CSE 604 Artificial Intelligence

Chapter 5: Adversarial Search

Adapted from slides available in Russell & Norvig's textbook webpage

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Outline

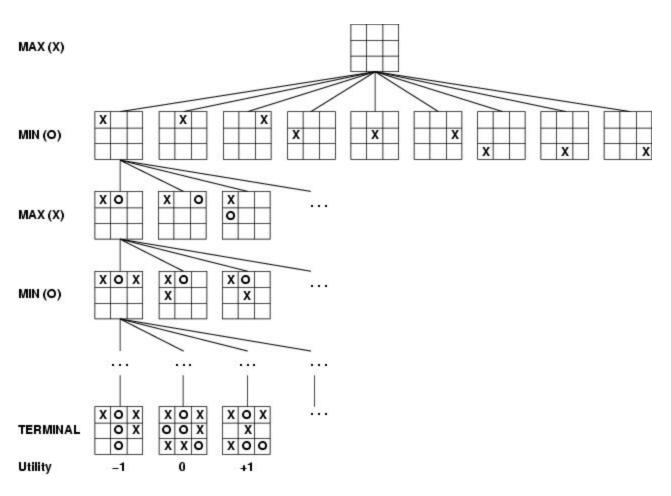
- Optimal decisions
- α-β pruning
- Imperfect, real-time decisions

Games vs. search problems

• "Unpredictable" opponent specifying a move for every possible opponent reply

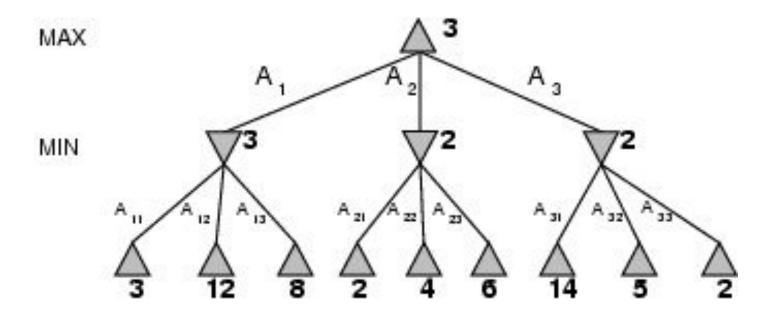
• Time limits unlikely to find goal, must approximate

Game tree (2-player, deterministic, turns)



Minimax

- Perfect play for deterministic games
- Idea: choose move to position with highest minimax value = best achievable payoff against best play
- E.g., 2-ply game:

