



Samueli
Computer Science



CS32: Introduction to Computer Science II

Discussion Week 9

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May 31, 2019

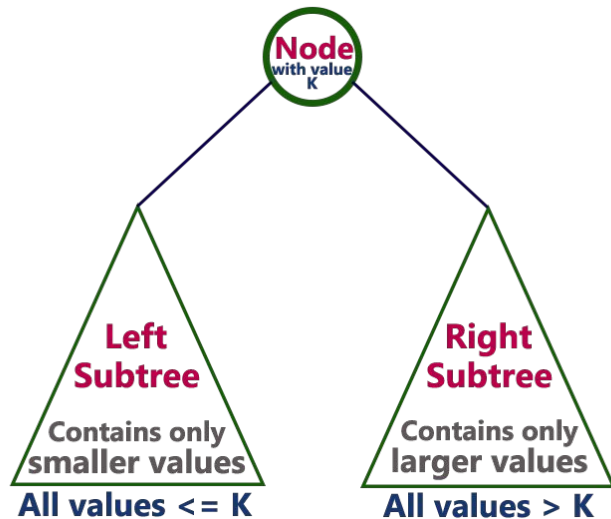
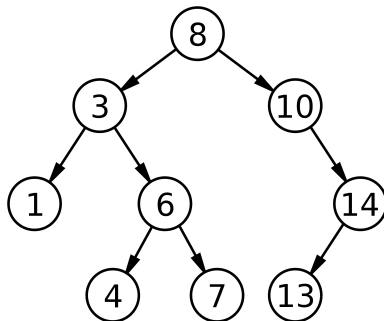
- Homework 5 is due 11PM next Thursday, June 6. (@A@)
- Project 4 is also due 11PM next Thursday, June 6. (@A@ >_< @A@)

- Binary Search Tree
- Heap (Preview)
- Hash Tables
- Project 4 Guidelines

Binary Search Tree

Definition, Properties

- The left subtree of a node contains only nodes with keys lesser than the node's key.
- The right subtree of a node contains only nodes with keys greater than the node's key.
- The left and right subtree each must also be a binary search tree.



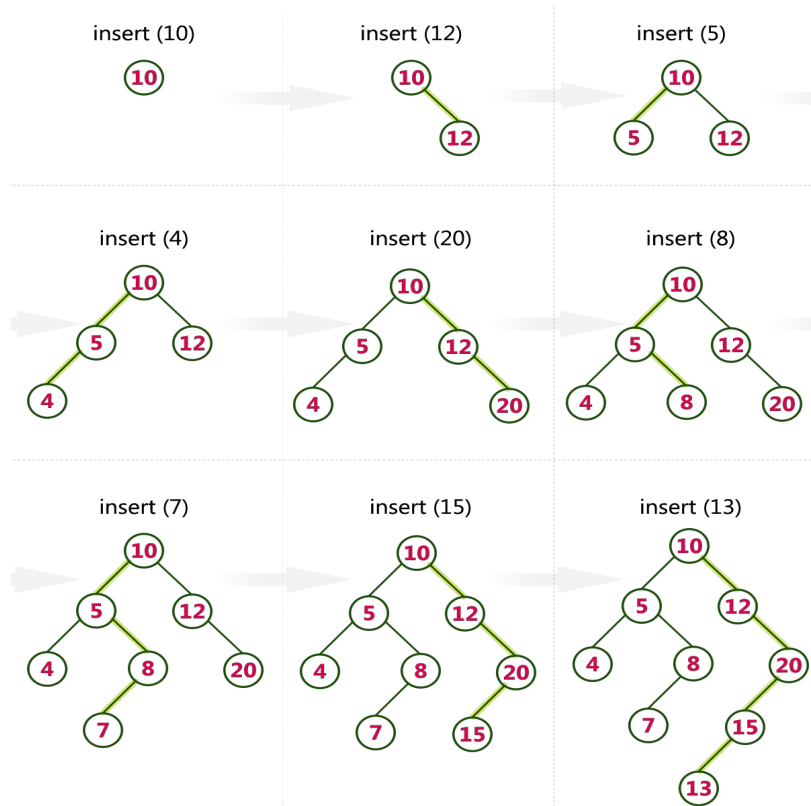
Binary Search Tree

Insertion

```
node* insert(node* node, ItemType key)
{
    /* If the tree is empty, return a new node */
    if (node == NULL) return newNode(key);

    /* Otherwise, recur down the tree */
    if (key < node->key)
        node->left = insert(node->left, key);
    else if (key > node->key)
        node->right = insert(node->right, key);

    /* return the (unchanged) node pointer */
    return node;
}
```

