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COMP 314: Algorithms and Complexity  
Lab work 2: Sorting

Github Link:

[https://github.com/Junth19/COMP-314-Algorithms-and-Complexity-Lab/tree/master/CE\\_III\\_05\\_Lab2](https://github.com/Junth19/COMP-314-Algorithms-and-Complexity-Lab/tree/master/CE_III_05_Lab2)

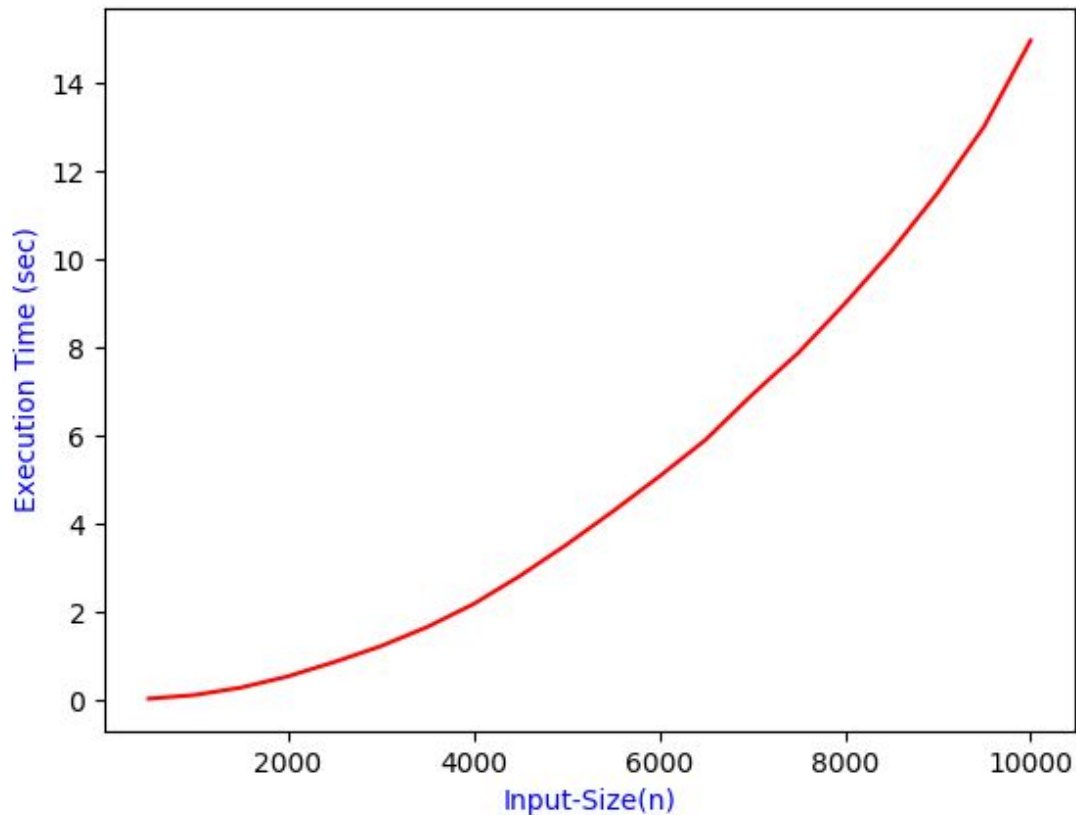
## Insertion Sort

The insertion sort algorithm iterates through an input array and removes one element per iteration, finds the place the element belongs in the array, and then places it there. This process grows a sorted list from left to right.

❖ Worst Case:  $O(n^2)$

| Input Size(n) | Execution Time(sec) |
|---------------|---------------------|
| 500           | 0.04690384864807129 |
| 1000          | 0.1315937042236328  |
| 1500          | 0.3013467788696289  |
| 2000          | 0.5533123016357422  |
| 2500          | 0.8784968852996826  |
| 3000          | 1.2374794483184814  |
| 3500          | 1.6719088554382324  |
| 4000          | 2.1939640045166016  |
| 4500          | 2.8319811820983887  |
| 5000          | 3.53788685798645    |
| 5500          | 4.29442286491394    |
| 6000          | 5.085555791854858   |
| 6500          | 5.9187257289886475  |
| 7000          | 6.938089370727539   |
| 7500          | 7.893252849578857   |
| 8000          | 9.006293535232544   |
| 8500          | 10.196868896484375  |
| 9000          | 11.520684957504272  |
| 9500          | 13.025882720947266  |

