

Echo's notebook

```
1  FLAGS=-Wall -Wextra -Wshadow -Wno-unused-result -D_GLIBCXX_DEBUG -fsanitize=address -fsanitize=undefined
   ↪ -fno-sanitize-recover
2
3  @g++ A.cpp $(FLAGS) -DJUNCO_DEBUG && ./a.out < z.in
```

```
1  // Iterate over all submasks of a mask. CONSIDER SUBMASK = 0 APART.
2  for(submask = mask; submask > 0; submask = (submask-1)&mask) {}
```

LIS

```
1  vll v_LIS(vll &v) {
2      int i, j, n = v.size();
3      vll lis, lis_time(n), ans;
4      if(!n) return ans;
5      lis.pb(v[0]); lis_time[0] = 1;
6      for(i = 1; i < n; i++) {
7          if(v[i] > lis.back()) {lis.pb(v[i]); lis_time[i] = lis.size(); continue;}
8          int pos = upper_bound(lis.begin(), lis.end(), v[i]) - lis.begin();
9          // if(pos > 0 && lis[pos-1] == v[i]) continue; // USE IF YOU WANT STRICTLY INCREASING.
10         lis[pos] = v[i];
11         lis_time[i] = pos+1;
12     }
13     j = lis.size();
14     for(i = n-1; i >= 0; i--) {
15         if(lis_time[i] == j && (ans.empty() || v[i] <= ans.back())) {ans.pb(v[i]); j--;} // <= or <.
16     }
17     reverse(ans.begin(), ans.end());
18     return ans;
19 }
```

IO

```
1  ios::sync_with_stdio(false); cin.tie(nullptr); cout.tie(nullptr);
2
3  stringstream ss;
4  ss << "Hello world";
5  ss.str("Hello world");
6  while(ss >> s) cout << s << endl;
7  ss.clear();
```

Dates

<pre>1 // Change here and date_to_num. 2 ll is_leap_year(ll y) { 3 // if(y%4 (y%100==0 && y%400)) return 0; // ↪ Complete leap year. 4 if(y%4 != 0) return 0; // Restricted leap year. 5 return 1; 6 } 7 ll days_month[12] = {31, 28, 31, 30, 31, 30, 31, ↪ 31, 30, 31, 30, 31}; 8 ll days_month_accumulate[12] = {31, 59, 90, 120, ↪ 151, 181, 212, 243, 273, 304, 334, 365}; 9 10 // d 1-index, m 1-index. 11 ll date_to_num(ll d, ll m, ll y) { 12 ll sum = d; 13 m -= 2; 14 if(m >= 1) sum += is_leap_year(y); 15 y--; 16 if(m >= 0) sum += days_month_accumulate[m]; 17 if(y >= 0) {</pre>	<pre>18 sum += 365*y; 19 // sum += y/4 -y/100 + y/400; // Complete ↪ leap year. 20 sum += y/4; // Restricted leap year. 21 } 22 return sum; 23 } 24 25 // Tiny optimization, binary search the year, month ↪ and day. 26 void num_to_date(ll num, ll &d, ll &m, ll &y) { 27 d = 1; m = 1; y = 0; // The date searched is >= ↪ this date. 28 while(date_to_num(d, m, y) <= num) y++; 29 y--; 30 while(date_to_num(d, m, y) <= num) m++; 31 m--; 32 while(date_to_num(d, m, y) <= num) d++; 33 d--; 34 }</pre>
--	---

