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Computer Science



# CS32: Introduction to Computer Science

## **Discussion Week 7**

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Feb 22, 2019

# Announcement

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- Midterm Part 2 is due on Tuesday, Feb 26.
- Project 3 Part 2 is due on Thursday, February 28

- 
- STL (continued)
  - Algorithmic Efficiency (Big-O notation)
  - Sorting (1)

# STL Table list

Check out when you need



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- C++ Containers library
  - Sequence containers: array, vector, deque, list, forward\_list
  - Container adaptors: set, map, multiset, multimap
  - Associate containers: unordered\_set, unordered\_map
  - Unordered associative containers: stack, queue, priority\_queue
- Check link
  - [https://en.cppreference.com/w/cpp/container#Sequence\\_containers](https://en.cppreference.com/w/cpp/container#Sequence_containers)

# Project 3 warmup

Why fail on test 3?

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- Pitfalls in modifying STL containers while you're traversing them, given the various iterator invalidation rules.
- **Another reminder:** For a container of raw pointers to dynamically allocated objects, the container operations know nothing about that: erasing an item does **NOT** call delete on the pointer.

Category	Container	After <b>insertion</b> , are...		After <b>erasure</b> , are...		Conditionally
		iterators valid?	references valid?	iterators valid?	references valid?	
Sequence containers	array	N/A		N/A		
	vector	No		N/A		Insertion changed capacity
		Yes		Yes		Before modified element(s)
		No		No		At or after modified element(s)
	deque	No	Yes	Yes, except erased element(s)		Modified first or last element
			No	No		Modified middle only
	list	Yes		Yes, except erased element(s)		
	forward_list	Yes		Yes, except erased element(s)		
Associative containers	set multiset map multimap	Yes		Yes, except erased element(s)		
Unordered associative containers	unordered_set unordered_multiset unordered_map unordered_multimap	No	Yes	N/A		Insertion caused rehash
		Yes		Yes, except erased element(s)		No rehash

