

# **Chapter 3:**

## **Assembly Language Programming II**

### **Topics**

- **Condition Codes/Flags**
  - **CF, ZF, OF, SF**
- **Conditional Control Flow**
  - **Conditional Jumps & Branching**
  - **Conditional Move**

# Announcements

- **Data Lab grading this week**
  - Sign up for grading time slots with your TA
- **Bomb Lab available on moodle, due Friday Oct 3**
  - TA reviewed this in recitation on Monday
- **Prof. Han traveling today**
  - No office hours today, will hold Wed office hours 3-4 pm
- **First midterm probably the week of Oct 6**
- **Essential that you read the textbook in detail & do the practice problems**
  - Read Chapter 3.1-3.14, skip 3.12 for now

# Recap...

- **X86 Assembly language**

- Eight 32-bit CPU registers: %eax, %ebx, %ecx, %edx, %edi, %esi, %esp, %ebp
- `movl Source, Dest`
  - Move between memory and CPU registers
  - Covered a swap example
  - Complex indexed addressing mode:
    - » `movl 8(%eax,%edx,4), dst`
- `addl Src, Dest`
  - Also `subl`, `imull`, `sall`, `sarl`, `andl`, etc.
- `leal Src, Dest`
  - Helpful for array arithmetic
- Covered an arithmetic assembly example

# Another Example

```
int logical(int x, int y)
{
    int t1 = x^y;
    int t2 = t1 >> 17;
    int mask = (1<<13) - 7;
    int rval = t2 & mask;
    return rval;
}
```

logical:

```
pushl %ebp
movl %esp,%ebp
```

} Set Up

```
movl 8(%ebp),%eax
xorl 12(%ebp),%eax
sarl $17,%eax
andl $8185,%eax
```

} Body

```
movl %ebp,%esp
popl %ebp
ret
```

} Finish

```
movl 8(%ebp),%eax
xorl 12(%ebp),%eax
sarl $17,%eax
andl $8185,%eax
```

```
# eax = x
# eax = x^y
# eax = t1>>17
# eax = t2 & 8185
```

# Another Example

```
int logical(int x, int y)
{
    int t1 = x^y;
    int t2 = t1 >> 17;
    int mask = (1<<13) - 7;
    int rval = t2 & mask;
    return rval;
}
```

logical:

```
pushl %ebp
movl %esp,%ebp
```

} Set Up

```
movl 8(%ebp),%eax
xorl 12(%ebp),%eax
sarl $17,%eax
andl $8185,%eax
```

} Body

```
movl %ebp,%esp
popl %ebp
ret
```

} Finish

```
movl 8(%ebp),%eax
xorl 12(%ebp),%eax
sarl $17,%eax
andl $8185,%eax
```

```
# eax = x
# eax = x^y    (t1)
# eax = t1>>17    (t2)
# eax = t2 & 8185
```

# Another Example

```
int logical(int x, int y)
{
    int t1 = x^y;
    int t2 = t1 >> 17;
    int mask = (1<<13) - 7;
    int rval = t2 & mask;
    return rval;
}
```

logical:

```
pushl %ebp
movl %esp,%ebp
```

} Set  
Up

```
movl 8(%ebp),%eax
xorl 12(%ebp),%eax
sarl $17,%eax
andl $8185,%eax
```

} Body

```
movl %ebp,%esp
popl %ebp
ret
```

} Finish

```
movl 8(%ebp),%eax
xorl 12(%ebp),%eax
sarl $17,%eax
andl $8185,%eax
```

```
# eax = x
# eax = x^y    (t1)
# eax = t1>>17  (t2)
# eax = t2 & 8185
```

# Another Example

```
int logical(int x, int y)
{
    int t1 = x^y;
    int t2 = t1 >> 17;
    int mask = (1<<13) - 7;
    int rval = t2 & mask;
    return rval;
}
```

$2^{13} = 8192, 2^{13} - 7 = 8185$

logical:

```
pushl %ebp
movl %esp,%ebp
```

} Set  
Up

```
movl 8(%ebp),%eax
xorl 12(%ebp),%eax
sarl $17,%eax
andl $8185,%eax
```

