

CS 381
HW 7
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1.
(a)

1. Construct current residual network // $O(E)$
2. Use Edmonds Karp algorithm to construct max flow and final residual network // $O(VE^2)$
3. Do DFS on residual network, make reachable vertices from source // $O(V+E)$
4. A minimum cut consists of all edges from reachable vertex to non-reachable vertex.
// $O(E)$

Time: Edmonds Karp algorithm dominates the time.
 $O(VE^2)$

(b)

1. Construct current residual network // $O(E)$
2. Use Edmonds Karp algorithm to construct max flow and final residual network // $O(VE^2)$
3. Do DFS on residual network, make reachable vertices from source // $O(V+E)$
4. A minimum cut C consists of all edges from reachable vertex to non-reachable vertex.
// $O(E)$
5. For all edges (u_i, v_i) in G // $O(E)$
 If (u_i, v_i) does not exist in residual network && minimum cut C
 // New saturated path from source to sink
 Print first minimum cut which includes (u_i, v_i) , second minimum cut C
Print the only minimum cut C // $O(1)$

Time: Edmonds Karp algorithm dominates the time.
 $O(VE^2)$

2.

Given:

S: set of n points

Input line: $y = mx + b$

Distance from point to line:

$$\text{distance}(ax + by + c = 0, (x_0, y_0)) = \frac{|ax_0 + by_0 + c|}{\sqrt{a^2 + b^2}}$$

Pseudocode:

Initial array Distance []

// determine the point is in left or right side of $y = mx + b$

For each point (x_i, y_i) // O(n)

 If $mx_i + b > y_i$

 Distance[i] = -1 // left side

 If $mx_i + b = y_i$

 Distance[i] = 0

 Else

 Distance[i] = 1 // right side

For each point (x_i, y_i)

 Distance[i] = Distance[i] * Distance from point to line $((x_i, y_i), y = mx + b)$

Find the median k of Distance array

// O (n), use the method that introduced in class to

// separate array in group of five entries to find

// median of an array in O(n) time

Take line $y = mx + n$ with (x_k, y_k) to get the number of n // O (1)

Return $y = mx + n$

Time: The algorithm to find the median of Distance array takes O (n) time dominating the whole algorithm.

So, time is O(n)

