


Chapter

6

*Lists Plus*



*Third Edition*

# C<sup>++</sup> *Plus* Data Structures

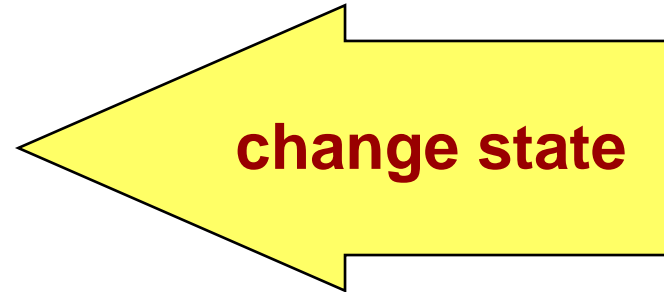
*Nell Dale*



# ADT Sorted List Operations

## Transformers

MakeEmpty  
InsertItem  
DeleteItem



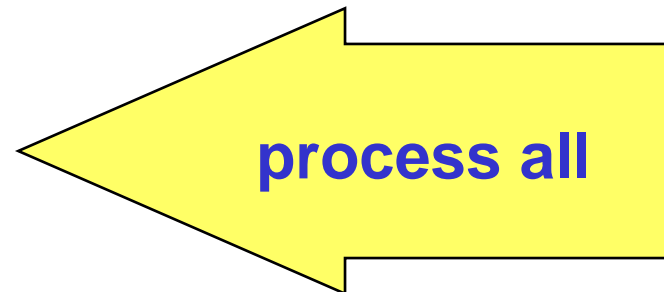
## Observers

IsFull  
LengthIs  
RetrieveItem



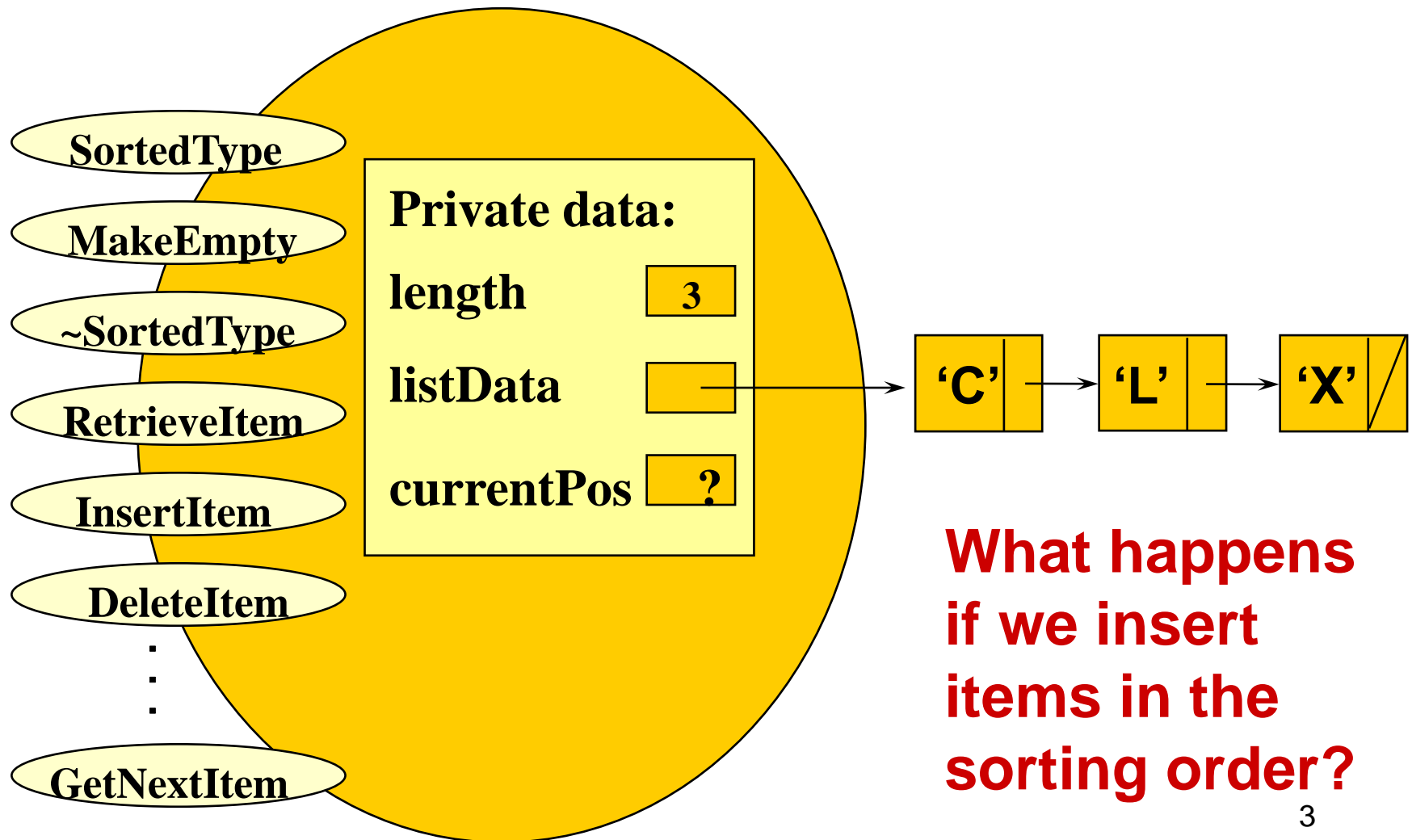
## Iterators

ResetList  
GetNextItem





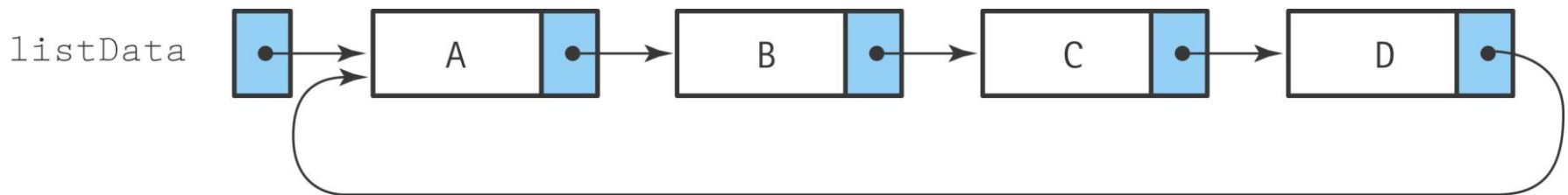
# class SortedType<char>



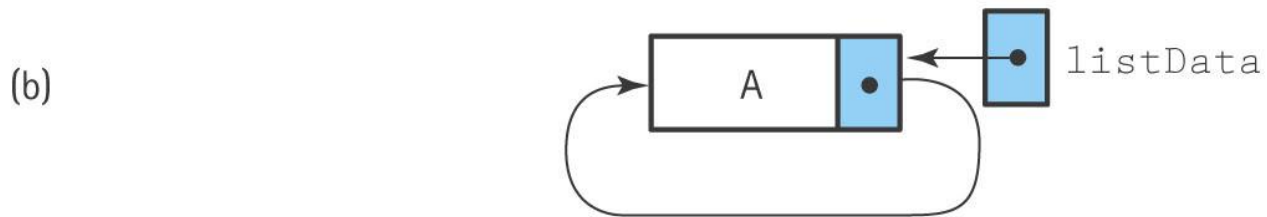
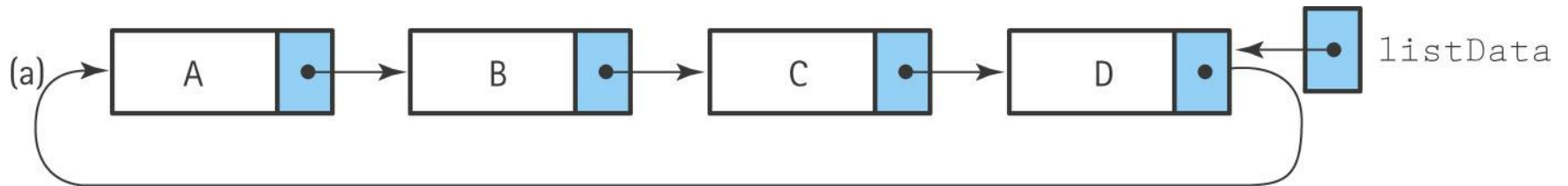


# What is a Circular Linked List?

A circular linked list is a list in which **every node has a successor**; the “last” element is succeeded by the “first” element.



