪ factor*dato.v[i][k];
121         }
122     }
123     // Solving Ux = dato.
124     // For every column of dato.
125     for(k = 0; k < dato.ncol; k++) {
126         for(i = mat.nrow-1; i >= 0; i--) {
127             val = dato.v[i][k];
128             for(j = i+1; j < mat.ncol; j++) {
129                 val -= mat.v[i][j] *
130                     ↪ dato.v[j][k];
131             }
132             dato.v[i][k] = val / mat.v[i][i];
133         }
134     }
135     return dato.delete_negative_cero();
136 }
137 // If you are going to *, it loses a lot of
138 ↪ precision.
139 static Matrix<T> inverse(Matrix<T> mat) {
140     Matrix<T> id(mat.nrow, mat.ncol);
141     id.convert_to_identity();
142     return gaussian_elimination(mat, id);
143 }
144 };

```

Discrete logarithm/root

```

1  ll elevate(ll a, ll b, ll mod) { // b >= 0.
2      ll ans = 1;
3      while(b) {
4          if(b & 1) ans = ans * a % mod;
5          b >>= 1;
6          a = a * a % mod;
7      }
8      return ans;
9  }
10
11  // phi of Euler. O(sqrt(n)).
12  ll get_phi(ll n) {
13      ll ans = n, i;
14      for(i = 2; i*i <= n; i++) {
15          if(n%i == 0) {
16              while(n%i == 0) n /= i;
17              ans -= ans/i;
18          }
19      }
20      if(n > 1) ans -= ans/n;
21      return ans;
22  }
23
24  // Return g such that for all x coprime with mod
25  ↪ exists k : (g^k == x)% mod.
26  // g^k generate all the elements.
27  // If g is the primitive root of mod, you can take
28  ↪ log_g{x} in both sides.
29  // Exists iff mod is 1, 2, 4, (odd p)^k, 2*(odd
30  ↪ p)^k.
31  // Complexity O(mod log(mod)).
32  ll primitive_root(ll mod) {
33      ll phi = get_phi(mod);
34      vll factors; // Factorize phi.
35      ll i, num = phi;
36      bool ok;
37      if(mod == 1) return 0;
38      if(mod == 2) return 1;
39      for(i = 2; i*i <= num; i++) {
40          if(num%i == 0) {
41              factors.pb(i);
42              while(num%i == 0) num /= i;
43          }
44      }
45      if(num > 1) factors.pb(num);
46      // Try every coprime number.
47      for(i = 2; i < mod; i++) {
48          if(__gcd(i, mod) != 1) continue;
49          ok = true;
50          for(auto p : factors) {
51              if(elevate(i, phi/p, mod) == 1) {
52                  ok = false;
53                  break;
54              }
55          }
56          if(ok) return i;
57      }
58      return -1;
59  }
60
61  // Return the smallest x such as a^x == b % mod, or
62  ↪ -1 if no answer exists.
63  // x = sqrt(mod)*p - q. Baby step - Giant step
64  ↪ algorithm.
65  // Complexity O(sqrt(mod)).
66  ll discrete_logarithm(ll a, ll b, ll mod) {
67      a = (a%mod + mod)%mod; b = (b%mod + mod)%mod;
68      if(a == 0 && b == 0) return 1;
69      if(a == 0) return -1; // 0^0 sometimes is 1.
70      if(b == 1) return 0;
71      ll k = 1, sq = sqrt(mod) + 1, q, p, g, asq = 1,
72      ↪ aq = 1, ap = 1, add = 0;
73      unordered_map<ll, ll> value;
74      // if a is not coprime con mod then transform
75      ↪ it to k*a^x == b%mod.
76      while((g = __gcd(a, mod)) != 1) {
77          if(b == k) return add; // Stop decreasing
78          ↪ x.
79          if(b%g) return -1;
80          b /= g;
81          mod /= g;
82          add++;
83          k = (k * (a / g))%mod;
84      }
85      // Meet in the middle, smallest x is high q and
86      ↪ low p.
87      for(q = 0; q <= sq; q++) {
88          value[(b*aq)%mod] = q;
89          aq = (aq*a)%mod;
90      }
91      for(p = 1; p <= sq; p++) asq = (asq*a)%mod;
92      for(p = 1; p <= sq; p++) {
93          ap = (ap*asq)%mod;
94          if(value.count((k*ap)%mod))
95              return sq*p - value[(k*ap)%mod] + add;
96      }
97      return -1;
98  }
99
100  ll p_root = -1;

```



```

92 // Get x such that  $x^k \equiv b \pmod{mod}$ .
93 ll discrete_root(ll k, ll b, ll mod) {
94     // If you change the mod then calculate again
95     // the primitive root.
96     if(p_root == -1) p_root = primitive_root(mod);
97     ll num = elevate(p_root, k, mod);
98     ll y = discrete_logarithm(num, b, mod);
99     return elevate(p_root, y, mod);
100 }

```

Mobius inversion

```

1 // You can do Mobius inversion, that is  $\sum_{d|n}$ 
2 //  $\mu[d] = [n == 1]$ . (Maybe  $n = \gcd(a_i)$  usually).
3 const int MAX = 1e5;
4 int mu[MAX]; //  $\mu(n) = 0$  if  $n$  is square prime and
5 //  $(-1)^t$  if  $n = p_1 \dots p_t$ .
6 void mobius_ini() {
7     vi prime;
8     ll i;
9     bool is_composite[MAX];
10    fill(is_composite, is_composite+MAX, false);
11    mu[1] = 1;
12    for(i = 2; i < MAX; i++) {
13        if(!is_composite[i]) prime.pb(i), mu[i] = -1;
14        for(auto p : prime) {
15            if(i * p >= MAX) break;
16            is_composite[i * p] = true;
17            if(i % p == 0) { // if p divides i.
18                //  $\mu[i * p] = 0$ ; // already 0.
19                break;
20            } else mu[i * p] = mu[i] * mu[p];
21        }
22    }
23 }

```

Lazy Segment Tree