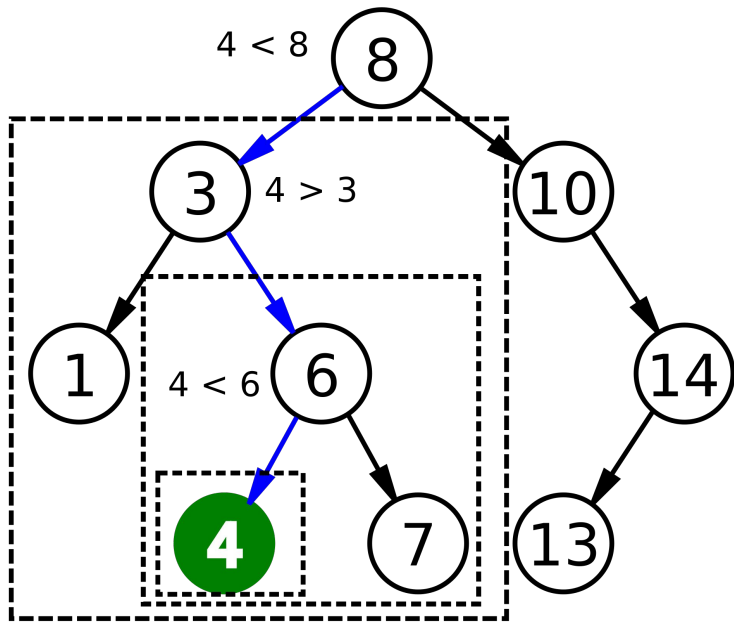


Binary Search Tree

Search

```
node* search(node* node, ItemType key)
{
    /* If the tree is empty, return null pointer */
    if (node == nullptr) return nullptr;

    /* compare with current node and decide*/
    if (key == node->key)
        return node;
    else if (key < node->key)
        return search(node->left, key);
    else
        return (node->right, key);
}
```



Binary Search Tree

Deletion

- Deletion is the most tricky one.

- 3 cases:

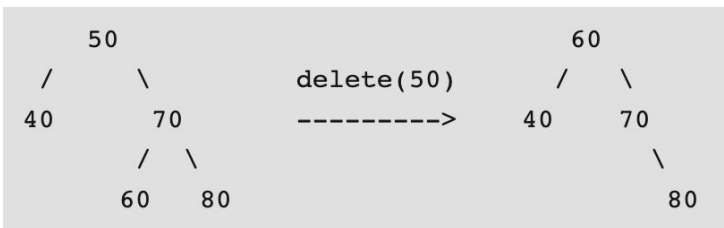
- The node is a leaf node.
- The node has one child.
- The node have two children.

Super easy! Just delete that node!

Not difficult! Copy the child to the node and delete the child.

Find inorder successor of the node. Copy contents of the inorder successor to the node and delete the inorder successor.

Note that inorder predecessor can also be used.



Binary Search Tree

Deletion

```
node* delete(node* node, ItemType key)
{
    if (node == nullptr) return nullptr;
    if (key < node->key) { node->left = delete(node->left, key); }
    else if (key > node->key) { node->right = delete(node->right, key); }
    else{
        /* case 1 & case 2*/
        if (node->left == nullptr) { node *temp = node->right; delete(node); return temp; }
        else if (node->right == nullptr) {node *temp = node->left; delete(node); return temp;}
        /* case 3 */
        node* temp = minValueNode(node->right);
        node->key = temp->key;
        root->right = deleteNode(root->right, temp->key);
    }
}
```


Binary Search Tree

Analysis of BST

- Insertion

- The (worst) case time complexity of search and insert operations is $O(h)$ where h is height of Binary Search Tree. In worst case, we may have to travel from root to the deepest leaf node.
- The height of a skewed tree may become n and the time complexity of search and insert operation may become $O(n)$.
- Average time complexity: $O(\log n)$

- Deletion

- Similar to insertion for complexity analysis

Binary Search Tree

FindMin and FindMax by using recursion

How to implement **findMaxKey** and **findMinKey** function?