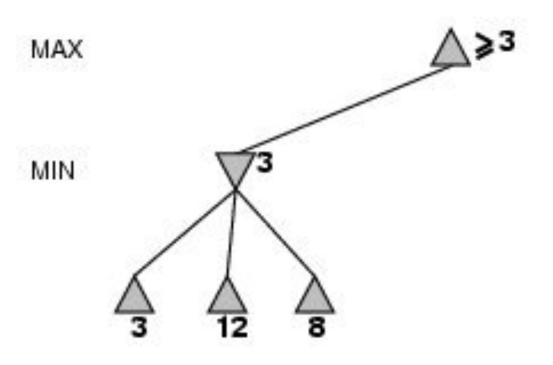


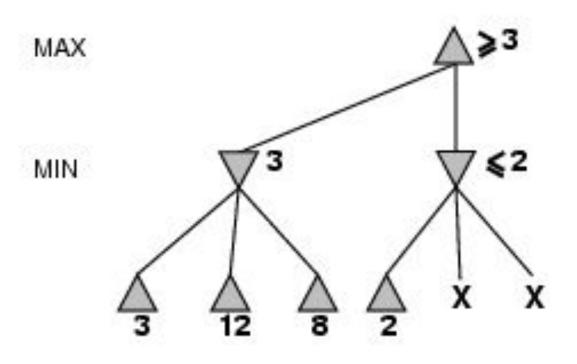
Minimax algorithm

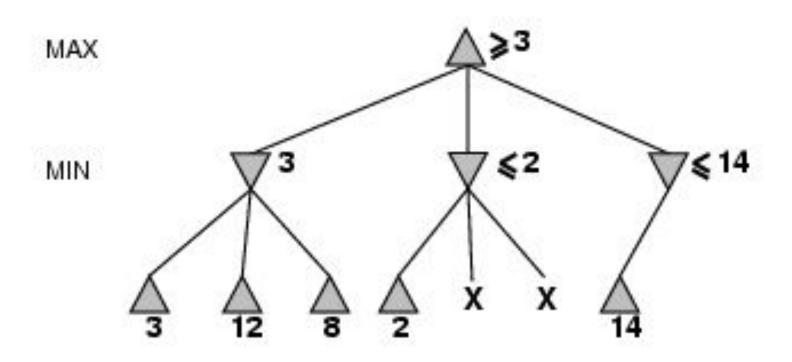
```
function Minimax-Decision(state) returns an action
   v \leftarrow \text{Max-Value}(state)
   return the action in Successors(state) with value v
function Max-Value(state) returns a utility value
   if Terminal-Test(state) then return Utility(state)
   v \leftarrow -\infty
   for a, s in Successors(state) do
      v \leftarrow \text{Max}(v, \text{Min-Value}(s))
   return v
function Min-Value(state) returns a utility value
   if Terminal-Test(state) then return Utility(state)
   v \leftarrow \infty
   for a, s in Successors(state) do
      v \leftarrow \text{MIN}(v, \text{MAX-VALUE}(s))
   return v
```

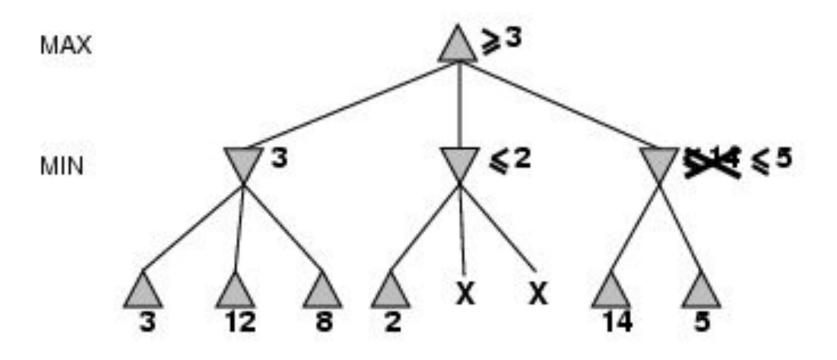
Properties of minimax

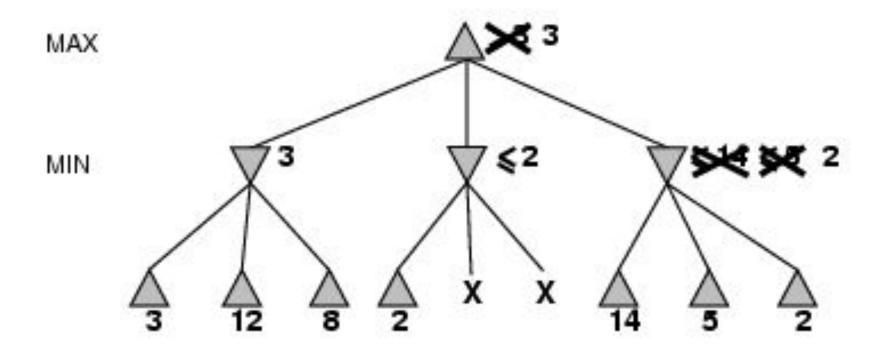
- <u>Complete?</u> Yes (if tree is finite)
- Optimal? Yes (against an optimal opponent)
- <u>Time complexity?</u> O(b^m)
- Space complexity? O(bm) (depth-first exploration)
- For chess, $b \approx 35$, $m \approx 100$ for "reasonable" games exact solution completely infeasible











Properties of α - β

- Pruning does not affect final result
- Good move ordering improves effectiveness of pruning
 - Try the moves that are "likely to be best" first
 - E.g., in chess, try captures, threats, forward moves, backward moves in that order
- With "perfect ordering," time complexity $\approx O(b^{m/2})$ doubles depth of search
- A simple example of the value of reasoning about which computations are relevant (a form of metareasoning)

Why is it called α - β ?

