# H# Official Documentation

The official specification for H#

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### Abstract

This document is the official document for the H-Sharp (H#) programming language. The H stands for hybrid - as this is a multi-paradigm programming language heavily inspired by C# and Scala. This document will only contain the grammar of the language, the operational semantics as well as the type semantics of the language. The operational semantics may be explained with code samples - but obvious uses will not be explained.

This document will also contain the byte-instruction semantics. The official compiler, is written in C# for productivity purposes while the Virtual Machine is implemented using C++. At no point will this document be documenting the internal processes of those applications.

#### Grammar

The offical H# grammar. Note:  $Id \in VARENV$  and  $TypeId \in TYPEENV$ . Both refer to the same grammatical definition, defined by the regular expression:

$$(|[a-Z]|^+(|[0-9]|[a-Z])^*$$

If multiple elements can occur - and it'd be convenient, the element be sufficed with 'n' to show it may be contain n of such elements.  $n \in \{0, 1, 2, \dots\}$ . The notation  $e_0, \dots, e_n$  represents the 1 to nth element with a specific separator. The full grammar is defined as follows:

```
CompileUnit
                    ::= CompileUnitElement
CompileUnitElement ::= CompileUnitElement CompileUnitElement
                     | Directive | Scope | ScopeElement | Declaration
                    ::= \{ ScopeElement \}
Scope
ScopeElement
                    ::= ScopeElement\ ScopeElement
                     | Expr; | Statement | VarDeclaration;
                    ::= Expr \mid (Expr) \mid Scope \mid LambdaExpr
Expr
                        Expr BinaryOp Expr | UnaryOp Expr | Id UnaryOp
                        Expr ? Expr : Expr
                        Id \mid Name \mid this | base
                        Expr(Argument)
                        new TypeId(Argument) | new TypeId[Argument]
                        Expr.Id \mid Expr?.Id
                        Expr.Id(Argument) \mid Expr?.Id(Argument)
                        Expr is TypeID | Expr is null Expr is TypeID Id
                        Expr is TypeID\ Id where Expr
                        Expr is not TypeID \mid Expr is not null
                        Expr is not TypeID\ Id where Expr
                        Expr as TypeId
                        (TypeID) Expr
                        sizeof(TypeId) \mid addressof(Id) \mid typeof(Expr)
                        Assignment
                        Literal
LambdaExpr
                    ::= (Param) \Rightarrow Expr
Directive
                    ::= type Id = TypeId;
                        using Name; | using TypeId from Name;
                        Name space Decl
Name space Decl
                    ::= namespace Name\ Scope
Statement
                    ::= Assignment; | ControlStatement | MatchStatement;
                        TryCatchStatement
ControlStatement
                    ::= if Expr\ Scope
                        if Expr\ Scope else Scope
                        if Expr Scope else if Expr Scope
                        if Expr\ Scope else if Expr\ Scope else Scope
                        while Expr\ Scope
                        do Scope while Expr;
                        for (Assignment; Expr; Expr) Scope
                        for (VarDeclaration; Expr; Expr) Scope
                        foreach (TypeId\ Id\ in\ Expr) Scope
                        throw TypeID(Argument)
                        return Expr;
                        break;
```

```
::= Expr match \{ MatchCase \}
MatchStatement
MatchCase
                     ::= MatchCase, MatchCase
                         \verb|case| \textit{Literal} \Rightarrow \textit{Expr} \mid \verb|case| \textit{TypeId} \Rightarrow \textit{Expr}
                         case TypeId\ Id \Rightarrow Expr \mid case TypeId\ Id\ when Expr \Rightarrow Expr
                         case (MatchCaseId) \Rightarrow Expr
                         case (MatchCaseId) when Expr \Rightarrow Expr
                         case TypeId (MatchCaseId) => Expr
                         case TypeId (MatchCaseId) when Expr \Rightarrow Expr
                         case \Rightarrow Expr
MatchCaseId
                     ::= MatchCaseId, MatchCaseId
                      \mid Id \mid
TryCatchStatement ::= try Scope catch (TypeId Id) Scope
                      try Scope catch (TypeId Id) when Expr Scope
                         try Scope catch (TypeId Id) Scope finally Scope
                         try Scope catch (TypeId Id) when Expr Scope finally Scope
                     ::= Id = Expr
Assignment
                       Id += Expr \mid Id -= Expr
                         Id *= Expr \mid Id /= Expr
                         Id \&= Expr \mid Id \mid= Expr
                         Id %= Expr
Declaration
                     ::= VarDecl \mid FuncDecl
                         ClassDecl \mid StaticClassDecl \mid InterfaceDecl
                         UnionDecl \mid EnumDecl \mid StructDecl
ClassDecl
                     ::= class Id \{ ClassMember \}
                         AccessMod class Id { ClassMember }
                         StorageMod class Id { ClassMember }
                         AccessMod StorageMod class Id { ClassMember }
                         class Id(Param) { ClassMember }
                         AccessMod class Id(Param) { ClassMember }
                         StorageMod class Id(Param) { ClassMember }
                         AccessMod StorageMod class Id(Param) { ClassMember }
StructDecl
                     ::= struct Id { ClassMember }
                         AccessMod struct Id { ClassMember }
                         StorageMod struct Id { ClassMember }
                         AccessMod StorageMod struct Id { ClassMember }
                         struct Id(Param) { ClassMember }
                         AccessMod struct Id(Param) { ClassMember }
                         StorageMod struct Id(Param) { ClassMember }
                         AccessMod StorageMod struct Id(Param) { ClassMember }
StaticClassDecl\\
                     ::= object Id { ClassMember }
                         AccessMod object Id { ClassMember }
                         object Id(Param) { ClassMember }
                         AccessMod object Id(Param) { ClassMember }
                     ::= interface Id { ClassMember }
Interface Decl
                      | AccessMod interface Id { ClassMember }
ClassMember
                     ::= ClassMember\ ClassMember
                         Id(Param) FuncBody | AcessMod Id(Param) FuncBody
                         VarDecl; | AcessMod VarDecl;
                         FuncDecl \mid ClassDecl \mid UnionDecl \mid EnumDecl
                         event TypeId id; | AccessMod event TypeId id;
```