```

1 template<typename T>
2 class Node { // Only modify this class.
3 public:
4     T value = numeric_limits<T>::max(); // max for
5     // MIN query.
6     static const T lazy_default = 0; // Default
7     // value for lazy.
8     T lazy = lazy_default;
9     Node(T _value) {value = _value;}
10    // Merge nodes.
11    Node(Node<T> a, Node<T> b) {value =
12        min(a.value, b.value);} // MIN query.
13    Node() = default;
14    void actualize_update(T x) {
15        value += x; // MIN query + (= SET update),
16        // (+= SUM update).
17        lazy += x; // MIN query + (= SET update),
18        // (+= SUM update).
19    }
20 };
21
22 template<typename T>
23 class Lazy_SegTree {
24 public:
25     vector<Node<T>> tree;
26     vector<T> v_input;
27     int v_size;
28     // Value is the real value, and lazy is only
29     // for its children.
30     void push_lazy(int k, int l, int r) {
31         if(l != r) {
32             tree[k<<1].\
33                 actualize_update(tree[k].lazy);
34             tree[k<<1|1].\
35                 actualize_update(tree[k].lazy);
36             tree[k] = Node<T>(tree[k<<1],
37                             tree[k<<1|1]);
38         }
39         tree[k].lazy = tree[k].lazy_default;
40     }
41     void build(int k, int l, int r) {
42         if(l == r) {tree[k] = Node<T>(v_input[l]);
43             return;}
44         int mid = (l + r) >> 1;
45         build(k<<1, l, mid);
46         build(k<<1|1, mid+1, r);
47         tree[k] = Node<T>(tree[k<<1],
48                             tree[k<<1|1]);
49     }
50     void update(int k, int l, int r, int ql, int
51         qr, T x) {
52         push_lazy(k, l, r);
53         if(ql < l || r < qr) return;
54         if(ql <= l && r <= qr) {
55             tree[k].actualize_update(x);
56         } else {
57             int mid = (l + r) >> 1;
58             update(k<<1, l, mid, ql, qr, x);
59             update(k<<1|1, mid+1, r, ql, qr, x);
60         }
61         push_lazy(k, l, r);
62     }
63     Node<T> query(int k, int l, int r, int ql, int
64         qr) {
65         push_lazy(k, l, r);
66         if(ql <= l && r <= qr) return tree[k];
67         int mid = (l + r) >> 1;
68         if(qr <= mid) return query(k<<1, l, mid,
69             ql, qr);
70         if(mid+1 <= ql) return query(k<<1|1, mid+1,
71             r, ql, qr);
72         Node<T> a = query(k<<1, l, mid, ql, qr);
73         Node<T> b = query(k<<1|1, mid+1, r, ql,
74             qr);
75         return Node<T>(a, b);
76     }
77 public:
78     Lazy_SegTree(vector<T> v) {
79         v_input = v;
80         v_size = v_input.size();
81         tree.assign(4*v_size, {});
82         build(1, 0, v_size-1);
83     }
84     void update(int ql, int qr, T x) { // [ql, qr].
85         update(1, 0, v_size-1, ql, qr, x);
86     }
87     T query(int ql, int qr) { // [ql, qr].
88         Node<T> ans = query(1, 0, v_size-1, ql,
89             qr);
90         return ans.value;
91     }
92 };

```

Sparse Table

```

1 // Sparse Table, table[i][j] = covers [i, i + 2^j - 1], range 2^j.
  ↳ 1], range 2^j.
2 // CAN'T UPDATE VALUES.
3 const ll MAX = 1e5;
4 const int LOG2_MAX = 22; // log2(MAX).
5 ll table[MAX][LOG2_MAX]; // Outside class.
6 template<typename T>
7 class SparseTable {
8     int n;
9     T f(T a, T b) {
10         return min(a, b);
11     }
12 public:
13     SparseTable(vector<T> &v) {
14         int i, j;
15         n = v.size();
16         for(i = 0; i < n; ++i) table[i][0] = v[i];
17
18         for(j = 1; j < LOG2_MAX; ++j){
19
20             for(i = 0; i < n; ++i){
21                 if(i + (1ll << (j - 1)) >= n)
22                     ↳ break;
23                 table[i][j] = f(table[i][j - 1],
24                               ↳ table[i + (1ll << (j - 1))][j -
25                               ↳ 1]);
26             }
27         }
28     }
29     // [ql..qr], [0..n-1].
30     T query(int ql, int qr) {
31         int lg2_dif = -1, num = qr - ql;
32         if(ql == qr) return table[ql][0];
33         while(num) lg2_dif++, num >>= 1;
34         return f(table[ql][lg2_dif], table[qr -
35                               ↳ (1ll << lg2_dif) + 1][lg2_dif]);
36     }
37 };

```

MO's algorithm

