

Assignment 5

EMPLOYEE

Fname	Minit	Lname	<u>Ssn</u>	Bdate	Address	Sex	Salary	Super_ssn	Dno
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DEPARTMENT

<u>Dname</u>	<u>Dnumber</u>	Mgr_ssn	Mgr_start_date
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DEPT_LOCATIONS

<u>Dnumber</u>	<u>Dlocation</u>
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PROJECT

Pname	<u>Pnumber</u>	Plocation	Dnum
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WORKS_ON

<u>Essn</u>	<u>Pno</u>	Hours
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DEPENDENT

<u>Essn</u>	<u>Dependent_name</u>	Sex	Bdate	Relationship
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Figure 5.5

Schema diagram for the COMPANY relational database schema.

7.7. In SQL, specify the following queries on the database in Figure 5.5 using the concept of nested queries and other concepts described in this chapter.

- a. Retrieve the names of all employees who work in the department that has the employee with the highest salary among all employees.
- b. Retrieve the names of all employees whose supervisor's supervisor has '888665555' for Ssn.
- c. Retrieve the names of employees who make at least \$10,000 more than the employee who is paid the least in the company.

a. **SELECT** Fname, Minit, Lname
FROM EMPLOYEE
WHERE Dno = (**SELECT** Dno
FROM EMPLOYEE
WHERE Salary = (**SELECT** MAX(Salary)
FROM EMPLOYEE))

b. **SELECT** Fname, Minit, Lname
FROM EMPLOYEE
WHERE Super_ssn = (**SELECT** Ssn
FROM EMPLOYEE
WHERE Super_ssn = '888665555')

c. **SELECT** Fname, Minit, Lname
FROM EMPLOYEE
WHERE Salary - 10000 >= (**SELECT** MIN(Salary)
FROM EMPLOYEE)

7.8. Specify the following views in SQL on the COMPANY database schema shown in Figure 5.5.

- a. A view that has the department name, manager name, and manager salary for every department
- b. A view that has the employee name, supervisor name, and employee salary for each employee who works in the 'Research' department
- c. A view that has the project name, controlling department name, number of employees, and total hours worked per week on the project for each project
- d. A view that has the project name, controlling department name, number of employees, and total hours worked per week on the project for each project *with more than one employee working on it*

a. **CREATE VIEW** DEPT_MGR(Dept_name, Mgr_fname, Mgr_minit, Mgr_lname
Mgr_sal)
AS SELECT Dname, Fname, Minit, Lname, Salary
FROM DEPARTMENT, EMPLOYEE
WHERE Mgr_ssn = Ssn;

b. **CREATE VIEW** RESEARCH_DEPT(E.Fname, E.Minit, E.Lname S.Fname,
S.Minit, S.Lname, E.salary)
AS SELECT E.Fname, E.Minit, Lname, S.Fname, S.Minit, S.Lname, E.Salary
FROM EMPLOYEE E, EMPLOYEE S, Department D
WHERE E.Super_ssn = S.ssn **AND** E.Dno = D.Dnumber **AND**
D.Dname = 'Research';

c. **CREATE VIEW** PROJECT_INFO(Project_name, Dept_name, Num_employees,
Total_hrs)
AS SELECT Pname, Dname, COUNT(*), SUM(Hours)
FROM PROJECT, DEPARTMENT, WORKS_ON
WHERE Pnumber = Pno **AND** Dnumber = Dnum
GROUP BY Pnumber, Pname;

d. **CREATE VIEW** PROJECT_INFO2(Project_name, Dept_name, Num_employees,
Total_hrs)
AS SELECT Pname, Dname, COUNT(*), SUM(Hours)
FROM PROJECT, DEPARTMENT, WORKS_ON
WHERE Pnumber = Pno **AND** Dnumber = Dnum
GROUP BY Pnumber, Pname
HAVING COUNT(*) > 1;

AIRPORT

<u>Airport_code</u>	Name	City	State
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FLIGHT

<u>Flight_number</u>	Airline	Weekdays
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FLIGHT_LEG

<u>Flight_number</u>	<u>Leg_number</u>	Departure_airport_code	Scheduled_departure_time
		Arrival_airport_code	Scheduled_arrival_time

LEG_INSTANCE

<u>Flight_number</u>	<u>Leg_number</u>	<u>Date</u>	Number_of_available_seats	Airplane_id	
		Departure_airport_code	Departure_time	Arrival_airport_code	Arrival_time

FARE

<u>Flight_number</u>	<u>Fare_code</u>	Amount	Restrictions
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AIRPLANE_TYPE

<u>Airplane_type_name</u>	Max_seats	Company
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CAN LAND

<u>Airplane_type_name</u>	<u>Airport_code</u>
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AIRPLANE

<u>Airplane_id</u>	Total_number_of_seats	Airplane_type
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SEAT_RESERVATION

<u>Flight_number</u>	<u>Leg_number</u>	<u>Date</u>	<u>Seat_number</u>	Customer_name	Customer_phone
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Figure 5.8

The AIRLINE relational database schema.

