HOMEWORK I

Due day: 3:00pm Oct. 14 (Wednesday), 2020

Introduction

This homework is to let you be familiar with SystemVerilog and CPU design.

After this homework, you will complete a simplified pipeline CPU with <u>33</u> instructions.

General rules for deliverables

- This homework needs to be completed by **INDIVIDUAL** student.
- Compress all files described in the problem statements into one tar file.
- Submit the compressed file to the course website before the due day.
 Warning! AVOID submitting in the last minute. Late submission is not accepted.

Grading Notes

- Important! DO remember to include your SystemVerilog code. NO code, NO grades. Also, if your code can not be recompiled by TA successfully using tools in SoC Lab and commands in Appendix B, you will receive NO credit.
- Write your report seriously and professionally. Incomplete description and information will reduce your chances to get more credits.
- If extra works (like synthesis, post-simulation or additional instructions) are done, please describe them in your final report clearly for bonus points.
- Please follow course policy.
- Verilog and SystemVerilog generators aren't allowed in this course.

Deliverables

- 1. All SystemVerilog codes including components, testbenches and machine codes for each problem. NOTE: Please DO NOT include source codes in the report!
- 2. Write a homework report in MS word and follow the convention for the file name of your report: *N260XXXXX.docx*. Please save as docx file format and replace N260XXXXX with your student ID number. (Let the letter be uppercase.)
- 3. Organize your files as the hierarchy in Appendix A.

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Report Writing Format

- a. Use the submission cover which is already in provided N260XXXXX.docx.
- b. A summary in the beginning to state what has been done.
- c. Report requirements from each problem.
- d. Describe the major problems you encountered and your resolutions.
- e. Lessons learned from this homework.

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Problem 1 (100/100)

1.1 Problem Description

In this course, we use the RISC-V instruction set architecture (ISA) to implement a pipelined CPU. The RISC-V ISA is composed of one base integer instruction set and some standard extensions. For simplicity of implementation, the homework of this course focuses on the RV32I base integer instruction set.

You need to implement a simplified pipeline CPU with the following features:

- a. The RISC-V ISA with the specified 33 instructions.
- b. The number of pipeline stage is 5.
- c. Register file size: 32×32-bit.
- d. Program counter with 32-bit.
- e. Mechanism to solve data hazard, control hazard and structural hazard.

You also need to use two memories outside the CPU with specified size:

- a. Instruction memory size: 64KB.
- b. Data memory size: 64KB.

Your RTL code should comply with Superlint within <u>85%</u> of your code. Besides, you should use programs listed in Section 1.5 to verify your design. Note that you <u>DO NOT</u> need to synthesize your design. A more detailed description of this problem can be found in Section 1.4.

1.2 Block Overview

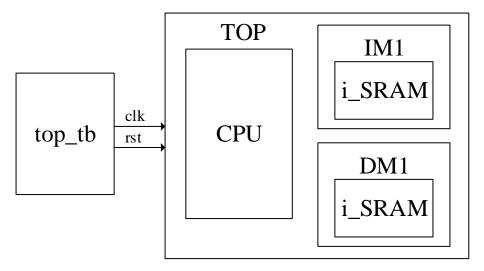


Fig. 1-1: System block diagram

Program Counter Instruction Memory (Outside CPU) Immediate Generator HOMEWORK I Register File ALLU ALU Outside CPU) Outside CPU)

Fig. 1-2: A CPU architecture for reference

1.3 Module Specification

Name Category File Module **SDF Instance RTL** TOP top.sv top **RTL** SRAM_wrapper.sv SRAM_wrapper IM1 RTL DM1 SRAM wrapper.sv SRAM_wrapper RTL SRAM_rtl.sv **SRAM** i SRAM