Geometry

```

1  template<typename T>
2  class Point {
3      public:
4          static const int LEFT_TURN = 1;
5          static const int RIGHT_TURN = -1;
6          T x = 0, y = 0;
7          Point() = default;
8          Point(T _x, T _y) {
9              x = _x;
10             y = _y;
11         }
12         friend ostream &operator << (ostream &os,
13             ↪ Point<T> &p) {
14             os << "(" << p.x << " " << p.y << ")";
15             return os;
16         }
17         bool operator == (const Point<T> other) const {
18             return x == other.x && y == other.y;
19         }
20         // Get the (1º) bottom (2º) left point.
21         bool operator < (const Point<T> other) const {
22             if(y != other.y) return y < other.y;
23             return x < other.x;
24         }
25         T euclidean_distance(Point<T> other) {
26             T dx = x - other.x;
27             T dy = y - other.y;
28             return sqrt(dx*dx + dy*dy);
29         }
30         T euclidean_distance_squared(Point<T> other) {
31             T dx = x - other.x;
32             T dy = y - other.y;
33             return dx*dx + dy*dy;
34         }
35     };
36
37     T manhatan_distance(Point<T> other) {
38         return abs(other.x - x) + abs(other.y - y);
39     }
40
41     // Get the height of the triangle with base b1,
42     ↪ b2.
43     T height_triangle(Point<T> b1, Point<T> b2) {
44         if(b1 == b2 || *this == b1 || *this == b2)
45             ↪ return 0; // It's not a triangle.
46         T a = euclidean_distance(b1);
47         T b = b1.euclidean_distance(b2);
48         T c = euclidean_distance(b2);
49         T d = (c*c-b*b-a*a)/(2*b);
50         return sqrt(a*a - d*d);
51     }
52
53     int get_quadrant() {
54         if(x > 0 && y >= 0) return 1;
55         if(x <= 0 && y > 0) return 2;
56         if(x < 0 && y <= 0) return 3;
57         if(x >= 0 && y < 0) return 4;
58         return 0; // Point (0, 0).
59     }
60
61     // Relative quadrant respect the point other,
62     ↪ not the origin.
63     int get_relative_quadrant(Point<T> other) {
64         Point<T> p(other.x - x, other.y - y);
65         return p.get_quadrant();
66     }
67
68     // Orientation of points *this -> a -> b.
69     int get_orientation(Point<T> a, Point<T> b) {
70         T prod = (a.x - x)*(b.y - a.y) - (a.y -
71             ↪ y)*(b.x - a.x);
72         if(prod == 0) return 0;
73         return prod > 0? LEFT_TURN : RIGHT_TURN;
74     }
75
76     // True if a have less angle than b, if *this->a->b is a left turn.
77     bool angle_cmp(Point<T> a, Point<T> b) {
78         if(get_relative_quadrant(a) != get_relative_quadrant(b))
79             return get_relative_quadrant(a) < get_relative_quadrant(b);
80         int ori = get_orientation(a, b);
81         if(ori == 0) return euclidean_distance_squared(a) < euclidean_distance_squared(b);
82         return ori == LEFT_TURN;
83     }
84
85     // Anticlockwise sort starting at 1º quadrant, respect to *this point.
86     void polar_sort(vector<Point<T>> &v) {
87         sort(v.begin(), v.end(), [&](Point<T> a, Point<T> b) {return angle_cmp(a, b);});
88     }
89
90     // Convert v to its convex hull, Do a Graham Scan. O(n log n).
91     void convert_convex_hull(vector<Point<T>> &v) {
92         if(v.size() < 3) return;
93         Point<T> bottom_left = v[0], p2;
94         for(auto p : v) bottom_left = min(bottom_left, p);
95         bottom_left.polar_sort(v);
96         vector<Point<T>> v_input = v; v.clear();
97         for(auto p : v_input) {
98             while(v.size() >= 2) {
99                 p2 = v.back(); v.pop_back();
100                 if(v.back().get_orientation(p2, p) == LEFT_TURN) {
101                     v.pb(p2);
102                     break;
103                 }
104             }
105             v.pb(p);
106         }
107     }
108 };

```

Graphs

Articulation points and bridges

```

1 vector<vi> adyList; // Graph
2 vi num, low; // num and low for DFS
3 int cnt; // Counter for DFS
4 int root, rchild; // Root and number of (DFS)
  ↳ children
5 vi artic; // Contains the articulation
  ↳ points at the end
6 set<pii> bridges; // Contains the bridges at the
  ↳ end
7
8 void dfs(int nparent, int nnode) {
9     num[nnode] = low[nnode] = cnt++;
10    rchild += (nparent == root);
11
12    for (auto a : adyList[nnode]) {
13        if (num[a] == -1) { // Tree edge
14            dfs(nnode, a);
15            low[nnode] = min(low[nnode], low[a]);
16            if (low[a] >= num[nnode]) {
17                artic[nnode] = true;
18            }
19            if (low[a] > num[nnode]) {
20                bridges.insert((nnode < a) ?
  ↳ mp(nnode, a) : mp(a, nnode));
21
22        } else if (a != nparent) { // Back edge
23            low[nnode] = min(low[nnode], num[a]);
24        }
25    }
26 }
27 void findArticulations(int n) {
28     cnt = 0;
29     low = num = vi(n, -1);
30     artic = vi(n, 0);
31     bridges.clear();
32
33     for (int i = 0; i < n; ++i) {
34         if (num[i] != -1) {
35             continue;
36         }
37         root = i;
38         rchild = 0;
39         dfs(-1, i);
40         artic[root] = rchild > 1; //Special case
41     }
42 }

```

Max Flow: Edmond Karp's $\mathcal{O}(VE^2)$

```

1 vector<vector<ll>> adjList;
2 vector<vector<ll>> adjMat;
3
4 void initialize(int n) {
5     adjList = decltype(adjList)(n);
6     adjMat = decltype(adjMat)(n, vector<ll>(n, 0));
7 }
8
9 map<int, int> p;
10 bool bfs(int source, int sink) {
11     queue<int> q;
12     vi visited(adjList.size(), 0);
13     q.push(source);
14     visited[source] = 1;
15     while (!q.empty()) {
16         int u = q.front();
17         q.pop();
18         if (u == sink)
19             return true;
20         for (auto v : adjList[u]) {
21             if (adjMat[u][v] > 0 && !visited[v]) {
22                 visited[v] = true;
23                 q.push(v);
24                 p[v] = u;
25             }
26         }
27     }
28     return false;
29 }
30 int max_flow(int source, int sink) {
31     ll max_flow = 0;
32     while (bfs(source, sink)) {
33         ll flow = inf;
34         for (int v = sink; v != source; v = p[v]) {
35             flow = min(flow, adjMat[p[v]][v]);
36         }
37         for (int v = sink; v != source; v = p[v]) {
38             adjMat[p[v]][v] -= flow; // Decrease
  ↳ capacity forward edge
39             adjMat[v][p[v]] += flow; // Increase
  ↳ capacity backward edge
40         }
41         max_flow += flow;
42     }
43     return max_flow;
44 }
45 void addEdgeUni(int orig, int dest, ll flow) {
46     adjList[orig].pb(dest);
47     adjMat[orig][dest] = flow;
48     adjList[dest].pb(orig); //Add edge for residual
  ↳ flow
49 }
50 void addEdgeBi(int orig, int dest, ll flow) {
51     adjList[orig].pb(dest);
52     adjList[dest].pb(orig);
53     adjMat[orig][dest] = flow;
54     adjMat[dest][orig] = flow;
55 }