We assume the root is not nullptr.

```
node* findMaxKey(const node* node)
{
    /* only need to check the right subtree*/
    if (node->right == nullptr) return node->key;
    return findMaxKey(node->right);
}
```

```
node* findMinKey(const node* node)
{
    /* only need to check the left subtree*/
    if (node->left == nullptr) return node->key;
    return findMinKey(node->left);
}
```

Binary Search Tree

FindMin and FindMax by using recursion

Question: How to test a tree is a valid BST?

One possible recursion solution by using `findMinKey` and `findMaxKey`.

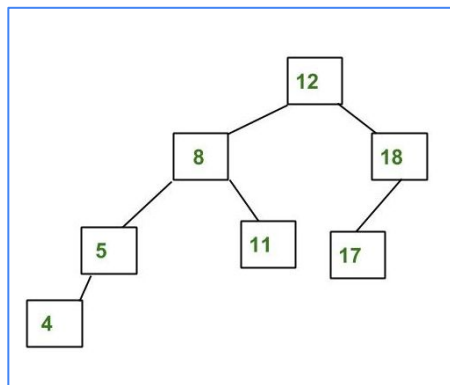
```
bool isValidBST(const node* node)
{
    if (node == nullptr) return true;
    /* check left subtree and right subtree condition*/
    if (node->left != nullptr && findMaxKey(node->left) > node->key)
        return false;
    if (node->right != nullptr && findMinKey(node->right) < node->key)
        return false;
    /* further check subtree with left child and right child */
    return isValidBST(node->left) && isValidBST(node->right)
}
```

What is the complexity of this `isValidBST()` function?

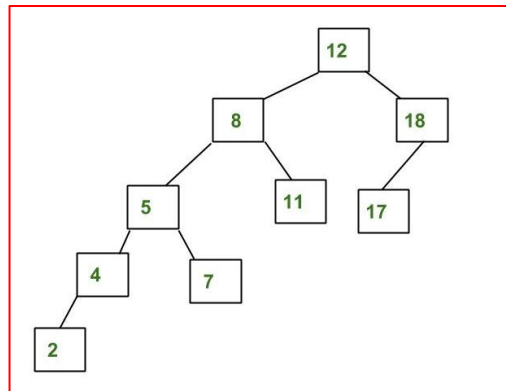
Beyond Binary Search Tree

AVL Tree

- What is the drawback of naive BST? → It can be skewed! Not good!
- AVL (Adelson-Velsky and Landis) Tree is a self-balancing BST.



This is OK!



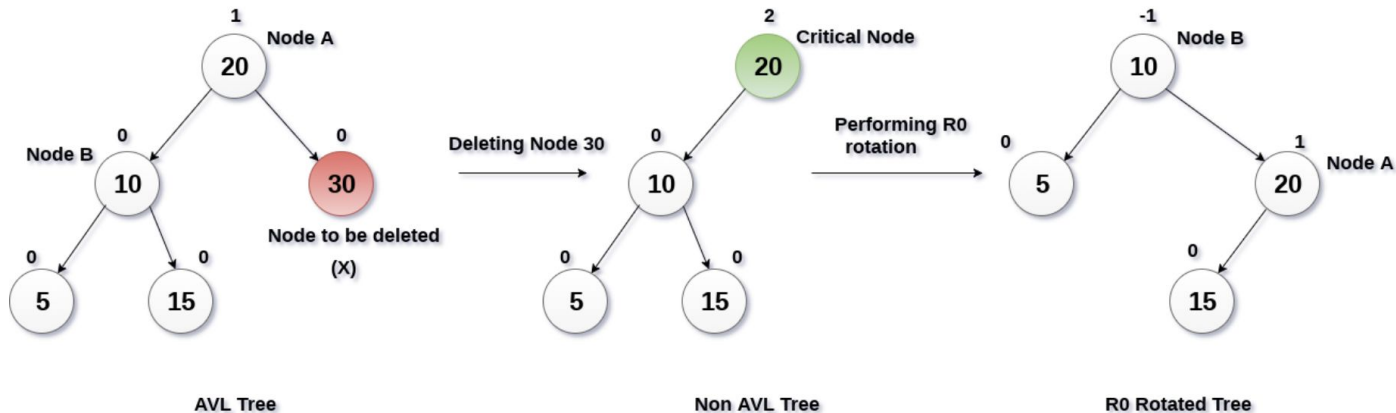
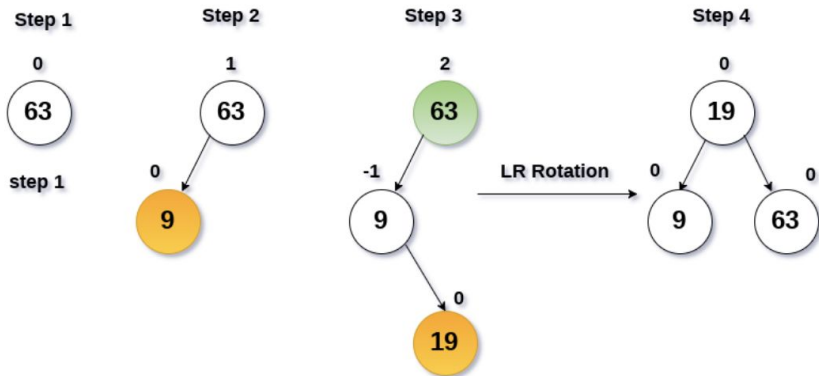
Just OK is not OK!