# STL Table list

## Member function table

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Header Container	Sequence containers					Associative containers				Unordered associative containers				Container adaptors			
	<array> array (constructor) (destructor) operator=	<vector> vector (implicit) ~vector operator=	<deque> deque (implicit) ~deque operator=	<forward_list> forward_list (implicit) ~forward_list operator=	<list> list (implicit) ~list operator=	<set> set (implicit) ~set operator=	multiset multiset (implicit) ~multiset operator=	map map (implicit) ~map operator=	multimap multimap (implicit) ~multimap operator=	unordered_set unordered_set (implicit) ~unordered_set operator=	unordered_multiset unordered_multiset (implicit) ~unordered_multiset operator=	unordered_map unordered_map (implicit) ~unordered_map operator=	unordered_multimap unordered_multimap (implicit) ~unordered_multimap operator=	<stack> stack (implicit) ~stack operator=	queue queue (implicit) ~queue operator=	priority_queue priority_queue (implicit) ~priority_queue operator=	
Iterators	begin	begin	begin	begin	begin	begin	begin	begin	begin	begin	begin	begin	begin				
	cbegin	cbegin	cbegin	cbegin	cbegin	cbegin	cbegin	cbegin	cbegin	cbegin	cbegin	cbegin	cbegin				
	end	end	end	end	end	end	end	end	end	end	end	end	end				
	cend	cend	cend	cend	cend	cend	cend	cend	cend	cend	cend	cend	cend				
	rbegin	rbegin	rbegin	rbegin	rbegin	rbegin	rbegin	rbegin	rbegin	rbegin	rbegin	rbegin	rbegin				
Element access	cbegin	cbegin	cbegin	cbegin	cbegin	cbegin	cbegin	cbegin	cbegin	cbegin	cbegin	cbegin	cbegin				
	rend	rend	rend	rend	rend	rend	rend	rend	rend	rend	rend	rend	rend				
	crend	crend	crend	crend	crend	crend	crend	crend	crend	crend	crend	crend	crend				
	at	at	at	at				at				at					
	operator[]	operator[]	operator[]	operator[]				operator[]				operator[]					
Capacity	data	data	data														
	front	front	front	front	front	front									front	top	
	back	back	back	back	back	back									back		
	empty	empty	empty	empty	empty	empty	empty	empty	empty	empty	empty	empty	empty		empty	empty	
	size	size	size	size	size	size	size	size	size	size	size	size	size		size	size	
Modifiers	max_size	max_size	max_size	max_size	max_size	max_size	max_size	max_size	max_size	max_size	max_size	max_size	max_size				
	resize	resize	resize	resize	resize												
	capacity	capacity	capacity	capacity	capacity												
	reserve	reserve	reserve	reserve	reserve												
	shrink_to_fit	shrink_to_fit	shrink_to_fit	shrink_to_fit	shrink_to_fit					bucket_count	bucket_count	bucket_count	bucket_count				
List operations	clear	clear	clear	clear	clear	clear	clear	clear	clear	clear	clear	clear	clear				
	insert	insert	insert	insert	insert	insert	insert	insert	insert	insert	insert	insert	insert				
	insert or assign	insert or assign	insert or assign	insert or assign	insert or assign	insert or assign	insert or assign	insert or assign	insert or assign	insert or assign	insert or assign	insert or assign	insert or assign				
	emplace	emplace	emplace	emplace	emplace	emplace	emplace	emplace	emplace	emplace	emplace	emplace	emplace				
	emplace_hint	emplace_hint	emplace_hint	emplace_hint	emplace_hint	emplace_hint	emplace_hint	emplace_hint	emplace_hint	emplace_hint	emplace_hint	emplace_hint	emplace_hint				
Lookup	try_emplace	try_emplace	try_emplace	try_emplace	try_emplace	try_emplace	try_emplace	try_emplace	try_emplace	try_emplace	try_emplace	try_emplace	try_emplace				
	erase	erase	erase	erase	erase	erase	erase	erase	erase	erase	erase	erase	erase				