```

Bellman Ford's

```

1 for(i = 0; i < n - 1; i++) { // Iterate n - 1 times.
2     for(auto e : edge) {
3         if(dist[e.fi.fi] != inf)
4             dist[e.fi.se] = min(dist[e.fi.se], dist[e.fi.fi] + e.se);
5     }
6 }

```

Floyd cycle detection

```

1 void floyd_detection() {
2     ll pslow = f(F_0), pfast = f(f(F_0)), iteration = 0;
3     while(pslow != pfast) pslow = f(pslow), pfast = f(f(pfast));
4     pslow = F_0;
5     while(pslow != pfast) pslow = f(pslow), pfast = f(pfast), iteration++;
6     cout << "In " << iteration << " coincide with value: " << pslow << endl;
7     pfast = f(pfast), iteration++;
8     while(pslow != pfast) pfast = f(pfast), iteration++;
9     cout << "In " << iteration << " coincide with value: " << pfast << endl;
10 }

```

Max Flow: Dinic's $\mathcal{O}(V^2E)$

```

1 //  $\mathcal{O}(V^2E)$  max flow algorithm. For bipartite
2   ↳ matching  $\mathcal{O}(\sqrt{V} * E)$ , always faster than
3   ↳ Edmond-Karp.
4 // Creates layer's graph with a BFS and then it
5   ↳ tries all possibles DFS, branching while the
6   ↳ path doesn't reach the sink
7 struct EdgeFlow {
8     ll u, v;
9     ll cap, flow = 0; //capacity and current flow
10    EdgeFlow(ll _u, ll _v, ll _cap) : u(_u), v(_v),
11      ↳ cap(_cap) {}
12 };
13
14 struct Dinic {
15     vector<EdgeFlow> edge; //keep the edges
16     vector<vll> graph; //graph[u] is the list of
17       ↳ their edges
18     ll n, n_edges = 0;
19     ll source, sink, inf_flow = inf;
20     vll lvl; //lvl of the node to the source
21     vll ptr; //ptr[u] is the next edge you have to
22       ↳ take in order to branch the DFS
23     queue<ll> q;
24
25     Dinic(ll _n, ll _source, ll _sink) : n(_n),
26       ↳ source(_source), sink(_sink) { //n nodes
27         graph.assign(_n, vll());
28     }
29
30     void add_edge(ll u, ll v, ll flow) { //u->v
31       ↳ with cost x
32         EdgeFlow uv(u, v, flow), vu(v, u, 0);
33         edge.pb(uv);
34         edge.pb(vu);
35         graph[u].pb(n_edges);
36         graph[v].pb(n_edges+1);
37         n_edges += 2;
38     }
39
40     bool BFS() {
41         ll u;
42         while(q.empty() == false) {
43             u = q.front(); q.pop();
44             for(auto el : graph[u]) {
45                 if(lvl[edge[el].v] != -1) {
46                     continue;
47                 }
48                 if(edge[el].cap - edge[el].flow <=
49                   ↳ 0) {
50                     continue;
51                 }
52                 lvl[edge[el].v] = lvl[edge[el].u] +
53                   ↳ 1;
54                 q.push(edge[el].v);
55             }
56         }
57         return lvl[sink] != -1;
58     }
59
60     ll dfs(ll u, ll min_flow) {
61         if(u == sink) return min_flow;
62         ll pushed, el;
63         for(; ptr[u] < (int)graph[u].size();
64           ↳ ptr[u]++) { //if you can pick ok, else
65           ↳ you crop that edge for the current bfs
66           ↳ layer
67             el = graph[u][ptr[u]];
68             if(lvl[edge[el].v] != lvl[edge[el].u] +
69               ↳ 1 || edge[el].cap - edge[el].flow
70               ↳ <= 0) {
71                 continue;
72             }
73             pushed = dfs(edge[el].v, min(min_flow,
74               ↳ edge[el].cap - edge[el].flow));
75             if(pushed > 0) {
76                 edge[el].flow += pushed;
77                 edge[el^1].flow -= pushed;
78                 return pushed;
79             }
80         }
81         return 0;
82     }
83
84     ll max_flow() {
85         ll flow = 0, pushed;
86         while(true) {
87             lvl.assign(n, -1);
88             lvl[source] = 0;
89             q.push(source);
90             if(!BFS()) {
91                 break;
92             }
93             ptr.assign(n, 0);
94             while(true) {
95                 pushed = dfs(source, inf_flow);
96                 if(!pushed) break;
97                 flow += pushed;
98             }
99         }
100        return flow;
101    }
102 };