Beyond Binary Search Tree

AVL Tree: Operations

- Insertion
- Deletion

Please check this interesting [demo!](#)



Beyond Binary Search Tree

The tree family (1)

There are many interesting tree structures such as:

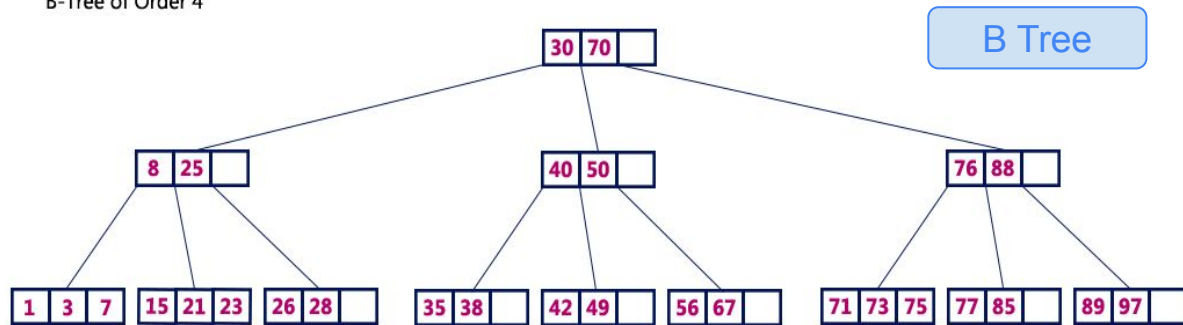
- B tree and B+ tree (I did not hear about A+ tree though)
- 2-3-4 tree
- R tree (spatial index tree)
- **Red-black tree**
- K-D tree

Please check the given links for more details!

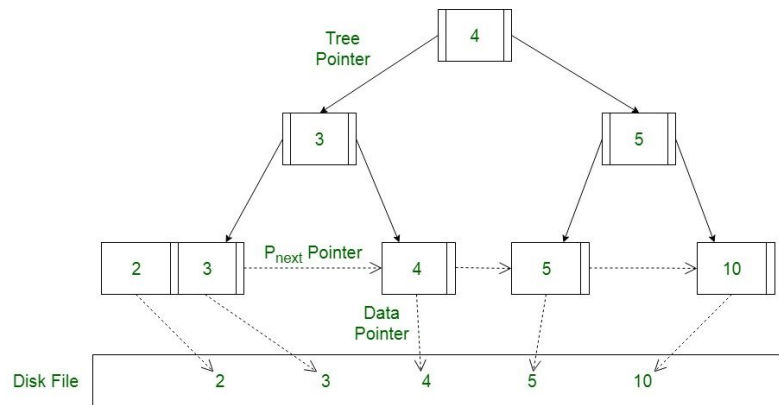
Beyond Binary Search Tree

The tree family (2)