VarDecl $::= TypeId \ Id = Expr \mid StorageMod \ TypeId \ Id = Expr$  $var\ Id = Expr \mid StorageMod\ var\ Id = Expr$ TypeId Id | StorageMod TypeId Id  $LambdaType\ Id = LambdaExpr$ TupleDecl::= (ParamType) Id = (Argument)TupleDecl| (Param) = (Argument)FuncDecl::= Id(FuncParam): TypeId FuncBody AccessMod Id(FuncParam): TypeId FuncBody Id = (FuncParam): TypeId FuncBodyAccessMod Id = (FuncParam): TypeId FuncBody AccessMod const Id = (FuncParam): TypeId FuncBodyFuncBody $::= Scope \mid \Rightarrow Expr$ FuncParam::= FuncParam, FuncParam | Param  $| TypeId\ Id = Literal$ UnionDecl::= union  $Id \{ UnionMember \}$ AccessMod union Id { UnionMember } AccessMod static union Id { UnionMember }  $::= \ Union Member \ Union Member$ UnionMember| TypeId Id;EnumDecl $::= \text{enum } Id \{ EnumBodyMember \}$  $AccessMod enum Id \{ EnumBodyMember \}$ enum  $Id(EnumMember) \{ EnumBodyMember \}$ AccessMod enum Id (EnumMember) { EnumBodyMember } EnumBody Member $::= EnumMember \mid FuncDecl \mid FuncDecl \mid EnumBodyMember$ ::= EnumMember, EnumMemberEnumMember $\mid Id \mid Id = LiteralNoNull$ StorageMod::= const | constexpr | static | abstract | override | virtual | final | lazy AccessMod::= public | private | protected | internal | external ::= Param , ParamParam $| TypeId\ Id\ |$  const  $TypeId\ Id$  $ParamType\ Id$  $::= TypeId \mid TypeId$ , TypeIdParamType $::= Expr \mid Expr$ , ExprArgument

 $::= Id \mid Name.Name$ 

Name

LambdaType ::= (ParamType): TypeId

| TypeId : TypeId

BinaryOp ::= + | - | \* | / | % | < | > | <= | == | !=

UnaryOp ::= - | ! | # | ++ | --

 $\textit{LiteralNoNull} \qquad ::= IntLit \mid FloatLit \mid DoubleLit \mid BoolLit \mid CharLit \mid StringLit$ 