```

Hungarian Algorithm

```

1 // The rows are jobs, the columns are workers
2 pair<ll, vl> hungarian(vector<vl> &matrix) {
3     int n = matrix.size(), m = matrix[0].size();
4     vl jobP(n), workerP(m + 1), matched(m + 1, -1);
5
6     vl dist(m + 1, inf);
7     vi from(m + 1, -1), seen(m + 1, 0);
8
9     for (int i = 0; i < n; ++i) {
10         int cWorker = m;
11         matched[cWorker] = i;
12         std::fill(all(dist), inf);
13         std::fill(all(from), -1);
14         std::fill(all(seen), false);
15
16         while (matched[cWorker] != -1) {
17             seen[cWorker] = true;
18             int i0 = matched[cWorker];
19             int nextWorker = -1;
20             ll delta = inf;
21
22             for (int worker = 0; worker < m;
23                  ↪ ++worker) {
24                 if (seen[worker])
25                     continue;
26                 ll candidateDistance =
27                     ↪ matrix[i0][worker];
28                 candidateDistance += -jobP[i0] -
29                     ↪ workerP[worker];
30
31                 if (candidateDistance < dist[worker]) {
32                     dist[worker] =
33                         ↪ candidateDistance;
34                     from[worker] = cWorker;
35                 }
36                 if (dist[worker] < delta) {
37                     delta = dist[worker];
38                     nextWorker = worker;
39                 }
40             }
41             for (int j = 0; j <= m; ++j) {
42                 if (seen[j]) {
43                     jobP[matched[j]] += delta;
44                     workerP[j] -= delta;
45                 } else {
46                     dist[j] -= delta;
47                 }
48             }
49             cWorker = nextWorker;
50         }
51         while (cWorker != m) {
52             int prevWorker = from[cWorker];
53             matched[cWorker] = matched[prevWorker];
54             cWorker = prevWorker;
55         }
56     }
57     ll ans = -workerP[m];
58     vl rowMatchesWith(n);
59     for (int j = 0; j < m; ++j) {
60         if (matched[j] != -1) {
61             rowMatchesWith[matched[j]] = j;
62         }
63     }
64     return {ans, std::move(rowMatchesWith)};
65 }

```

Floyd - Warshall: k->i->j
Kosaraju

```

1 vector<vi> adyList; // Graph
2 vector<int> visited; // Visited for DFS
3 vector<vi> sccs; // Contains the SCCs at the
4 ↪ end
5 void dfs(int nnode, vector<int> &v, vector<vi>
6 ↪ &adyList) {
7     if (visited[nnode]) {
8         return;
9     }
10    visited[nnode] = true;
11    for (auto a : adyList[nnode]) {
12        dfs(a, v, adyList);
13    }
14    v.push_back(nnode);
15 }
16 void Kosaraju(int n) {
17     visited = vi(n, 0);
18     stack<int> s = stack<int>();
19     sccs = vector<vi>();
20
21     vector<int> postorder;
22
23     for (int i = 0; i < n; ++i) {
24         dfs(i, postorder, adyList);
25     }
26     reverse(all(postorder));
27
28     vector<vi> rAdyList = vector<vi>(n, vi());
29     for (int i = 0; i < n; ++i) {
30         for (auto v : adyList[i]) {
31             rAdyList[v].push_back(i);
32         }
33     }
34
35     visited = vi(n, 0);
36     vi data;
37     for (auto a : postorder) {
38         if (!visited[a]) {
39             data = vi();
40             dfs(a, data, rAdyList);
41             if (!data.empty())
42                 sccs.pb(data);
43         }
44     }
45 }

```

LCA tree

```

1  const int MAX_N = 1e5 + 5;
2  const int MAX_LOG_N = 18;
3  int n;
4  vector<vi> graph; // Directed graph, allways
   ↳ reserve memory for it.
5  vector<vi> bigraph; // Undirected graph, reserve
   ↳ memory only if needed.
6
7  int level[MAX_N]; // level of the node rooted.
8  int parent[MAX_N][MAX_LOG_N]; // parent[i][j] is
   ↳ the parent 2^j of the node i.
9
10 vector<bool> visited_bigraph;
11 // root_graph(u, -1) roots the bigraph at node u.
12 void root_graph(int u, int p) {
13     if(p == -1) visited_bigraph.assign(n, false);
14     for(auto v : bigraph[u]) {
15         if(v == p) continue;
16         graph[u].pb(v);
17         root_graph(v, u);
18     }
19 }
20
21 // Calculate the level and parent 1. Don't call.
22 void dfs_level(int u, int p) {
23     parent[u][0] = p;
24     level[u] = level[p] + 1;
25     for(auto v : graph[u]) {
26         if(v == p) continue;
27         dfs_level(v, u);
28     }
29 }
30 // Builds the LCA.
31
32 void build_lca(int root) {
33     int i, j;
34     level[root] = -1;
35     dfs_level(root, root); // The parent of the
   ↳ root is itself.
36     for(j = 1; j < MAX_LOG_N; j++) {
37         for(i = 0; i < MAX_N; i++) {
38             parent[i][j] = parent[parent[i][j -
   ↳ 1]][j - 1];
39         }
40     }
41     // Calculates the LCA(u, v) in O(log n).
42     int lca(int u, int v) {
43         if(level[u] > level[v]) swap(u, v);
44         int i, d = level[v] - level[u];
45         for(i = MAX_LOG_N - 1; i >= 0; i--) {
46             if(is_set(d, i)) v = parent[v][i];
47         }
48         if(u == v) return u;
49         for(i = MAX_LOG_N - 1; i >= 0; i--) {
50             if(parent[u][i] != parent[v][i])
51                 u = parent[u][i], v = parent[v][i];
52         }
53         return parent[u][0];
54     }
55     // Calculates the distance(u, v) in a tree in O(log
   ↳ n).
56     int dist(int u, int v) {
57         return level[u] + level[v] - 2 * level[lca(u,
   ↳ v)];
58     }