Table 1-1: Module naming rule

CD 11	1 0	3 / 1	1	•	1
Table	1-7.	Modu	II e	ciona	IC
Table	1-4.	MOGU	-	SIZHA	ıο

Module	Specifications				
	Name	Signal	Bits	Function explanation	
top	clk	input	1	System clock	
	rst	input	1	System reset (active high)	
	CK	input	1	System clock	
	CS	input	1	Chip select (active high)	
	OE	input	1	Output enable (active high)	
SRAM_wrapper	WEB	input	4	Write enable (active low)	
	A	input	14	Address	
	DI	DI input		Data input	
	DO	output	32	Data output	

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Module	Name	Signal	Bits	Function explanation			
	Memory Space						
	Memory_byte0	logic	8	Size: [16384]			
SRAM	Memory_byte1	logic	8	Size: [16384]			
	Memory_byte2	logic	8	Size: [16384]			
	Memory_byte3	logic	8	Size: [16384]			

1.4 Detailed Description

Fig. 1-1 is a system overview of this problem. You SHOULD NOT modify any port declarations or your design may have error when TA runs the hidden testbench.

Fig. 1-2 is a block diagram of a simplified 5-stage pipeline CPU. Note that this diagram shows only some possible components and signals of a pipeline CPU, and you may need to add others, e.g. a controller and its signals, for your design. You can also develop your own architecture. The only restriction is that your architecture SHOULD have 5 stages of pipeline.

You should implement the instructions in Table 1-3. The detailed instruction types and the immediate formats can be found in Appendix C. You can also study *The RISC-V Instruction Set Manual* posted on the course website.

Table 1-3: Instruction lists

R-type

31 25	24 20	19 15	14 12	11 7	6 0		
funct7	rs2	rs1	funct3	rd	opcode	Mnemonic	Description
0000000	rs2	rs1	000	rd	0110011	ADD	rd = rs1 + rs2
0100000	rs2	rs1	000	rd	0110011	SUB	rd = rs1 - rs2
0000000	rs2	rs1	001	rd	0110011	SLL	$rd = rs1_u \le rs2[4:0]$
0000000	rs2	rs1	010	rd	0110011	SLT	$rd = (rs1_s < rs2_s)$? 1:0
0000000	rs2	rs1	011	rd	0110011	SLTU	$rd = (rs1_u < rs2_u)$? 1:0
0000000	rs2	rs1	100	rd	0110011	XOR	$rd = rs1 ^ rs2$
0000000	rs2	rs1	101	rd	0110011	SRL	$rd = rs1_u >> rs2[4:0]$
0100000	rs2	rs1	101	rd	0110011	SRA	$rd = rs1_s >> rs2[4:0]$
0000000	rs2	rs1	110	rd	0110011	OR	$rd = rs1 \mid rs2$
0000000	rs2	rs1	111	rd	0110011	AND	rd = rs1 & rs2

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F I-type

31	20	19 15	14 12	11 7	6 0		
imm[11:	0]	rs1	funct3	rd	opcode	Mnemonic	Description
imm[11:	0]	rs1	010	rd	0000011	LW	rd = M[rs1+imm]
imm[11:	0]	rs1	000	rd	0010011	ADDI	rd = rs1 + imm
imm[11:	0]	rs1	010	rd	0010011	SLTI	$rd = (rs1_s < imm_s)? 1:0$
imm[11:	0]	rs1	011	rd	0010011	SLTIU	$rd = (rs1_u < imm_u)? 1:0$
imm[11:	0]	rs1	100	rd	0010011	XORI	rd = rs1 ^ imm
imm[11:	0]	rs1	110	rd	0010011	ORI	rd = rs1 imm
imm[11:	0]	rs1	111	rd	0010011	ANDI	rd = rs1 & imm
imm[11:	0]	rs1	000	rd	0000011	LB	$rd = M[rs1+imm]_{bs}$
0000000	shamt	rs1	001	rd	0010011	SLLI	$rd = rs1_u \ll shamt$
0000000	shamt	rs1	101	rd	0010011	SRLI	$rd = rs1_u >> shamt$
0100000	shamt	rs1	101	rd	0010011	SRAI	$rd = rs1_s >> shamt$
							rd = PC + 4
imm[11:	0]	rs1	000	rd	1100111	JALR	PC = imm + rs1
							(Set LSB of PC to 0)