	push_front	push_front	push_front	push_front	push_front	push_front	push_front	push_front	push_front	push_front	push_front	push_front	push_front				
	emplace_front	emplace_front	emplace_front	emplace_front	emplace_front	emplace_front	emplace_front	emplace_front	emplace_front	emplace_front	emplace_front	emplace_front	emplace_front				
	pop_front	pop_front	pop_front	pop_front	pop_front	pop_front	pop_front	pop_front	pop_front	pop_front	pop_front	pop_front	pop_front		pop	pop	
Observers	push_back	push_back	push_back	push_back	push_back	push_back	push_back	push_back	push_back	push_back	push_back	push_back	push_back		push	push	
	emplace_back	emplace_back	emplace_back	emplace_back	emplace_back	emplace_back	emplace_back	emplace_back	emplace_back	emplace_back	emplace_back	emplace_back	emplace_back		emplace	emplace	
	pop_back	pop_back	pop_back	pop_back	pop_back	pop_back	pop_back	pop_back	pop_back	pop_back	pop_back	pop_back	pop_back		pop	pop	
	swap	swap	swap	swap	swap	swap	swap	swap	swap	swap	swap	swap	swap		swap	swap	
	merge	merge	merge	merge	merge	merge	merge	merge	merge	merge	merge	merge	merge				
Allocator	extract	extract	extract	extract	extract	extract	extract	extract	extract	extract	extract	extract	extract				
	splice	splice	splice	splice	splice	splice	splice	splice	splice	splice	splice	splice	splice				
	remove	remove	remove	remove	remove	remove	remove	remove	remove	remove	remove	remove	remove				
	remove_if	remove_if	remove_if	remove_if	remove_if	remove_if	remove_if	remove_if	remove_if	remove_if	remove_if	remove_if	remove_if				
	reverse	reverse	reverse	reverse	reverse	reverse	reverse	reverse	reverse	reverse	reverse	reverse	reverse				
Container	unique	unique	unique	unique	unique	unique	unique	unique	unique	unique	unique	unique	unique				
	sort	sort	sort	sort	sort	sort	sort	sort	sort	sort	sort	sort	sort				
	count	count	count	count	count	count	count	count	count	count	count	count	count				
	find	find	find	find	find	find	find	find	find	find	find	find	find				
	contains	contains	contains	contains	contains	contains	contains	contains	contains	contains	contains	contains	contains				
Container	lower_bound	lower_bound	lower_bound	lower_bound	lower_bound	lower_bound	lower_bound	lower_bound	lower_bound	lower_bound	lower_bound	lower_bound	lower_bound				
	upper_bound	upper_bound	upper_bound	upper_bound	upper_bound	upper_bound	upper_bound	upper_bound	upper_bound	upper_bound	upper_bound	upper_bound	upper_bound				
	equal_range	equal_range	equal_range	equal_range	equal_range	equal_range	equal_range	equal_range	equal_range	equal_range	equal_range	equal_range	equal_range				
	key_comp	key_comp	key_comp	key_comp	key_comp	key_comp	key_comp	key_comp	key_comp	key_comp	key_comp	key_comp	key_comp				
	value_comp	value_comp	value_comp	value_comp	value_comp	value_comp	value_comp	value_comp	value_comp	value_comp	value_comp	value_comp	value_comp				
Container	hash_function	hash_function	hash_function	hash_function	hash_function	hash_function	hash_function	hash_function	hash_function	hash_function	hash_function	hash_function	hash_function				
	key_eq	key_eq	key_eq	key_eq	key_eq	key_eq	key_eq	key_eq	key_eq	key_eq	key_eq	key_eq	key_eq				
	get_allocator	get_allocator	get_allocator	get_allocator	get_allocator	get_allocator	get_allocator	get_allocator	get_allocator	get_allocator	get_allocator	get_allocator	get_allocator				
	Container	array	vector	deque	forward_list	list	set	multiset	map	multimap	unordered_set	unordered_multiset	unordered_map	unordered_multimap	stack	queue	priority_queue

