

1 (https://www.interviewcake.com/question/javascript/coin)

## Imagine you landed a new job as a cashier...

Your quirky boss found out that you're a programmer and has a weird request about something they've been wondering for a long time.

Write a function that, given:

1. an amount of money
2. an array of coin denominations

computes the number of ways to make amount of money with coins of the available denominations.

**Example:** for amount=4 (4¢) and denominations=[1, 2, 3] (1¢, 2¢ and 3¢), your program would output 4—the number of ways to make 4¢ with those denominations:

1. 1¢, 1¢, 1¢, 1¢
2. 1¢, 1¢, 2¢
3. 1¢, 3¢
4. 2¢, 2¢

## Gotchas

We can do this in  $O(n * m)$  time and  $O(n)$  space, where  $n$  is the amount of money and  $m$  is the number of denominations.

A simple recursive approach works, but you'll find that your function gets called more than once with the same inputs. We can do better.

We could avoid the duplicate function calls by memoizing<sup>1</sup>, but there's a cleaner bottom-up<sup>1</sup> approach.

## Breakdown

We need to find some way to break this problem down into subproblems.

Here's one way: for **each denomination**, we can use it once, or twice, or...as many times as it takes to reach or overshoot the amount with coins of that denomination alone.

For each of those choices of how many times to include coins of each denomination, we're left with the subproblem of seeing how many ways we can get the remaining amount from the remaining denominations.

Here's that approach in pseudocode:

```
function numWays(amount, denominations) {
  denominations.forEach(denomination) {
    possibleNumTimesToUseDenominationWithoutOvershootingAmount.forEach(numTimesToUseDenomination) {
      answer += numWays(amountRemaining, otherDenominations);
    }
  }
}
```

The answer for some of those subproblems will of course be 0. For example, there's no way to get 1¢ with only 2¢ coins.

As a recursive function, we could formalize this as:

```
function changePossibilitiesTopDown(amountLeft, denominationsLeft) {
  // base cases:
  // we hit the amount spot on. yes!
  if (amountLeft === 0) return 1;
  // we overshoot the amount left (used too many coins)
  if (amountLeft < 0) return 0;
  // we're out of denominations
  if (denominationsLeft.length === 0) return 0;

  console.log('checking ways to make ' + amountLeft + ' with ' + denominationsLeft);

  // choose a current coin
  var currentCoin = denominationsLeft[0],
      restOfCoins = denominationsLeft.slice(1);

  // see how many possibilities we can get
  // for each number of times to use currentCoin
  var numPossibilities = 0;
  while (amountLeft >= 0) {
    numPossibilities += changePossibilitiesTopDown(amountLeft, restOfCoins);
    amountLeft -= currentCoin;
  }

  return numPossibilities;
}
```

But there's a problem—we'll often **duplicate** the work of checking remaining change possibilities. Note the duplicate calls with the input 4, [1,2,3]:

```
> changePossibilitiesTopDown(4, [1, 2, 3]);
checking ways to make 4 with 1,2,3
checking ways to make 4 with 2,3
checking ways to make 4 with 3
checking ways to make 2 with 3
checking ways to make 3 with 2,3
checking ways to make 3 with 3
checking ways to make 1 with 3
checking ways to make 2 with 2,3
checking ways to make 2 with 3
checking ways to make 1 with 2,3
checking ways to make 1 with 3
< 4
```

For example, we check ways to make 2 with [3] twice.

We can do better. How do we avoid this duplicate work and bring down the time cost?

One way is to memoize<sup>1</sup>.

Here's what the memoization might look like:

```
function Change() {  
  this.memo = {};  
}  
  
Change.prototype.changePossibilitiesTopDown = function(amountLeft, denominationsLeft) {  
  
  // check our memo and short-circuit if we've already solved this one  
  var memoKey = String([amountLeft, denominationsLeft]);  
  if (this.memo.hasOwnProperty(memoKey)) {  
    console.log('grabbing memo[' + memoKey + ']');  
    return this.memo[memoKey];  
  }  
  
  // base cases:  
  // we hit the amount spot on. yes!  
  if (amountLeft === 0) return 1;  
  
  // we overshot the amount left (used too many coins)  
  if (amountLeft < 0) return 0;  
  
  // we're out of denominations  
  if (denominationsLeft.length === 0) return 0;  
  
  console.log('checking ways to make ' + amountLeft + ' with ' + denominationsLeft);  
  
  // choose a current coin  
  var currentCoin = denominationsLeft[0];  
  var restOfCoins = denominationsLeft.slice(1);  
  
  // see how many possibilities we can get  
  // for each number of times to use currentCoin  
  var numPossibilities = 0;  
  while (amountLeft >= 0) {  
    numPossibilities += this.changePossibilitiesTopDown(amountLeft, restOfCoins);  
    amountLeft -= currentCoin;  
  }  
  
  // save the answer in our memo so we don't compute it again  
  this.memo[memoKey] = numPossibilities;  
  return numPossibilities;  
};
```

And now our checking has no duplication:

```

> new Change().changePossibilitiesTopDown(4, [1, 2, 3]);
checking ways to make 4 with 1,2,3
checking ways to make 4 with 2,3
checking ways to make 4 with 3
checking ways to make 2 with 3
checking ways to make 3 with 2,3
checking ways to make 3 with 3
checking ways to make 1 with 3
checking ways to make 2 with 2,3
grabbing memo[2,3]
checking ways to make 1 with 2,3
grabbing memo[1,3]
< 4

```

This answer is quite good. It certainly solves our duplicate work problem. It takes  $O(n * m)$  time and  $O(n * m)$  space, where  $n$  is the size of amount and  $m$  is the number of items in denominations.

