

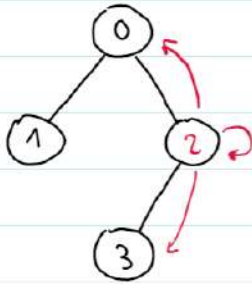
## 968. Binary Tree Cameras

Rules: A camera at a node covers:

- The node itself
- Its parent
- Its two immediate children

Goal: Minimize the amount of cameras needed to cover entire tree

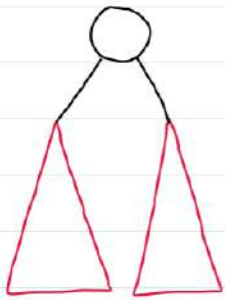
Example:



A camera on node 2 could cover nodes: 0, 2 and 3

Cases: For any given node (A.K.A. root)

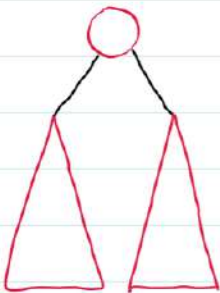
0:



Both left and right subtrees are covered except root node

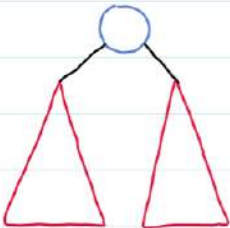
All nodes below current are covered

1:



Entire subtree is covered including its root

2:



All nodes are covered plus root has a camera

Base: null nodes are considered as covered and can't have cameras on them.

So leaf nodes can meet any of three cases above.

## 968. Binary Tree Cameras

Possible States:

$\forall u$  node, we define  $f(u)$  that returns  $d[3]$ :

$d[0]$ : Min cameras needed to cover entire subtree excluding  $u$

$d[1]$ : Min cameras needed to cover entire subtree including  $u$  with no camera in  $u$

$d[2]$ : Min cameras needed to cover entire subtree with a camera on  $u$

Recurrence

$$L = f(u.\text{left}) \quad R = f(u.\text{right})$$

$$dp[0] = L[1] + R[1] \quad \text{Children don't have camera}$$

$$dp[1] = \min(L[1] + R[2], L[2] + R[2], L[2] + R[2]) \quad \text{At least 1 camera}$$

$$dp[2] = 1 + \min(L) + \min(R) \quad \text{Min of each}$$

Base Case:

$$f(\text{null}) = (0, 0, \infty)$$

Cameras not needed for a nonexistent node

A camera cannot be put on a null node

Final Result:

$$dp = f(\text{root}) \Rightarrow \min(dp[1], dp[2])$$

Root must be covered, thus  $dp[0]$  cannot be used