Graph:  $O(n^2)$

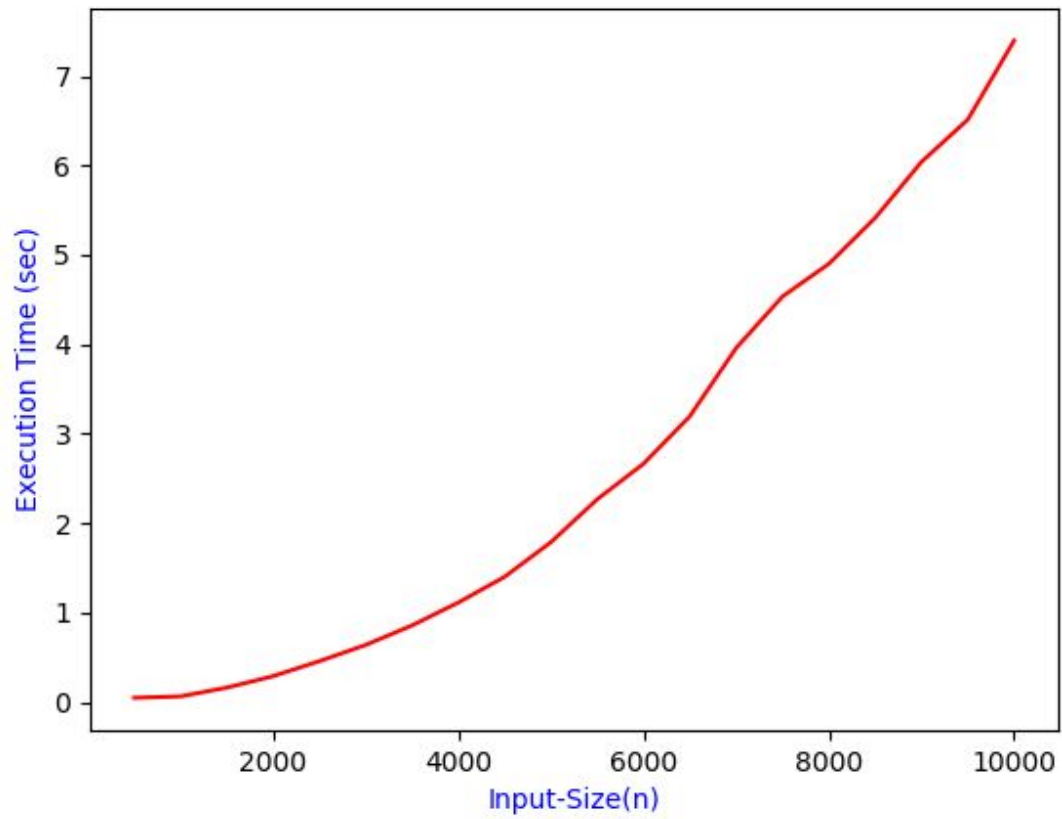


The simplest worst-case input is an array sorted in reverse order. The set of all worst-case inputs consists of all arrays where each element is the smallest or second-smallest of the elements before it. In these cases, every iteration of the inner loop will scan and shift the entire sorted subsection of the array before inserting the next element. This gives insertion sort a quadratic running time (i.e.,  $O(n^2)$ ).

❖ Average Case:  $O(n^2)$

| Input Size(n) | Execution Time(sec) |
|---------------|---------------------|
| 500           | 0.04681134223937988 |
| 1000          | 0.06245088577270508 |
| 1500          | 0.16046404838562012 |
| 2000          | 0.2882354259490967  |
| 2500          | 0.4565401077270508  |
| 3000          | 0.6378993988037109  |
| 3500          | 0.8562748432159424  |
| 4000          | 1.1122260093688965  |
| 4500          | 1.3997390270233154  |
| 5000          | 1.7873189449310303  |
| 5500          | 2.2673544883728027  |
| 6000          | 2.664799213409424   |
| 6500          | 3.196070432662964   |
| 7000          | 3.9591143131256104  |
| 7500          | 4.534657716751099   |
| 8000          | 4.899816513061523   |
| 8500          | 5.41332221031189    |
| 9000          | 6.0402185916900635  |
| 9500          | 6.512476205825806   |