```

1 //MO's algorithm, similar than sqrt decomposition.
  ↳ First sort the queries and then keep adding and
2 //removing elements until your current interval is
  ↳ the query interval and report the answer
3 //Usefull when you can compose the answer with a
  ↳ smaller or bigger interval. O((Q+N)sqrt(N))
4 const int BLOCK = //3; //sqrt(max v.size)
5
6 struct Query{
7     int l, r, id;
8 };
9
10 //Sort first by block, second by R
11 bool Query_cmp(Query a, Query b) {
12     if(a.l / BLOCK != b.l / BLOCK) {
13         return a.l / BLOCK < b.l / BLOCK;
14     }
15     return a.r < b.r;
16 }
17
18 const int MAX = 1e5+4;
19 vi v, freq(MAX, 0);
20 int answer = 0;
21
22 //add data to the answer
23 void add(int i) {
24     ++freq[v[i]];
25     if(freq[v[i]] == 1) ++answer;
26 }
27
28 //remove data to the answer
29 void remove(int i) {
30     --freq[v[i]];
31     if(freq[v[i]] == 0) --answer;
32 }
33
34 void MO() {
35     int i, currL = 0, currR = 0;
36     v = {2, 3, 1, 1, 2, 1, 2, 3};
37     vector<Query> vq = {{0, 5, 0}, {6, 7, 1}, {0,
38                               ↳ 3, 2}};
39     //Sort the queries
40     sort(vq.begin(), vq.end(), Query_cmp);
41
42     //The answer contains data of the interval
43     ↳ [L..R)
44     for(i = 0; i < (int)vq.size(); i++) {
45         while(currL < vq[i].l) {
46             remove(currL);
47             ++currL;
48         }
49         while(currL > vq[i].l) {
50             --currL;
51             add(currL);
52         }
53
54         while(currR <= vq[i].r) {
55             add(currR);
56             currR++;
57         }
58         while(currR - 1 > vq[i].r) {
59             currR--;
60             remove(currR);
61         }
62
63         cout << "[" << vq[i].l << " " << vq[i].r <<
64             ↳ ":" << answer << endl;
65         //ans[vq[i].id] = answer //to sort the
66         ↳ answer
67     }
68 }

```

Square Root Decomposition

```

1 template<typename T>
2 class SquareRootDecomposition {
3     const int B = //3; //size of the bucket, ~
4     ↳ sqrt(N)
5     vector<T> bucket; //(N/B + 1, 0T);
6
7     vector<T> v;
8 public:
9     SquareRootDecomposition(vector<T> &_v) {
10         bucket.assign((int)_v.size()/B + 1, 0);
11         v = _v;

```

```

10     for(int i = 0; i < (int)v.size(); i++) {
11         bucket[i/B] += v[i];
12     }
13 }
14 // [l..r].
15 T query(int l, int r) {
16     T ans = 0;
17     int i;
18     if(l/B == r/B) {
19         for(i = l; i <= r; i++) ans += v[i]; //
20         ↪ Same block.
21         return ans;
22     }
23     for(i = l/B + 1; i <= r/B - 1; i++) ans +=
24         ↪ bucket[i]; //middle blocks
25
26     for(i = l; i/B == l/B; i++) ans += v[i];
27     ↪ //left block
28     for(i = r; i/B == r/B; i--) ans += v[i];
29     ↪ //right block
30
31     return ans;
32 }
33 // Replace v[x] by dx.
34 void update(int x, T dx) {
35     bucket[x/B] += dx - v[x];
36     v[x] = dx;
37 }
38 };

```

Aho Corasick

```

1 //construct trie O(m) + automaton O(mk), O(mk)
2 ↪ memory, m = sum(len(word_i))
3 #define next asdfa
4 //size of alphabet, 26 lowercase
5 const int k = 26;
6
7 struct vertex{
8     vi next;
9     //number of words ending at current vertex
10    int leaf;
11    //ancestor p and ch is the transition of p->v
12    int p;
13    char pch;
14    //proper suffix link of the vertex
15    int link;
16    vi go;
17    //how many suffixes there are in the tree;
18    int count;
19
20    vertex(int _p, char _pch) {
21        next.assign(k, -1);
22        leaf = 0;
23        this->p = _p;
24        this->pch = _pch;
25        link = -1;
26        go.assign(k, -1);
27        count = -1;
28    }
29 };
30
31 vector<vertex> t = {{-1, '$'}};
32 int t_size = 1;
33
34 //add string to the trie t
35 void add_string(string &s) {
36     int c, p = 0;
37     for(char ch : s) {
38         c = ch - 'a';
39         if(t[p].next[c] == -1) {
40             t.pb({p, ch});
41             t[p].next[c] = t_size++;
42         }
43         p = t[p].next[c];
44     }
45     t[p].leaf++;
46 }
47
48 //Search for any proper suffix of v that has
49 ↪ next[c] transition
50 //call go(v, ch) for move the automaton from the
51 ↪ vertex v using transition ch
52
53 int go(int v, char ch);
54
55 //get the proper suffix link of v. Once called,
56 ↪ don't call anymore add_strings
57 int get_link(int v) {
58     if(t[v].link == -1) {
59         if(v == 0 || t[v].p == 0) t[v].link = 0;
60         else t[v].link = go(get_link(t[v].p),
61             ↪ t[v].pch);
62     }
63     return t[v].link;
64 }
65
66 int go(int v, char ch) {
67     int c = ch - 'a';
68     if(t[v].go[c] == -1) {
69         if(t[v].next[c] != -1) t[v].go[c] =
70             ↪ t[v].next[c];
71         //The root doesn't have next[c]
72         else if(v == 0) t[v].go[c] = 0;
73         else {
74             t[v].go[c] = go(get_link(v), ch);
75         }
76     }
77     return t[v].go[c];
78 }
79
80 //get the count of v
81 int count(int v) {
82     if(t[v].count == -1) {
83         t[v].count = t[v].leaf;
84         if(v != 0) t[v].count +=
85             ↪ count(get_link(v));
86     }
87     return t[v].count;
88 }
89
90 //search the number of the strings in the automaton
91 ↪ that are in the text
92 int search_num_string(string &text) {
93     int p = 0, ans=0;
94
95     for(auto ch : text) {
96         ans += count(p);
97         p = go(p, ch);
98     }
99     ans += count(p);
100    return ans;
101 }

```

LPS