# External Pointer to the Last Node



# What is a Doubly Linked List?

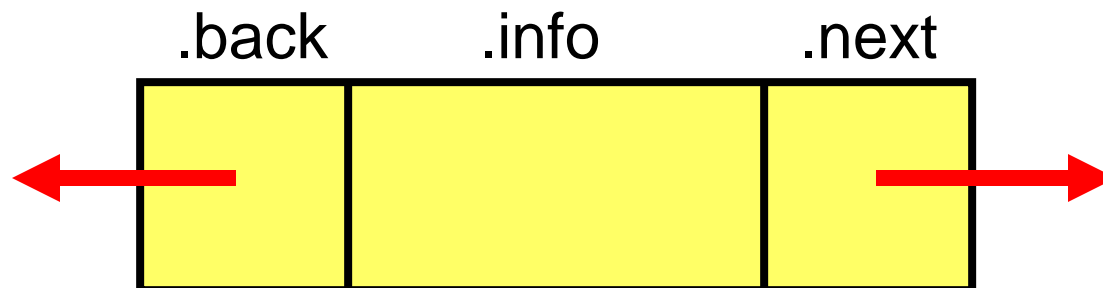
A doubly linked list is a list in which **each node is linked to both its successor and its predecessor.**



# Node data

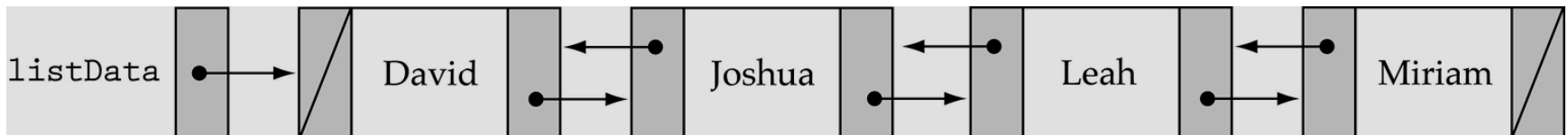
info: the user's data

next, back: the address of the next and previous node in the list



# Node data (cont.)

```
template<class ItemType>
struct NodeType {
    ItemType info;
    NodeType<ItemType>* next;
    NodeType<ItemType>* back;
};
```