# \*Smart Pointer

A good tool in modern C++



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- A smart pointer is an abstract data type that simulates a pointer while providing added features, such as automatic memory management or bounds checking.
- C++ libraries provide implementations of smart pointers in the form of `unique_ptr`, `shared_ptr` and `weak_ptr`
- Trade-off by using smart pointers: may increase memory usage (for example in `list`)
- More info: [\[Smart pointer tutorial\]](#)

```
// normal pointers
void UseNormalPointer{
    MyClass *ptr = new MyClass();
    ptr->doSomething();
}
// We must delete ptr to avoid memory leak!
```

```
// smart pointers, defined in std
void UseSmartPointer{
    unique_ptr<MyClass> ptr(new MyClass());
    ptr->doSomething();
}
// ptr is deleted automatically here!
// unique_ptr: encapsulated pointer as only data member
```



# \*Smart Pointer

## unique\_ptr and shared\_ptr

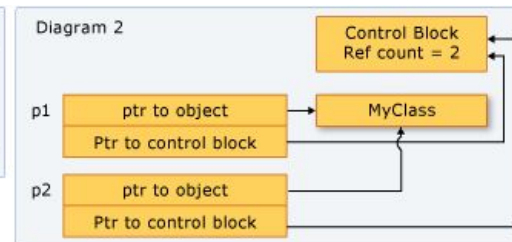
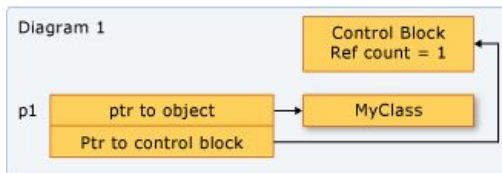


- **unique\_ptr**
  - Allows exactly one owner of the underlying pointer.
  - Can be moved to a new owner, but not copied or shared.
  - Small and efficient (the size is one pointer as data member)
  - More about unique\_ptr: [\[unique\\_ptr tutorial\]](#)
- **shared\_ptr**
  - Reference-counted smart pointer. Use when you want to assign one raw pointer to multiple owners.
  - The size is two pointers; one for the object and one for the shared control block that contains the reference count.
  - More about shared\_ptr: [\[shared\\_ptr tutorial\]](#)

```
auto ptrA = make_unique<Song>(L"Diana Krall", L"The Look of Love");
```



```
auto ptrB = std::move(ptrA);
```



# Exercise

Try to implement a `unique_ptr`?