8.17. Consider the AIRLINE relational database schema shown in Figure 5.8, which was described in Exercise 5.12. Specify the following queries in relational algebra:

- For each flight, list the flight number, the departure airport for the first leg of the flight, and the arrival airport for the last leg of the flight.
- List the flight numbers and weekdays of all flights or flight legs that depart from Houston Intercontinental Airport (airport code 'iah') and arrive in Los Angeles International Airport (airport code 'lax').
- List the flight number, departure airport code, scheduled departure time, arrival airport code, scheduled arrival time, and weekdays of all flights or flight legs that depart from some airport in the city of Houston and arrive at some airport in the city of Los Angeles.
- List all fare information for flight number 'co197'.
- Retrieve the number of available seats for flight number 'co197' on '2009-10-09'.

$\sigma_{\langle \text{selection condition} \rangle}(\mathbf{R})$

$\pi_{\langle \text{attribute list} \rangle}(\mathbf{R})$

$\rho_{\mathbf{S}(\mathbf{B}_1, \mathbf{B}_2, \dots, \mathbf{B}_n)}(\mathbf{R})$ or $\rho_{\mathbf{S}}(\mathbf{R})$ or $\rho_{(\mathbf{B}_1, \mathbf{B}_2, \dots, \mathbf{B}_n)}(\mathbf{R})$

$\mathbf{R}_1 \times \mathbf{R}_2$

$\mathbf{R}_1 \bowtie_{\langle \text{join condition} \rangle} \mathbf{R}_2$

$\mathbf{R}_1 *_{\langle \text{join condition} \rangle} \mathbf{R}_2$

a. $\text{INT1} \leftarrow \sigma_{\langle \text{Leg_number} = 1 \rangle}(\text{FLIGHT_LEG})$

$\text{INT2} \leftarrow \pi_{\langle \text{Flight_number}, \text{Departure_airport_code} \rangle}(\text{INT1})$

$\text{INT3} \leftarrow \sigma_{\langle \text{Leg_number} = \text{MAX}(\text{Leg_number}) \rangle}(\text{FLIGHT_LEG})$

$\text{INT4} \leftarrow \pi_{\langle \text{Flight_number}, \text{Arrival_airport_code} \rangle}(\text{INT3})$

$\text{RESULT} \leftarrow (\text{INT2} * \text{INT4})$

b. $\text{INT1} \leftarrow \sigma_{\langle \text{Departure_airport_code} = \text{'iah'} \text{ AND } \text{Arrival_airport_code} = \text{'lax'} \rangle}(\text{FLIGHT_LEG})$

$\text{INT2} \leftarrow \pi_{\langle \text{Flight_number} \rangle}(\text{INT1})$

$\text{INT3} \leftarrow (\text{INT2} * \text{FLIGHT})$

$\text{RESULT} \leftarrow \pi_{\langle \text{Flight_number}, \text{Weekdays} \rangle}(\text{INT3})$

c. $\text{INT1} \leftarrow \sigma_{\langle \text{Departure_airport_code} = \text{'iah'} \text{ AND } \text{Arrival_airport_code} = \text{'lax'} \rangle}(\text{FLIGHT_LEG})$

$\text{INT2} \leftarrow \pi_{\langle \text{Flight_number}, \text{Departure_airport_code}, \text{Departure_time}, \text{Arrival_airport_code}, \text{Arrival_time} \rangle}(\text{INT1})$

$\text{INT3} \leftarrow (\text{INT2} * \text{FLIGHT})$

$\text{RESULT} \leftarrow \pi_{\langle \text{Flight_number}, \text{Departure_airport_code}, \text{Departure_time}, \text{Arrival_airport_code}, \text{Arrival_time}, \text{Weekdays} \rangle}(\text{INT3})$

d. $\text{RESULT} \leftarrow \sigma_{\langle \text{Flight_number} = \text{'co197'} \rangle}(\text{FARE})$

e. $\text{INT1} \leftarrow \sigma_{\langle \text{Flight_number} = \text{'co197'} \text{ AND } \text{Date} = \text{'2009-10-09'} \rangle}(\text{LEG_INSTANCE})$

$\text{RESULT} \leftarrow \pi_{\langle \text{Number_of_available_seats} \rangle}(\text{INT1})$