```

Mathematics

Binary operations

```

1  ll elevate(ll a, ll b) { // b >= 0.
2      ll ans = 1;
3      while(b) {
4          if(b & 1) ans = ans * a % mod;
5          b >>= 1;
6          a = a * a % mod;
7      }
8      return ans;
9  }
10 // a^(mod - 1) = 1, Euler.
11 ll inv(ll a) {
12     return elevate(((a%mod) + mod)%mod, mod - 2);
13 }
14
15 ll mul(ll a, ll b) {
16     ll ans = 0, neg = (a < 0) ^ (b < 0);
17     a = abs(a); b = abs(b);
18     while(b) {
19         if(b & 1) ans = (ans + a) % mod;
20         b >>= 1;
21         a = (a + a) % mod;
22     }
23     if(neg) return -ans;
24     return ans;
25 }

```

$$\text{Catalan numbers: } C_n = \frac{1}{n+1} \binom{2n}{n}$$

Combinatoric numbers

```

1  const int MAX_C = 1+66; // 66 is the for long
   ↳ long, C(66, x)
2  ll Comb[MAX_C][MAX_C];
3
4  void calc() {
5      int i, j;
6      for(i = 0; i < MAX_C; i++) {
7          Comb[i][0] = 1;
8          Comb[0][i] = 1;
9      }
10     for(i = 1; i < MAX_C; i++) {
11         for(j = 1; j < MAX_C; j++) {
12             if(i+j >= MAX_C) continue;
13             Comb[i][j] = Comb[i-1][j] +
   ↳ Comb[i][j-1];
14         }
15     }
16 }
17 ll C(ll i, ll j) {
18     return Comb[i-j][j];
19 }

```

Chinese Remainder

```

1  const ll MAX = 10;
2  ll a[MAX], p[MAX], n;
3  // Given  $n$   $x \equiv a[i] \pmod{p[i]}$ , find  $x$ ,
4  // or -1 if it doesn't exist.
5  // Let  $q[i] = (\prod_{i=0}^{n-1} p[j]) / p[i]$ .
6  //  $x$  will be  $\sum_{i=0}^{n-1} a[i] * q[i]$ 
7  //  $*inv(q[i], \pmod{p[i]})$ 
8  ll chinese_remainder() {
9      ll i, j, g, ans = 0, inv1, inv2;
10     mod = 1;
11     for(i = 0; i < n; i++) {
12         // If the  $p[i]$  are not coprimes, do them
13         //  $\hookrightarrow$  coprimes.
14         a[i] %= p[i]; a[i] += p[i]; a[i] %= p[i];
15         for(j = 0; j < i; j++) {
16             g = __gcd(p[i], p[j]);
17             if((a[i]%g + g)%g != (a[j]%g + g)%g)
18                 return -1;
19         }
20         // Delete the repeated factor at the
21         //  $\hookrightarrow$  correct side.
22         if (__gcd(p[i]/g, p[j]) == 1) {p[i] /=
23             g; a[i] %= p[i];}
24         else {p[j] /= g; a[j] %= p[j];}
25     }
26     // If you have a supermod, take  $P = \min(P,$ 
27     //  $\hookrightarrow$  supermod);
28     for(i = 0; i < n; i++) {
29         mod *= p[i];
30     }
31     for(i = 0; i < n; i++) {
32         gcdEx(mod/p[i], p[i], &inv1, &inv2);
33         ans += mul(a[i], mul(mod/p[i], inv1));
34         ans %= mod;
35     }
36     return (ans%mod + mod) % mod;
37 }

```

Euclides

```

1  ll gcdEx(ll a, ll b, ll *x1, ll *y1) {
2      if(a == 0) {
3          *x1 = 0;
4          *y1 = 1;
5          return b;
6      }
7      ll x0, y0, g;
8      g = gcdEx(b%a, a, &x0, &y0);
9      *x1 = y0 - (b/a)*x0;
10     *y1 = x0;
11     return g;
12 }