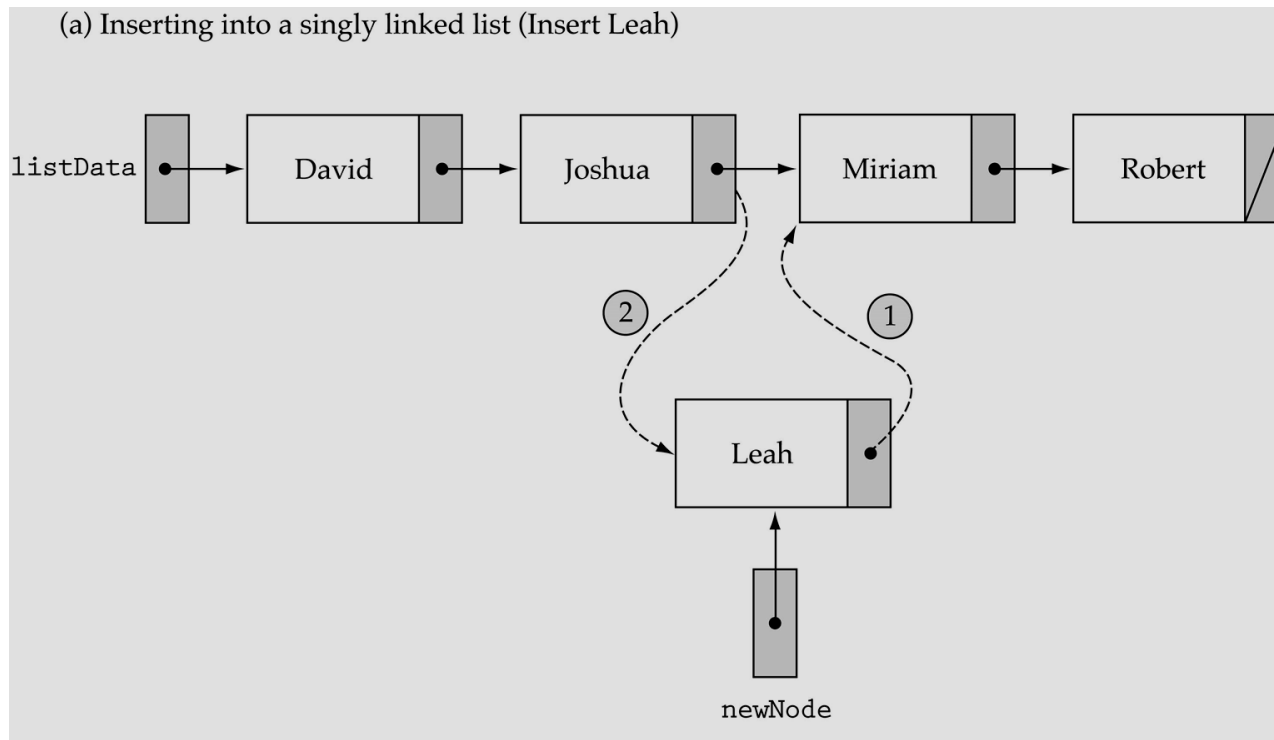




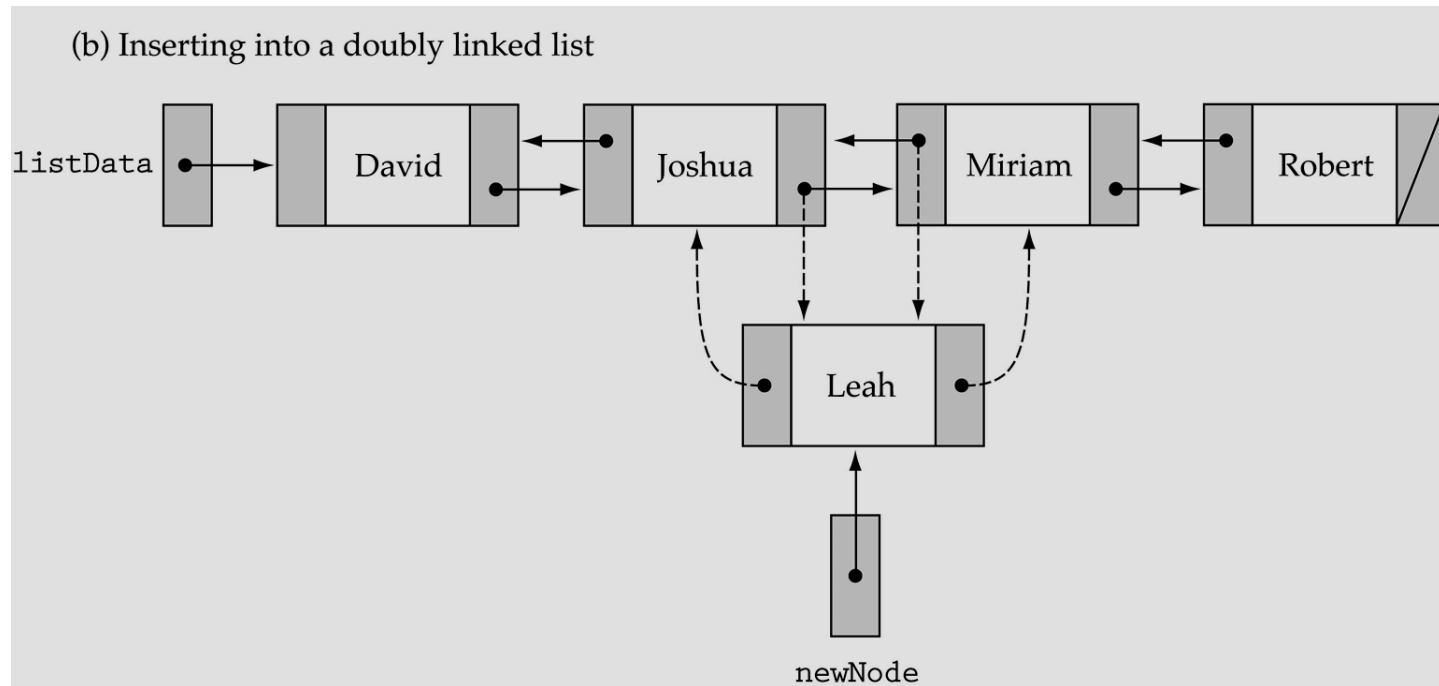


# Finding a List Item

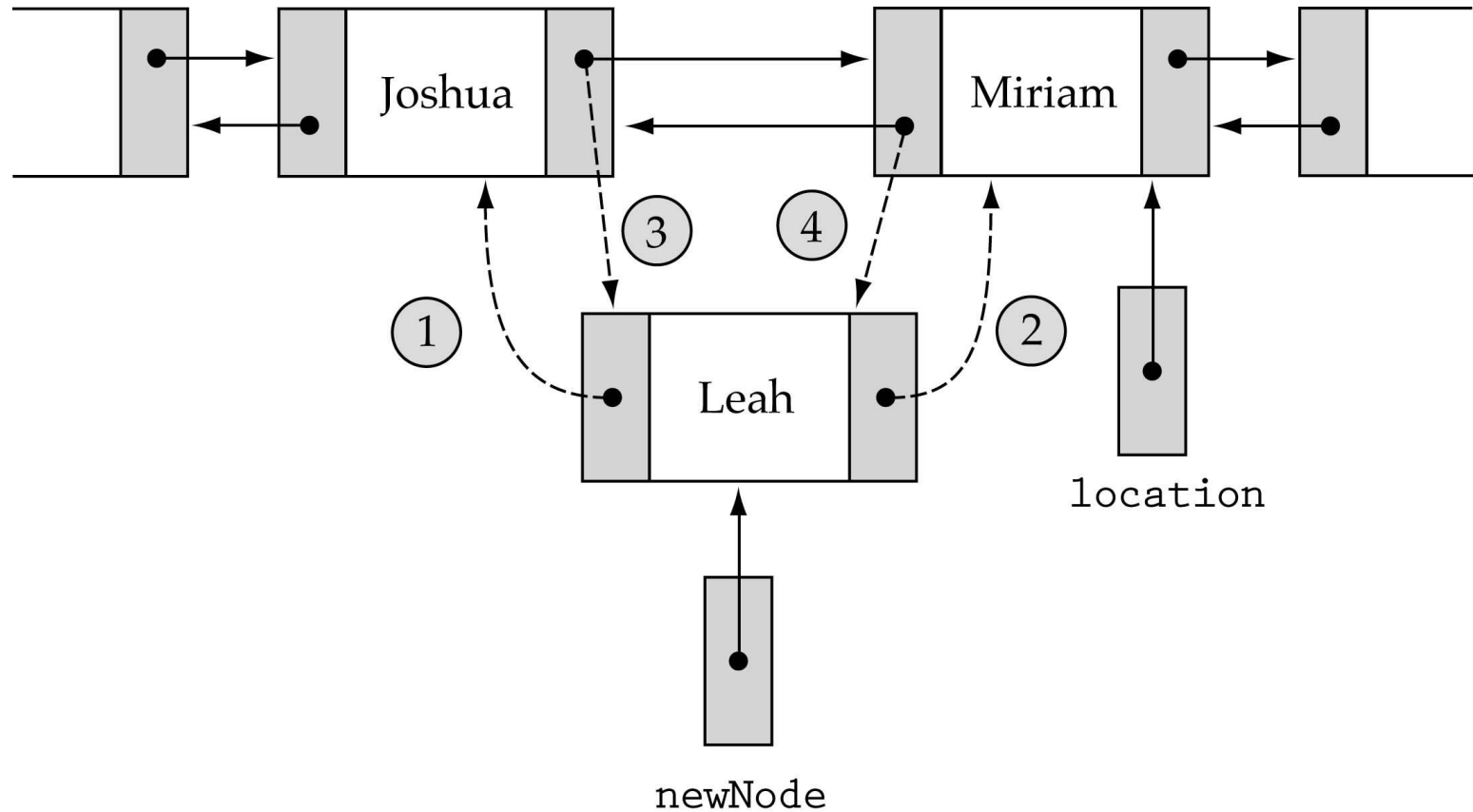
We no longer need to use *prevLocation* (we can get the predecessor of a node using its *back* member)



# Finding a List Item (cont.)



# Inserting into a Doubly Linked List



1.  $\text{newNode} \rightarrow \text{back} = \text{location} \rightarrow \text{back}$ ;
2.  $\text{newNode} \rightarrow \text{next} = \text{location}$
3.  $\text{location} \rightarrow \text{back} \rightarrow \text{next} = \text{newNode}$ ;
4.  $\text{location} \rightarrow \text{back} = \text{newNode}$ ;



## FindItem(listData, item, location, found)

RetrieveItem, InsertItem, and DeleteItem all require a search !

Write a general non-member function

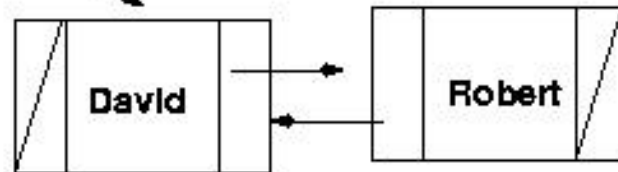
FindItem that takes *item* as a parameter and returns location and found.

InsertItem and DeleteItem need *location* (ignore *found*)

RetrieveItem needs *found* (ignores *location*)

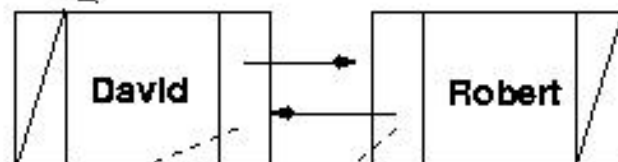


l1stData



Retrieve: **Robert**

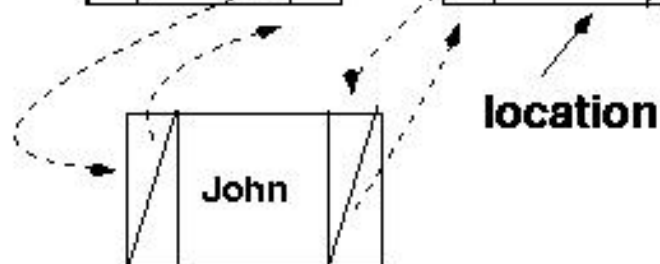
l1stData



location

Insert: **John**

(similar for Delete)



location



# Finding a List Item (cont.)

```
template<class ItemType>
void FindItem(NodeType<ItemType>* listData, ItemType item,
    NodeType<ItemType>* &location, bool &found)
{
    // precondition: list is not empty

    bool moreToSearch = true;

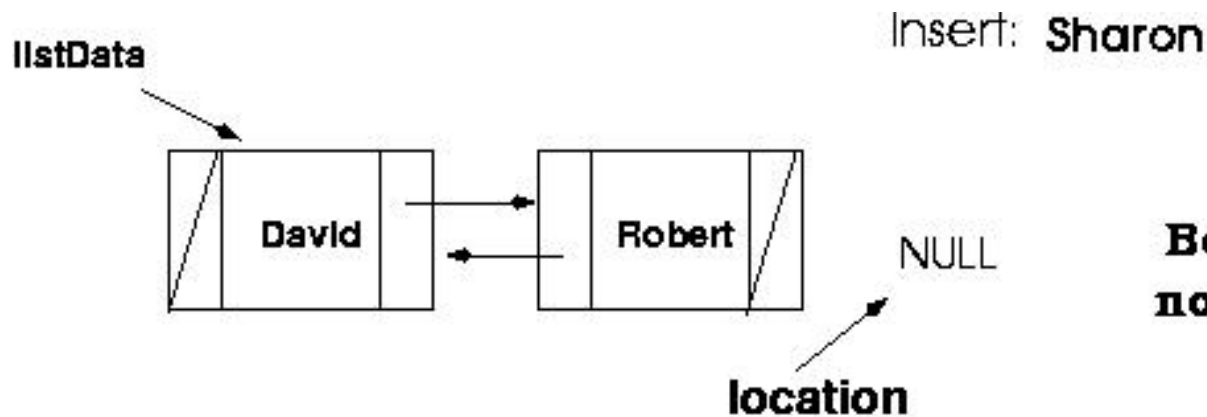
    location = listData;
    found = false;

    while( moreToSearch && !found) {

        if(item < location->info)
            moreToSearch = false;
        else if(item == location->info)
            found = true;

        else {
            location = location->next;
            moreToSearch =
                (location != NULL);
        }
    }
}
```