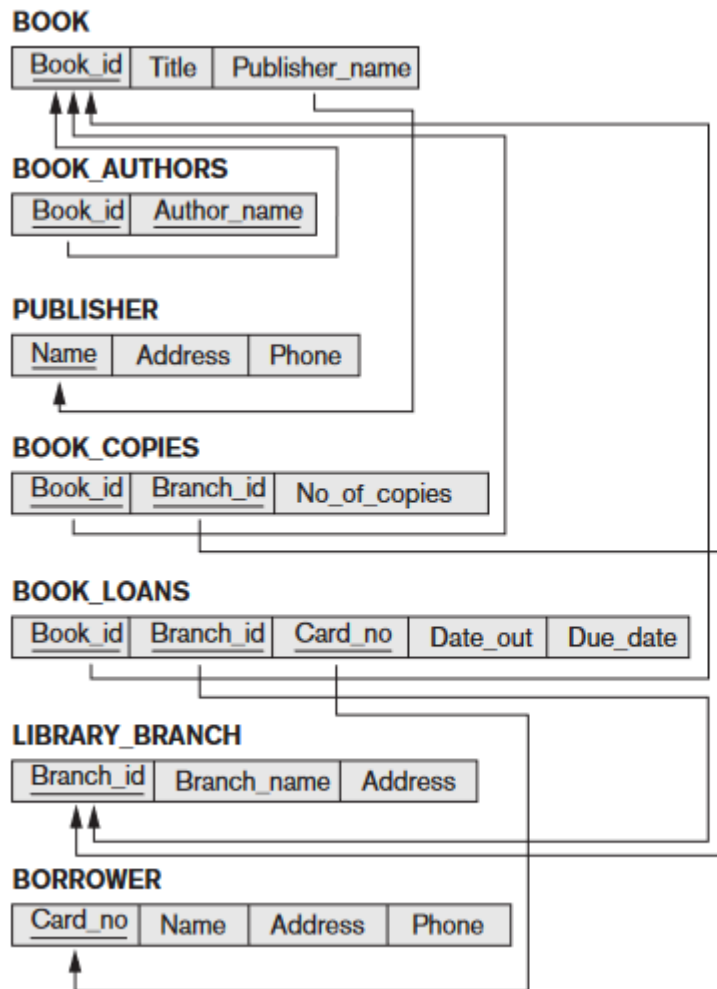


Figure 8.14
A relational database
schema for a LIBRARY
database.

- 8.18. Consider the LIBRARY relational database schema shown in Figure 8.14, which is used to keep track of books, borrowers, and book loans. Referential integrity constraints are shown as directed arcs in Figure 8.14, as in the notation of Figure 5.7. Write down relational expressions for the following queries:
- How many copies of the book titled *The Lost Tribe* are owned by the library branch whose name is 'Sharpstown'?
 - How many copies of the book titled *The Lost Tribe* are owned by each library branch?
 - Retrieve the names of all borrowers who do not have any books checked out.
 - For each book that is loaned out from the Sharpstown branch and whose Due_date is today, retrieve the book title, the borrower's name, and the borrower's address.
 - For each library branch, retrieve the branch name and the total number of books loaned out from that branch.
 - Retrieve the names, addresses, and number of books checked out for all borrowers who have more than five books checked out.
 - For each book authored (or coauthored) by Stephen King, retrieve the title and the number of copies owned by the library branch whose name is Central.

a. $INT1 \leftarrow \sigma_{\langle \text{Branch_name} = \text{'Sharpstown'} \rangle}(\text{LIBRARY_BRANCH})$
 $INT2 \leftarrow (INT1 * \text{BOOK_COPIES})$
 $INT3 \leftarrow \sigma_{\langle \text{Title} = \text{'The Lost Tribe'} \rangle}(\text{BOOK})$
 $INT4 \leftarrow (INT3 * INT2)$
 $RESULT \leftarrow \pi_{\langle \text{No_of_copies} \rangle}(INT4)$

b. $INT1 \leftarrow \sigma_{\langle \text{Title} = \text{'The Lost Tribe'} \rangle}(\text{BOOK})$
 $INT2 \leftarrow \pi_{\langle \text{Book_id} \rangle}(INT1)$
 $INT3 \leftarrow (INT2 * \text{BOOK_COPIES})$
 $INT4 \leftarrow \pi_{\langle \text{Branch_id}, \text{No_of_copies} \rangle}(INT3)$
 $INT5 \leftarrow (INT4 * \text{LIBRARY_BRANCH})$
 $RESULT \leftarrow \pi_{\langle \text{Branch_name}, \text{No_of_copies} \rangle}(INT5)$

c. $INT1 \leftarrow (\text{BORROWER} * \text{BOOK_LOANS})$
 $INT2 \leftarrow \pi_{\langle \text{Card_no}, \text{Name}, \text{Address}, \text{Phone} \rangle}(INT1)$
 $INT3 \leftarrow (\text{BORROWER} - INT2)$
 $RESULT \leftarrow \pi_{\langle \text{Name} \rangle}(INT3)$