```

1 // LPS Longest Palindromic Substring, O(n).
2 void Manacher(string &str) {
3     char ch = '#'; // '#' a char not contained in
4     ↪ str.
5     string s(1, ch), ans;
6     for(auto c : str) {s += c; s += ch;}
7     int i, n = s.length(), c = 0, r = 0;
8     vi lps(n, 0);
9     for(i = 1; i < n; i++) {
10         // lps[i] >= it's mirror, but falling in
11         ↪ the interval [L..R]. L = c - (R - c).
12         if(i < r) lps[i] = min(r - i, lps[c - (i -
13         ↪ c)]);
14         // Try to increase.
15
16         while(i-lps[i]-1 >= 0 && i+lps[i]+1 < n &&
17         ↪ s[i-lps[i]-1] == s[i+lps[i]+1])
18         ↪ lps[i]++;
19         // Update the interval [L..R].
20         if(i + lps[i] > r) c = i, r = i + lps[i];
21     }
22     // Get the longest palindrome in ans.
23     int pos = max_element(lps.begin(), lps.end()) -
24     ↪ lps.begin();
25     for(i = pos - lps[pos]; i <= pos + lps[pos];
26     ↪ i++) {
27         if(s[i] != ch) ans += s[i];
28     }
29     //cout << ans.size() << "\n";
30 }

```

Suffix Automaton

```

1 #define next _42_
2 //Suffix Automaton, save a directed acyclic graph
3 ↪ and a suffix link tree with all the suffix of a
4 ↪ word
5 struct state {
6     //length of the longest string in the
7     ↪ equivalence classes
8     int len;
9     //suffix link
10    int link = -1;
11    map<char, int> next;
12    state(int _len) {
13        len = _len;
14    }
15 };
16
17 vector<state> t = {{0}};
18 int t_size = 1, last = 0;
19
20 //add a character to the automaton
21 //last is the state of the last char c added, p is
22 ↪ the head of the automaton
23 //q is the state to duplicate
24 void sa_extend(char c) {
25     int p = last, q;
26     t.pb({t[last].len + 1});
27     last = t_size; t_size++;
28     //add c to the previous suffixes
29     while(p != -1 && t[p].next.find(c) ==
30     ↪ t[p].next.end()) {
31         t[p].next[c] = last;
32         p = t[p].link;
33     }
34     //first time of c in the string
35     if(p == -1) {
36         t[last].link = 0;
37         return;
38     }
39     q = t[p].next[c];
40     if(t[p].len + 1 == t[q].len) {
41         t[last].link = q;
42         return;
43     }
44     //clone state q
45     t.pb({t[p].len + 1});
46     t_size++;
47     t[t_size - 1].next = t[q].next;
48     t[t_size - 1].link = t[q].link;
49
50     //add links of last and q
51     t[last].link = t_size - 1;
52     t[q].link = t_size - 1;
53
54     //point the last suffixes to q cloned
55     while(p != -1 && t[p].next.find(c) !=
56     ↪ t[p].next.end()) {
57         t[p].next[c] = t_size - 1;
58         p = t[p].link;
59     }
60
61     //O(s.length()) to create the automaton. Be careful
62     ↪ adding any char once called another function
63     void sa_ini(string &s) {
64         for(char c : s) sa_extend(c);
65     }
66
67     //A path from root to a terminal node is a suffix
68     ↪ of the automaton string
69     vector<bool> terminal;
70     void sa_terminal() {
71         int p = last;
72         if(terminal.empty() == false) return;
73         ↪ //previously calculated
74         terminal.assign(t_size, false);
75         while(p != -1) {
76             terminal[p] = true;
77             p = t[p].link;
78         }
79     }
80
81     //true if w is a substring of the automaton string
82     //Also s is the longest prefix of w that is in s
83     //w is a suffix if the last p is a terminal state
84     bool sa_is_substr(string &w) {
85         int p = 0; //string s;
86         for(char ch : w) {
87             if(t[p].next.find(ch) == t[p].next.end())
88             ↪ return false;
89             p = t[p].next[ch];
90             //s += c;
91         }
92         return true;
93     }
94 }
95
96 vll dp_num_substr;

```

```

87 ll num_substr_rec(int i) {
88     ll sum = 1;
89     if(dp_num_substr[i] != -1) return
        ↳ dp_num_substr[i];
90     for(auto el : t[i].next) sum +=
        ↳ num_substr_rec(el.se);
91     return dp_num_substr[i] = sum;
92 }
93 //Number of different substrings of the automaton
    ↳ string (Is the number of different paths in the
    ↳ automaton)
94 //For the number of the length of all different
    ↳ substring the recursive formula is
95 // sum of dp_num_substr[i] + dp_num_len_substr[i],
    ↳ the previous answer + 1*number of different
    ↳ substrings
96 ll sa_num_substr() {
97     if(dp_num_substr.empty() == false) return
        ↳ dp_num_substr[0]; //previously calculated
98     dp_num_substr.assign(t_size, -1);
99     num_substr_rec(0);
100    return dp_num_substr[0]; // -1 if you don't
        ↳ want the empty substring
101 }
102
103 //k-th string in the sorted substrings set of the
    ↳ automaton string. It's the k-th path in the
    ↳ graph
104 //k is [0..sa_num_substr()-1]
105 string sa_k_substr(int k) {
106     int p = 0;
107     char prev = '$';
108     string ans = "";
109     if(k > sa_num_substr()) return ans; //not
        ↳ exists that k-th string, error
110     while(k > 0) {
111         prev = '$';
112         for(auto el : t[p].next) {
113             prev = el.fi;
114             if(dp_num_substr[el.se] >= k) break;
115             k -= dp_num_substr[el.se];
116         }
117         if(prev == '$') break; //error
118         ans += prev;
119         p = t[p].next[prev];
120         k--;
121     }
122     return ans;
123 }
124 //lexicographically smallest cyclic shift of the
    ↳ string s
125 string sa_small_cyclic_shift(string &s) {
126     int p = 0, cnt = s.length();
127     string ans = "";
128     sa_ini(s + s); //initialize sa with s+s, the
        ↳ ans is greedy the first path with length
        ↳ s.length()
129     while(cnt-->0) {
130         auto el = *(t[p].next.begin()); //take
        ↳ greedy the first edge
131         ans += el.fi;
132         p = el.se;
133     }
134     return ans;
135 }
136
137 //int sa_num_ocurrences(string w); //Better use
    ↳ Aho-Corasick
138
139 //Test of the automaton string, the number of the
    ↳ substrings and the substrings, sorted
140 void sa_test1() {
141     ll i, n;
142     sa_ini("test");
143     n = sa_num_substr();
144     cout << n << endl;
145     for(i = 0; i < n; i++)
146         cout << sa_k_substr(i) << endl;
147 }

```

Python Template

