

# Mathematical Foundations of Artificial Intelligence and Machine Learning (NCM-CEP Course)

Guest Lecture on Multi-agent AI

(a.k.a. **Game Theory and Mechanism Design**)

Swaprava Nath



CSE, IIT Bombay

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# Let us play a game: Neighboring King(Queen)dom's Dilemma

Each kingdom can invest either in Agriculture or War – but not both

- If both choose Agri – happiness is 5 for each
- If both choose War – happiness is 1 for each
- If one chooses Agri, but the other War – the Agri kingdom stand to lose everything and War kingdom gets happiness more than 5

What should the King / Queen do?  
What is collectively best?  
(Agriculture, Agriculture)

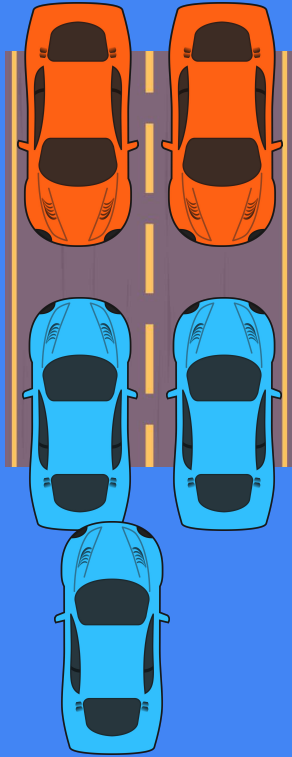


	Agriculture	War
Agriculture	<div style="border: 2px solid green; padding: 5px; display: inline-block;">○</div>	○
War	<div style="text-align: center;">↓</div> <div style="display: inline-block;">○</div>	<div style="text-align: center;">↓</div> <div style="display: inline-block;">○</div>

Decisions are simultaneous

War is a **Dominant Strategy** for Queen as well as King  
(War, War) is a **Dominant Strategy Equilibrium**

## Another game: Traffic Movement



Does this game have a dominant strategy equilibrium?