d. $INT1 \leftarrow \sigma_{\langle \text{Branch_name} = \text{'Sharpstown'} \rangle}(\text{LIBRARY_BRANCH})$
 $INT2 \leftarrow \pi_{\langle \text{Branch_id} \rangle}(INT1)$
 $INT3 \leftarrow (INT2 * \text{BOOK_LOANS})$
 $INT4 \leftarrow \sigma_{\langle \text{Due_date} = \text{Current Date} \rangle}(INT3)$
 $INT5 \leftarrow (INT4 * \text{BOOK})$
 $INT6 \leftarrow (INT5 * \text{BORROWER})$
 $RESULT \leftarrow \pi_{\langle \text{Title}, \text{Name}, \text{Address} \rangle}(INT6)$

e. $INT1 \leftarrow \pi_{\langle \text{Branch_id}, \text{Branch_name} \rangle}(\text{LIBRARY_BRANCH})$
 $INT2 \leftarrow \sigma_{\langle \text{Branch_id} \rangle \mathfrak{J} \langle \text{COUNT Book_id} \rangle}(\text{BOOK_LOANS})$
 $INT3 \leftarrow (INT1 * INT2)$
 $RESULT \leftarrow \pi_{\langle \text{Branch_name}, \text{COUNT Book_id} \rangle}(INT3)$

f. $INT1 \leftarrow \sigma_{\langle \text{Card_no} \rangle \mathfrak{J} \langle \text{COUNT Book_id} \rangle}(\text{BOOK_LOANS})$
 $INT2 \leftarrow (INT1 * \sigma_{\langle \text{COUNT Book_id} \rangle > 5}(\text{BORROWER}))$
 $RESULT \leftarrow \pi_{\langle \text{Name}, \text{Address}, \text{COUNT Book_id} \rangle}(INT2)$

g. $INT1 \leftarrow \sigma_{\langle \text{Author_name} = \text{'Steven King'} \rangle}(\text{BOOK_AUTHORS})$
 $INT2 \leftarrow (INT1 * \text{BOOK})$
 $INT3 \leftarrow (INT2 * \text{BOOK_COPIES})$
 $INT4 \leftarrow (INT3 * \text{LIBRARY_BRANCH})$
 $INT5 \leftarrow \sigma_{\langle \text{Branch_name} = \text{'Central'} \rangle}(INT4)$
 $RESULT \leftarrow \pi_{\langle \text{Title}, \text{No_of_copies} \rangle}(INT5)$

8.32. A nested query is a query within a query. More specifically, a nested query is a parenthesized query whose result can be used as a value in a number of places, such as instead of a relation. Specify the following queries on the database specified in Figure 5.5 using the concept of nested queries and the relational operators discussed in this chapter. Also show the result of each query as it would apply to the database state in Figure 5.6.

- List the names of all employees who work in the department that has the employee with the highest salary among all employees.
- List the names of all employees whose supervisor's supervisor has '888665555' for Ssn.
- List the names of employees who make at least \$10,000 more than the employee who is paid the least in the company.

a. $RESULT \leftarrow \pi_{\text{Fname}, \text{Lname}} (\sigma_{\text{Dno} = \pi_{\text{Dno}}(\sigma_{\text{Salary} = \text{MAX}(\text{Salary})(\text{EMPLOYEE}))}(\text{EMPLOYEE}))$

b. $RESULT \leftarrow \pi_{\text{Fname}, \text{Lname}} (\sigma_{\text{Super_ssn} = \pi_{\text{Ssn}}(\sigma_{\text{Super_ssn} = \text{'888665555'}(\text{EMPLOYEE}))}(\text{EMPLOYEE}))$

c. $RESULT \leftarrow \pi_{\text{Fname}, \text{Lname}} (\sigma_{\text{Salary} \geq 10000 + (\pi_{\text{Salary}}(\sigma_{\text{MIN}(\text{Salary})(\text{EMPLOYEE}))}(\text{EMPLOYEE}))}$

EMPLOYEE

Fname	Minit	Lname	<u>Ssn</u>	Bdate	Address	Sex	Salary	Super_ssn	Dno
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DEPARTMENT

Dname	<u>Dnumber</u>	Mgr_ssn	Mgr_start_date
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DEPT_LOCATIONS

<u>Dnumber</u>	<u>Dlocation</u>
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PROJECT

Pname	<u>Pnumber</u>	Plocation	Dnum
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WORKS_ON

<u>Essn</u>	<u>Pno</u>	Hours
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DEPENDENT

<u>Essn</u>	<u>Dependent_name</u>	Sex	Bdate	Relationship
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Figure 5.5

Schema diagram for the COMPANY relational database schema.