Graph:  $O(n^2)$

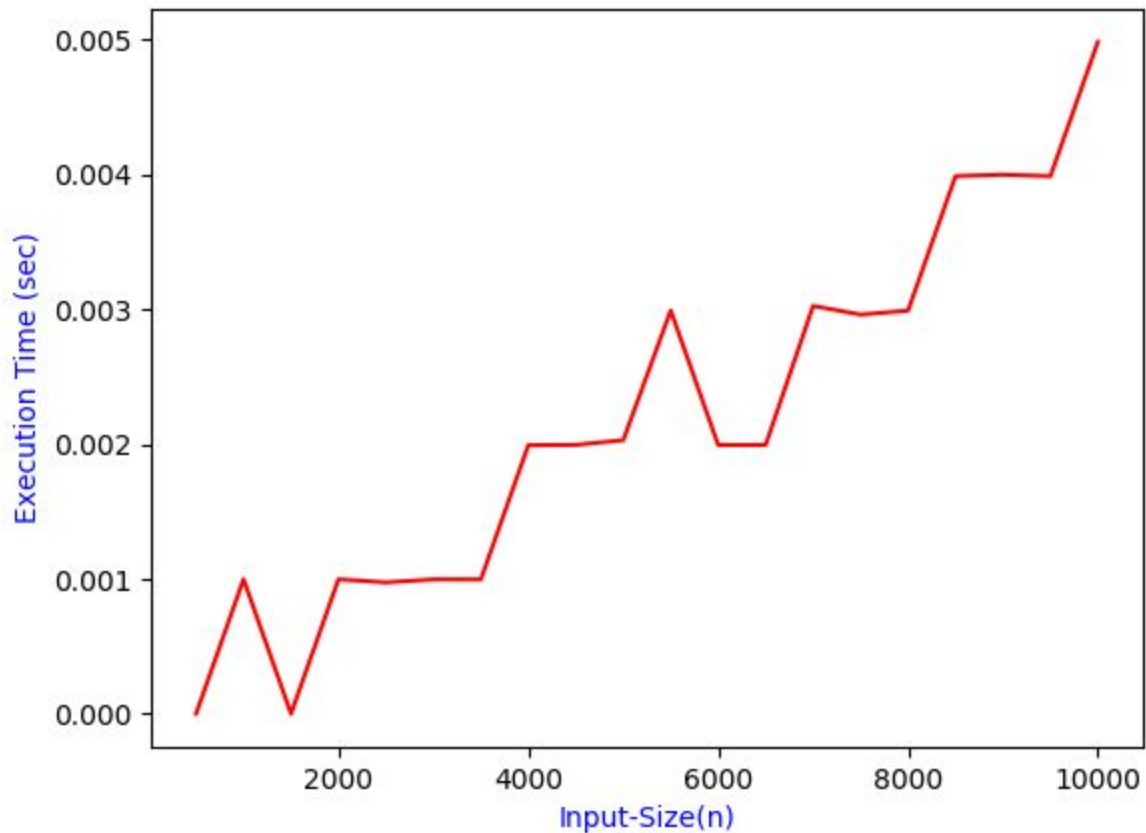


The average case is also quadratic, which makes insertion sort impractical for sorting large arrays.

❖ Best Case :  $O(n)$

| Input Size(n) | Execution Time(sec)   |
|---------------|-----------------------|
| 500           | 0.0                   |
| 1000          | 0.0009982585906982422 |
| 1500          | 0.0                   |
| 2000          | 0.000997781753540039  |
| 2500          | 0.0009748935699462891 |
| 3000          | 0.0009970664978027344 |
| 3500          | 0.0009970664978027344 |
| 4000          | 0.0019927024841308594 |
| 4500          | 0.0019948482513427734 |
| 5000          | 0.002029895782470703  |
| 5500          | 0.002991914749145508  |
| 6000          | 0.001993894577026367  |
| 6500          | 0.0019943714141845703 |
| 7000          | 0.0030260086059570312 |
| 7500          | 0.002961397171020508  |
| 8000          | 0.0029909610748291016 |
| 8500          | 0.003988742828369141  |
| 9000          | 0.003999471664428711  |
| 9500          | 0.003988504409790039  |

Graph:  $O(n)$



The best case input is an array that is already sorted. In this case insertion sort has a linear running time (i.e.,  $O(n)$ ). During each iteration, the first remaining element of the input is only compared with the right-most element of the sorted subsection of the array.