```

1 import math, sys
2 input = sys.stdin.readline
3 # try:
4 #     x = input() # until EOF.
5 # except:
6 #     exit(0)
7 # v = [k for k in map(int, s.split(' '))]
8
9 #define mp make_pair
10 #define pb push_back
11 #define fi first
12 #define se second
13 #define LSB(x) ((x) & (-x))
14 #define is_set(x, i) (((x)>>(i))&1)
15
16 #define set_bit(x, i) {(x) |= 1ll<<(i);}
17 #define unset_bit(x, i) {(x) = ((x) | (1ll<<(i))) ^
    ↳ (1ll<<(i))};
18
19 const long double PI = acos(-1);
20 const long double eps = 1e-9;
21 const long long inf = LLONG_MAX / 10;

```

2-SAT

```

1 // 2-SAT. Check values (xi or xj) and ... and (xk
    ↳ or xz).
2 // xi will be element 2*i and not xi will be 2*i+1.
    ↳ Change them with xi xor 1.
3 vector<vi> graph; // size of graph will be 2*(the
    ↳ number of xi).
4
5 int get_element(int n) {return 2*n;} // Get pos of
    ↳ xi.
6 int get_not_element(int n) {return 2*n + 1;} // Get
    ↳ pos of not xi.
7 // Add (xi or xj), two edges: (not xj => xi) and
    ↳ (not xi => xj).
8 // inclusive or: 1 or 1 = 1. For exclusive use 2
    ↳ clauses. (xi or xj) and (not xi or not xj).
9 void add_or_clause(int i, int j) {
10     int neg_i = i^1, neg_j = j^1;
11     graph[neg_i].pb(j);
12     graph[neg_j].pb(i);

```

```

13 }
14
15 // Use Kosaraju to find the SCCs.
16 vector<vi> graphRev;
17 stack<int> s;
18 vector<bool> visited; // It will be reutilized in
    ↪ SAT.
19 vector<vi> components;
20
21 void dfs1(ll u){
22     visited[u] = true;
23     for(auto v : graph[u]){
24         if(!visited[v]) dfs1(v);
25     }
26     s.push(u);
27 }
28 void dfs2(ll u){
29     visited[u] = true;
30     for(auto v : graphRev[u]){
31         if(!visited[v]) dfs2(v);
32     }
33     components.back().pb(u); // One element more to
    ↪ the current component.
34 }
35 void Kosaraju(){
36     ll i, n = graph.size();
37     graphRev.assign(n, vi());
38     s = stack<int>();
39     //transpose graph to graphRev
40     for(i = 0; i < n; ++i){
41         for(auto v : graph[i]){
42             graphRev[v].pb(i);
43         }
44     }
45     visited.assign(n, false);
46     for(i = 0; i < n; i++){
47         if(!visited[i])
48             dfs1(i);
49     }
50     visited.assign(n, false);
51     components.pb(vi());
52     while(true) {
53         while(!s.empty() && visited[s.top()] ==
            ↪ true) s.pop();
54         if(s.empty()) break;
55         dfs2(s.top());
56         components.pb(vi()); // End of the current
            ↪ component.
57     }
58 }
59 // Do a topoSort of the SCCs.
60 vector<vi> graph_topo;
61 vi vSorted;
62 void topo_rec(int u) {
63     if(visited[u]) return;
64     visited[u] = true;
65     for(auto _v : graph_topo[u]) topo_rec(_v);
66     vSorted.pb(u);
67 }

```

```

68 // Topo sort the n_sz first values of graph.
69 void topo_sort(int n_sz) {
70     vSorted.clear();
71     visited.assign(n_sz, false);
72     for(int i = 0; i < n_sz; i++) topo_rec(i);
73     reverse(vSorted.begin(), vSorted.end());
74 }
75 // xi_value[i] is 0 if xi is false, 1 if true.
76 vi xi_value;
77 // If you know in advance elements of xi add a
    ↪ clause (xi or xi).
78 // Return true if the base is satisfiable, false
    ↪ otherwise.
79 // All calls are O(n).
80 vi node2component; // Index of the component of the
    ↪ node.
81 vi component2order; // Order of the component in
    ↪ the topological sort.
82 bool SAT() {
83     Kosaraju();
84     int n_components = components.size(), i, n =
        ↪ graph.size();
85     node2component.assign(n, 0);
86     graph_topo.assign(n_components, vi());
87     component2order.assign(n_components, 0);
88     // All components in the same SCC will have the
        ↪ same truth value.
89     for(i = 0; i < n_components; i++) {
90         for(auto u : components[i])
91             ↪ node2component[u] = i;
92     }
93     // If xi and not xi are in the same component
        ↪ is UNSAT.
94     for(i = 0; i < n; i += 2) {
95         if(node2component[i] ==
            ↪ node2component[i+1]) return false;
96     }
97     for(i = 0; i < n; i++) {
98         for(auto u : graph[i]) {
99             if(node2component[i] !=
                ↪ node2component[u])
100             ↪ graph_topo[node2component[i]].pb(node2com
101         }
102     }
103     topo_sort(n_components);
104     for(i = 0; i < n_components; i++) {
105         component2order[vSorted[i]] = i;
106     }
107     xi_value.assign(n/2, -1);
108     for(i = 0; i < n; i += 2) {
109         if(component2order[node2component[i]] >
            ↪ component2order[node2component[i+1]])
110             ↪ xi_value[i/2] = true;
111         else xi_value[i/2] = false;
112     }
113     return true;
114 }

```

Convex Hull Trick