} Body

```
movl %ebp,%esp
popl %ebp
ret
```

} Finish

```
movl 8(%ebp),%eax
xorl 12(%ebp),%eax
sarl $17,%eax
andl $8185,%eax
```

```
# eax = x
# eax = x^y    (t1)
# eax = t1>>17  (t2)
# eax = t2 & 8185 ←
```

Note how  
compiler  
combines  
2 source code  
lines into 1

# Processor State (IA32, Partial)

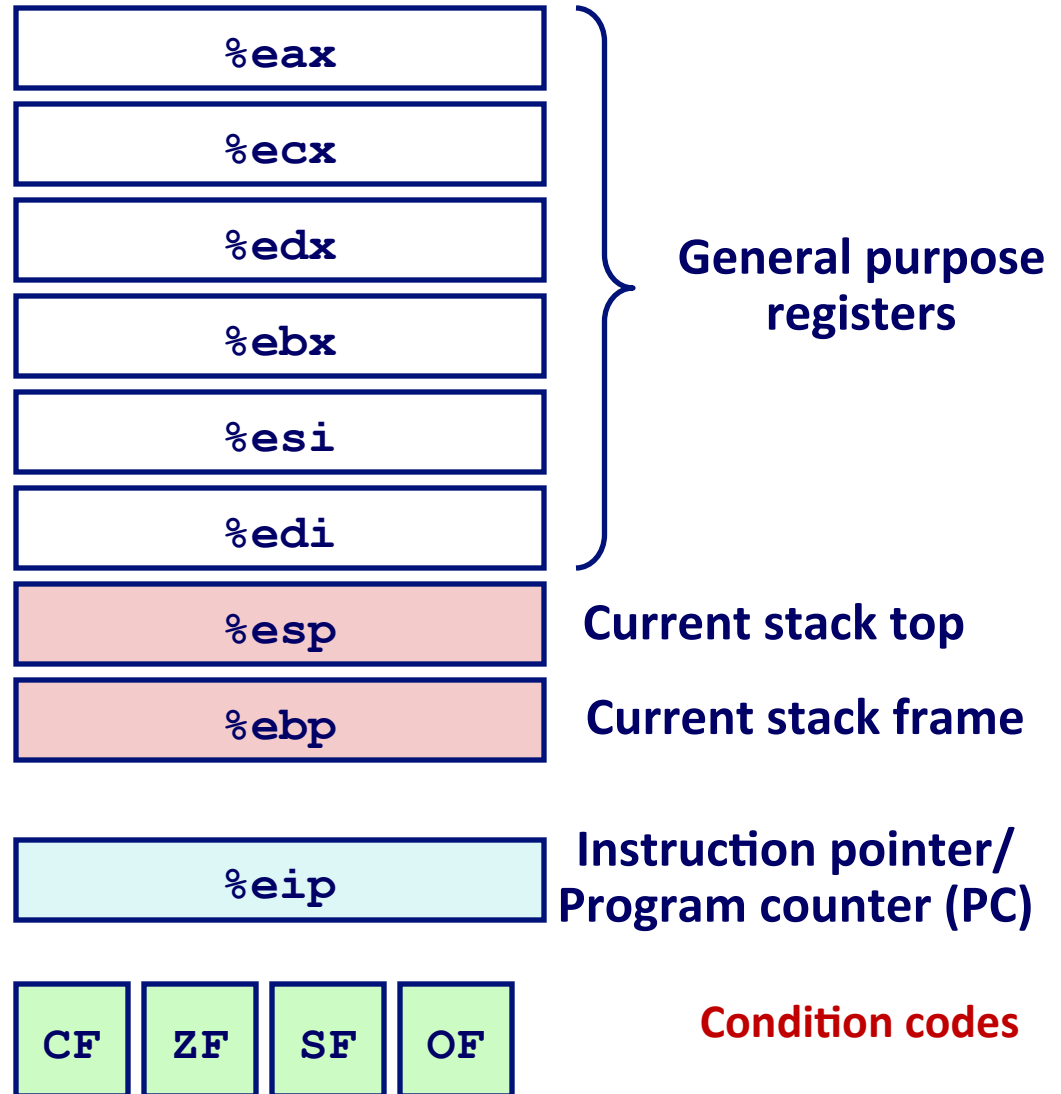
- Information about currently executing program