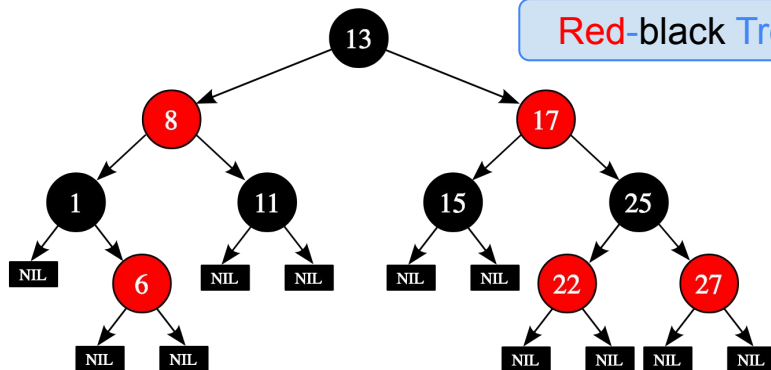
B-Tree of Order 4



B+ Tree



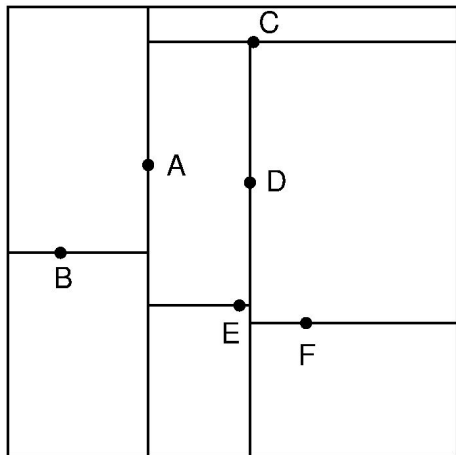
Red-black Tree



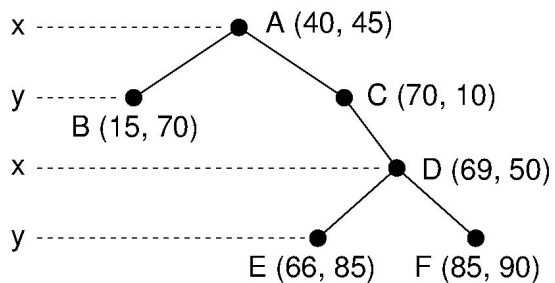
Beyond Binary Search Tree

The tree family (3)

KD Tree

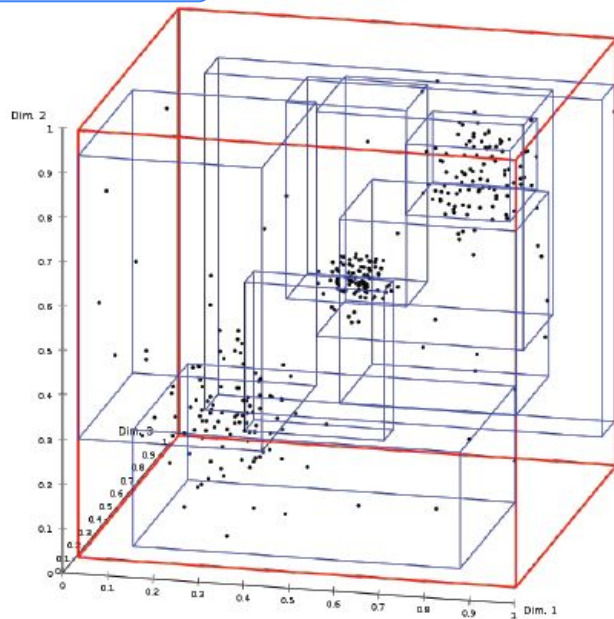


(a)



(b)

R Tree



Visualization of an R*-tree for 3D cubes using ELKI

Heap (preview)