S-type

31 25	24 20	19 15	14 12	11 7	6 0		
imm[11:5]	rs2	rs1	funct3	imm[4:0]	opcode	Mnemonic	Description
imm[11:5]	rs2	rs1	010	imm[4:0]	0100011	SW	M[rs1+imm] = rs2
imm[11:5]	rs2	rs1	000	imm[4:0]	0100011	SB	$M[rs1+imm]_b = rs2_b$

[®] B-type

31 25	24 20	19 15	14 12	11 7	6 0		
imm[12 10:5]	rs2	rs1	funct3	imm[4:1 11]	opcode	Mnemonic	Description
imm[12 10:5]	rs2	rs1	000	imm[4:1 11]	1100011	BEQ	PC = (rs1 == rs2)? PC + imm: PC + 4
imm[12 10:5]	rs2	rs1	001	imm[4:1 11]	1100011	BNE	PC = (rs1 != rs2)? PC + imm: PC + 4
imm[12 10:5]	rs2	rs1	100	imm[4:1 11]	1100011	BLT	$PC = (rs1_s < rs2_s)?$ $PC + imm: PC + 4$
imm[12 10:5]	rs2	rs1	101	imm[4:1 11]	1100011	BGE	$PC = (rs1_s \ge rs2_s)?$ $PC + imm: PC + 4$
imm[12 10:5]	rs2	rs1	110	imm[4:1 11]	1100011	BLTU	$PC = (rs1_u < rs2_u)?$ $PC + imm: PC + 4$
imm[12 10:5]	rs2	rs1	111	imm[4:1 11]	1100011	BGEU	$PC = (rs1_u \ge rs2_u)?$ $PC + imm: PC + 4$

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U-type

31 12	11 7	6 0		
imm[31:12]	rd	opcode	Mnemonic	Description
imm[31:12]	rd	0010111	AUIPC	rd = PC + imm
imm[31:12]	rd	0110111	LUI	rd = imm

J-type

31 12	11 7	6 0		
imm[20 10:1 11 19:12]	rd	opcode	Mnemonic	Description
imm[20 10:1 11 19:12]	rd	1101111	JAL	rd = PC + 4
mm[20 10.1 11 19.12]	Iu	1101111	JAL	PC = PC + imm

The data size of your design SHOULD BE 32 BITS. You need to implement a register file with 32 registers. There are 31 general-purpose registers x1–x31. Note that the register x0 is hardwire to the constant 0. In addition to the register file, you also need to implement a 32-bit program counter. Every memory address should be four-byte aligned, i.e., the program counter and the load/store address should be multiple of 4. You should implement the mechanism to avoid hazards, e.g. forwarding or hazard detection.

The size of the instruction memory is 16K×4-byte (64KB), and the size of the data memory is 16K×4-byte (64KB). You should use SRAM_wrapper as instruction memory (IM1) and data memory (DM1). You SHOULDN'T modify Verilog code in SRAM wrapper and SRAM.

Your RTL code needs to comply with Superlint within <u>85%</u> of your code, i.e., the number of errors & warnings in total shall not exceed 15% of the number of lines in your code. HINT: You can use the command in Appendix B to get the number of lines in *src* and *include* directories.

1.5 Verification

You should complete following programs and use the commands in Appendix B to verify your design.

- a. Use $prog\theta$ to perform verification for the functionality of instructions.
- b. Write a program defined as *prog1* to perform a sort algorithm. The number of sorting elements is stored at the address named *array_size* in ".rodata" section defined in *data.S*. The first element is stored at the address named *array_addr* in ".rodata" section defined in *data.S*, others are stored at adjacent addresses. The maximum number of elements is 64. All elements are **signed 4-byte integers** and you should sort them in **ascending order**.

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Rearranged data should be stored at the address named _test_start in "_test" section defined in link.ld.