- Temporary data ( `%eax`, ... )

- Location of runtime stack ( `%ebp`, `%esp` )

- Location of current code control point ( `%eip`, ... )

- Status of recent tests ( `CF`, `ZF`, `SF`, `OF` )





# Condition Codes (Implicit Setting)

- Single bit registers

CF Carry Flag (for unsigned)

SF Sign Flag (for signed)

ZF Zero Flag

OF Overflow Flag (for signed)

- Implicitly set (think of it as side effect) by arithmetic operations

`addl Src, Dest`

C analog: `t = a+b`

- **CF set** if carry out from most significant bit

- Used to detect unsigned overflow

- **ZF set** if `t == 0`

- **SF set** if `t < 0` (as signed)

- **OF set** if two's complement (signed) overflow

`(a>0 && b>0 && t<0) || (a<0 && b<0 && t>=0)`

- Codes set differently depending on instructions, and in some cases not set at all, e.g. `leal` instruction

# Example: Left Shifting <<

- `shl` or `sal`
- **How are condition flags set?**
  - `shl/sal` sets carry flag to last shifted out bit.
    - For unsigned, if  $CF=1$ , then overflow occurred.
  - for overflow flag  $OF$ ,
    - if shift by one,
      - » if  $CF$  &  $MSbit$  identical after shift,  $OF = 0$
      - » else  $OF = 1$  ( $MS$  bit changes from  $1 \rightarrow 0$  or  $0 \rightarrow 1$ ),  
i.e.  $OF = CF \wedge MSbit$
    - else if shift by  $> 1$ ,  $OF$  undefined.
    - For signed, if  $OF=1$ , then overflow occurred.

# Condition Codes (Explicit Setting: Compare)

- **Explicit Setting by Compare Instruction**

`cmpl Src2,Src1`

`cmpl b,a` like computing `a-b` without setting destination

- **CF set** if carry out from most significant bit (used for unsigned comparisons)
- **ZF set** if `a == b`
- **SF set** if `(a-b) < 0` (as signed)
- **OF set** if two's complement (signed) overflow  
`(a>0 && b<0 && (a-b)<0) || (a<0 && b>0 && (a-b)>0)`

# Condition Codes (Explicit Setting: Test)

- **Explicit Setting by Test instruction**

`testl/testq Src2,Src1`

`testl b,a` like computing `a&b` without setting destination

- Sets condition codes based on value of *Src1* & *Src2*
- Useful to have one of the operands be a mask
- ZF set when `a&b == 0`
- SF set when `a&b < 0`

# Reading Condition Codes

- **SetX Instructions:**

Set single byte based on combination of condition codes

- **One of 8 addressable byte registers**

- Does not alter remaining 3 bytes
- Typically use `movzbl` to finish job

```
int gt (int x, int y)
{
    return x > y;
}
```

Body

```
movl 12(%ebp),%eax
cmpl %eax,8(%ebp)
setg %al
movzbl %al,%eax
```

%eax	%ah	%al
%ecx	%ch	%cl
%edx	%dh	%dl
%ebx	%bh	%bl
%esi		
%edi		
%esp		
%ebp		

# Reading Condition Codes (Cont.)

- **SetX Instructions:**

Set single byte based on combination of condition codes

- **One of 8 addressable byte registers**

- Does not alter remaining 3 bytes
- Typically use `movzbl` to finish job

```
int gt (int x, int y)
{
    return x > y;
}
```

Body

```
movl 12(%ebp),%eax    # eax = y
cmpl %eax,8(%ebp)      # Compare x and y
setg %al              # al = x > y
movzbl %al,%eax        # Zero rest of %eax
```

Note  
inverted  
ordering!