```

Hash Set

```

1  const int MAX = 2*1e5+5;
2  ll val[MAX]; // For random numbers and not index
3  //  $\hookrightarrow$  use  $f$  with random xor.
4  void ini() { // CALL ME ONCE.
5      srand(time(0));
6      for(int i = 0; i < MAX; i++) val[i] = rand();
7  }
8  // Hash_set contains a set of indices  $[0..MAX-1]$ 
9  //  $\hookrightarrow$  with duplicates.
10 //  $a[i] = \sum_x \{val_x\} \% \pmod{p[i]}$ .
11 class Hash_set {
12 public:
13     vll p = {1237273, 1806803, 3279209}; // Prime
14     //  $\hookrightarrow$  numbers.
15     vll a = {0, 0, 0};
16     int n = 3; //  $n = p.size()$ ;
17 }
18 void insert(int x) { // Insert index  $x$ .
19     for(int i = 0; i < n; i++) a[i] = (a[i] +
20         //  $\hookrightarrow$  val[x]) % p[i];
21     }
22 // Insert all the elements of  $hs$ .
23 void insert (Hash_set hs) {
24     for(int i = 0; i < n; i++) a[i] = (a[i] +
25         //  $\hookrightarrow$  hs.a[i]) % p[i];
26     }
27 bool operator == (Hash_set hs) {
28     for(int i = 0; i < n; i++) if(a[i] !=
29         //  $\hookrightarrow$  hs.a[i]) return false;
30     return true;
31 }
32 };

```

Hash of pairs

```

1  // Use unordered_set<pii, pair_hash> us or
2  //  $\hookrightarrow$  unordered_map<pii, int, pair_hash> um;
3  struct pair_hash
4  {
5      template <class T1, class T2>
6      size_t operator () (pair<T1, T2> const &pair)
7          //  $\hookrightarrow$  const
8      {
9          size_t h1 = hash<T1>()(pair.first);
10         size_t h2 = hash<T2>()(pair.second);
11         return (h1 ^ 0b11001001011001101) +
12             //  $\hookrightarrow$  (0b011001010011100111 ^ h2);
13     }
14 };

```

Linear Sieve

```

1  const int MAX_PRIME = 1e6+5;
2  bool num[MAX_PRIME]; // If num[i] = false  $\Rightarrow$   $i$  is prime.
3  int num_div[MAX_PRIME]; // Number of divisors of  $i$ .
4  int min_div[MAX_PRIME]; // The smallest prime that divide  $i$ .
5  vector<int> prime;
6

```



```

7 void linear_sieve(){
8     int i, j, prime_size = 0;
9     min_div[1] = 1;
10    for(i = 2; i < MAX_PRIME; ++i){
11        if(num[i] == false) {prime.push_back(i); ++prime_size; num_div[i] = 1; min_div[i] = i;}
12
13        for(j = 0; j < prime_size && i * prime[j] < MAX_PRIME; ++j){
14            num[i * prime[j]] = true;
15            num_div[i * prime[j]] = num_div[i] + 1;
16            min_div[i * prime[j]] = min(min_div[i], prime[j]);
17            if(i % prime[j] == 0) break;
18        }
19    }
20 }
21 bool is_prime(ll n) {
22     for(auto el : prime) {
23         if(n == el) return true;
24         if(n%el == 0) return false;
25     }
26     return true;
27 }
28 vll fact, nfact; // The factors of n and their
    ↪ exponent.
29 void factorize(int n) { // Up to
    ↪ MAX_PRIME*MAX_PRIME.
30     ll cont, prev_p;
31     fact.clear(); nfact.clear();
32     for(auto p : prime) {
33         if(n < MAX_PRIME) break;
34         if(n%p == 0) {
35             fact.pb(p);
36             cont = 0;
37             while(n%p == 0) n /= p, cont++;
38             nfact.pb(cont);
39         }
40     }
41     if(n >= MAX_PRIME) {
42         fact.pb(n);
43         nfact.pb(1);
44         return;
45     }
46     while(n != 1) { // When n < MAX_PRIME,
    ↪ factorization in almost O(1).
47         prev_p = min_div[n];
48         cont = 0;
49         while(n%prev_p == 0) n /= prev_p, cont++;
50         fact.pb(prev_p);
51         nfact.pb(cont);
52     }
53 }