- c. Write a program defined as *prog2* to perform multiplication. The multiplicand is stored at the address named *mul1* in ".rodata" section defined in *data.S*. The multiplier is stored at the address named *mul2* in ".rodata" section defined in *data.S*. The multiplicand and the multiplier are **signed 4-byte integers**. Their product is **signed 8-byte integers** and should be stored at the address named *_test_start* in "_test" section defined in *link.ld*.
- d. Write a program defined as *prog3* to perform greatest common divisor(GCD). The first number is stored at the address named *div1* in ".rodata" section defined in *data.S*. The second data is stored at the address named *div2* in ".rodata" section defined in *data.S*. These two numbers are **unsigned 4-byte integers**. The result should be stored at the address named _*test_start* in "_test" section defined in *link.ld*. The values of the quotient and the remainder should follow C99 specification. "When integers are divided, the result of the / operator is the algebraic quotient with any fractional part discarded. If the quotient **a/b** is representable, the expression (**a/b**)***b** + **a**%**b** shall equal **a**.

Don't forget to return from *main* function to finish the simulation in each program. Save your assembly code or C code as *main.S* or *main.c* respectively. You should also explain the result of this program in the report. In addition to these verification, **TA will use another program to verify your design.** Please make sure that your design can execute the listed instructions correctly.

1.6 Report Requirements

Your report should have the following features:

- a. Proper explanation of your design is required for full credits.
- b. Block diagrams shall be drawn to depict your designs.
- c. Show your snapshots of the waveforms and the simulation results on the terminal for the different test cases in your report and illustrate the correctness of your results.
- d. Report the number of lines of your RTL code, the final results of running Superlint and 3~5 most frequent warning/errors in your code. Describe how you modify your code to comply with the Superlint.

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Appendix

A. File Hierarchy Requirements

All homework SHOULD be uploaded and follow the file hierarchy and the naming rules, especially letter case, specified below. You should create a main folder named your student ID number. It contains your homework report and other files. The names of the files and the folders are labeled in red color, and the specifications are labeled in black color. Filenames with * suffix in the same folder indicate that you should provide one of them. Before you submit your homework, you can use Makefile macros in Appendix B to check correctness of the file structure.

Fig. A-1 File hierarchy

N260XXXXX.tar (Don't add version text in filename, e.g. N260XXXXX v1.tar) *N260XXXXX* (Main folder of this homework) *N260XXXXX.docx* (Your homework report) StudentID (Specify your student ID number in this file) *Makefile* (You shouldn't modify it) *Src* (Your RTL code with *sv* format) top.sv SRAM wrapper.sv Other submodules (*.sv) include (Your RTL definition with svh format, optional) Definition files (*.svh) www. (Your synthesized code and timing file, optional) top syn.v top syn.sdf script (Any scripts of verification, synthesis or place and route) 鄶 Script files (*.sdc, *.tcl or *.setup) *sim* (Testbenches and memory libraries) top tb.sv (Main testbench. You shouldn't modify it) **CYCLE** (Specify your clock cycle time in this file) MAX (Specify max clock cycle number in this file) SRAM (SRAM libraries and behavior models) Library files (*.lib, *.db, *.lef or *.gds) *SRAM.ds* (SRAM datasheet) **SRAM** rtl.sv (SRAM RTL model) *SRAM.v* (SRAM behavior model)

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prog	g0 (Subfolder for Program 0)
	Makefile (Compile and generate memory content)
	main.S (Assembly code for verification)
	setup.S (Assembly code for testing environment setup)
	link.ld (Linker script for testing environment)
	golden.dat (Golden hexadecimal data)
prog	gl (Subfolder for Program 1)
	Makefile (Compile and generate memory content)
	main.S * (Assembly code for verification)
	main.c * (C code for verification)
	data.S (Assembly code for testing data)
	setup.S (Assembly code for testing environment setup)
	<i>link.ld</i> (Linker script for testing environment)
	golden.dat (Golden hexadecimal data)
prog	g2 (Subfolder for Program 2)
	Makefile (Compile and generate memory content)
	main.S * (Assembly code for verification)
	main.c * (C code for verification)
	data.S (Assembly code for testing data)
	setup.S (Assembly code for testing environment setup)
	<i>link.ld</i> (Linker script for testing environment)
	golden.dat (Golden hexadecimal data)
prog	g3 (Subfolder for Program 3)
	<i>Makefile</i> (Compile and generate memory content)
	main.S * (Assembly code for verification)
	main.c * (C code for verification)
	data.S (Assembly code for testing data)
	setup.S (Assembly code for testing environment setup)
	<i>link.ld</i> (Linker script for testing environment)