## Merge Sort

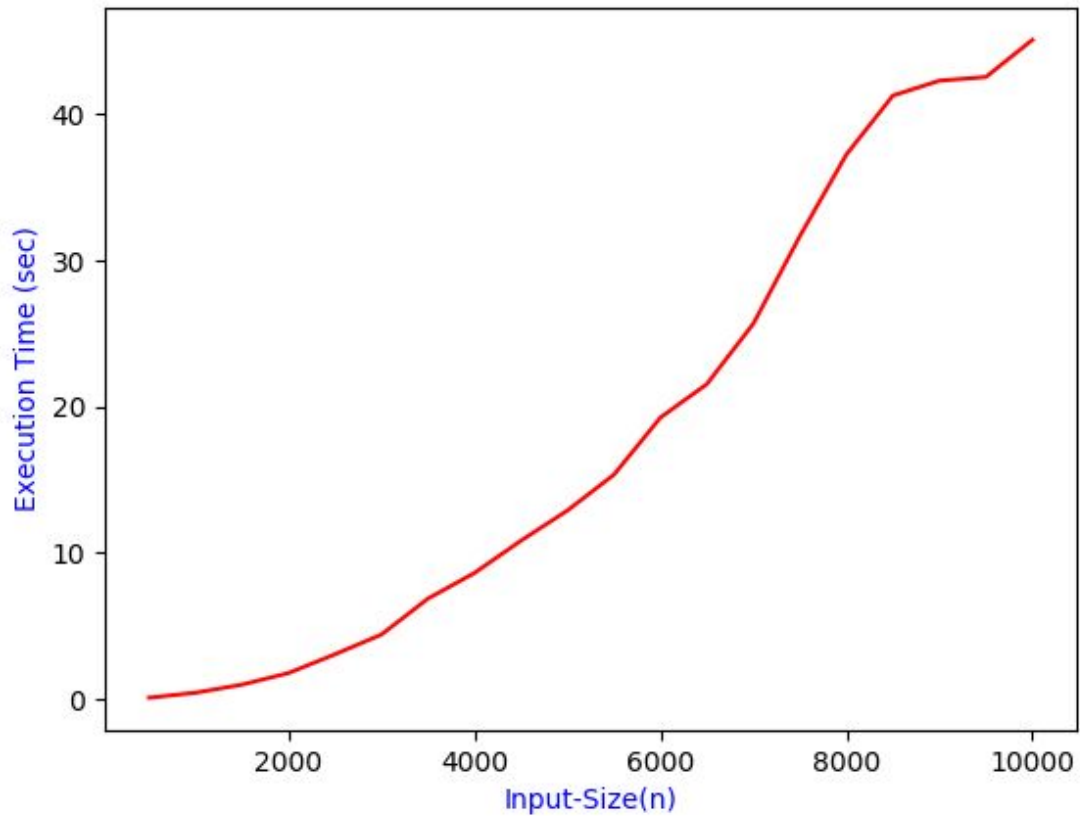
Merge Sort is a Divide and Conquer algorithm. It divides the input array into two halves, calls itself for the two halves and then merges the two sorted halves.

❖ Worst Case:  $O(n \log n)$

| Input Size(n) | Execution Time(sec) |
|---------------|---------------------|
| 500           | 0.10771346092224121 |
| 1000          | 0.4497954845428467  |
| 1500          | 1.0042991638183594  |
| 2000          | 1.794447422027588   |
| 2500          | 3.0780818462371826  |
| 3000          | 4.4286789894104     |
| 3500          | 6.883840084075928   |
| 4000          | 8.637100458145142   |
| 4500          | 10.845529794692993  |
| 5000          | 12.909912824630737  |
| 5500          | 15.355231046676636  |
| 6000          | 19.261536359786987  |
| 6500          | 21.552732467651367  |
| 7000          | 25.685041427612305  |
| 7500          | 31.6725409030914    |
| 8000          | 37.24225950241089   |
| 8500          | 41.27059197425842   |
| 9000          | 42.287206411361694  |
| 9500          | 42.54072284698486   |



