

Assignment : 3

Calculate time complexity : - Recursive Recurrence Relation

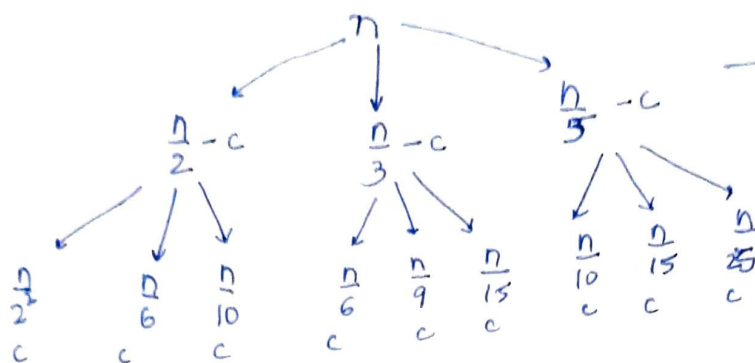
$$T(n) = T(n/2) + T(n/3) + T(n/5) + c$$

Different Cost

$$c \rightarrow 3^0 c$$

$$\text{Level 1} \\ 3c = 3^1 c$$

$$\text{Level 2} \\ 9c = 3^2 c$$



we get an idea that we have a GP series for cost.
 $3^0 c + 3^1 c + 3^2 c + \dots + 3^K c = \text{Overall time complexity Cost}$

$K \rightarrow \text{Order of Magnitude}$
 (No. of times recursion)

Let's talk about base condition : stopping criteria.

$$\frac{n}{2^{K_L}} = 1$$

$$n = 2^{K_L}$$

$$\Rightarrow K_L = \log_2 n$$

Lower BASE
 Higher the value.

$$\frac{n}{3^{K_M}} = 1$$

$$n = 3^{K_M}$$

$$\Rightarrow K_M = \log_3 n$$

$$\frac{n}{5^{K_R}} = 1$$

$$n = 5^{K_R}$$

$$\Rightarrow K_R = \log_5 n$$

$$\begin{aligned} \text{Overall time complexity} &= O(\text{cost} \times K_E) \\ &= O(3^0 c + 3^1 c + 3^2 c + \dots + 3^{K_L} c) \end{aligned}$$

$$\left\{ \begin{aligned} \text{GP series: sum} &= \frac{a(r^n - 1)}{r - 1} \quad r > 1 \\ &= \frac{1(3^{\log_2 n} - 1)}{3 - 1} = \frac{(n^{\log_2 3} - 1)}{3 - 1} \\ &\approx (n^{1.58}) \end{aligned} \right.$$

$$\text{Value } \log_2 3 \\ = 1.58$$

$$\text{Time Complexity} = O(n^{1.58}) \quad \underline{\underline{\text{Ans.}}}$$

(2) find the time complexity of reverse one relation by Substitution method.

$$T(n) = \begin{cases} 1 & n=1 \\ 2T(\frac{n}{2}) + n & n>1 \end{cases}$$

stopping criteria $T(1) = 1$

$$T(n) = 2T(\frac{n}{2}) + n \longrightarrow \text{level 1}$$

$$= 2 \left[2T(\frac{n}{2^2}) + \frac{n}{2} \right] + n$$

$$= 2^2 T(\frac{n}{2^2}) + 2n \longrightarrow \text{level 2}$$

$$= 2^2 \left[2T(\frac{n}{2^3}) + \frac{n}{2^2} \right] + 2n$$

$$= 2^3 T(\frac{n}{2^3}) + 3n \longrightarrow \text{level 3}$$

\downarrow k times

$$= 2^k T(\frac{n}{2^k}) + kn \longrightarrow k^{\text{th}} \text{ level.} \quad \text{Order of Magnitude.}$$

base condition $\rightarrow T(1) = 1$

$$\begin{aligned} T\left(\frac{n}{2^k}\right) &= 1 \quad \Rightarrow \quad \frac{n}{2^k} = 1 \\ \frac{n}{2^k} &= 1 \end{aligned} \quad \boxed{k = \log_2 n}$$

$$T(n) = 2^{\log_2 n} T\left(\frac{n}{2^{\log_2 n}}\right) + n \log_2 n$$

$$= n^{\log_2 2} T\left(\frac{n}{n^{\log_2 2}}\right) + n \log_2 n$$

$$= n T(1) + n \log_2 n$$

$$= n + n \log_2 n$$

$$O(n) < O(n \log_2 n)$$

$$\underline{\underline{\text{Time Complexity} = O(n \log_2 n)}}$$

⑧ Calculate time complexity of recursive relation by substitution method.

$$T(n) = \begin{cases} 1 & n \geq 1 \\ 8T\left(\frac{n}{2}\right) + n^2 & n > 1 \end{cases}$$

stopping criteria $\rightarrow T(1) = 1$

$$T(n) = 8T\left(\frac{n}{2}\right) + n^2 \quad \rightarrow \text{level 1}$$

$$= 8 \left[8T\left(\frac{n}{2^2}\right) + \left(\frac{n}{2^2}\right)^2 \right] + n^2$$

$$= 8^2 T\left(\frac{n}{2^2}\right) + 3n^2 \quad \rightarrow \text{level 2}$$

$$= 8^2 \left[8T\left(\frac{n}{2^3}\right) + \left(\frac{n}{2^3}\right)^2 \right] + 3n^2$$

$$= 8^3 T\left(\frac{n}{2^3}\right) + 8^2 \frac{n^2}{2^4} + 3n^2$$

$$= 8^3 T\left(\frac{n}{2^3}\right) + 8 \cdot 2^2 n^2 + 3n^2$$

$$= 8^3 T\left(\frac{n}{2^3}\right) + 7n^2 \quad \rightarrow \text{level 3}$$

\downarrow k times.

$$T(n) = 8^k T\left(\frac{n}{2^k}\right) + (2^k - 1) n^2 \quad \rightarrow k^{\text{th}} \text{ level.}$$

1, 3, 7, \dots k

$(2^k - 1)_{k=0,1,2,\dots}$

stopping criteria $T(1) = 1$

$$\text{base condition} \Rightarrow T\left(\frac{n}{2^k}\right) = 1 \quad \Rightarrow \quad \frac{n}{2^k} = 1 \quad \Rightarrow \quad \boxed{k = \log_2 n}$$

substitute value of k in $T(n)$

$$T(n) = 8^{\log_2 n} T\left(\frac{n}{2^{\log_2 n}}\right) + (2^{\log_2 n} - 1) n^2$$

$$= n^{\log_2 8} T\left(\frac{n}{n}\right) + (n^{\log_2 2} - 1) n^2$$

$$= n^3 T\left(\frac{n}{n}\right) + (n - 1) n^2$$

$$= n^3 + n^3 - n^2$$

$$O(n^2) < O(n^3)$$

$$\text{Time Complexity} = O(n^3)$$