%eax	%ah	%al
%ecx	%ch	%cl
%edx	%dh	%dl
%ebx	%bh	%bl
%esi		
%edi		
%esp		
%ebp		

# Reading Condition Codes

- **setX dest // e.g. setl %al**
  - Set single byte of dest to 0 or 1 based on combinations of condition codes (expressions below are only relevant if previous instruction was a compare `cmp`)

SetX	Condition	Description
<code>sete</code>	<code>ZF</code>	Equal / Zero
<code>setne</code>	<code>~ZF</code>	Not Equal / Not Zero
<code>sets</code>	<code>SF</code>	Negative
<code>setns</code>	<code>~SF</code>	Nonnegative
<code>setg</code>	<code>~(SF^OF) &amp; ~ZF</code>	Greater (Signed)
<code>setge</code>	<code>~(SF^OF)</code>	Greater or Equal (Signed)
<code>setl</code>	<code>(SF^OF)</code>	Less (Signed) – derivation in text
<code>setle</code>	<code>(SF^OF)   ZF</code>	Less or Equal (Signed)
<code>seta</code>	<code>~CF &amp; ~ZF</code>	Above (unsigned)
<code>setb</code>	<code>CF</code>	Below (unsigned)

# Jumping – Conditional or not

- **jX Instructions**

- **Unconditional jump:** `jmp Label` or `jmp *%eax` or `jmp *(%eax)`
- **Conditional jumps to different part of code depending on condition codes, e.g. `jle Label`**

jX	Condition	Description
<code>jmp</code>	1	Unconditional
<code>je</code>	ZF	Equal / Zero
<code>jne</code>	$\sim$ ZF	Not Equal / Not Zero
<code>js</code>	SF	Negative
<code>jns</code>	$\sim$ SF	Nonnegative
<code>jg</code>	$\sim (SF \wedge OF) \ \& \ \sim ZF$	Greater (Signed)
<code>jge</code>	$\sim (SF \wedge OF)$	Greater or Equal (Signed)
<code>j1</code>	$(SF \wedge OF)$	Less (Signed)
<code>jle</code>	$(SF \wedge OF) \   \ ZF$	Less or Equal (Signed)
<code>ja</code>	$\sim CF \ \& \ \sim ZF$	Above (unsigned)
<code>jb</code>	CF	Below (unsigned)



# Conditional Branch Example

```
int absdiff(int x, int y)
{
    int result;
    if (x > y) {
        result = x-y;
    } else {
        result = y-x;
    }
    return result;
}
```

```
absdiff:
    pushl    %ebp
    movl     %esp, %ebp
    movl     8(%ebp), %edx
    movl     12(%ebp), %eax
    cmpl     %eax, %edx
    jle      .L7
    subl     %eax, %edx
    movl     %edx, %eax
.L8:
    leave
    ret
.L7:
    subl     %edx, %eax
    jmp      .L8
```

Diagram illustrating the assembly code for the `absdiff` function, grouped into sections:

- Setup**: `pushl %ebp`, `movl %esp, %ebp`
- Body1**: `movl 8(%ebp), %edx`, `movl 12(%ebp), %eax`, `cmpl %eax, %edx`, `jle .L7`, `subl %eax, %edx`, `movl %edx, %eax`
- Finish**: `.L8:`, `leave`, `ret`
- Body2**: `.L7:`, `subl %edx, %eax`, `jmp .L8`

# Conditional Branch Example (Cont.)

```
int absdiff(int x, int y)
{
    int result;
    if (x > y) {
        result = x-y;
    } else {
        result = y-x;
    }
    return result;
}
```

```
int goto_ad(int x, int y)
{
    int result;
    if (x <= y) goto Else;
    result = x-y;
Exit:
    return result;
Else:
    result = y-x;
    goto Exit;
}
```

- C allows “goto” as means of transferring control
  - Closer to machine-level programming style
- Generally considered bad coding style

# Conditional Branch Example (Cont.)

```
int goto_ad(int x, int y)
{
    int result;
    if (x <= y) goto Else;
    result = x-y;
Exit:
    return result;
Else:
    result = y-x;
    goto Exit;
}
```

