

# TP5 Report

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## First Part

In this part we had to solve 2 systems of congruences a) and b)

### Notes:

- We'll consider the definition of a function  $\text{gcd}(a, b)$  that gives the greater common divisor of two values  $a$  and  $b$ .
- $k, l$  and  $m$  are positive integers
- The congruences will be convertible to equations in the form:

$$x \equiv a \pmod{b} \rightarrow x = b \cdot k + a, k \text{ being a positive integer}$$

### a)

$$\begin{aligned} (1): x &\equiv 48 \pmod{13} \\ (2): x &\equiv 57 \pmod{23} \\ (3): x &\equiv 39 \pmod{27} \end{aligned}$$

Firstly we begin with the congruency with largest modulus that is (3)  $x \equiv 39 \pmod{27}$

Then we substitute this congruences expression for  $X$  into the congruence with the next largest modulus (2):

$$27 \cdot m + 39 \equiv 57 \pmod{23}$$

Solving this Linear Congruency ...

$$27 \cdot m \equiv 18 \pmod{23}$$

note:  $\text{gcd}(27, 23) = 1$  so there is a solution

$$\text{gcd}(27, 18) = 9$$

$$3 \cdot m \equiv 2 \pmod{23}$$

$$3 \cdot m \equiv -21 \pmod{23}$$

$$m \equiv -7 \pmod{23}$$

$$m \equiv 16 \pmod{23}$$

$$m = 23 \cdot k + 16$$

Replacing the expression form of this result in the expression for  $X$  we get:

$$X = 27*(23*l + 16) + 39$$

$$X = 621*l + 471$$

Then we replace this expression in the last congruency (1) and solve this Linear Congruency

$$621*l + 471 \equiv 48 \pmod{13}$$

note:  $\gcd(48, 13) = 1$  so there is a solution

$$621*l \equiv -423 \pmod{13}$$

$$621*l \equiv 6 \pmod{13}$$

$$207*l \equiv 2 \pmod{13}$$

$$207*l \equiv -24 \pmod{13}$$

$$\gcd(24, 207) = 3$$

$$69*l \equiv -8 \pmod{13}$$

...

$$l \equiv 11 \pmod{13}$$

Finally we replace this in the expression obtained previously and get the solution

$$X = 621*(13*m + 11) + 471$$

$$X = 8073*m + 7302$$

$$X = 7302$$

The smallest solution for this system is 7302

**b)**

For the second system the congruences we first simplified the each one to remove the coefficient

$$19*X \equiv 21 \pmod{16}$$

$$37*X \equiv 100 \pmod{15}$$

Solving this Linear Congruences ...

$$(1): X \equiv 7 \pmod{16}$$

$$(2): X \equiv 10 \pmod{15}$$

With the same intension as before, we begin with the congruency with largest modulus that is (1)  $X \equiv 7 \pmod{16}$

Then we substitute this congruences expression for X into the congruence with the next largest modulus (2):

$$16*k + 7 \equiv 10 \pmod{15}$$

Solving this Linear Congruence ...

$$k \equiv 3 \pmod{15}$$

$$k = 15 \cdot l + 3$$

Replacing the expression form of this result in the expression for X we get:

$$X = 16 \cdot (15 \cdot l + 3) + 7$$

$$X = 240 \cdot l + 55$$

$$X = 55$$

The smallest solution for this system is 55

## Second Part

For this part we had the objective to decrypt a cipher which was encrypted with RSA with  $e = 17$  and  $n = 213271$ , beyond this, the ciphered integer was a result of an encoding of three letters in a row in the form:

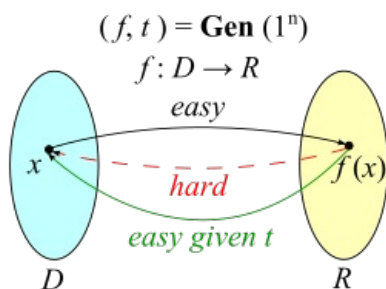
$$27^2 L_1 + 27 L_2 + L_3$$

For this encoding we developd two functions to make conversions in this scheme:

$$\begin{aligned} n &= \text{encode}(l_1, l_2, l_3) \\ l_1, l_2, l_3 &= \text{decode}(n) \end{aligned}$$

Firstly, since we know the public key, and that the value of the modulus is small, we can compute the prime factorization in relevant time since this value isn't bigger enough.

The key concept in RSA is that it is based on a trapdoor function, a function that is easy to compute one way but difficult in the opposite direction unless we know the value  $t$ .



For the RSA this function is:

$$f(x) = x^e \pmod{n}$$

More specifcly, we can compute the private key  $d$ :

$$(1) d = e^{-1} \pmod{t}$$

The private key can be computed knowing  $e$ , and the prime factorization of  $n$  as  $p$  and  $q$

$t$  is the trapdoor value which can be computed as the Least Common Multiple of  $p$  and  $q$  ( $\text{lcm}(p, q)$ ).

$$\text{lcm}(a, b) = \frac{|a \cdot b|}{\text{gcd}(a, b)}$$

To obtain this value we developed the following auxiliar functions:

```
gcd(a,b)
t = lcm(a,b)
```

For the prime factorization we provided a function `prime_factorization(n)` that returns a list of all prime factors

Finally, we apply (1) with all the data and obtain the value  $d$  that is the private/decryption key (together with  $n$ )

Once we obtain this value  $d$  we can decrypt the ciphertext and obtain the plaintext by decoding by the specified scheme.

```
for c in cipher:
    l1, l2, l3 = decode((c**d) % n)
```

After joining every letter decripted we obtain the full plaintext, this text can be found in the file `plaintext.txt` provided