Definition and properties

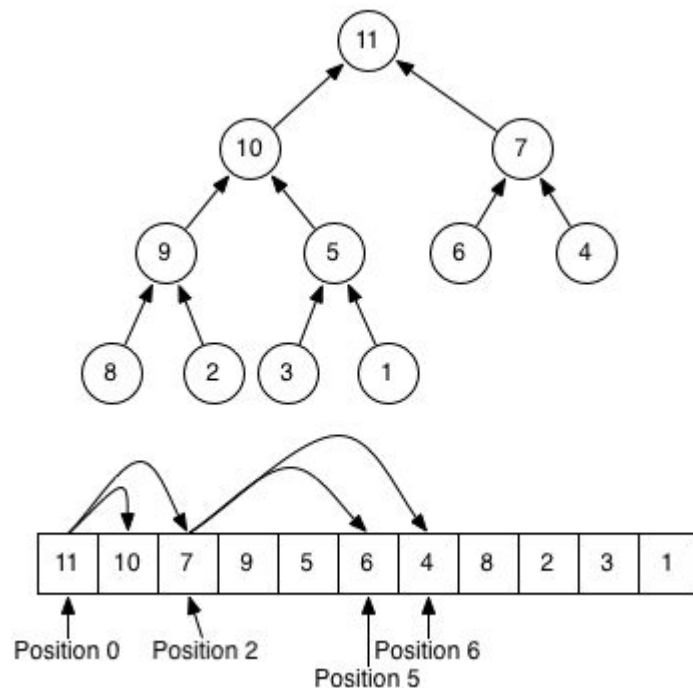
- About heap
 - Heap is considered as complete binary tree.
 - Every nodes carries a value greater than or equal to its children (for MaxHeap).
 - Often implemented as an array.
 - Body structure of priority queue.



Stack



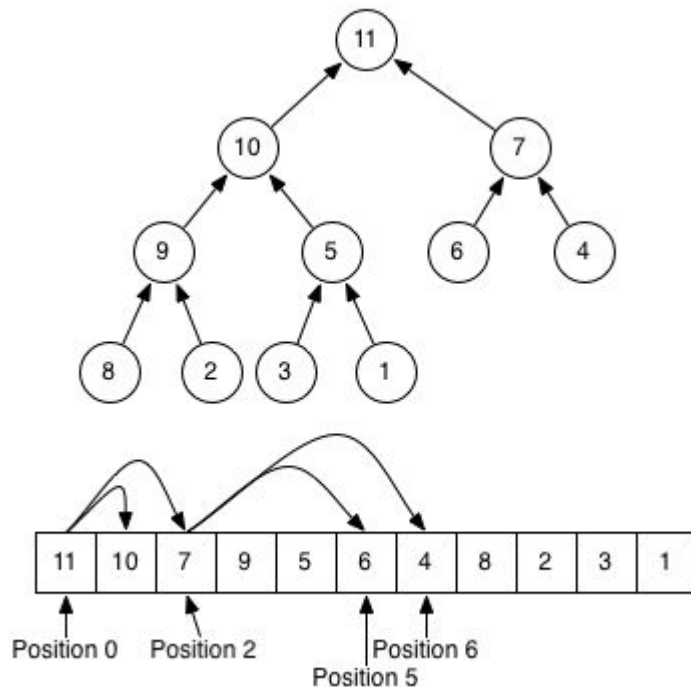
Heap



Heap (preview)

Standard operations

- Three operations of heaps
 - Find Max (search)
 - Insert Node (insert)
 - Delete Max (delete)
- How to implement FindMax() function of a heap?
 - Well, that is just too obvious!



Heap & Heapsort (preview)

Complexity of heap operations

- Find Max $\rightarrow O(1)$
- Insert $\rightarrow O(\log n)$
- Delete Max Node $\rightarrow O(\log n)$

6 5 3 1 8 7 2 4

- Bonus: How can you sort based on heap?
 - Insert all elements into a heap.
 - Extract the maximum element from the heap one by one.
 - Check the example in [Wikipedia](#)
- What is the complexity of heapsort?
 - $O(n \log n)$

Hash Tables

Start from hashing and hash functions

- Hash functions: Take a “key” and map it to a number
- Requirement for hash function: should return the same value for the same key
- Good hash functions:
 - Spreads out the values: two different key are likely to results in different hash values. → Avoid confliction
 - Compute each value quickly.
- Example: FNV-1