- **Rewritten absdiff with labels closely reflects assembly language**

```
absdiff:
    pushl    %ebp
    movl     %esp, %ebp
    movl     8(%ebp), %edx
    movl     12(%ebp), %eax
    cmpl     %eax, %edx
    jle      .L7
    subl     %eax, %edx
    movl     %edx, %eax
.L8:
    leave
    ret
.L7:
    subl     %edx, %eax
    jmp      .L8
```

# Conditional Branch Example (Cont.)

```
int goto_ad(int x, int y)
{
    int result;
    if (x <= y) goto Else;
    result = x-y;
Exit:
    return result;
Else:
    result = y-x;
    goto Exit;
}
```

```
absdiff:
    pushl    %ebp
    movl     %esp, %ebp
    movl     8(%ebp), %edx
    movl     12(%ebp), %eax
    cmpl     %eax, %edx
    jle      .L7
    subl     %eax, %edx
    movl     %edx, %eax
.L8:
    leave
    ret
.L7:
    subl     %edx, %eax
    jmp      .L8
```

# Conditional Branch Example (Cont.)

```
int goto_ad(int x, int y)
{
    int result;
    if (x <= y) goto Else;
    result = x-y;
Exit:
    return result;
Else:
    result = y-x;
    goto Exit;
}
```

```
absdiff:
    pushl    %ebp
    movl     %esp, %ebp
    movl     8(%ebp), %edx
    movl     12(%ebp), %eax
    cmpl     %eax, %edx
    jle      .L7
    subl     %eax, %edx
    movl     %edx, %eax
.L8:
    leave
    ret
.L7:
    subl     %edx, %eax
    jmp      .L8
```

# Conditional Branch Example (Cont.)

```
int goto_ad(int x, int y)
{
    int result;
    if (x <= y) goto Else;
    result = x-y;
Exit:
    return result;
Else:
    result = y-x;
    goto Exit;
}
```

```
absdiff:
    pushl    %ebp
    movl     %esp, %ebp
    movl     8(%ebp), %edx
    movl     12(%ebp), %eax
    cmpl     %eax, %edx
    jle      .L7
    subl     %eax, %edx
    movl     %edx, %eax
.L8:
    leave
    ret
.L7:
    subl     %edx, %eax
    jmp      .L8
```

# Conditional Branch Example (Cont.)

```
int goto_ad(int x, int y)
{
    int result;
    if (x <= y) goto Else;
    result = x-y;
Exit:
    return result;
Else:
    result = y-x;
    goto Exit;
}
```

```
absdiff:
    pushl    %ebp
    movl     %esp, %ebp
    movl     8(%ebp), %edx
    movl     12(%ebp), %eax
    cmpl     %eax, %edx
    jle      .L7
    subl     %eax, %edx
    movl     %edx, %eax
.L8:
    leave
    ret
.L7:
    subl     %edx, %eax
    jmp      .L8
```

# General Conditional Expression Translation

## C Code

```
val = Test ? Then-Expr : Else-Expr;
```

```
val = x > y ? x - y : y - x;
```

## Goto Version

```
nt = !Test;  
if (nt) goto Else;  
val = Then-Expr;  
Done:  
. . .  
Else:  
val = Else-Expr;  
goto Done;
```

- *Test* is expression returning integer
  - = 0 interpreted as false
  - ≠ 0 interpreted as true
- Create separate code regions for then & else expressions
- Execute appropriate one



# Conditional Move

- **cmovX src, dest**
  - Set dest=src only if condition X holds
  - More efficient than conditional branching for highly pipelined processors – easier to guess the next instruction to execute
  - But overhead: both branches are evaluated

cmovX	Condition	Description
cmove	ZF	Equal / Zero
cmovne	~ZF	Not Equal / Not Zero
cmovs	SF	Negative
cmovns	~SF	Nonnegative
cmovg	~ (SF^OF) & ~ZF	Greater (Signed)
cmovge	~ (SF^OF)	Greater or Equal (Signed)
cmovl	(SF^OF)	Less (Signed)
cmovle	(SF^OF)   ZF	Less or Equal (Signed)
cmova	~CF & ~ZF	Above (unsigned)
cmovb	CF	Below (unsigned)