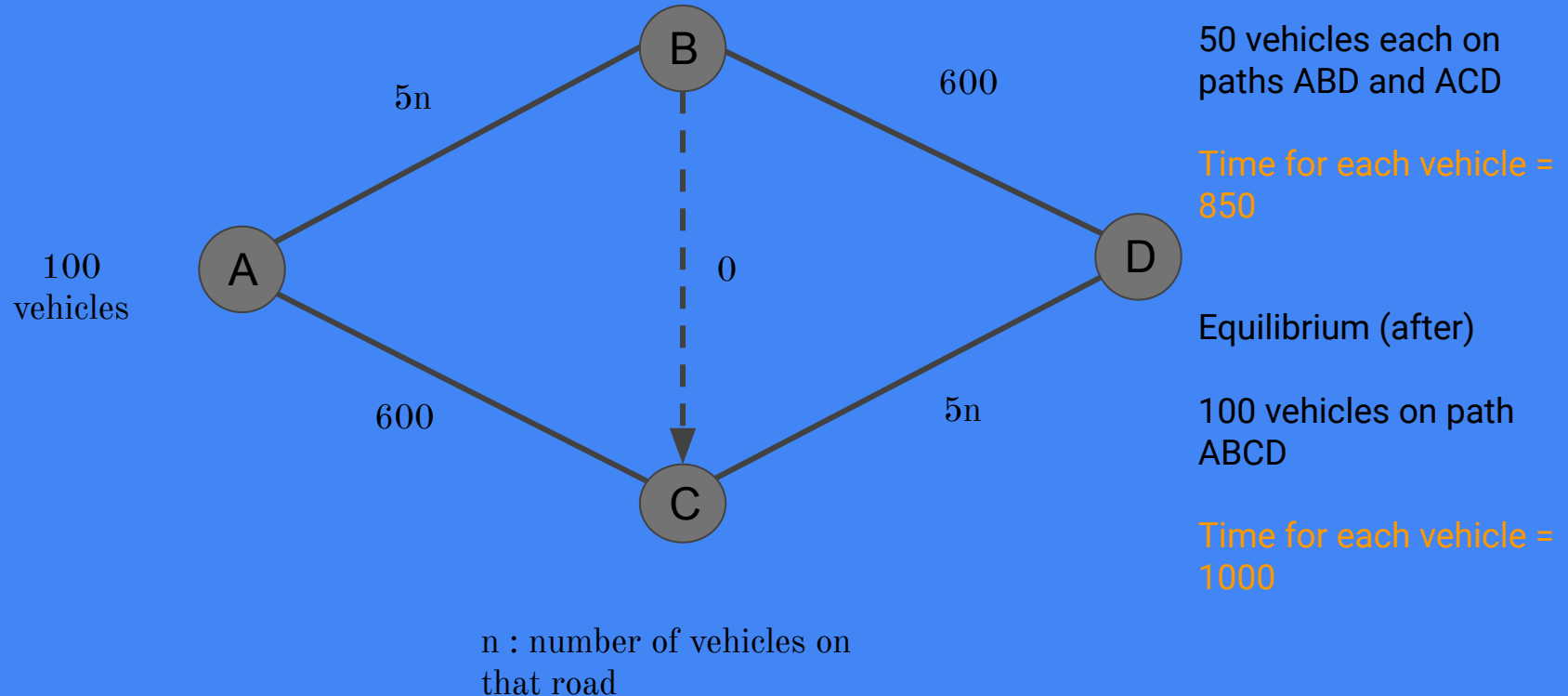


The Nash Equilibrium (John Nash, 1951)

	Left	Right
Left	<div></div>	
Right		<div></div>

Equilibrium here is a **strategy profile** from where no player wants to **unilaterally** deviate

# Adding resources (blindly) does not improve the society

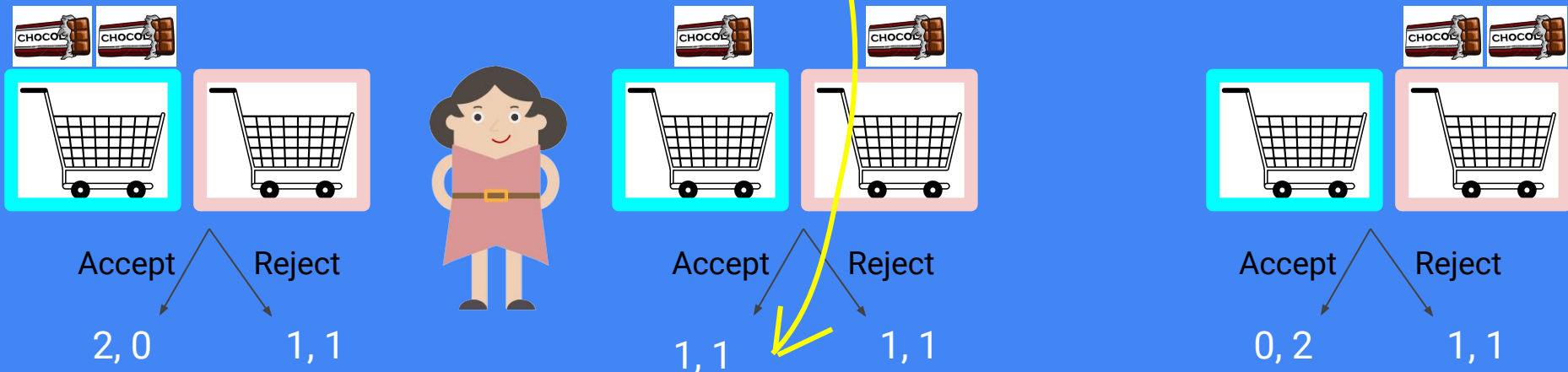


# All games do not end in the same round

Subgame Perfect  
Nash Equilibrium

Found via Backward  
Induction

Settle  
yourselves  
amicably or I'll  
give one to  
each of you



# Can we design algorithms for a better society?

## Fair division

Equal amount may  
have different values  
for an agent

Heterogeneous

Divisible

Differing preferences

Any fractional  
allocation is  
feasible

Different agents may have  
different preferences for the  
same piece



# Proportional division

For each agent  $i$

$$v_i(A_i) \geq 1/n$$

Each agent gets at least the average share

**Normalization:** for each  $i$ ,  $v_i([0,1]) = 1$



**“I cut, you choose” algorithm**

# Envy free division

For each pair of agents  $i, j$

$$v_i(A_i) \geq v_i(A_j)$$

Each agent likes her own share at least as much as others

Proportional?