```

Suffix Array

```

1 class SuffixArray {
2     public:
3     int n;
4     string s;
5     vi p; // p[i] is the position in the order
    ↪ array of the ith suffix (s[i..n-1]).
6     vi c; // c[i] is the equivalence class of the
    ↪ ith suffix. When build, c[p[i]] = i,
    ↪ inverse.
7     // dont use lcp[0] = 0.
8     vi lcp; // lcp[i] is the longest common prefix
    ↪ in s[p[i-1]..n-1] and s[p[i]..n-1].
9     // To get lcp(s[i..n-1], s[j..n-1]) is
    ↪ min(lcp[c[i]+1], lcp[c[j]]) (use SegTree).
10    void radix_sort(vector<pair<pii, int>> &v) { //
    ↪ O(n).
11        vector<pair<pii, int>> v2(n);
12        vi freq(n, 0); // first frequency and then
    ↪ the index of the next item.
13        int i, sum = 0, temp;
14        for(i = 0; i < n; i++) freq[v[i].fi.se]++;
    ↪ // Sort by second component.
15        for(i = 0; i < n; i++) {temp = freq[i];
    ↪ freq[i] = sum; sum += temp;}
16        for(i = 0; i < n; i++)
    ↪ {v2[freq[v[i].fi.se]] = v[i];
    ↪ freq[v[i].fi.se]++;}
17        freq.assign(n, 0); sum = 0;
18        for(i = 0; i < n; i++) freq[v2[i].fi.fi]++;
    ↪ // Sort by first component.
19        for(i = 0; i < n; i++) {temp = freq[i];
    ↪ freq[i] = sum; sum += temp;}
20
21        for(i = 0; i < n; i++)
    ↪ {v[freq[v2[i].fi.fi]] = v2[i];
    ↪ freq[v2[i].fi.fi]++;}
22    }
23    SuffixArray() = default;
24    SuffixArray(string &s) {
25        s = _s;
26        s += "$"; // smaller char to end the
    ↪ string.
27        n = s.size();
28        int i, k;
29        p.assign(n, 0);
30        c.assign(n, 0);
31        vector<pii> v1(n); // temporal vector to
    ↪ sort.
32        for(i = 0; i < n; i++) v1[i] = mp(s[i], i);
33        sort(v1.begin(), v1.end());
34        for(i = 0; i < n; i++) p[i] = v1[i].se;
35        c[p[0]] = 0;
36        for(i = 1; i < n; i++) {
37            if(v1[i].fi == v1[i-1].fi) c[p[i]] =
    ↪ c[p[i-1]];
38            else c[p[i]] = c[p[i-1]] + 1;
39        }
40        k = 0; // in k+1 iterations sort strings of
    ↪ length 2^(k+1).
41        while(c[p[n-1]] != n-1) { // At most
    ↪ ceil(log2(n)).
42            vector<pair<pii, int>> v2(n); //
    ↪ temporal vector to sort.
43            for(i = 0; i < n; i++) v2[i] =
    ↪ mp(mp(c[i], c[(i + (1 << k)) % n]),
    ↪ i);

```



```

43         radix_sort(v2);
44         for(i = 0; i < n; i++) p[i] = v2[i].se;
45         c[p[0]] = 0;
46         for(i = 1; i < n; i++) {
47             if(v2[i].fi == v2[i - 1].fi)
48                 c[p[i]] = c[p[i - 1]];
49             else c[p[i]] = c[p[i - 1]] + 1;
50         }
51         k++;
52     }
53     void show_suffixes() { // IMPORTANT use this to
54         debug.
55         for(int i = 0; i < n; i++) cout << i << " "
56         << p[i] << " " << s.substr(p[i]) <<
57         endl;
58         if(!lcp.empty()) cout << "LCP: " << lcp <<
59         endl;
60     }
61     // cmp s with t. return -1 if s < t, 1 if s >
62     // t, 0 if s == t.
63     int cmp_string(int pos, string &t) {
64         for(int i = p[pos], j = 0; j < (int)
65         t.size(); i++, j++) {
66             if(s[i] < t[j]) return -1; // i < n
67             // because s[n-1] = '$'.
68             if(s[i] > t[j]) return 1;
69         }
70         return 0;
71     }
72     // Count the number of times t appears in s.
73     int count_substring(string &t) {
74         int l = -1, r = n, mid, L, R;
75         while(l + 1 < r) { //
76             // -1,...,-1=L,0,...,0,1=R...1.
77             mid = (l + r) / 2;
78             if(cmp_string(mid, t) < 0) l = mid;
79             else r = mid;
80         }
81         L = l;
82         l = -1; r = n;
83         while(l + 1 < r) {
84             mid = (l + r) / 2;
85             if(cmp_string(mid, t) <= 0) l = mid;
86             else r = mid;
87         }
88         R = r;
89         return R - L - 1;
90     }
91     // 0(n) build. At most 2n lcp++ and n lcp--;
92     void build_lcp() {
93         lcp.assign(n, 0);
94         for(int i = 0; i < n - 1; i++) {
95             if(i > 0) lcp[c[i]] = max(lcp[c[i - 1]]
96             <- - 1, 0);
97             while(s[i + lcp[c[i]]] == s[p[c[i] - 1]
98             <- + lcp[c[i]]) lcp[c[i]]++;
99         }
100     }
101     ll number_substrings() {
102         ll ans = 0, i;
103         for(i = 1; i < n; i++) {
104             ans += n - p[i-1] - lcp[i]; // Length
105             // of the suffix - lcp with the next
106             // suffix.
107         }
108         ans += n - p[n - 1]; // Plus the last
109         // suffix.
110         return ans - n; // Remove the '$' symbol on
111         // n substrings.
112     }
113     string LCS(string s, string &t) {
114         int mx = 0, mxi = 0, i, n2 = t.length();
115         string ans = "";
116         s += "@" + t; // Concatenate with a special
117         // char.
118         SuffixArray sa(s);
119         sa.build_lcp();
120         for(i = 1; i < sa.n; i++) {
121             // Suffix of s and before suffix of t.
122             if(sa.n - sa.p[i] > n2 + 2 && sa.n -
123             sa.p[i-1] <= n2 + 1) {
124                 if(sa.lcp[i] > mx) mx = sa.lcp[i], mxi
125                 = i;
126             }
127             // Suffix of t and before suffix of s.
128             if(sa.n - sa.p[i] <= n2 + 1 && sa.n -
129             sa.p[i-1] > n2 + 2) {
130                 if(sa.lcp[i] > mx) mx = sa.lcp[i], mxi
131                 = i;
132             }
133         }
134         return sa.s.substr(sa.p[mxi], mx);
135     }
136 }