# Conditional Move Example

- Rewrite the absdiff example using conditional moves instead of conditional branching

```
int absdiff(  
    int x, int y)  
{  
    int result;  
    if (x > y) {  
        result = x-y;  
    } else {  
        result = y-x;  
    }  
    return result;  
}
```

```
absdiff: # x in %edi, y in %esi  
    movl    %edi, %eax    # eax = x  
    movl    %esi, %edx    # edx = y  
    subl    %esi, %eax    # eax = x-y  
    subl    %edi, %edx    # edx = y-x  
    cmpl    %esi, %edi    # x<y?  
    cmovl    %edx, %eax    # eax=edx if <=  
    ret
```

- Note how the control flow is much easier to predict than all the jumping around with labels, e.g. `jle`, in the conditionally branched version of `absdiff`

# Conditional Move Example

```
int absdiff(  
    int x, int y)  
{  
    int result;  
    if (x > y) {  
        result = x-y;  
    } else {  
        result = y-x;  
    }  
    return result;  
}
```



```
int cmovdiff(int x, int  
y)  
{  
    int rval = x-y;  
    int tval = y-x;  
    int test = x < y;  
    if (test) rval = tval;  
    return rval;  
}
```

# Conditional Move Example

```
int cmovdiff(int x, int
y)
{
    int rval = x-y;
    int tval = y-x;
    int test = x < y;
    if (test) rval = tval;
    return rval;
}
```

```
absdiff: # x in %edi, y in %esi
    movl    %edi, %eax    # eax = x
    movl    %esi, %edx    # edx = y
    subl    %esi, %eax    # eax = x-y
    subl    %edi, %edx    # edx = y-x
    cmpl    %esi, %edi    # x<y?
    cmovl    %edx, %eax    # eax=edx if <=
    ret
```

# Conditional Move Example

```
int cmovdiff(int x, int
y)
{
    int rval = x-y;
    int tval = y-x;
    int test = x < y;
    if (test) rval = tval;
    return rval;
}
```

```
absdiff: # x in %edi, y in %esi
    movl    %edi, %eax    # eax = x
    movl    %esi, %edx    # edx = y
    subl    %esi, %eax    # eax = x-y
    subl    %edi, %edx    # edx = y-x
    cmpl    %esi, %edi    # x<y?
    cmovl    %edx, %eax    # eax=edx if <=
    ret
```

# Conditional Move Example

```
int cmovdiff(int x, int
y)
{
    int rval = x-y;
    int tval = y-x;
    int test = x < y;
    if (test) rval = tval;
    return rval;
}
```

```
absdiff: # x in %edi, y in %esi
    movl    %edi, %eax    # eax = x
    movl    %esi, %edx    # edx = y
    subl    %esi, %eax    # eax = x-y
    subl    %edi, %edx    # edx = y-x
    cmpl    %esi, %edi    # x<y?
    cmovl    %edx, %eax    # eax=edx if <=
    ret
```

# Conditional Move Example

```
int cmovdiff(int x, int
y)
{
    int rval = x-y;
    int tval = y-x;
    int test = x < y;
    if (test) rval = tval;
    return rval;
}
```

```
absdiff: # x in %edi, y in %esi
    movl    %edi, %eax    # eax = x
    movl    %esi, %edx    # edx = y
    subl    %esi, %eax    # eax = x-y
    subl    %edi, %edx    # edx = y-x
    cmpl    %esi, %edi    # x<y?
    cmovl    %edx, %eax    # eax=edx if <=
    ret
```

- Control flow is more predictable, but both branches must be evaluated

# General Form with Conditional Move

## C Code

```
val = Test ? Then-Expr : Else-Expr;
```

## Conditional Move Version

```
val1 = Then-Expr;  
val2 = Else-Expr;  
val1 = val2 if !Test;
```

- Both values get computed
- Overwrite then-value with else-value if condition doesn't hold
- **Don't use when:**
  - Then or else expression have side effects, like dereferencing a null pointer or incrementing a global variable (then & else expressions always evaluated)
  - Then and else expressions are too expensive