CPSC 457 T03/T04

Week 4 Day 1

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Agenda

- Condition variable
- unique_lock
- lock_guard
- Semaphores
- Assignment 3

Condition variables

- A synchronization primitive
- Used together with mutexes
- When to use: critical sections w/ loops waiting for some 'condition' to happen
- This 'condition' can only become true if another thread runs its critical section

```
// critical section protected with mutex m
lock(m)
...
while(!some_condition) { wait(cv); }
...
unlock(m)

// some other thread ...
some_condition = TRUE;
signal(cv);
...
unlock(m)
```

Download CPSC 457 – Synchronization – Code.pdf from D2L under W4D1

• Exercise + Demo

<u>CV</u> in C++

unique lock in C++

lock guard in C++

Semaphores

- Semaphores can be used to provide mutual exclusion.
- A semaphore has integer value.
- A locked semaphore can be unlocked by any thread
- Three operations: initialization, decrement (wait), increment (signal)
- decrement and increment must be executed atomically

Mutex	Semaphore
Binary operation/value (lock/unlock)	Integer value
Lock/unlock done by the same thread	Locked threads can be unlocked by any thread

Binary semaphore

- Behave very similarly to mutex locks.
- On systems that do not support mutex locks, binary semaphores can be used instead to provide mutual exclusion.
- Note: A locked thread can be unlocked by any thread using a semaphore.

Counting semaphores

- Can be used to control access to some given resource consisting of a finite number of instances.
 - Initialize the semaphore to the number of resources available
 - Each process that uses a resource performs a wait() operation (aka. decrement resource count)
 - When the count for the semaphore goes to 0, all resources are being used.
 - Additional processes that wish to use a resource will now be blocked until count > 0.

Assignment 3

Ignore 0,1

Input:	25	4012009	165
Factors:	1,5	1,2003,401209 * Typo on A3	1, 3, 5,,165
		Jr	

Ignore 0,1

Input:	25	4012009	165
Factors:	1,5	1,2003 4012009 * Typo on A3	1,35,,165
		* Typo on A3	
smallest non-trivial factor:	5	2003	3

Ignore 0,1

Input:	25	4012009	165	
Factors:	1,5	1,2003)4012009 * Typo on A3	1,35,,165	
		* Typo on A3		
smallest non-trivial factor:	5	2003	3	
$\Sigma = 5 + 2003 + 3 = 2011$ Output				

Solution 1: Let N=2 5 10 33 1123 Thread #1 3 100000 Thread #2

```
Solution 1: Let N=2
            Not Good.

Thread #1

Thread #2

Thread #2

132456799817

#1 10 33

1123

Thread #2

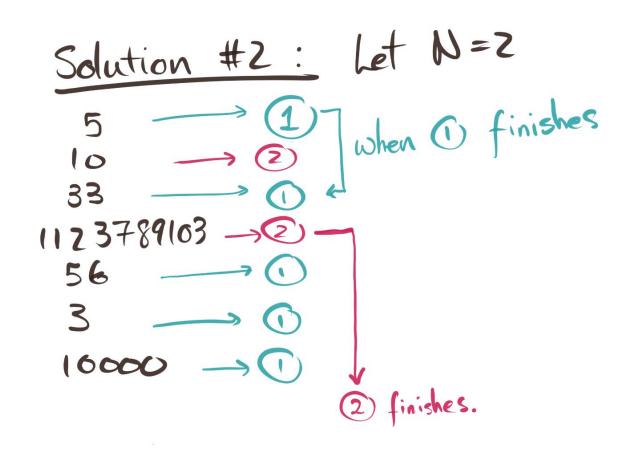
123

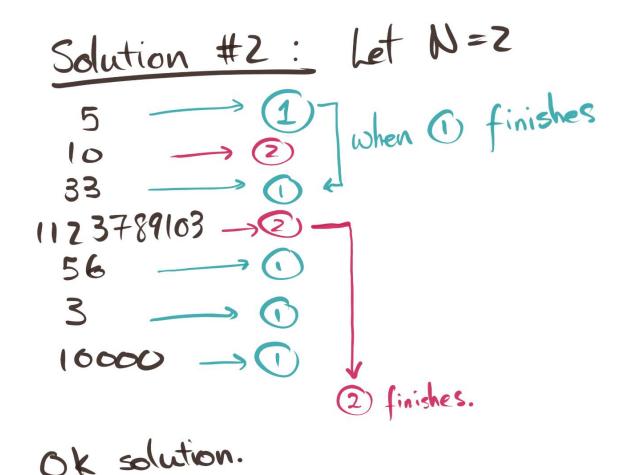
Thread #2

123

Thread #2
                                             bottleneck.

>runtime will
                                            be similar to single thread code.
```





However, consider input: 1123789103

Runtime?

Solution #3: Let N=2 33 1123789103 -> 112 + Parallelize get Smallest Divisor ().

Assignment 3 suggestions

- 1. Implement solution #2 first.
- 2. Use <u>mutex</u>, <u>condition variable</u>, <u>semaphores</u>, and <u>pthread barrier</u> to try and implement solution #3.

Next Time

- pthread versions of everything
- Assignment 3 help (?)