```

BIT Fenwick tree

```

1  template<typename T>
2  class BIT{
3      vector<T> bit;
4      int n;
5  public:
6      BIT(int _n) {
7          n = _n;
8          bit.assign(n+1, 0);
9      }
10     BIT(vector<T> v) {
11         n = v.size();
12         bit.assign(n+1, 0);
13         for(int i = 0; i < n; i++) update(i, v[i]);
14     }
15     // Point update.
16     void update(int i, T dx) {
17         for(i++; i < n+1; i += LSB(i)) bit[i] +=
18         dx;
19     }
20     T query(int r) { // query [0, r].
21         T ans = 0;
22         for(r++; r > 0; r -= LSB(r)) ans += bit[r];
23         return ans;
24     }
25     T query(int l, int r) { // query [l, r].
26         return query(r) - query(l-1);
27     }
28     // k-th smallest element inserted.
29     int k_element(ll k) { // k > 0 (1-indexed).
30         int l = 0, r = n+1, mid;
31         if(query(0) >= k) return 0;
32         while(l + 1 < r) {
33             mid = (l + r)/2;
34             if(query(mid) >= k) r = mid;
35             else l = mid;
36         }
37         return r;
38     }
39 };

```

Strings: KMP

```

1  template <typename T>
2  vi prefixFun(const T &s, int n) {
3      vi res(n);
4      for (int i = 1; i < n; ++i) {
5          int j = res[i - 1];
6          while (j > 0 && s[i] != s[j]) {
7              j = res[j - 1];
8          }
9          res[i] = j + (s[i] == s[j]);
10     }
11     return res;
12 }
13
14 template <typename T>
15 int kmpSearch(const T &text, int n,
16               const T &pattern, int m,
17               const vi &patternPre) {
18
19     int count = 0;
20     int j = 0;
21     for (int i = 0; i < n; ++i) {
22         while (j > 0 && text[i] != pattern[j]) {
23             j = max(0, patternPre[j] - 1);
24         }
25         j += (text[i] == pattern[j]);
26         if (j == m) {
27             count++;
28             j = patternPre[j - 1];
29         }
30     }
31     return count;
32 }

```

Longest Palindromic Substring

```

1  // LPS Longest Palindromic Substring, O(n).
2  void Manacher(string &str) {
3      char ch = '#'; // '#' a char not contained in str.
4      string s(1, ch), ans;
5      for(auto c : str) {s += c; s += ch;}
6      int i, n = s.length(), c = 0, r = 0;
7      vi lps(n, 0);
8      for(i = 1; i < n; i++) {
9          // lps[i] >= it's mirror, but falling in the interval [L..R]. L = c - (R - c).
10         if(i < r) lps[i] = min(r - i, lps[c - (i - c)]);
11         // Try to increase.
12         while(i-lps[i]-1 >= 0 && i+lps[i]+1 < n && s[i-lps[i]-1] == s[i+lps[i]+1]) lps[i]++;
13         // Update the interval [L..R].
14         if(i + lps[i] > r) c = i, r = i + lps[i];
15     }
16     // Get the longest palindrome in ans.
17     int pos = max_element(lps.begin(), lps.end()) - lps.begin();
18     for(i = pos - lps[pos]; i <= pos + lps[pos]; i++) {
19         if(s[i] != ch) ans += s[i];
20     }
21     //cout << ans.size() << "\n";
22 }

```

Z-algorithm

```

1  // Search the occurrences of t (pattern to search)
2  // in s (the text).
3  // O(n + m). It increases R at most 2n times
4  // and decreases at most n times.
5  // z[i] is the longest string s[i..i+z[i]-1]
6  // that is a prefix = s[0..z[i]-1].
7  void z_algorithm(string &s, string &t) {
8      s = t + "$" + s;
9      // "$" is a char not present in s nor t.
10     int n = s.length(), m = t.length(), i;
11     int L = 0, R = 0;
12     vi z(n, 0);
13     // s[L..R] = s[0..R-L], [L, R]
14     // is the current window.
15     for(i = 1; i < n; i++) {
16         if(i > R) { // Old window, recalculate.
17             L = R = i;
18             while(R < n && s[R] == s[R-L]) R++;
19             R--;
20             z[i] = R - L + 1;
21         } else {
22             // z[i] will fall in the window.
23             if(z[i-L] < R - i) z[i] = z[i-L];
24             // z[i] can fall outside the window,
25             // try to increase the window.
26             else {
27                 L = i;
28                 while(R < n && s[R] == s[R-L]) R++;
29                 R--;
30                 z[i] = R - L + 1;
31             }
32         }
33         if(z[i] == m) { // Match found.
34             //echo("Pattern found at: ", i-m-1);
35         }
36     }
37 }

```