golden.dat (Golden hexadecimal data)

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B. Simulation Setting Requirements

You **SHOULD** make sure that your code can be simulated with specified commands in Table B-1. **TA will use the same command to check your design under SoC Lab environment.** If your code can't be recompiled by **TA successfully, you receive NO credit.**

Table B-1: Simulation commands

Simulation Level	Command	
Problem1		
RTL	make rtl_all	
Post-synthesis (optional)	make syn_all	

TA also provide some useful Makefile macros listed in Table B-2. Braces {} means that you can choose one of items in the braces. X stands for 0,1,2,3..., depend on which verification program is selected.

Table B-2: Makefile macros

Situation	Command
RTL simulation for progX	make rtlX
Post-synthesis simulation for progX	make synX
Dump waveform (no array)	make {rtlX,synX} FSDB=1
Dump waveform (with array)	make {rtlX,synX} FSDB=2
Open nWave without file pollution	make nWave
Open Superlint without file pollution	make superlint
Open DesignVision without file pollution	make dv
Synthesize your RTL code (You need write	make synthesize
synthesis.tcl in script folder by yourself)	make synthesize
Delete built files for simulation, synthesis	make clean
or verification	make clean
Check correctness of your file structure	make check
Compress your homework to tar format	make tar

You can use the following command to get the number of lines:

wc -1 src/* include/*	
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C. RISC-V Instruction Format

Tu-immediate

31

Inst[31]

30

20 19

inst[30:20]

Table C-1: Instruction type													
R-type											1		
31	25	24	20	19	15	14	12	11 7			6	0	
funct7 rs2			rs1 funct3			ct3	rd				opcode		
F I-type	☞ I-type								1				
31 20 1				19	15	14	12	11 7			6	0	
	imm[31:20]			rs1 funct3			ct3	rd				opcode	
S-type		ı				ı		_					
31	25	24	20	19	15	14	12	11		7	6	0	
im	m[11:5]	rs	s2	rs1 funct3			ct3	i	nm[4	:0]	opcode		
B-type						1				T	1		
31	30 25	24	20	19	15	14	12	11	8	7	6	0	
imm[12]	imm[10:5]	rs	s2	rs	s1	fun	ct3	imm[4	1]	imm[11]	opcode		
TU-type	;							I					
31	31 12 11 7									6	0		
	imm[31:12] rd									opcode			
J-type		-		1				<u> </u>			ı		
31	30	21	20	19 12			11 7				0		
imm[20]	imm[10:1]	im	ım[11]	imm[19:12]				rd				ode	
			Table (C-2:	Imn	nedia	te ty	pe					
T-imme	ediate												
31							11	10	5 4	1	C)	
							inst[30:25] inst[24:21]			inst[20]			
S-imm	ediate												
31	31 11 10 5 4 1							1	0)			
$-\inf[31]$ $-\inf[30:25]$ $\inf[11:8]$							nst[11:8]	inst	[7]				
F B-immediate													
31	31 12							10	5 4	1	C)	
— inst[31] —							7]	inst[30:25	0)			

1	2	/	13	

12 11

inst[19:12]

0

-0-

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J-immediate

31	20	19	12	11	10	5	4	1	0
— inst[31] —		inst[19:12]		inst[20]	inst[30:25]		inst[2	4:21]	0

[&]quot;— X —" indicates that all the bits in this range is filled with X.