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```
template<class T>
class unique_ptr {
public:
    unique_ptr(T* p) : ptr_(p) {}
    ~unique_ptr() {
        delete ptr_;
    }
private:
    T* ptr_;
};
```

# Pointers vs Smart Pointers

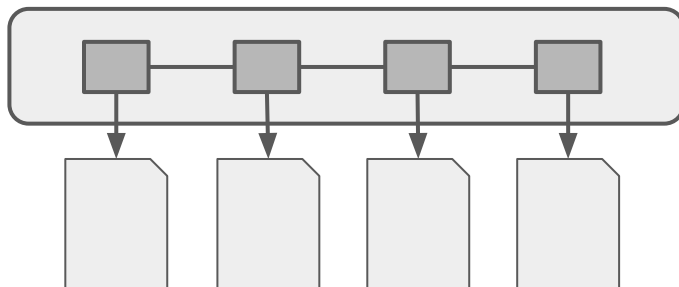
Example: Container of pointers



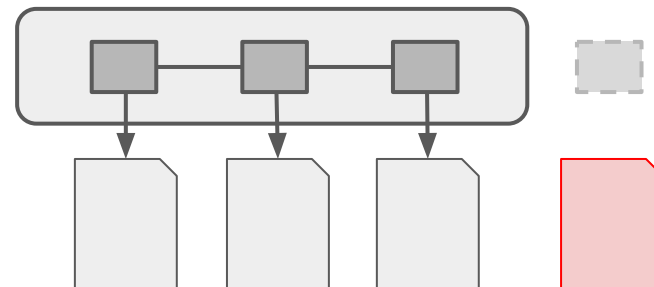
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## Normal Pointers

vector 1

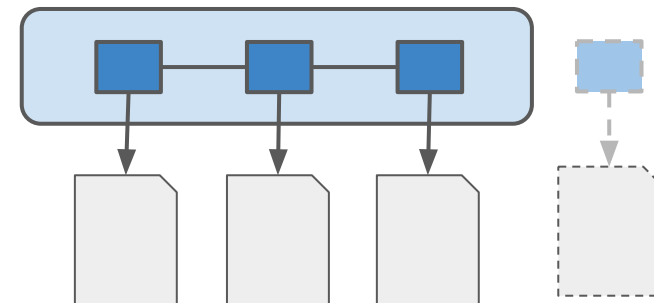
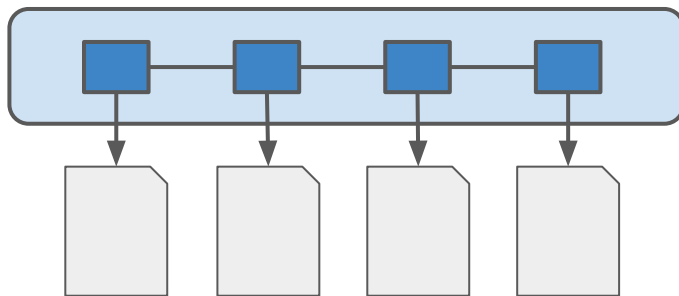


pop\_back()



## Smart Pointers

vector 2



# Algorithm Efficiency

Note: Complexity of a program



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- Quantify the efficiency of a program.
- The magnitude of time and space cost for an algorithm given certain size of input.
  - Time complexity: quantifies the run time.
  - Space complexity: quantifies the usage of the memory (or sometimes hard disk drives, cloud disk drives, etc.).
- Naturally, the size of input determines how long a program runs.
  - Often, the larger the size of input, the longer the run time. But not always that case.
  - Consider: sort an array of 1,000 items and 1,000,000 items vs get size of an array of 1,000 items and 1,000,000 items
- Big-O notation