```

1 typedef long double ftype; //NOT USE LONG LONG,
    ↪ complex cast to minor precision
2 typedef complex<ftype> point;
3 #define x real
4 #define y imag
5

```

```

6 ftype dot(point a, point b) {
7     return (conj(a)*b).x();
8 }
9
10 ftype cross(point a, point b) {
11     return (conj(a)*b).y();

```



```

12 }
13 //get min{k_i * x + b_i}.
14 //Insert k_i in ascending order. max {} = -min{-()}
15 //Decreasing k_i then add -k_i and query -x
16 class ConvexHullTrick {
17     point im = {0, 1};
18     vector<point> hull, normal;
19 public:
20     void add_line(ftype k, ftype b) {
21         point nw = {k, b};
22         //Create lower convex hull, with increasing
23         ↪ k (add lines to the right only)
24         while(!normal.empty() && dot(normal.back(),
25         ↪ (nw - hull.back())) < 0) {
26             hull.pop_back();
27             normal.pop_back();
28         }
29         if(!hull.empty()) { //add the normal of
30             ↪ vector (nw - hull.back())
31
32         normal.pb(im * (nw - hull.back()));
33     }
34     hull.pb(nw);
35 }
36 // normal anti-clockwise, hull[it], normal
37 ↪ clockwise
38 ftype query(ftype x) {
39     point px = {x, 1}; //query is min dot
40     ↪ product {k, b} * {x, 1}
41
42     int pos = lower_bound(normal.begin(),
43     ↪ normal.end(), px, [](point a, point b)
44     ↪ {
45         return cross(a, b) > 0;
46     }) - normal.begin();
47     return dot(px, hull[pos]);
48 }
49 };

```

Find Centroid

```

1 int subtree_sz[MAX]; // Number of nodes in the
2 ↪ subtree u, rooted at 1.
3 int mn = MAX; // Min {Max{sz(T_i)}} and T_i are the
4 ↪ trees created when cutting node i.
5 vi centroid; // The centroids. Always there are 1
6 ↪ centroid or 2: x and y, and edge x-y exist.
7 vector<pii> subtree[MAX]; // subtree[u] = (fi, se)
8 ↪ fi is one child of u and se is the size of that
9 ↪ subtree.
10 void fill_sz(int u, int p) { // Recursive fill
11 ↪ subtree array.
12     subtree_sz[u] = 1;
13     for(auto v : graph[u]) {
14         if(v == p) continue;
15         fill_sz(v, u);
16         subtree_sz[u] += subtree_sz[v];
17         subtree[u].pb(mp(v, subtree_sz[v]));
18     }
19     if(p != -1)
20         subtree[u].pb(mp(p, n-subtree_sz[u])); // n
21         ↪ is the number of nodes of the graph.
22
23 int mx = 0;
24 for(auto el : subtree[u]) mx = max(mx, el.se);
25 mn = min(mn, mx);
26 }
27 // 1-indexed!
28 void find_centroid() { // fill vi centroid.
29     int mx = 0, i;
30     mn = MAX;
31     centroid.clear();
32     for(i = 1; i <= n; i++) subtree[i].clear();
33     fill_sz(1, -1);
34     for(i = 1; i <= n; i++) {
35         mx = 0;
36         for(auto el : subtree[i]) {
37             mx = max(mx, el.se);
38         }
39         if(mx == mn) centroid.pb(i);
40     }
41 }

```

Java Template

```

1 import java.io.*;
2 import java.math.*;
3 // .setScale(m, RoundingMode.FLOOR);
4 // .divide(denom, m, RoundingMode.FLOOR);
5 @SuppressWarnings("unused")
6 public class Main {
7     private final static int MAX = 100_005;
8     public static void solve() {
9     }
10    public static void main(String[] args) {
11        out = new PrintWriter(new BufferedOutputStream(System.out));
12        sc = new MyScanner();
13        int zz = 1; //sc.nextInt();
14        while (zz-- > 0) solve();
15        out.close();
16    }
17    private static PrintWriter out;
18    private static MyScanner sc;
19    private static class MyScanner {
20        private static final int BUF_SIZE = 2048;
21        BufferedReader br;
22        private MyScanner() {
23            br = new BufferedReader(new InputStreamReader(System.in));

```



```

24     }
25     private boolean isSpace(char c) {
26         return c == '\n' || c == '\r' || c == ' ';
27     }
28     String next() {
29         try {
30             StringBuilder sb = new StringBuilder();
31             int r;
32             while ((r = br.read()) != -1 && isSpace((char) r));
33             if (r == -1) {
34                 return null;
35             }
36             sb.append((char) r);
37             while ((r = br.read()) != -1 && !isSpace((char) r)) {
38                 sb.append((char) r);
39             }
40             return sb.toString();
41         } catch (IOException e) {
42             e.printStackTrace();
43         }
44         return null;
45     }
46     int nextInt() {
47         return Integer.parseInt(next());
48     }
49     long nextLong() {
50         return Long.parseLong(next());
51     }
52     double nextDouble() {
53         return Double.parseDouble(next());
54     }
55 }
56 }

```

Treap