# Problem: Finding a List Item



**Book's solution will  
not work in this case !**

**fix the problem:**

```
if(location->next != NULL)
    location = location->next;
```



# Correction: Finding a List Item

```
template<class ItemType>
void FindItem(NodeType<ItemType>* listData, ItemType item,
    NodeType<ItemType>* &location, bool &found)
{
    // precondition: list is not empty

    bool moreToSearch = true;

    location = listData;
    found = false;

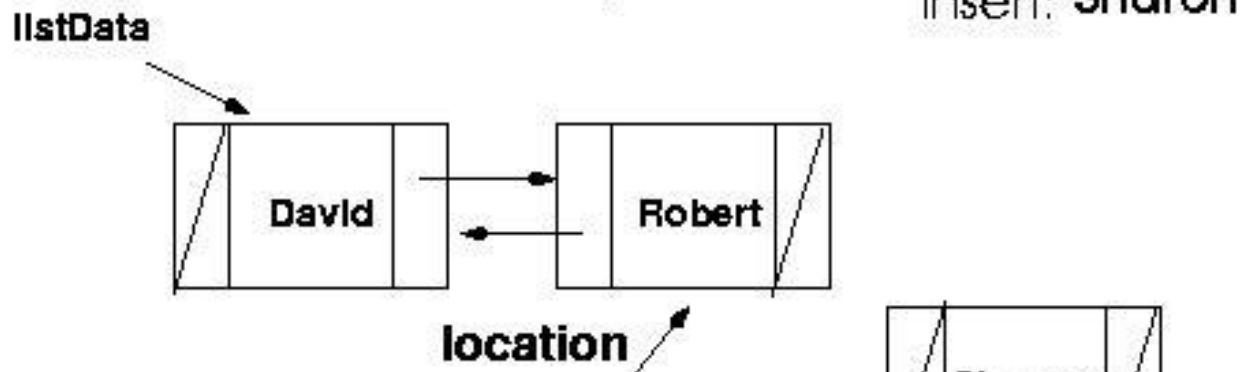
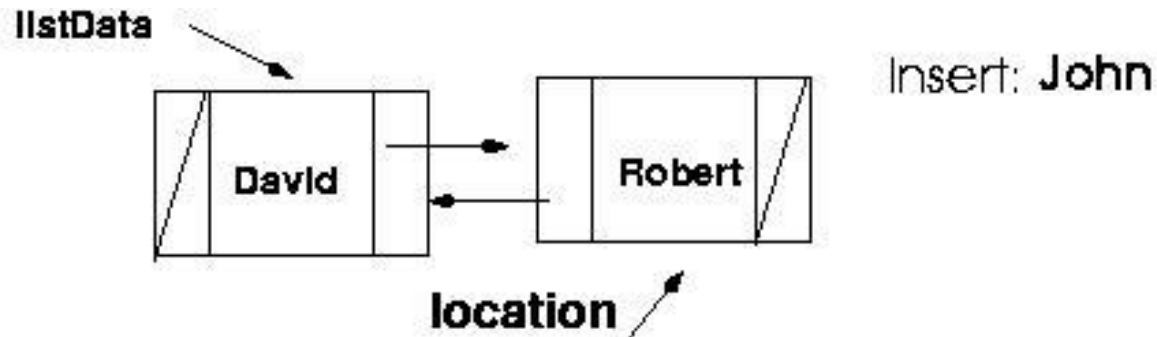
    while( moreToSearch && !found) {

        if(item < location->info)
            moreToSearch = false;
        else if(item == location->info)
            found = true;

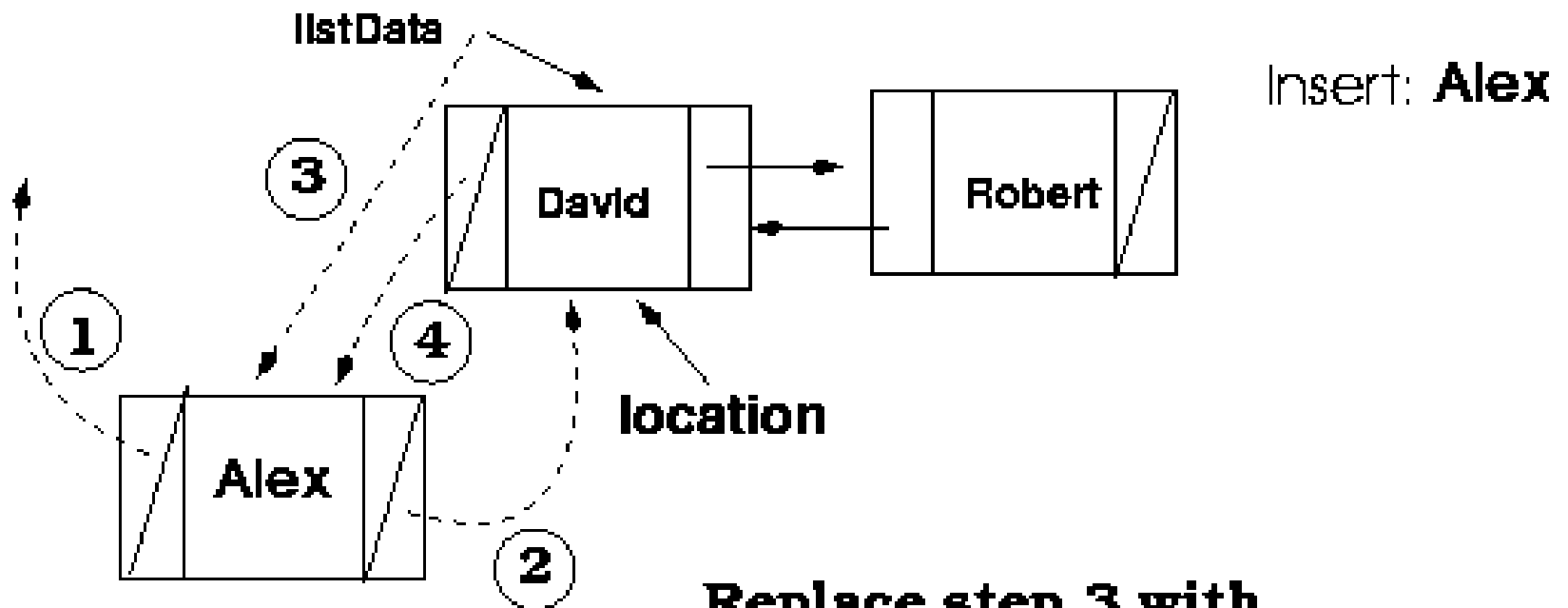
        else {
            if(location->next == NULL)
                moreToSearch = false;
            else
                location = location->next;
        }
    }
}
```



# How can we distinguish between the following two cases?



## Special case: inserting in the beginning



**Replace step 3 with**

`listData = newNode;`

**(See Page 11)**



# Inserting into a Doubly Linked List

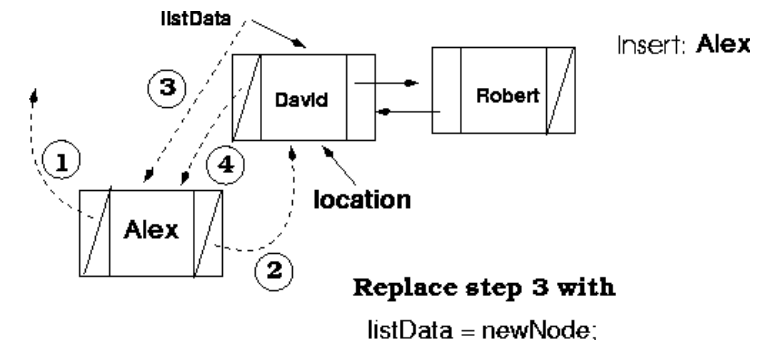
```
template<class ItemType>
void SortedType<ItemType>::InsertItem(ItemType item)
{
```

```
    NodeType<ItemType>* newNode;
    NodeType<ItemType>* location;
    bool found;
```

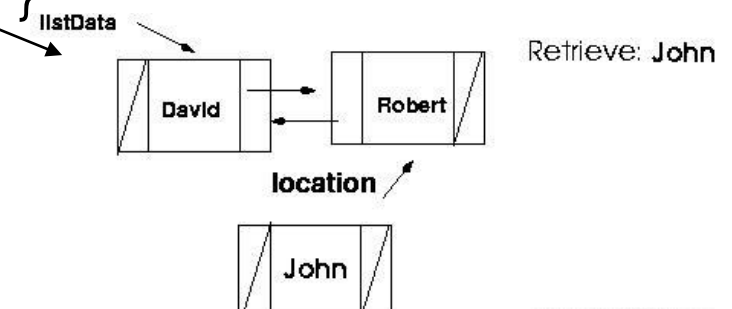
```
    newNode = new NodeType<ItemType>;
    newNode->info = item;
    if (listData != NULL) {
```

```
        FindItem(listData, item, location, found);
```

```
        if (location->info > item) {
            newNode->back = location->back; (1)
            newNode->next = location; (2)
            if (location != listData) // special case
                (location->back)->next = newNode; (3)
```