# Big-O Notation

## Formal definition



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If you are interested in formal definition, check [here](#).

Well, you can simply understand as how many operations given input size of  $n$  regardless of the constant.

No need to memorize definitions.

Example: if your program takes,

- about  $n$  steps  $\rightarrow O(n)$
- about  $2n$  steps  $\rightarrow O(n)$
- about  $n^2$  steps  $\rightarrow O(n^2)$
- about  $3n^2+10n$  steps  $\rightarrow O(n^2)$
- about  $2^n$  steps  $\rightarrow O(2^n)$

Question: What is the speed of growth for typical function?

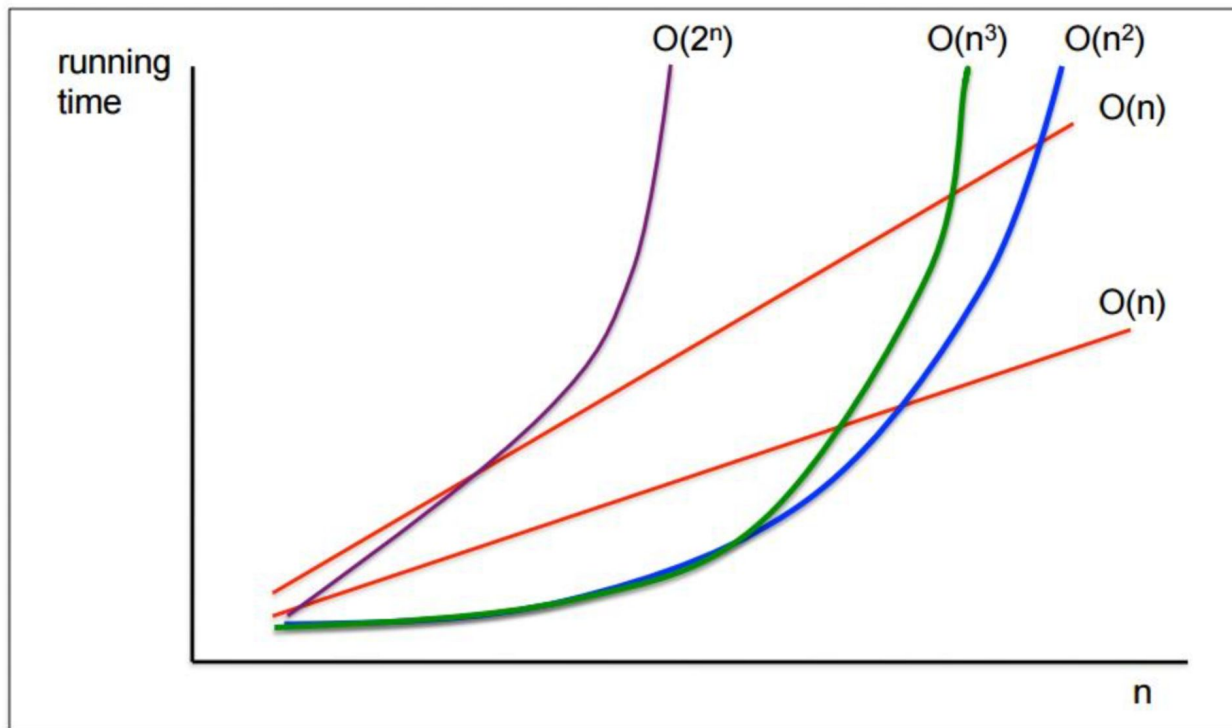
$f(n) = \log(n) / n / n^2 / 2^n / n!$

# Big-O Notation

Growth speed

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# Big-O Arithmetic

How to determine the entire program?

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Generally,

- If things happen sequentially, we add Big-Os;
- If one thing happen with another, then we multiply Big-Os.
- Watch the **LOOPS** in your programs!

Rules:

$$O(f(n)) + O(g(n)) = O(\max(f(n), g(n)))$$

$$O(f(n)) \times O(g(n)) = O(f(n) \times g(n))$$

# Efficiency Analysis

## Example 1: Linear Search



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- Linear search: Look for one item in an unsorted array
- Best cases? Average cases? Worst cases?
- What if the array is ordered?