```

1  add srand
2  //https://cp-
   ↳ algorithms.com/data_structures/treap.html
3  //not implemented: find_by_order, order_of_key,
   ↳ find/search
4  template<typename K, typename V> //key (unique) and
   ↳ value (data)
5  class Treap {
6      const pair<V, K> MINVK =
   ↳ mp(numeric_limits<V>::min(),
   ↳ numeric_limits<K>::min());
7      const pair<V, K> MAXVK =
   ↳ mp(numeric_limits<V>::max(),
   ↳ numeric_limits<K>::max());
8      struct node {
9          K key; //unique, time or x-axis for example
10         V data; // f[key] = data
11         node *l = NULL, *r = NULL;
12         int priority;
13         pair<V, K> mx, mn; //maximum and minimum
   ↳ values over all subtrees, V 1° to be
   ↳ comparable <
14         node(K _key, V _data) {
15             key = _key;
16             data = _data;
17             mx = mn = mp(_data, _key);
18             priority =
   ↳ rand(); //((ll)rand() << 16) ~ (ll)rand() * 2
   ↳ //hope there is no collision...
19         }
20     };
21     typedef node* pnode;

22     pnode root = NULL;
23     pair<V, K> getMaximumVK(pnode t) { //O(1)
24         if(!t) return MINVK;
25         else return t->mx;
26     }
27     pair<V, K> getMinimumVK(pnode t) { //O(1)
28         if(!t) return MAXVK;
29         else return t->mn;
30     }
31     void update (pnode t) {
32         if(!t) return;
33         t->mx = max({mp(t->data, t->key),
   ↳ getMaximumVK(t->l),
   ↳ getMaximumVK(t->r)});
34         t->mn = min({mp(t->data, t->key),
   ↳ getMinimumVK(t->l),
   ↳ getMinimumVK(t->r)});
35     }
36     //return a subtree l and r such as key(l) < key
   ↳ < key(r), similar to rotations
37     void split(pnode t, pnode &l, pnode &r, K key)
   ↳ {
38         if(!t) l = r = NULL;
39         else if(key < t->key) split(t->l, l, t->l,
   ↳ key), r = t;
40         else split(t->r, t->r, r, key), l = t;
41         update(t);
42     }
43     //merge two trees l and r into one, t,
   ↳ allKey(l) < allKey(r)
44     void merge(pnode &t, pnode l, pnode r) {
45         if(!l || !r) t = l ? l : r;

```

```

46     else if(l->priority < r->priority)
47         ↪ merge(r->l, l, r->l), t = r;
48     else merge(l->r, l->r, r), t = l;
49     update(t);
50 }
51 void insert(pnode &t, pnode it) {
52     if(!t) t = it;
53     else if(t->priority < it->priority)
54         ↪ split(t, it->l, it->r, it->key), t =
55         ↪ it;
56     else insert(it->key < t->key ? t->l : t->r,
57         ↪ it);
58     update(t);
59 }
60 void erase(pnode &t, K key) { //only erase if
61     ↪ the element exists
62     if(!t) {echo("estas borrando pero no
63         ↪ esta:", key); exit(-1);}
64     if(t->key == key) merge(t, t->l, t->r);
65     else erase(key < t->key ? t->l : t->r,
66         ↪ key);
67     update(t);
68 }
69 void showTree(pnode t) { //preOrder, the
70     ↪ inOrder is the tree sorted
71     if(!t) return;
72     cout << "(" << t->key << "," << t->data <<
73     ↪ " ";
74     showTree(t->l); showTree(t->r);
75 }
76 // UP CONSTRUCTION, BELOW QUERIES
77 pair<V, K> getMaximumAfter(pnode t, K key) {
78     ↪ //max node with key >= t, O(log(n))
79     if(!t) return MINVK;
80     pair<V, K> mid = t->key >= key ?
81     ↪ mp(t->data, t->key) : MINVK;
82     if(key < t->key) return
83     ↪ max({getMaximumAfter(t->l, key),
84     ↪ getMaximumVK(t->r, mid)}); //maybe
85     ↪ equals also here
86     else return max({getMaximumAfter(t->r,
87     ↪ key), mid});
88 }
89 pair<V, K> getMinimumBefore(pnode t, K key) {
90     ↪ //min node with key <= t, O(log(n))
91     if(!t) return MAXVK;
92     pair<V, K> mid = t->key <= key ?
93     ↪ mp(t->data, t->key) : MAXVK;
94     if(key > t->key) return
95     ↪ min({getMinimumBefore(t->r, key),
96     ↪ getMinimumVK(t->l, mid)});
97     else return min({getMinimumBefore(t->l,
98     ↪ key), mid});
99 }
100 pair<V, K> getMinimumKAll(pnode t) { //is the
101     ↪ first node in the inOrder traversal
102     if(t->l) return getMinimumKAll(t->l);
103     return mp(t->data, t->key);
104 }
105 public:
106 void insert(K key, V data) { //O(log n)
107     pnode n = new node(key, data);
108     insert(root, n);
109 }
110 void erase(K key) { //O(log n)
111     erase(root, key);
112 }
113 void showTree() { //debug
114     showTree(root);
115     cout << endl;
116 }
117 pair<V, K> getMaximumAfter(K t) { //O(log n)
118     return getMaximumAfter(root, t);
119 }
120 pair<V, K> getMinimumBefore(K t) { //O(log n)
121     return getMinimumBefore(root, t);
122 }
123 pair<V, K> getMaximumVKAll() { //all the tree,
124     ↪ O(1)
125     return getMaximumVK(root);
126 }
127 pair<V, K> getMinimumVKAll() { //all the tree,
128     ↪ O(1)
129     return getMinimumVK(root);
130 }
131 pair<V, K> getMinimumKAll() { //top of the
132     ↪ tree, O(1), check root not null
133     return getMinimumKAll(root);
134 }
135 };

```

Factorizator