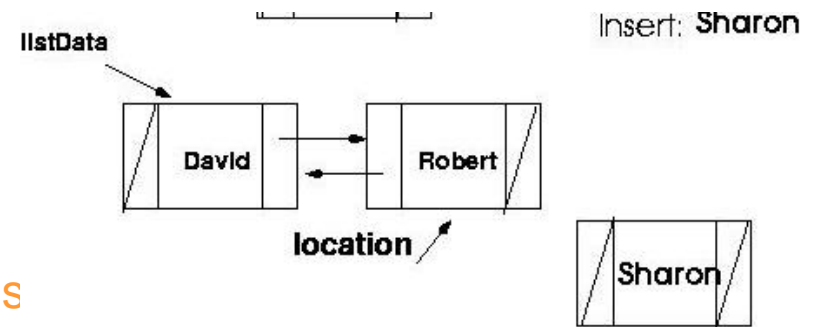


```
    else
        listData = newNode; (3)
        location->back = newNode; (4)
```

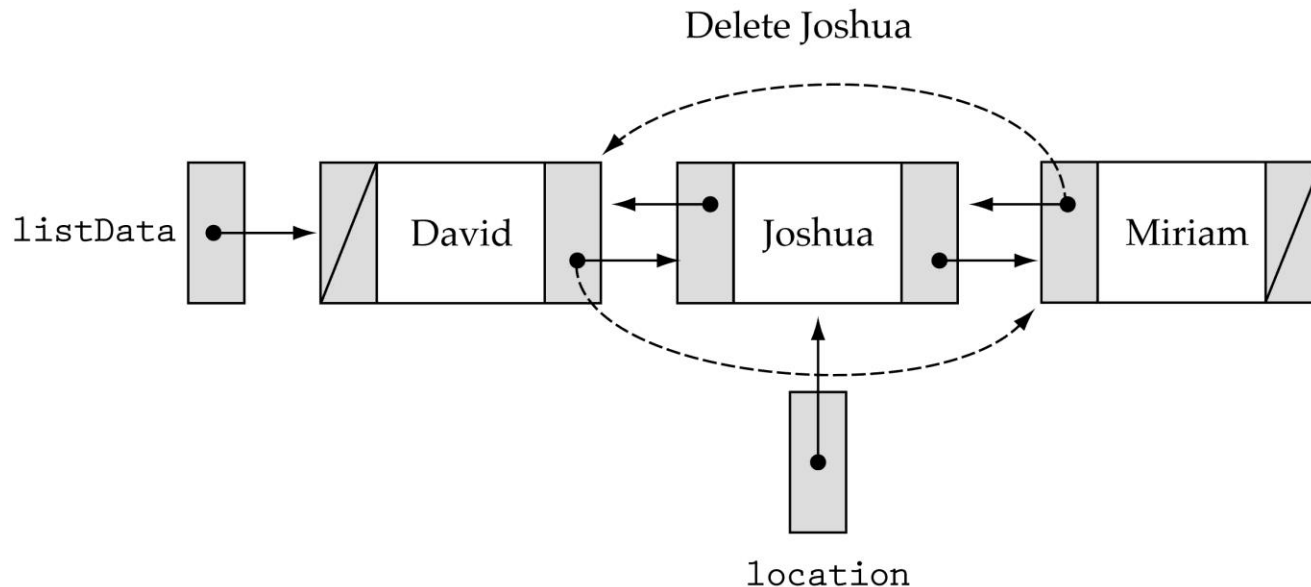


# Inserting into a Doubly Linked List (cont.)

```
else {                                // insert at the end
    newNode->back = location;
    location->next = newNode;
    newNode->next = NULL;
}
}
else {                                // insert into an empty list
    listData = newNode;
    newNode->next = NULL;
    newNode->back = NULL;
}
length++;
}
```



# Deleting from a Doubly Linked List



Be careful about the end cases!!



# Headers and Trailers

Special cases arise when we are dealing with the first or last nodes

How can we simplify the implementation?

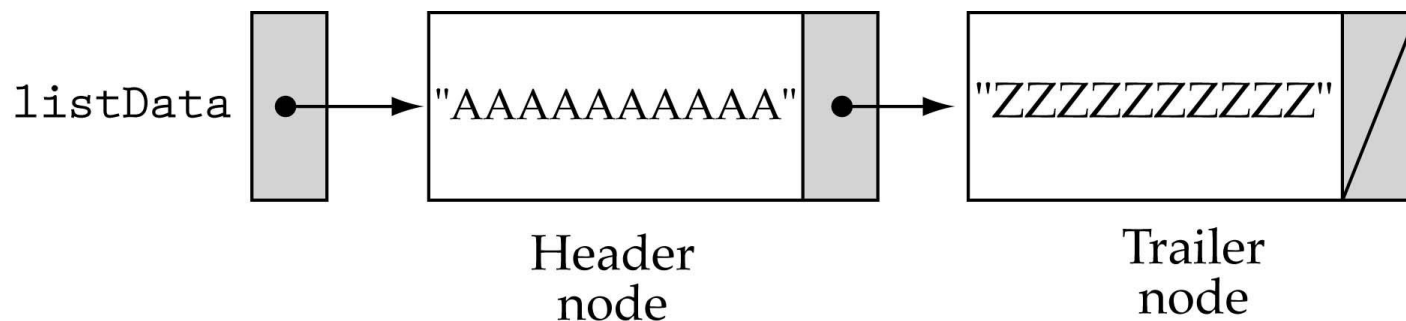
Idea: make sure that we never insert or delete the ends of the list

How? Set up dummy nodes with values outside of the range of possible values

# Headers and Trailers (cont.)

*Header Node*: contains a value smaller than any possible list element

*Trailer Node*: contains a value larger than any possible list element





# A linked list as an array of records

What are the advantages of using linked lists?

- (1) Dynamic memory allocation
- (2) Efficient insertion-deletion (for sorted lists)

Can we implement a linked list without dynamic memory allocation ?





# A linked list as an array of records (cont.)

nodes	.info	.next
[0]	David	4
[1]		
[2]	Miriam	6
[3]		
[4]	Joshua	7
[5]		
[6]	Robert	-1
[7]	Leah	2
[8]		
[9]		

list 0



# A Sorted list Stored in an Array of Nodes

nodes	.info	.next
[0]	David	4
[1]		
[2]	Miriam	6
[3]		
[4]	Joshua	7
[5]		
[6]	Robert	-1
[7]	Leah	2
[8]		
[9]		
list	0	

# An Array with Linked List of Values and Free Space

nodes	.info	.next
[0]	David	4
[1]		5
[2]	Miriam	6
[3]		8
[4]	Joshua	7
[5]		3
[6]	Robert	NUL
[7]	Leah	2
[8]		9
[9]		NUL

list	0
free	1

# An Array with Three Lists (Including the Free List)

		free	7
nodes	.info	.next	
[0]	John	4	
[1]	Mark	5	
[2]		3	
[3]		NUL	
[4]	Nell	8	
[5]	Naomi	6	
[6]	Robert	NUL	
[7]		2	
[8]	Susan	9	
[9]	Susanne	NUL	
list1	0		
list2	1		



# Recall Definition of Stack

***Logical (or ADT) level:*** A stack is an ordered group of **homogeneous items** (elements), in which the removal and addition of stack items can take place only at the top of the stack.

A stack is a **LIFO** “last in, first out” structure.