Graph:  $O(n \log n)$

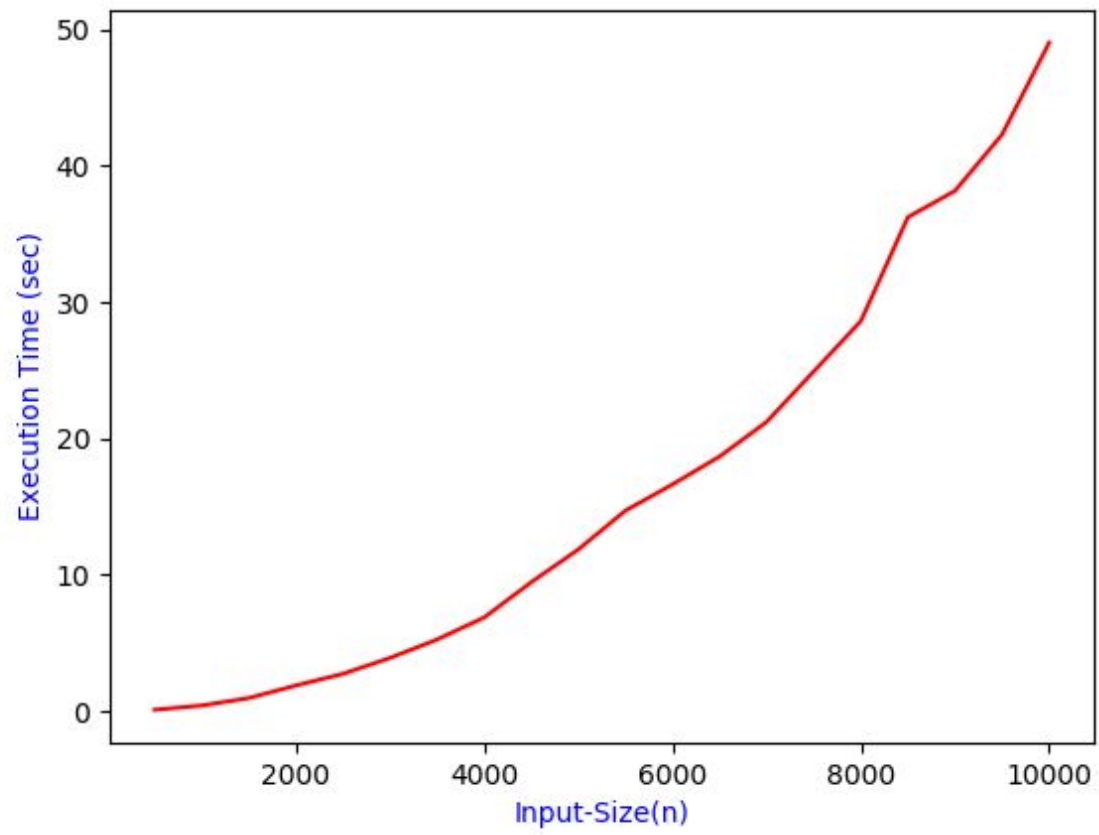


In sorting  $n$  objects, merge sort has an average and worst-case performance of  $O(n \log n)$ . If the running time of merge sort for a list of length  $n$  is  $T(n)$ , then the recurrence  $T(n) = 2T(n/2) + n$  follows from the definition of the algorithm.

❖ Average Case:  $O(n \log n)$

| Input Size(n) | Execution Time(sec) |
|---------------|---------------------|
| 500           | 0.10932183265686035 |
| 1000          | 0.42177653312683105 |
| 1500          | 0.9685547351837158  |
| 2000          | 1.8891375064849854  |
| 2500          | 2.7453551292419434  |
| 3000          | 3.9080986976623535  |
| 3500          | 5.260749340057373   |
| 4000          | 6.873384475708008   |
| 4500          | 9.459927558898926   |
| 5000          | 11.863524913787842  |
| 5500          | 14.70813512802124   |
| 6000          | 16.6305570602417    |
| 6500          | 18.677709102630615  |
| 7000          | 21.21852684020996   |
| 7500          | 24.914501190185547  |
| 8000          | 28.60091233253479   |
| 8500          | 36.22527456283569   |
| 9000          | 38.156410932540894  |
| 9500          | 42.286498069763184  |

Graph:  $O(n \log n)$



❖ Best Case:  $O(n \log n)$

| Input Size(n) | Execution Time(sec) |
|---------------|---------------------|
| 500           | 0.0997316837310791  |
| 1000          | 0.4200928211212158  |
| 1500          | 1.0856542587280273  |
| 2000          | 2.2834055423736572  |
| 2500          | 5.587977409362793   |
| 3000          | 5.988343000411987   |
| 3500          | 5.91169548034668    |
| 4000          | 7.348695516586304   |
| 4500          | 9.177350759506226   |
| 5000          | 11.173984050750732  |
| 5500          | 13.889242887496948  |
| 6000          | 18.092814922332764  |
| 6500          | 20.00065588951111   |
| 7000          | 23.773826599121094  |
| 7500          | 27.113330841064453  |
| 8000          | 28.393465995788574  |
| 8500          | 32.80443024635315   |
| 9000          | 37.42474818229675   |
| 9500          | 44.60211253166199   |

Graph:  $O(n \log n)$