```

1  add srand(time(0)); //ADD srand(time(0));
2  // USE THIS ONLY WHEN NO OTHER OPTION LEFT ...
3  namespace Factorizator {
4      vll primes; //Add primes manually
5      const vll fixed_primes = {2, 3, 5, 7, 11, 13,
6          ↪ 17, 19, 23, 29, 31, 37, 41}; //47, 53, 59
7      //return (a*b)%mod, with numbers up to 1e18
8      ↪ (LL)
9      ll mult(ll a, ll b, ll mod) {
10         ↪ //return ((_int128_t)a*b) % mod; //only
11         ↪ with 64 bits GCC
12         ll ans = 0;
13         while(b) {
14             if(b&1) ans = (ans+a) % mod;
15             b >>= 1;
16             a = (a+a)%mod;
17         }
18         return ans;
19     }
20     ll elevate(ll a, ll b, ll mod){
21         ll ans = 1;
22         while(b){
23             if(b & 1) ans = mult(ans, a, mod);
24             b >>= 1;
25             a = mult(a, a, mod);
26         }
27         return ans;
28     }
29     //a^(mod - 1) = 1, Euler
30     ll inv(ll a, ll mod){
31         return elevate(((a%mod) + mod)%mod, mod -
32             ↪ 2);
33     }
34     //a^{p-1} = 1 mod p => p divides some factor of
35     ↪ (a^{d2^s}+1)*(a^{d2^{s-
36     ↪ 1}}+1)*...*(a^{d+1})*(a^{d-1}
37     //return true if the number is composite, false
38     ↪ if it is not sure
39     bool check_composite(ll num, ll a, ll d, int s)
40     ↪ {

```

```

33     ll x = elevate(a, d, num);
34     if(x == 1 || x == num-1) return false;
35     int i;
36     for(i = 0; i < s; i++) {
37         x = mult(x, x, num);
38         if(x == num-1) return false;
39     }
40     return true;
41 }
42 //Miller_Rabin DETERMINISTIC Version for num up
43   ↳ to 1e18 (all LL)
44 bool isPrime(ll num) { //num-1 == d*2^s
45     bool flag = true;
46     ll d = num-1, s = 0;
47     if(num <= 1) return false;
48     for(auto p : fixed_primes) {
49         if(p == num) return true;
50         if(p%num == 0) return false;
51         ↳ //optimization
52     }
53     while(d%2 == 0) {
54         d /= 2;
55         ++s;
56     }
57     for(auto p : fixed_primes) flag &=
58     ↳ !check_composite(num, p, d, s);
59     return flag;
60 }
61 //a polynomail function modulo mod, it will
62   ↳ contain a cycle
63 ll f_pollard_rho(ll x, ll c, ll mod) {
64     return (mult(x, x, mod)+c) % mod;
65 }
66 //found a factor (maybe not prime) of num. x0
67   ↳ and c are random, change them if the return
68   ↳ is num
69 ll pollard_rho(ll num, ll x0, ll c) {
70     ll x1 = x0, x2 = x0;
71     ll g = 1;
72     if(num == 1) return 1;
73     if(num%2 == 0) return 2;
74     x0 %= num;
75     c %= num;
76     while(g == 1) { //Floyd cycle detection
77         x1 = f_pollard_rho(x1, c, num);
78         x2 = f_pollard_rho(x2, c, num);
79         x2 = f_pollard_rho(x2, c, num);
80         g = __gcd(abs(x1 - x2), num);
81     }
82     return g;
83 }
84 vector<pll> factors; // .fi is the prime, .se
85   ↳ is the exponent.
86 vector<pll> factorize(ll num) {
87     factors.clear();
88     ll y = num, cont;
89     while(num > 1) {
90         if(isPrime(num)) y = num;
91         else {
92             y = pollard_rho(num, rand(),
93             ↳ rand());
94             while(isPrime(y) == false) y =
95             ↳ pollard_rho(y, rand(), rand());
96         }
97         cont = 0;
98         while(num % y == 0) {
99             num /= y; cont++;
100         }
101         factors.pb(mp(y, cont));
102     }
103     sort(factors.begin(), factors.end());
104     return factors;
105 }
106 vll divisors; //will save all the divisors
107 void dfs_div(ll x, ll i) {
108     if(i == (int)factors.size())
109     ↳ {divisors.pb(x); return;}
110     dfs_div(x, i+1);
111     int j;
112     for(j = 0; j < factors[i].se; j++) {x*=
113     ↳ factors[i].fi; dfs_div(x, i+1);}
114 }
115 //NOT TESTED
116 vector<ll> get_divisors(ll num) { //1 and num
117   ↳ inclusive
118     if(factors.empty()) factorize(num);
119     divisors.clear();
120     dfs_div(1, 0);
121     sort(divisors.begin(), divisors.end());
122     return divisors;
123 }
124 //factorize knowing that its primes are in
125   ↳ primes vector
126 vector<pll> factorize_using_primes(ll num) {
127     vector<pll> ans;
128     ll i, cont;
129     for(i = 0; i < (int)primes.size() &&
130     ↳ primes[i] <= num/primes[i]; i++) {
131         if(num%primes[i] == 0) {
132             cont = 0;
133             while(num%primes[i] == 0) {
134                 num /= primes[i];
135                 cont++;
136             }
137             ans.pb(mp(primes[i], cont));
138         }
139     }
140     if(num > 1) ans.pb(mp(num, 1));
141     return ans;
142 }
143 //empty if gcd = 1
144 vector<pll> gcd(vector<pll> &va, vector<pll>
145   ↳ &vb) {
146     vector<pll> ans;
147     int l = 0, r = 0, va_sz = (int)va.size(),
148     ↳ vb_sz = (int)vb.size();
149     while(l < va_sz && r < vb_sz) {
150         if(va[l].fi == vb[r].fi)
151         ↳ {ans.pb(mp(va[l].fi, min(va[l].se,
152         ↳ vb[r].se))); ++l; ++r;}
153         else if(va[l].fi < vb[r].fi) ++l;
154         else ++r;
155     }
156     return ans;
157 }
158 }

```