# Stack ADT Operations

**IsEmpty** -- Determines whether the stack is currently empty.

**IsFull** -- Determines whether the stack is currently full.

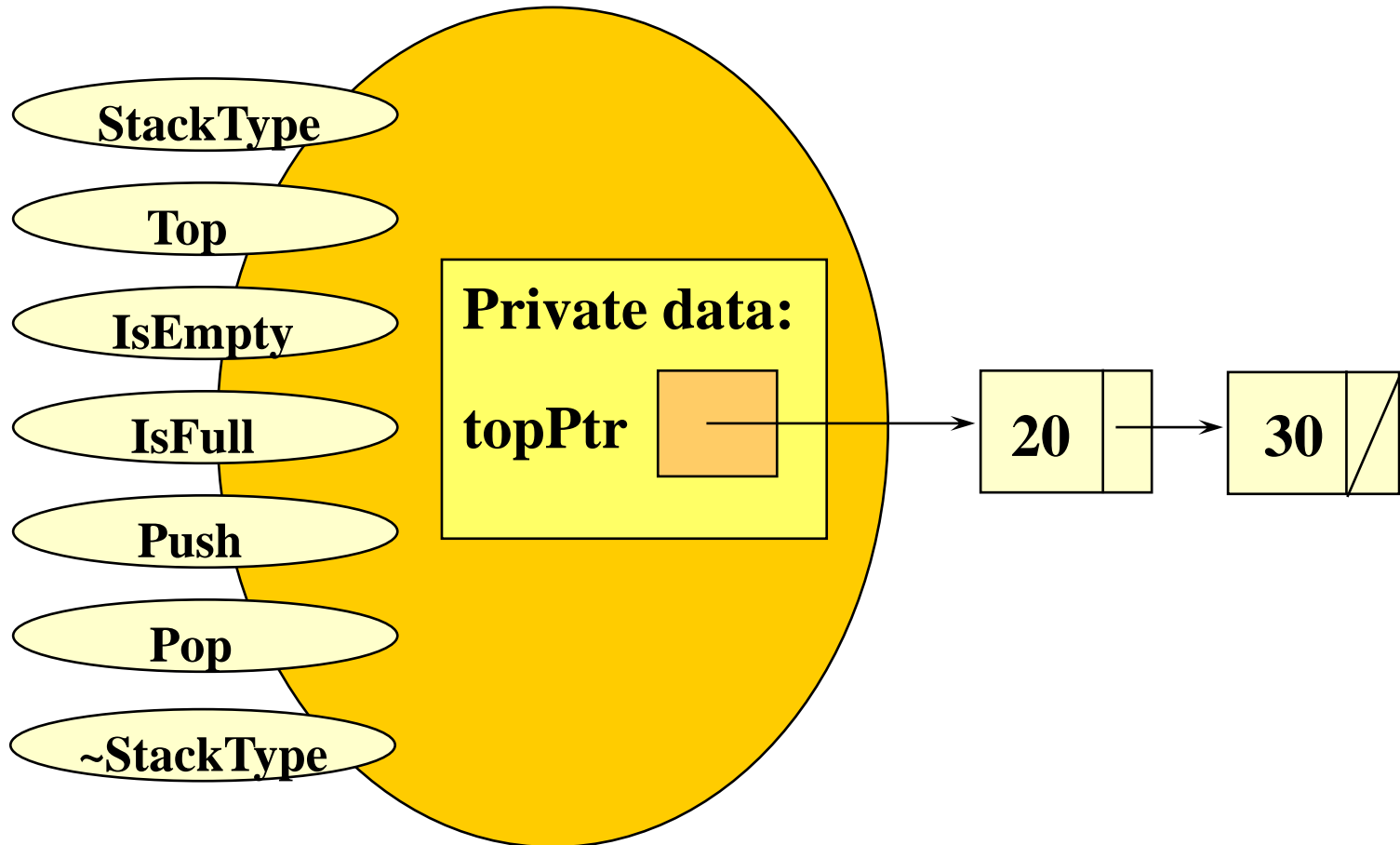
**Push (ItemType newItem)** -- Adds newItem to the top of the stack.

**Pop** -- Removes the item at the top of the stack and returns it in item.

**Top** -- Returns a copy of the top item

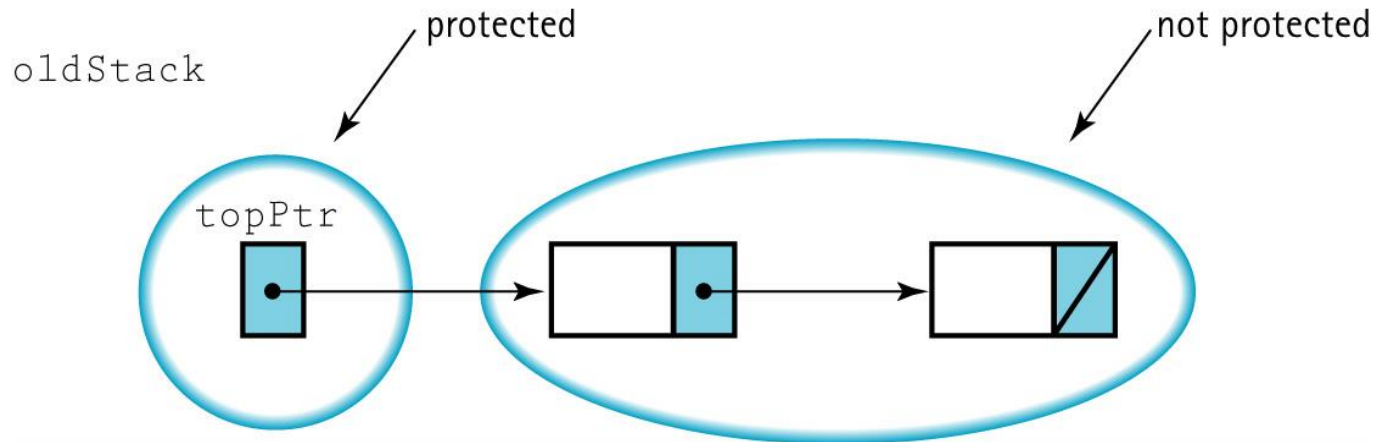


# class StackType<int>



# What happens . . .

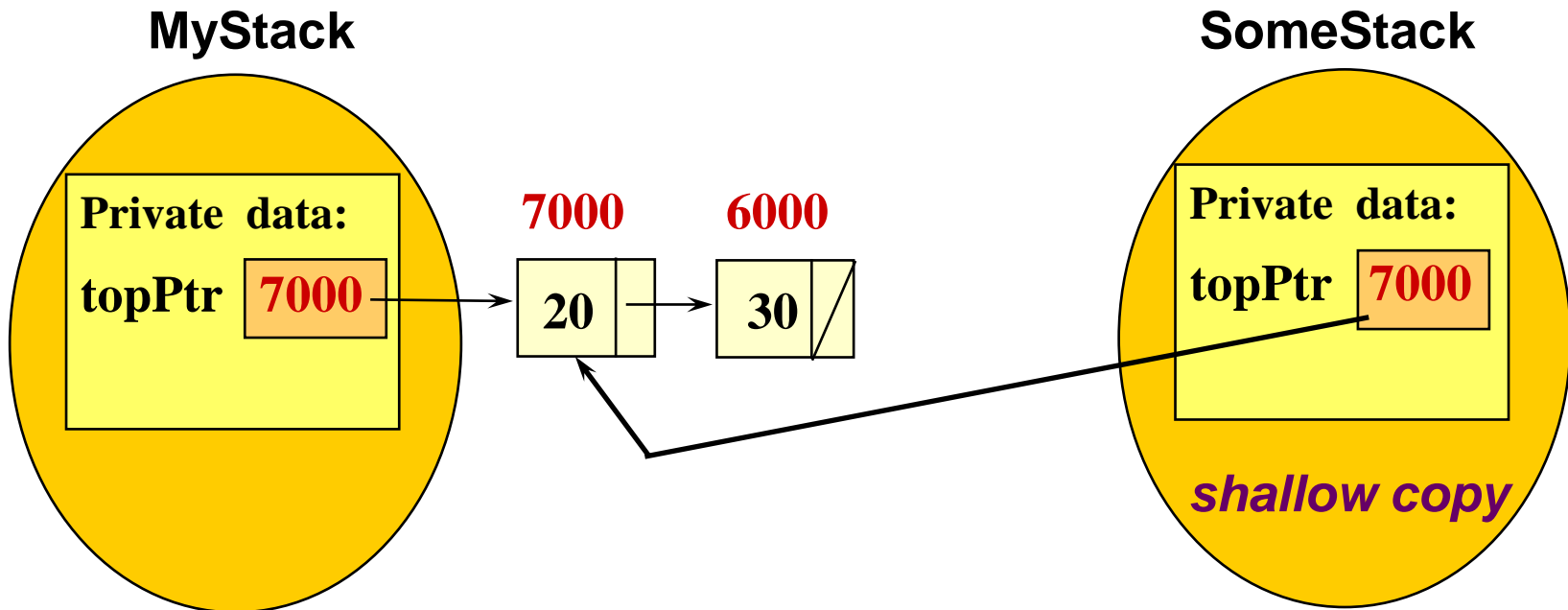
When a function is called that uses **pass by value** for a class object like our dynamically linked stack?