Literal ::=  $LiteralNoNull \mid NullLit$ 

IntLit ::=  $Digit^+$ 

 $FloatLit ::= Digit^+ . Digit^+ f$ 

 $DoubleLit ::= Digit^+ . Digit^+$ 

 $CharLit ::= `Letter', | `\Letter'$ 

StringLit ::= "(Letter|Digit)\*"

BoolLit ::= true | false

NullLit ::= null

## **Operational Semantics**

$$\begin{array}{ll} \text{VariableLookup} & \text{HeapObjectLookup} \\ \underline{\rho,\mu,\phi,\kappa,\sigma\vdash\rho(x)=v\neq(\ell,\omega,\sigma)} & \underline{\rho,\mu,\phi,\kappa,\sigma\vdash\rho(x)=(\ell,\omega,\sigma)} & \sigma(\ell)=v & \omega=0 \\ \hline \\ \rho,\mu,\phi,\kappa,\sigma\vdash x\Rightarrow v,\sigma & \underline{\rho,\mu,\phi,\kappa,\sigma\vdash\rho(x)=(\ell,\omega,\sigma)} & \sigma(\ell)=v & \omega=S \\ \hline \\ \underline{\rho,\mu,\phi,\kappa,\sigma\vdash\rho(x)=(\ell,\omega,\sigma)} & \sigma(\ell)=v & \omega=S \\ \hline \\ \rho,\mu,\phi,\kappa,\sigma\vdash x\Rightarrow v,\sigma & \\ \hline \\ \text{HeapArrayLookup} \\ \underline{\rho,\mu,\phi,\kappa,\sigma\vdash\rho(x)=(\ell,\omega,\sigma)} & \sigma(\ell)=v & \omega=A \\ \hline \\ \rho,\mu,\phi,\kappa,\sigma\vdash x\Rightarrow v,\sigma & \\ \hline \end{array}$$

## **Type Semantics**

 $\theta = TypeEnv$ 

 $\gamma = TypeLookupEnv$ 

 $\eta = ReferenceTypeEnv$ 

 $\theta, \gamma, \eta \vdash Expr : Type$ 

 $\theta, \gamma, \eta \vdash Decl : \theta, \gamma$ 

 $\theta, \gamma, \eta \vdash Decl : \theta, \gamma, \eta$ 

int <: float <: double <: INumeric</pre>

T-IntLit  $i\in\mathbb{N}$  $\theta, \overline{\gamma, \eta} \vdash i \, : \, \mathtt{int}$  T-Addition

 $\frac{\theta, \gamma, \eta \vdash e_1 : \tau_1 \qquad \theta, \gamma, \eta \vdash e_2 : \tau_2 \qquad \qquad \tau <: \texttt{INumeric}}{\theta, \gamma, \eta \vdash e_1 + e_2 :}$ 

T-DeclVar-NonRef

 $\frac{\theta, \gamma, \eta \vdash e : \tau \quad \tau \notin \eta \quad \theta' = \theta[x \mapsto \tau'] \quad \tau' = \mathsf{typeof}(t, \gamma) \quad \tau <: \tau'}{\theta, \gamma, \eta \vdash t \; x = e : \theta', \gamma, \eta}$ 

 $\operatorname{T-DeclVar-Ref}$ 

 $\frac{\theta, \gamma, \eta \vdash e : \tau \qquad \tau \in \eta \qquad \theta' = \theta[x \mapsto \mathsf{Ref}(\tau')] \qquad \tau' = \mathsf{typeof}(t, \gamma) \qquad \tau <: \tau'}{\theta, \gamma, \eta \vdash t \ x = e : \theta', \gamma, \eta}$ 

T-NewObject

 $\frac{\theta, \gamma, \eta \vdash e_1, \dots, e_n \ : \ \tau_1, \dots, \tau_n \qquad \tau = \gamma(t)}{\theta, \gamma, \eta \vdash \text{new } t(e_1, \dots, e_n) \ : \ \tau}$ 

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# Compilation Semantics

# Bytecode Semantics