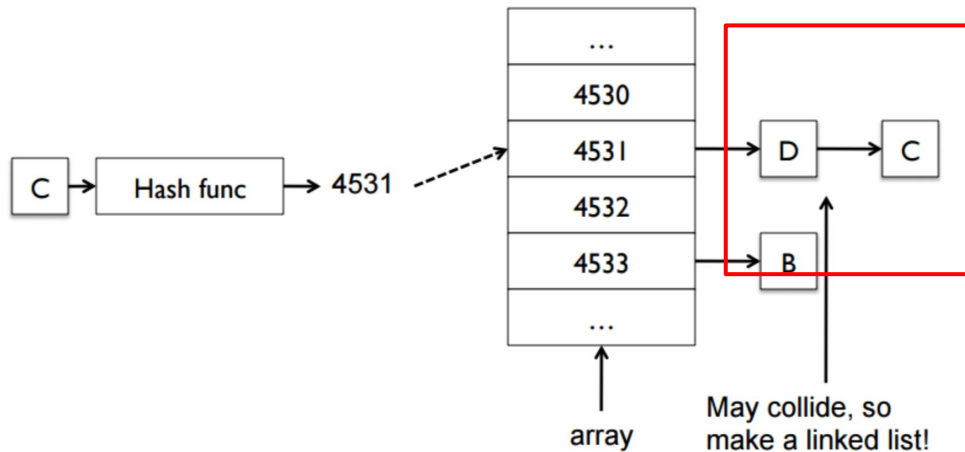
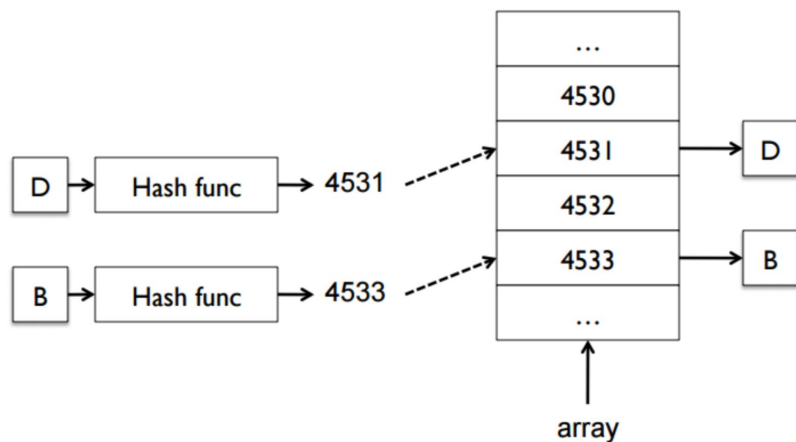
```
unsigned int FNV-1(string s) {  
    unsigned int h = 2166136261U;  
    for (int k = 0; k != s.size(); k++)  
    {  
        h += s[k];  
        h *= 16777619;  
    }  
    return h;  
}
```

Hash Tables

Examples

- Example: Use a hash table to store people.
- Use a linked list to collision in the hash function.

*"You should almost **NEVER** assume that collisions are impossible!!!" --David Smallberg*



- Insert
- Remove
- Search

- The complexity depends on your hash tables.
- Closed Hashing
 - Fixed number of buckets
 - All operations are $O(n)$ with a small constant of proportionality
- Open Hashing
 - Consider $\#entries / \#buckets$
 - Almost $O(1)$ for all operations

Hash Tables Problem 1

Word counting and sorting in a document

- Question: We have n words in a document, whose vocabulary size is v . Count top k frequent words in a document.

Two steps: counting + sorting

- The most efficient way to count the frequency for all words takes $O(\underline{n})$ time complexity.
- After getting the frequency of each word, the most efficient way to get the top k frequent words takes $O(\underline{v \log k})$ time complexity.
- Totally the entire procedure takes $O(\underline{n + v \log k})$.

Hash Tables Problem 2

The very first question in LeetCode - TwoSum

- Given an array of integers, return indices of the two numbers such that they add up to a specific target.
- You may assume that each input would have exactly one solution, and you may not use the same element twice.
- Example: Given `nums = [2, 7, 11, 15]` and `target = 9`. Because `nums[0] + nums[1] = 2 + 7 = 9`, return `[0, 1]`.