# Pass by value makes a shallow copy

```
StackType<int> MyStack;           // CLIENT CODE
:  
:  
MyFunction( MyStack );           // function call
```





# Shallow Copy vs. Deep Copy

**A shallow copy** copies only the class data members, and does not copy any pointed-to data.

**A deep copy** copies not only the class data members, but also makes separately stored copies of any pointed-to data.



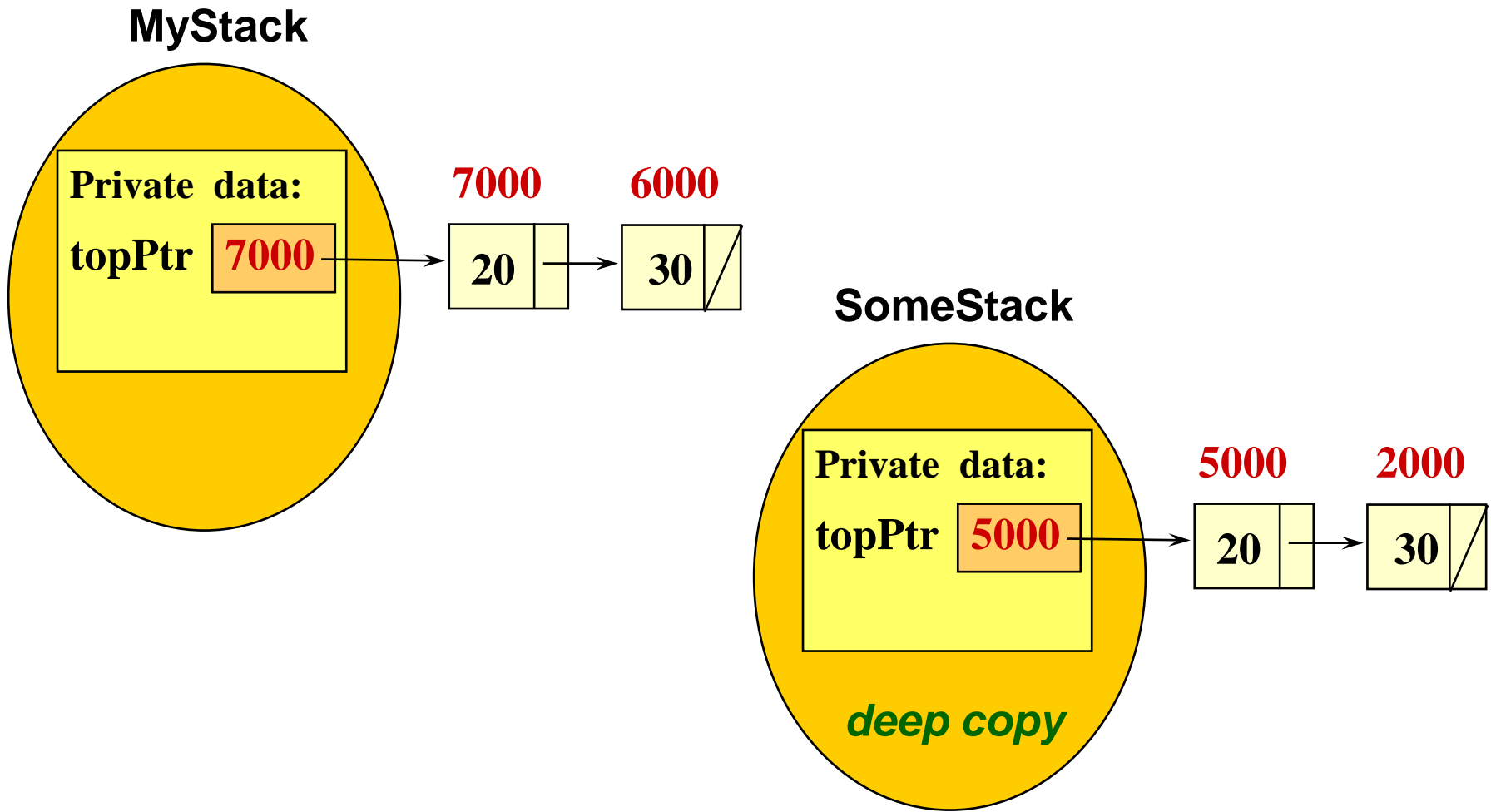
# What's the difference?

***A shallow copy*** shares the pointed to data with the original class object.

***A deep copy*** stores its own copy of the pointed to data at different locations than the data in the original class object.



# Making a deep copy





# Suppose MyFunction Uses Pop

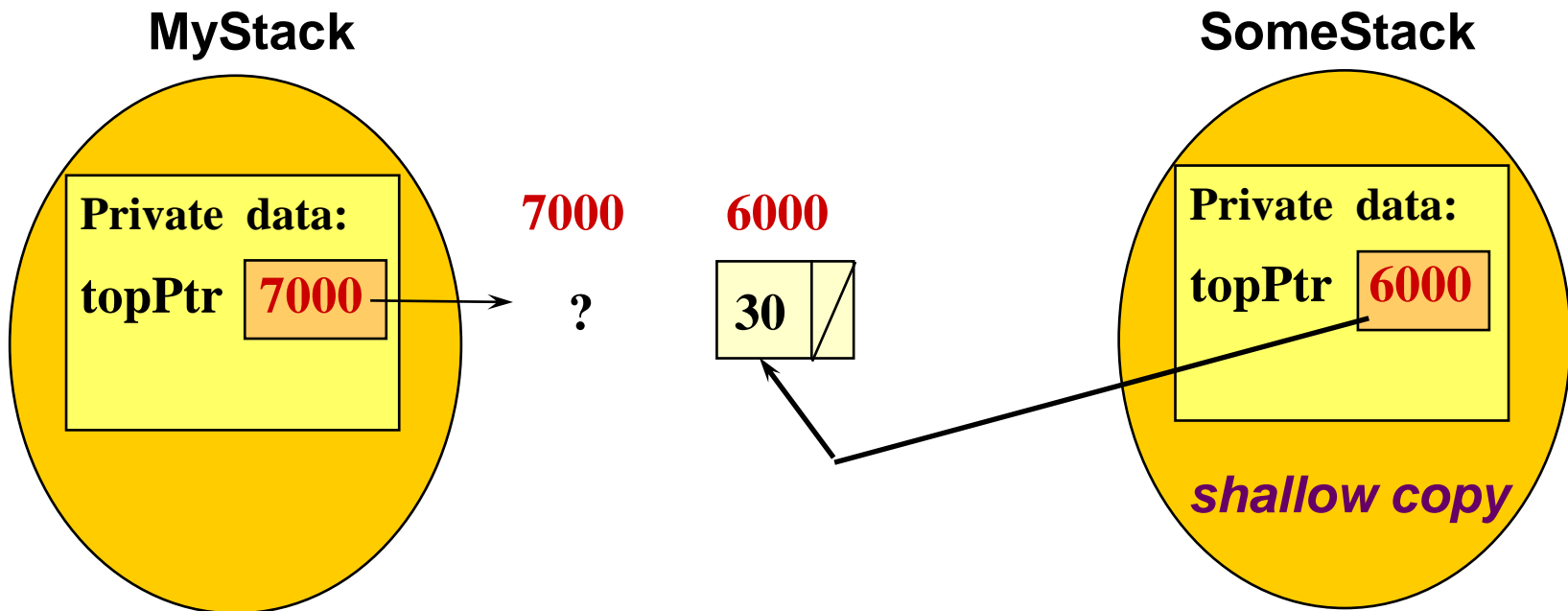
**// FUNCTION CODE**

```
template<class ItemType>
void MyFunction( StackType<ItemType> SomeStack )
    // Uses pass by value
{
    ItemType item;
    SomeStack.Pop(item) ;
    .
    .
    .
}
```

