

Tutorial3.3

Binary Arithmetic Assignment2

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Modulus Arithmetic

- Any carry out is ignored when using ordinary arithmetic instructions
 - Eg: $9 + 8$ should give 17
 - but actually gives $(9 + 8) \bmod 16$ on a 4-bit CPU

$$\begin{array}{r} 1001 \\ + \underline{1000} \\ 1\ 0001 \end{array}$$

- Instructions like **adds** or **subs** do set the carry flag
 - can be used to do extended precision arithmetic

Signed number Branching Conditions

- subs
 - alias for *cmp*
- ands
 - *ands Xzr, Xn, Xm*
 - alias for *tst*
 - The *tst* instruction performs a bitwise AND operation on the value in Xn and the value of Xm. This is the same as an ANDS instruction, except that the result is discarded, updates the N and Z flags according to the result.
 - eg
 - *tst x20, 0x8*
 - *b.eq bitclear*
- subs and ands set the flags

Signed number Branching Conditions

- Recall that signed branch instructions use the N, Z, V flags

Name	Meaning	C equivalent	Flags
eq	equal	==	Z == 1
ne	not equal	!=	Z == 0
gt	greater than	>	Z == 0 && N == V
ge	greater than or equal	>=	N == V
lt	less than	<	N != V
le	less than or equal	<=	!(Z == 0 && N == V)

ands (tst)

- `ands wzr, product, 0x1` or `tst product, 0x1`
 - bitwise AND of product and 0x1
- `b.eq productEndsWithZero`
 - `b.eq` doesn't mean true, but rather means the Z flag is set.
 - The **Z flag** is set here if the ANDS operation results in zero (0).
 - That means the Z flag is set (`b.eq`) when bit-0 of product is 0.
- `b.ne productEndsWithOne`
 - This is not necessary if you just let your code fall through
 - `b.ne` (**Z flag is not set**) when bit-0 of product is 1
 - If bit-0 of product is 1, then `ANDS product, 0x1` will product 1 (Z flag is not set).

Multiplication

- Can be done with an iterated procedure
 - Repeated for every bit in the source registers
- Each step consists of two sub-steps:
 - A conditional addition
 - An arithmetic Shift right
- The product may require twice as many bits as for the source registers
 - the result is put into 2 concatenated registers
- if the original multiplier is negative, an extra step is needed:
 - Subtract the multiplicand from the high-order part of the result(i.e.from the product register)

why extra step?

- Negating a number is done by:
 1. Taking the one's complement
 - Toggle all 0's to 1's, and vice versa
 2. Adding 1 to the result
- Eg: find the bit pattern for -5 in a 4-bit register
 - +5 is 0101
 - One's complement: 1010
 - Add 1: 1011
- The result equals 2^{N+1} -positive value
 - Eg:
 - $2^5(10000) - 5(0101) = 1011$
- $a * (\text{neg.}b) = a * (2^{n+1} - b) = a * 2^{n+1} - a * b = a * 2^{n+1} + a * (-b)$
so $a * (-b) = a * (\text{neg.}b) - a * 2^{n+1} = \text{product} - a * 2^{n+1}$

Assignment2

- Create an ARMv8 assembly language that implement the following integer multiplication program

conditional addition
using *tst* in assembly

shift the result to multiplier

arithmetic shift right(*asr*)

extra step for negative

extend the product to 64bit
using *sxtw* in assembly

```
#include <stdio.h>
#define FALSE 0
#define TRUE 1

int main()
{
    int multiplier, multiplicand, product, i, negative;
    long int result, temp1, temp2;

    // Initialize variables
    multiplicand = -16843010;
    multiplier = 70;
    product = 0;

    // Print out initial values of variables
    printf("multiplier = 0x%08x (%d)  multiplicand = 0x%08x (%d)\n\n",
        multiplier, multiplier, multiplicand, multiplicand);

    // Determine if multiplier is negative
    negative = multiplier < 0 ? TRUE : FALSE;

    // Do repeated add and shift
    for (i = 0; i < 32; i++) {
        if (multiplier & 0x1) {
            product = product + multiplicand;
        }

        // Arithmetic shift right the combined product and multiplier
        multiplier = multiplier >> 1;
        if (product & 0x1) {
            multiplier = multiplier | 0x80000000;
        } else {
            multiplier = multiplier & 0x7FFFFFFF;
        }
        product = product >> 1;
    }

    // Adjust product register if multiplier is negative
    if (negative) {
        product = product - multiplicand;
    }

    // Print out product and multiplier
    printf("product = 0x%08x  multiplier = 0x%08x\n",
        product, multiplier);

    // Combine product and multiplier together
    temp1 = (long int)product & 0xFFFFFFFF;
    temp1 = temp1 << 32;
    temp2 = (long int)multiplier & 0xFFFFFFFF;
    result = temp1 + temp2;

    // Print out 64-bit result
    printf("64-bit result = 0x%016lx (%ld)\n", result, result);

    return 0;
}
```


Assignment2 requirement

- use 32-bit registers for variables declared using `int`, 64-bit for `long int`
- use `m4` macros
- debug using `gdb`, capture it using script
- name requirement: `assign2a.asm`
`assign2b.asm`
`assign2c.asm`
- Don't use `gcc -S`