```
int linear_search(array arr, size n, value v)
{
    for (int i=0; i<n; i++)
    {
        if (arr[i] == v)
            return i;
    }
    return -1;
}
```



# Efficiency Analysis

## Example 2: Enumerate all pairs



- Task: Find all pairs from one array (Note: [1,2] and [2,1] are considered different pairs)

```
int all_pairs(array arr, size n, value v)
{
    for (int i=0; i<n; i++)
    {
        for (int j=0; j<n; j++)
        {
            if (i != j)
                cout << "Pair:" << arr[i] << "and" << arr[j] << endl;
        }
    }
    return -1;
}
```

A diagram illustrating the nested loops in the code. A red bracket on the left side groups the outer loop (the 'for' loop with 'i') and its closing brace. A blue bracket on the left side groups the inner loop (the 'for' loop with 'j') and its closing brace, showing it is nested within the outer loop.

# Efficiency Analysis

## Example 3: Binary search



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- Task: Look for one item in a sorted array

```
// this is pseudo code
int binary_search(array arr, value v, start_index s, end_index e)
{
    if (s > e) return -1
    find the middle point i=(s+e)/2
    if (arr[i] == v) return i
    else if (arr[i] < v) return binary_search(arr, v, i+1, e)
    else return binary_search(arr, v, s, i-1)
}
```

# Big-O and Complexity

Big O	Name	n = 128
$O(1)$	constant	1
$O(\log n)$	logarithmic	7
$O(n)$	linear	128
$O(n \log n)$	"n log n"	896
$O(n^2)$	quadratic	16192
$O(n^k), k \geq 1$	polynomial	
$O(2^n)$	exponential	$10^{40}$
$O(n!)$	factorial	$10^{214}$

Question: Can you find an algorithms with  $O(n!)$  complexity?

Most important algorithm ever!

Methods:

- Selection sort
- Bubble sort
- Insertion sort
- Merge sort
- Quick sort

Focus on:

1. Steps for each sorting algorithm
2. Runtime complexity for worst cases, best cases and average cases
3. Space complexity
4. How about additional assumptions, such as the array is “almost sorted” / “reversed” arrays

# Sorting

## Selection sort



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### Steps:



**Idea:** Find the smallest item in the unsorted portion and place it in the front.

### Runtime complexity:

Average:  $O(n^2)$

Worst:  $O(n^2)$

Best:  $O(n^2)$

**Space complexity:**  $O(1)$

# Sorting

## Insertion sort



### Steps:



**Idea:** Pick one from the unsorted part and place it in the right position.

### Runtime complexity:

Average:  $O(n^2)$

Worst:  $O(n^2)$

Best:  $O(n)$

**Space complexity:**  $O(1)$

# Sorting

## Bubble sort

### Steps:

4	3	1	5	2
3	4	1	5	2
3	1	4	5	2
3	1	4	2	5
1	3	2	4	5
1	2	3	4	5

**Idea:** Well, just “bubble” as its name

### Runtime complexity:

Average:  $O(n^2)$

Worst:  $O(n^2)$

Best:  $O(n)$

**Space complexity:**  $O(1)$

# Merge sort

## Idea: Divide and conquer



Average:  $O(n \log n)$

Worst:  $O(n \log n)$

Best:  $O(n \log n)$

**Space complexity:**  $O(n)$



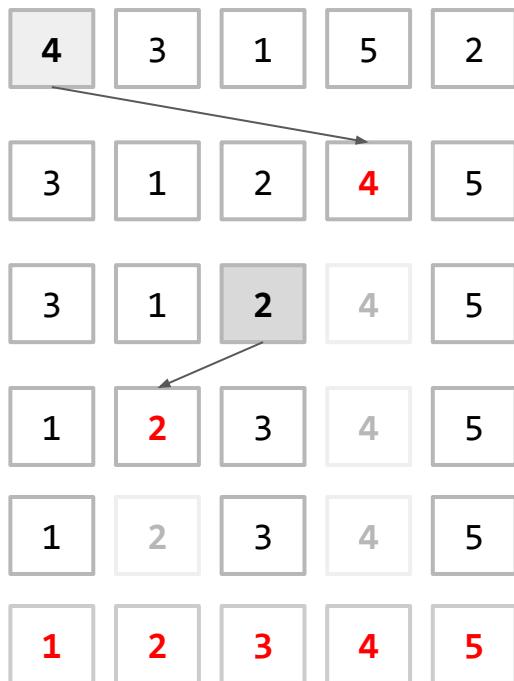
# Sorting

## Quicksort



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Steps:



**Idea:** Set a pivot. Numbers less than pivot are placed to front while others to end.

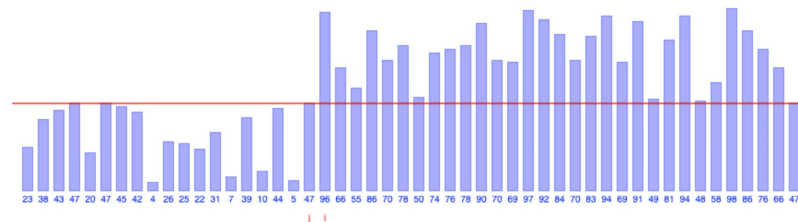
**Runtime complexity:**

Average:  $O(n \log n)$

Worst:  $O(n^2)$

Best:  $O(n \log n)$

**Space complexity:**  $O(\log n)$



# Sorting

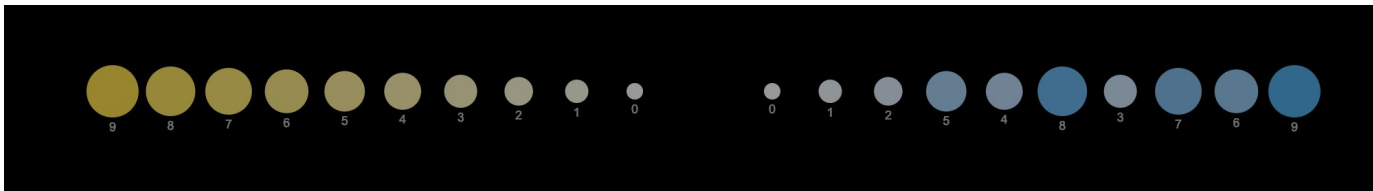
## Other methods and complexity?



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- $O(n \log n)$  is faster than  $O(n^2)$  → Merge sort is more efficient than selection, insertion and bubble sort in runtime.
- $O(n \log n)$  is best average complexity that a general sorting algorithm can achieve.
- With more information about the data provided, you can sometimes sort things almost linearly.

Question: What is the complexity of these sorting algorithms if you know the array is **reversed**? What if the array is **almost already sorted**?



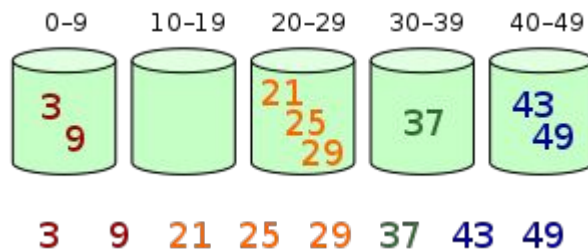
# Sorting

## Other methods and complexity?



There are many other sorting methods:

- Shell sort (shell 1959, Knuth 1973, Ciura 2001)
- Quicksort 3-way
- Heap sort
- Bucket sort



# Sorting

## Why sorting is important?

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Sorting is the most important and basic algorithm. Many other real-world problems are somewhat based on sorting, including:

Sorting Algorithms Animations: <https://www.toptal.com/developers/sorting-algorithms>

Other good demos:

<https://www.cs.usfca.edu/~galles/visualization/ComparisonSort.html>

<http://sorting.at/>

# Sorting

## Variant sorting problems

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Question: How about get the *K-th* largest numbers in one array?

[Leetcode question #215](#)

Hint:

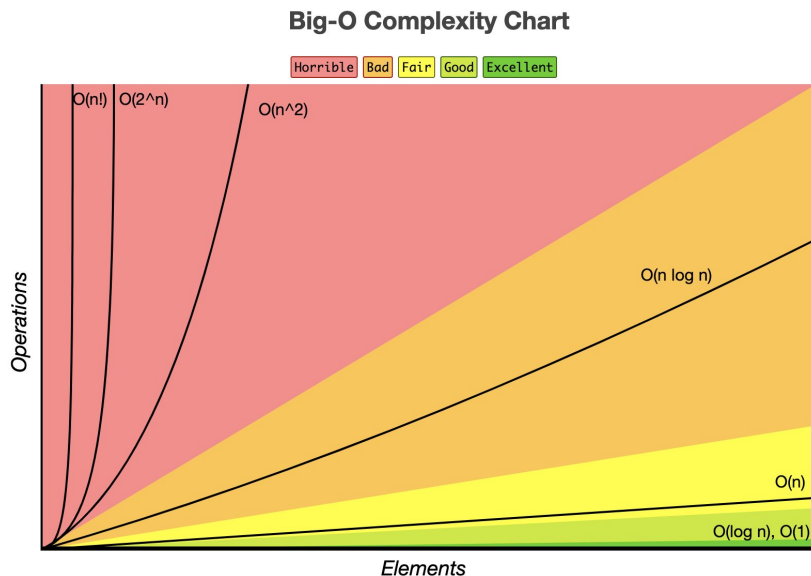
1. How to find the k-th largest numbers by merge sort and quicksort (or other sort methods)? What are the average and worst complexity?
2. What data structures is good to use?

# Big-O Notation

## Big-O Complexity Chart



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## Array Sorting Algorithms

Algorithm	Time Complexity			Space Complexity
	Best	Average	Worst	Worst
<u>Quicksort</u>	$\Omega(n \log(n))$	$\theta(n \log(n))$	$O(n^2)$	$O(\log(n))$
<u>Mergesort</u>	$\Omega(n \log(n))$	$\theta(n \log(n))$	$O(n \log(n))$	$O(n)$
<u>Timsort</u>	$\Omega(n)$	$\theta(n \log(n))$	$O(n \log(n))$	$O(n)$
<u>Heapsort</u>	$\Omega(n \log(n))$	$\theta(n \log(n))$	$O(n \log(n))$	$O(1)$
<u>Bubble Sort</u>	$\Omega(n)$	$\theta(n^2)$	$O(n^2)$	$O(1)$
<u>Insertion Sort</u>	$\Omega(n)$	$\theta(n^2)$	$O(n^2)$	$O(1)$
<u>Selection Sort</u>	$\Omega(n^2)$	$\theta(n^2)$	$O(n^2)$	$O(1)$
<u>Tree Sort</u>	$\Omega(n \log(n))$	$\theta(n \log(n))$	$O(n^2)$	$O(n)$
<u>Shell Sort</u>	$\Omega(n \log(n))$	$\theta(n(\log(n))^2)$	$O(n(\log(n))^2)$	$O(1)$
<u>Bucket Sort</u>	$\Omega(n+k)$	$\theta(n+k)$	$O(n^2)$	$O(n)$
<u>Radix Sort</u>	$\Omega(nk)$	$\theta(nk)$	$O(nk)$	$O(n+k)$
<u>Counting Sort</u>	$\Omega(n+k)$	$\theta(n+k)$	$O(n+k)$	$O(k)$
<u>Cubesort</u>	$\Omega(n)$	$\theta(n \log(n))$	$O(n \log(n))$	$O(n)$



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# Thank you!

Q & A