**WHAT HAPPENS IN THE SHALLOW COPY SCENARIO?**

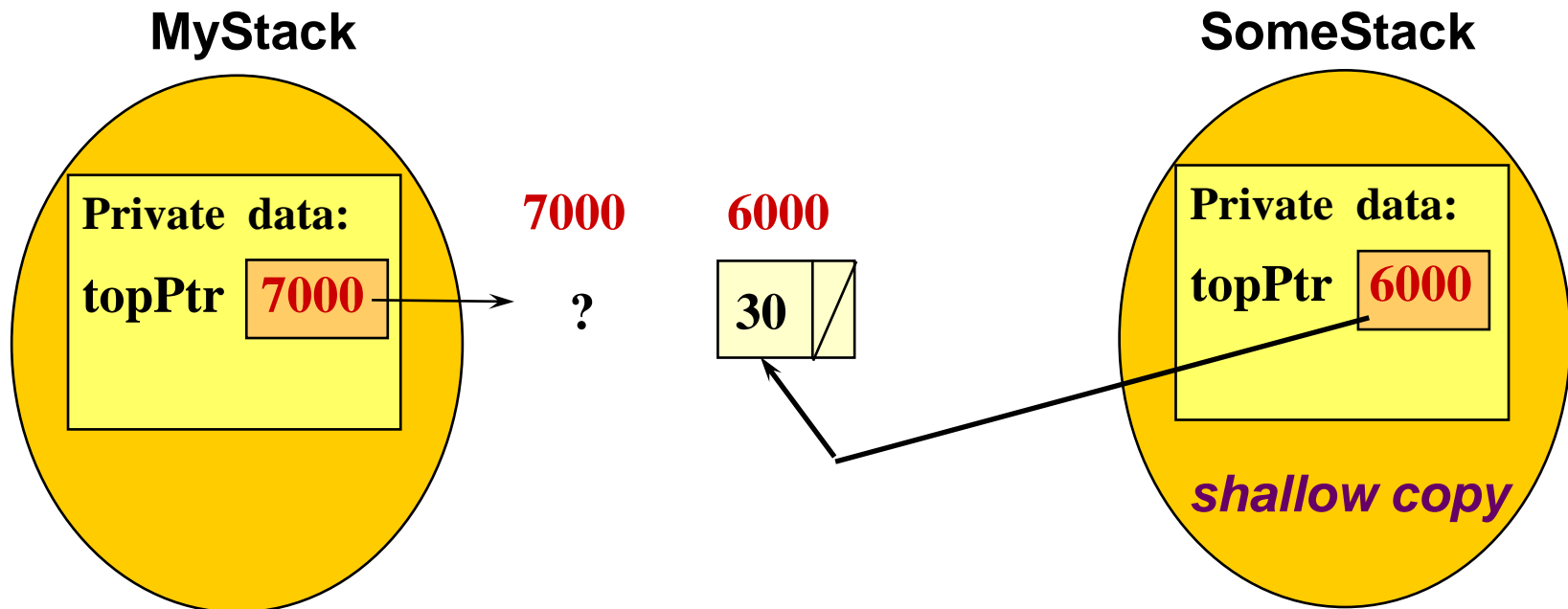
# MyStack.topPtr is left dangling

```
StackType<int> MyStack;    // CLIENT CODE
:  
:  
MyFunction( MyStack );
```



# MyStack.topPtr is left dangling

NOTICE THAT NOT JUST FOR THE SHALLOW COPY, BUT ALSO FOR ACTUAL PARAMETER MyStack, THE DYNAMIC DATA HAS CHANGED!





## As a result . . .

**This default method used for pass by value is not the best way when a data member pointer points to dynamic data.**

**Instead, you should write what is called a **copy constructor**, which makes a deep copy of the dynamic data in a different memory location.**

**Use **call by reference** in the case where you do not need a copy of the passing object.**





# More about copy constructors

**When there is a copy constructor provided for a class, the copy constructor is used to make copies for pass by value.**

**You do not call the copy constructor.**

**Like other constructors, it has no return type.**

**Because the **copy constructor** properly defines pass by value for your class, it **must use pass by reference in its definition.****



# Copy Constructor

**Copy constructor is a special member function of a class that is implicitly called in these three situations:**

**passing object parameters by value,  
initializing an object variable in a  
declaration,**

**returning an object as the return value of  
a function.**



```
// DYNAMICALLY LINKED IMPLEMENTATION OF STACK
template<class ItemType>
class StackType {
public:
    StackType( );
        // Default constructor.
        // POST: Stack is created and empty.
    StackType( const StackType<ItemType>& anotherStack );
        // Copy constructor.
        // Implicitly called for pass by value.
    :
    :
    :
    ~StackType( );
        // Destructor.
        // POST: Memory for nodes has been deallocated.
private:
    NodeType<ItemType>* topPtr ;
};
```



# Classes with Data Member Pointers Need

**CLASS CONSTRUCTOR**

**CLASS COPY CONSTRUCTOR**

**CLASS DESTRUCTOR**

```

template<class ItemType>                // COPY CONSTRUCTOR
StackType<ItemType>::
StackType( const StackType<ItemType>& anotherStack )
{
    NodeType<ItemType>* ptr1 ;
    NodeType<ItemType>* ptr2 ;
    if ( anotherStack.topPtr == NULL )
        topPtr = NULL ;
    else                                // allocate memory for first node
    {
        topPtr = new NodeType<ItemType> ;
        topPtr->info = anotherStack.topPtr->info ;
        ptr1 = anotherStack.topPtr->next ;
        ptr2 = topPtr ;
        while ( ptr1 != NULL )          // deep copy other nodes
        {
            ptr2->next = new NodeType<ItemType> ;
            ptr2 = ptr2->next ;
            ptr2->info = ptr1->info ;
            ptr1 = ptr1->next ;
        }
        ptr2->next = NULL ;
    }
}

```



# What About the Assignment Operator?

The **default method** used for assignment of class objects makes a **shallow copy**.

If your class has a data member pointer to dynamic data, you should write a member function to **overload the assignment operator to make a deep copy** of the dynamic data.



```
// DYNAMICALLY LINKED IMPLEMENTATION OF STACK
template<class ItemType>
class StackType {
public:
    StackType( );
        // Default constructor.
    StackType( const StackType<ItemType>& anotherStack );
        // Copy constructor.
    void operator= ( StackType<ItemType> );
        // Overloads assignment operator.
    :
    :
    :

    ~StackType( );
        // Destructor.
private:
    NodeType<ItemType>* topPtr ;
};
```



# C++ Operator Overloading Guides

- 1 All operators **except these** `::` `.` `sizeof` `?:` may be overloaded.
- 2 At least **one operand must be a class instance**.
- 3 You cannot change precedence, operator symbols, or number of operands.
- 4 Overloading `++` and `--` requires prefix form use by default, unless special mechanism is used.
- 5 To overload these operators `=` `()` `[]` member functions (not friend functions) must be used.
- 6 An operator can be given multiple meanings if the data types of operands differ.





# Using Overloaded Binary operator+

**When a Member Function was defined**

**myStack + yourStack**

**myStack.operator+(yourStack)**

**When a Friend Function was defined**

**myStack + yourStack**

**operator+(myStack, yourStack)**

A blue butterfly with black markings on its wings, perched on a green leaf.

## Case Study: Implementing a large integer ADT

The range of integer values varies from one computer to another

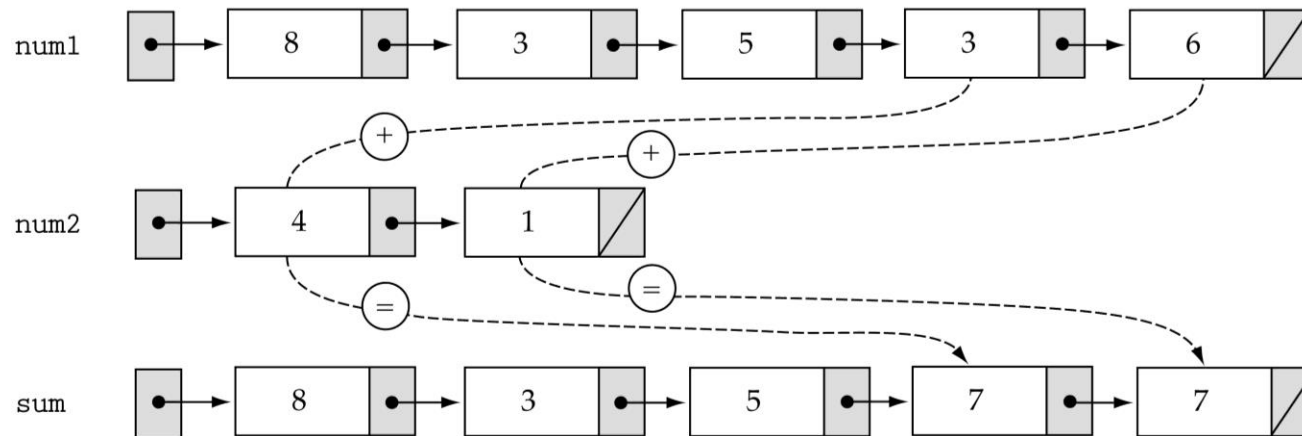
For *long* integers, the range is

$[-2,147,483,648 \text{ to } 2,147,483,647]$

How can we manipulate larger integers?

# Case Study: Implementing a large integer ADT (cont.)

(c)  $\text{sum} = 83536 + 41$



- A special list ADT