However, we can do better. Because our function is recursive it will build up a **large call stack**<sup>1</sup> of size  $O(m)$ . Of course, this cost is eclipsed by the memory cost of this `memo`, which is  $O(n * m)$ . But it's still best to avoid building up a large stack like this, because it can cause a **stack overflow** (yes, that means recursion is *usually* better to avoid for functions that might have arbitrarily large inputs).

It turns out we can get  $O(n)$  additional space.

A great way to avoid recursion is to go **bottom-up**<sup>1</sup>.

Our recursive approach was top-down because it started with the final value for amount and recursively broke the problem down into subproblems with smaller values for amount. What if instead we tried to **compute the answer for small values of amount first**, and use those answers to iteratively compute the answer for higher values until arriving at the final amount?

We can **start by making an array** `waysOfDoingNCents`, where the index is the amount and the value at each index is the number of ways of getting that amount.

This array will take  $O(n)$  space, where  $n$  is the size of amount.

To further simplify the problem, we can work with only the first coin in denominations, then add in the second coin, then the third, etc.

What would `waysOfDoingNCents` look like for just our first coin: 1¢? Let's call this `waysOfDoingNCents1`.

```

var waysOfDoingNCents1 = [
  1, // 0c: no coins
  1, // 1c: 1 1c coin
  1, // 2c: 2 1c coins
  1, // 3c: 3 1c coins
  1, // 4c: 4 1c coins
  1, // 5c: 5 1c coins
];

```

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Now what if we add a 2¢ coin?

```
var waysOfDoingNCents1And2 = [
  1,    // 0c: no change
  1,    // 1c: no change
  1+1,  // 2c: new [(2)]
  1+1,  // 3c: new [(2,1)]
  1+2,  // 4c: new [(2,1,1), (2,2)]
  1+2,  // 5c: new [(2,1,1,1), (2,2,1)]
];
```

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How do we formalize this process of going from waysOfDoingNCents1 to waysOfDoingNCents1And2?

Let's **suppose we're partway through already** (this is a classic dynamic programming approach). Say we're trying to calculate waysOfDoingNCents1And2[5]. Because we're going bottom-up, we know we already have:

1. waysOfDoingNCents1And2 for amounts less than 5
2. a fully-populated waysOfDoingNCents1

So how many *new* ways should we add to waysOfDoingNCents1[5] to get waysOfDoingNCents1And2[5]?

Well, if there are *any* new ways to get 5¢ now that we have 2¢ coins, those new ways must involve at least one 2¢ coin. So if we presuppose that we'll use one 2¢ coin, that leaves us with  $5 - 2 = 3$  left to come up with. We already know how many ways we can get 3¢ with 1¢ and 2¢ coins: waysOfDoingNCents1And2[3], which is 2.

So we can see that:

```
waysOfDoingNCents1And2[5] = waysOfDoingNCents1[5] + waysOfDoingNCents1And2[5-2]
```

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**Why don't we also need to check waysOfDoingNCents1And2[5 - 2 - 2] (two 2¢ coins)?** Because we already checked waysOfDoingNCents1And2[1] when calculating waysOfDoingNCents1And2[3]. We'd be counting some arrangements multiple times. In other words, waysOfDoingNCents1And2[k] already includes the full count of possibilities for getting  $k$ , including possibilities that use 2¢ any number of times. We're only interested in how many *more* possibilities we might get when we go from  $k$  to  $k + 2$  and thus have the ability to add one *more* 2¢ coin to each of the possibilities we have for  $k$ .

## Solution

We use a bottom-up<sup>1</sup> algorithm to build up a table waysOfDoingNCents such that waysOfDoingNCents[k] is how many ways we can get to  $k$  cents using our denominations. We start with the base case that there's **one way to create the amount zero**, and progressively add each of our denominations.

The number of new ways we can make a higherAmount when we account for a new coin is simply waysOfDoingNCents[higherAmount - coin], where we know that value already includes combinations involving coin (because we went bottom-up, we know smaller values have already been calculated).

```
function changePossibilitiesBottomUp(amount, denominations) {  
    // initialize an array of zeros with indices up to amount  
    var waysOfDoingNcents = [];  
    for (var i = 0; i <= amount; i++) {  
        waysOfDoingNcents[i] = 0;  
    }  
    waysOfDoingNcents[0] = 1;  
  
    for (var j = 0; j < denominations.length; j++) {  
        var coin = denominations[j];  
        for (var higherAmount = coin; higherAmount <= amount; higherAmount++) {  
            var higherAmountRemainder = higherAmount - coin;  
            waysOfDoingNcents[higherAmount] += waysOfDoingNcents[higherAmountRemainder];  
        }  
    }  
  
    return waysOfDoingNcents[amount];  
}
```

Here's how waysOfDoingNCents would look in successive iterations of our function for amount=5 and denominations=[1,3,5].

```

=====
key:
a = higherAmount
r = higherAmountRemainder
=====

=====
for coin = 1:
=====
[1, 1, 0, 0, 0, 0]
r a

[1, 1, 1, 0, 0, 0]
r a

[1, 1, 1, 1, 0, 0]
r a

[1, 1, 1, 1, 1, 0]
r a

[1, 1, 1, 1, 1, 1]
r a

function numWaysToMakeChange(a)
var numWaysChangeN

=====
for coin = 3:
=====
[1, 1, 1, 2, 1, 1]
r a

[1, 1, 1, 2, 2, 1]
r a

[1, 1, 1, 2, 2, 2]
r a

=====
for coin = 5:
=====
[1, 1, 1, 2, 2, 3]
r a

final answer: 3

```

## Complexity

$O(n * m)$  time and  $O(n)$  additional space, where  $n$  is the amount of money and  $m$  is the number of potential denominations.

Did you get it right?



Yes, I'm expert on this



Not quite, review later

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