Hash Table in STL

Map, multimap, unordered_map, HashMap

- Useful functions of STL map, multimap, unordered_map
 - `size()`, `begin()`, `end()`, `empty()`
 - `insert(keyvalue, mapvalue)`
 - `find()`
 - `operator[]`
- Internal implementation:
 - `map/multimap`: Red-black tree
 - `unordered_map`: Hash table
- Pros and cons between `map` and `unordered_map`

Note & Reminders:

1. Get File I/O done first.
2. Build BST or hashtable by vector, list, stack and queue (and string) instead of map or unordered_map.
 - Reason to build BST or hashtable: fast indexing the substrings
 - For complexity analysis
3. Algorithm improvements → Shorten the diff file



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Break Time! (5 minutes)

Q & A

- Exercise problems from **Worksheet 8** (see “LA worksheet” tab in CS32 website). Answers will be posted next week.
- Questions for today:
 - 1, 2, 4, 6

Question 1

Given an array `arr[0..n-1]` of distinct elements and a range `[low, high]`, use a hash table to find all numbers that are in the range, but not in the array. Print out the missing elements in sorted order. Use the function header:

```
void inRange(int arr[], int size, int low, int high)
```

For example:

Input: `arr[] = {10, 12, 11, 15}, low = 10, high = 15`

Output: `13, 14`

Input: `arr[] = {1, 14, 11, 51, 15}, low = 50, high = 55`

Output: `50, 52, 53, 54`

Question 1: Solution

```
#include <unordered_set>
#include <iostream>
using namespace std;

void inRange(int arr[], int size, int low, int high)
{
    // Insert all elements of arr[] in set
    unordered_set<int> set;
    for (int i=0; i<size; i++)
        set.insert(arr[i]);

    // Traverse through the range and print all
    // missing elements
    for (int x=low; x<=high; x++)
        if (set.find(x) == set.end()) //or if (set.count(x) == 0)
            cout << x << " ";
}
```

Question 2

Given an array of integers and a target sum, determine if two integers in the array can be added together to equal the sum. The time complexity of your solution should be $O(n)$ where n is the number of elements in the array. In other words, you cannot use the brute force method in which you compare each element with every other element using nested for loops. Example:

Array: 4 8 3 7 9 2 5
Target: 15

You can take 8 and 7 from the array, and their sum equals the target of 15. Thus, the function we will write will return true.

Use the following function header:

```
bool twoSum(const int arr[], int n, int target);
```

Question 2: Solution

```
bool twoSum(const int arr[], int n, int target) {  
    unordered_set<int> numsFound;  
  
    for (int i = 0; i < n; i++) {  
  
        int complement = target - arr[i];  
  
        if (numsFound.find(complement) != numsFound.end()) {  
            return true;  
        }  
        else {  
            numsFound.insert(arr[i]);  
        }  
    }  
    return false;  
}
```


Question 4

Given a string, find the first non-repeating character in it and return its index. If it doesn't exist, return -1. You may assume the string contain only lowercase letters. Use a hashtable to solve this problem.

Use the function header:

```
int firstUniqueChar(std::string s)
```

Example:

```
s = "hello"           // return 0  
s = "computer science" // return 1
```

Question 4: Solution

```
int firstUniqueChar(std::string s) {  
    // Map character to the frequency of occurrence  
    unordered_map<char, int> counter;  
    for(int i = 0; i < s.size(); i++) {  
        counter[s[i]]++;  
    }  
    for (int i = 0; i < s.size(); i++) {  
        if (counter[s[i]] == 1)  
            return i;  
    }  
    return -1;  
}
```

Question 6

Implement the following function, given the following data structure:

```
struct Node {  
    int val;  
    Node* left;  
    Node* right;  
};
```

```
bool isMinHeap(const Node* head);
```

This function takes in the head of a binary tree and returns whether or not that binary tree represents a binary min heap. In other words, this tree must follow the min heap property, where a parent is less than or equal to its children.

Question 6: Solution

```
// Precondition: head points to a complete tree
bool isMinHeap(const Node* head) {
    if (head == nullptr)
        return true;

    Node* left = head->left;
    Node* right = head->right;
    if (left != nullptr && head->val > left->val)
        return false;
    if (right != nullptr && head->val > right->val)
        return false;

    return isMinHeap(head->left) && isMinHeap(head->right);
}
```



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Thank you!

Q & A