

HW #8

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Exercises 2.11, 2.12, 2.14, 2.30, 2.31, 2.32, 2.35

2.11) $G_H = (V, \Sigma, R, \langle \text{EXPR} \rangle)$

$V = \langle \text{EXPR} \rangle$

V is $\{ \langle \text{EXPR} \rangle, \langle \text{TERM} \rangle, \langle \text{FACTOR} \rangle \}$

$T = \langle \text{TERM} \rangle$

$\Sigma = \{ a, +, \times, (,) \}$

$F = \langle \text{FACTOR} \rangle$

$E \rightarrow E + T \mid T$

$+ \rightarrow + \times F \mid T$

$F \rightarrow (E) \mid a$

$E, E \rightarrow E + T$

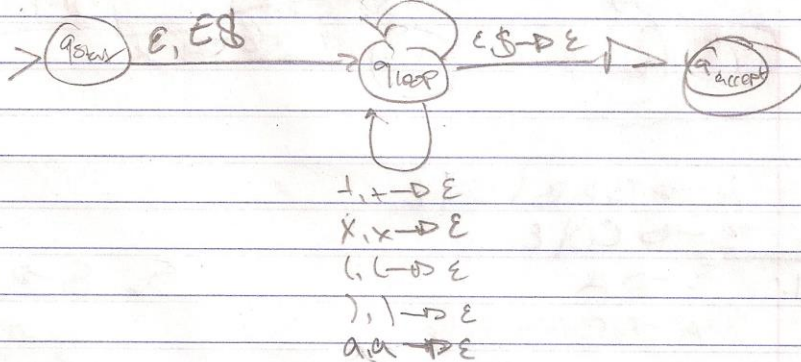
$E, E \rightarrow T$

$E, T \rightarrow + \times F$

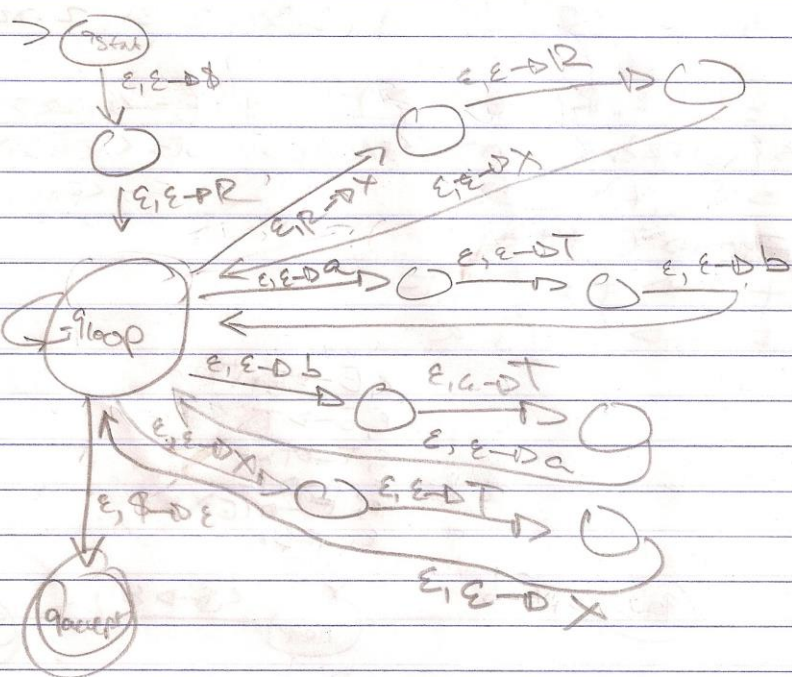
$E, T \rightarrow T$

$E, F \rightarrow (E)$

$E, F \rightarrow a$



2.12



R → X R X S
S → a T b b T a
T → X T X X X ε
X → a a b

ε, R → S
ε, T → X
ε, T → ε
ε, X → ε
ε, X → b
a, a → ε
b, b → ε

2.14 A → BAB | B | ε
B → ∅ | ε

1. S → A
A → BAB | B | ε
B → ∅ | ε
2. S → A
A → BAB | B | A B B A | A | ε
B → ∅
3. S → A | ε
A → BAB | B | A B | B A | A
B → ∅
4. S → BAB | A B | B A | B B | B | ε
A → BAB | A B | B A | B B | B
B → ∅
5. S → BAB | A B | B A | B B | ∅ | ε
A → BAB | A B | B A | B B | ∅
B → ∅