• α is the value of the best (i.e., highest-value) choice found so far at any choice point along the path for *max*

If v is worse than α, max
 will avoid it
 prune that branch

• Define β similarly for *min*

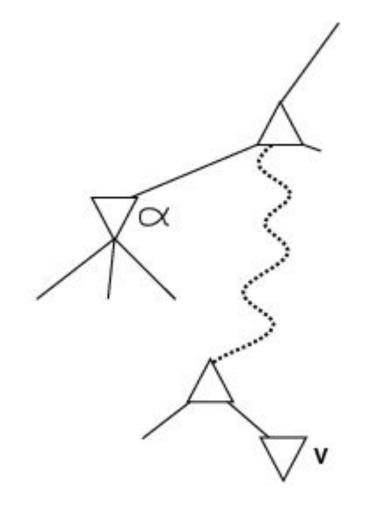
MAX

MIN

..

MAX

MIN



The α - β algorithm

```
function Alpha-Beta-Search(state) returns an action
   inputs: state, current state in game
   v \leftarrow \text{MAX-VALUE}(state, -\infty, +\infty)
   return the action in Successors(state) with value v
function MAX-VALUE(state, \alpha, \beta) returns a utility value
   inputs: state, current state in game
             \alpha, the value of the best alternative for MAX along the path to state
             eta, the value of the best alternative for MIN along the path to state
   if TERMINAL-TEST(state) then return UTILITY(state)
   v \leftarrow -\infty
   for a, s in Successors(state) do
      v \leftarrow \text{Max}(v, \text{Min-Value}(s, \alpha, \beta))
      if v \geq \beta then return v
      \alpha \leftarrow \text{Max}(\alpha, v)
   return v
```

The α - β algorithm

```
function MIN-VALUE(state, \alpha, \beta) returns a utility value inputs: state, current state in game \alpha, the value of the best alternative for MAX along the path to state \beta, the value of the best alternative for MIN along the path to state if Terminal-Test(state) then return Utility(state) v \leftarrow +\infty for a, s in Successors(state) do v \leftarrow \text{Min}(v, \text{Max-Value}(s, \alpha, \beta)) if v \leq \alpha then return v \beta \leftarrow \text{Min}(\beta, v) return v
```

Imperfect Real Time Decisions

Even with alpha-beta pruning, it is infeasible to grow the whole game tree!

Standard approach:

- evaluation function
 - = estimated desirability of position
- cut off search
 - e.g., depth limit or iterative deepening
- forward pruning
 - e.g., Beam search

Evaluation functions

• For chess, typically linear weighted sum of features

$$Eval(s) = w_1 f_1(s) + w_2 f_2(s) + ... + w_n f_n(s)$$

• e.g., $w_1 = 9$ with $f_1(s) = \text{(number of white queens)} - \text{(number of black queens)},$ $w_2 = 5$ with $f_2(s) = \text{(number of white rooks)} - \text{(number of black rooks)},$ etc.

Cutting off search

- We can use a modified algorithm *MinimaxCutoff*
- MinimaxCutoff is identical to MinimaxValue except
 - Terminal? is replaced by Cutoff?
 - *Utility* is replaced by *Eval*
- Does it work in practice?
 - Suppose we have 100 secs, explore 10⁴ nodes/sec
 10⁶ nodes per move

$$b^{m} = 10^{6}, b=35 m=4$$

- 4-ply lookahead is a hopeless chess player!
 - 4-ply ≈ human novice
 - 8-ply ≈ typical PC, human master
 - 12-ply ≈ Deep Blue, Kasparov

Deterministic games in practice

- **Checkers**: Chinook ended 40-year-reign of human world champion Marion Tinsley in 1994.
- Chess: Deep Blue defeated human world champion Garry Kasparov in a six-game match in 1997. Deep Blue searches 200 million positions per second, uses very sophisticated evaluation, and undisclosed methods for extending some lines of search up to 40 ply.
- Othello: human champions refuse to compete against computers, who are too good.
- Go: Until recently, human champions refused to compete against computers, who were too bad (in Go, b > 300).
 - But in 2016, Google's AlphaGo defeated human world champion Lee Sedol.
 - In 2017, AlphaGo Zero defeated the previous version of AlphaGo 100-0