Yes, agent 1 cuts  $1/2$ , and agent 2 picks the larger

Envy-free?

Yes, agent 1 gets  $1/2$ , which is the same as the other piece in his view  
agent 2 picks first, can't envy the other piece

# Fair division of indivisible objects

All are  
indivisible  
objects



Faculty retires and wants to give away  
his/her belongings to the department  
staff / existing faculty

Items:

1. Books
2. Shelves
3. Furnitures
4. Wall decor
5. Table decor
6. Electronic gadgets
7. Many more ...

Notice that Envy-free  
allocation is no longer  
possible

Consider a single item and  
two agents

Envy free upto one  
good allocation

For each pair of agents  $i, j$

$$v_i(A_i) \geq v_i(A_j \setminus x_j), \text{ for some } x_j \in A_j$$

Each agent likes her  
own share upto all but  
one item of every other  
agent



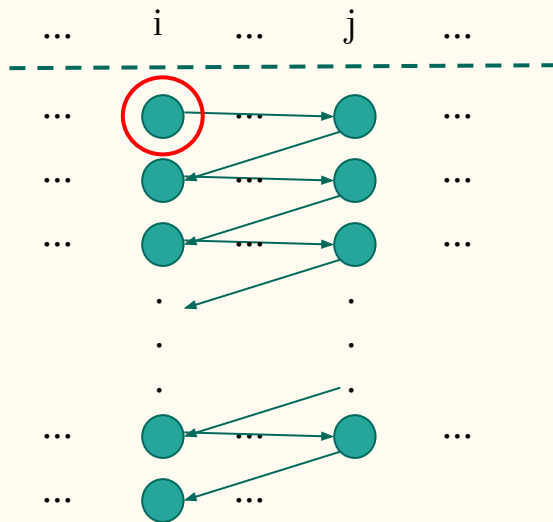
# Envy free upto one good (EF-1)

All are  
indivisible  
objects

Always exists and computable in polynomial time!

Example for additive valuations: Round-Robin Algorithm

Place the agents in any arbitrary order, and ask them to pick their favorite remaining item



i does not envy j, since  
it picks before j

j may envy i, but not if  
the first item i picked  
is dropped

Round-Robin achieves EF-1  
for additive valuations

## Envy free upto one good allocation

For each pair of agents i, j

$$v_i(A_i) \geq v_i(A_j \setminus x_j), \text{ for some } x_j \in A_j$$

Each agent likes her  
own share upto all but  
one item of every other  
agent

- A game is an interaction between agents who want to maximize their utilities
- **Game theory** predicts the outcome of a game
- This is a predictive approach
- **Mechanism design** tries to design the game with desirable outcomes
- This is a prescriptive approach

## References:

1. Lecture notes and videos from the modules 1, 2, 5, 6, 7 of course CS 6001 ([webpage](#)), and
2. For the fair division, lectures 7 and 8 of course CS 6002 ([webpage](#))

## General fun reading:

[Nash equilibrium](#), [Braess's paradox](#)

If you are interested in probing further, two courses in the CSE department are relevant – lecture materials (and videos for CS 6001) are available

1. CS 6001: Game Theory and Algorithmic Mechanism Design
2. CS 6002: Selected Areas of Mechanism Design

# Thanks!

[swaprava@cse.iitb.ac.in](mailto:swaprava@cse.iitb.ac.in)

<http://www.cse.iitb.ac.in/~swaprava>