5. S → B C | A B | B B | ∅ | ε
A → B C | A B | B A | B B | ∅
C → A ∅
B → ∅ ∅
∅ → ∅

2.30a | $A = \{0^n 1^n 0^n \mid n \geq 0\}$

Let p be the pumping length

$$S = 0^p 1^p 0^p$$

S can be split into 5 parts $uvxyz$

1. for $i \geq 0$, $uv^i xy^i z \in L$

2. $|xy| \geq 0$

3. $|vxy| \leq p$

vxy can only be like the following

(1) 0^p or 1^p

for each $i \geq 2$ $uv^i xy^i z$ will have more 0's or 1's in the 1st half or 2nd half of S

(2) $0^p 1^p$ in the 1st half of S

for each $i \geq 2$ $uv^i xy^i z$ will have more 0's or 1's in the 1st half of S

(3) $1^p 0^p$ in the middle of S

for each $i \geq 2$, $uv^i xy^i z$ will have more 0's or 1's in the middle of S

(4) $0^p 1^p$ at the 2nd half of S

for each $i \geq 2$ $uv^i xy^i z$ will have more 0's or 1's in the 2nd half of S

2.30b | $B = \{0^p \# 0^{2p} \# 0^{3p} \mid p \geq 0\}$

Let p be the pumping length

$$S = 0^p \# 0^{2p} \# 0^{3p}$$

$$S = uvxyz \quad |vxy| \leq p, |vy| \geq 0$$

(1) v or y can not have $\#$ otherwise it will break the pattern of the regex more than 2#.

(2) Dividing the $0^p, 0^{2p}, 0^{3p}$ into 3 segments it is not possible to keep them all in either v or y it would break the regex ratio.

Thus it is not possible to pump this language

2.30c) $L = \{w\#t \mid w \text{ is a substring of } t \text{ where } w, t \in \{a,b\}^*\}$

Let p be the pumping length

$S = a^p b^p \# a^p b^p$, $S = uvxy^2$, $|vxy| \leq p$, $|vy| > 0$

- (1) v and y can not have $\#$ in them since there is only $\#$
- (2) $\#$ if v and y are both on the left of the $\#$ then uv^2xy^2 would be longer than the right side of $\#$
- (3) If v and y are both on the right of the $\#$ then uv^0xy^0 would be shorter than the left side which is a substring
- (4) If either v or y is empty then it would be like either (2) or (3) on the same side.
- (5) If v and y are both nonempty and the $\#$ is in between then then the left side of the $\#$ will have more b 's than the right side making it not a substring

2.30d) $D = \{t_1 \# t_2 \# \dots \# t_k \mid k \geq 2, \text{ Each } t_i \in \{a,b\}^*, \text{ and } t_i = t_j \text{ for } i \neq j\}$

Let p be the pumping length

$S = a^p 1^p \# a^p 1^p \in L$, $S = uvxy^2$, $|vxy| \leq p$, $|vy| > 0$

- (1) If vxy is only on the left side of the $\#$, then uv^2xy^2 makes $t_1 \neq t_2$
- (2) If vxy contains the $\#$ then uv^0xy^0 does not contain the $\#$
- (3) If x is the $\#$ and v is a substring of a^p while y is a substring of 1^p Pumping would make the two sides not equal

2.31 Let B be the language of all palindromes over $\{0,1\}$ containing equal numbers of 0's and 1's. Show B is not context-free.
 Let p be the pumping lemma

$$S = 0^p 1^p 1^p 0^p, S = uvxy^2 \quad |vxy| \leq p, |vy| > 1$$

(1) $0^p 1^p$ or $1^p 0^p$ is vxy and vy must contain both 0's and 1's. But then uv^0xy^0z would provide the reverse since one side of S is changed making it a palindrome no longer.

2.32 Let $\Sigma = \{1, 2, 3, 4\}$ and $C = \{w \in \Sigma^+ \mid \text{the \# of 1's} = \text{\# of 2's and \# 3's} = \text{\# 4's}\}$

Let p be the pumping length

$$S = 1^p 4^p 2^p 3^p, S = uvxy^2 \quad |vy| > 0, |vxy| \leq p$$

(1) vxy contains only 1, only 2, only 3, or 4

for each $i > 1$ the $\#1's \neq \#2's$ or $\#3's \neq \#4's$ breaking the reverse

(2) vxy is in $1^p 4^p$ or $4^p 2^p$ or $2^p 3^p$

for each $i > 1$, the $\#1's \neq \#2's$ or $\#3's \neq \#4's$ since pumping would increase the numbers unevenly