

# Solving Linear Systems: Gaussian Elimination

Given:  $n$  equations in  $n$  variables,  $x_1, x_2, \dots, x_n$

Output: A ~~one~~ unique solution  $\vec{x} = (x_1, x_2, \dots, x_n)$  to the system

$$3x_1 - x_2 + 2x_3 = 1$$

$$4x_1 + 2x_2 - x_3 = 4$$

$$x_1 + 3x_2 + 3x_3 = 5$$

$$\begin{array}{c} A \\ \rightarrow \\ n \left\{ \begin{array}{ccc|c} x_1 & x_2 & x_3 & \\ \hline 3 & -1 & 2 & 1 \\ 4 & 2 & -1 & 4 \\ 1 & 3 & 3 & 5 \end{array} \right. \end{array}$$

$\underbrace{\hspace{10em}}_n$

rows = equations

cols = variables

$n \times (n+1)$  matrix

Augmented • Last column: scalar value

Goal: make  $A$  into an upper triangular matrix

$$\begin{bmatrix} \bullet & \# & \# & \# \\ & \# & \# & \# \\ -0- & & \# & \# \\ & & \cdot & \cdot \end{bmatrix}$$

$\underbrace{\hspace{10em}}_{\vec{a} \cdot \vec{x}_n = b_n}$

$$\begin{array}{c} x_n \\ x_{n-1} \\ a \end{array} \quad \begin{array}{c} a \square b \\ \uparrow \quad \curvearrowright \\ a \end{array}$$

valid elementary operations:

- multiply a row by a constant
- divide ~~by~~ row by a constant other than zero
- exchange rows
- exchange columns? yes as long as you keep track:  
exchanges variables
- you can add/subtract a row from another row

$$\begin{bmatrix} 3 & -1 & 2 & 1 \\ 4 & 2 & -1 & 4 \\ 1 & 3 & 3 & 5 \end{bmatrix} \begin{array}{l} \text{subtract } \frac{4}{3} \times \text{row 1 from row 2} \\ \text{subtract } \frac{1}{3} \times \text{row 1 from row 3} \end{array} \begin{bmatrix} 3 & -1 & 2 & 1 \\ 0 & \frac{10}{3} & -\frac{11}{3} & \frac{8}{3} \\ 0 & \frac{10}{3} & \frac{7}{3} & \frac{14}{3} \end{bmatrix}$$

→

Subtract 1 × row 2 from row 3 →

$$\begin{array}{c} \vec{x} \\ \begin{matrix} x_1 & x_2 & x_3 & b \end{matrix} \\ \left[ \begin{array}{ccc|c} 3 & -1 & 2 & 1 \\ 0 & \frac{10}{3} & -\frac{11}{3} & \frac{8}{3} \\ 0 & 0 & \frac{18}{3} & \frac{6}{3} \end{array} \right] \\ \begin{matrix} (6) & (2) \end{matrix} \end{array}$$

~~line~~ row 3:

$$\begin{aligned} 6x_3 &= b_3 \\ &= 2 \end{aligned}$$

$$x_3 = \frac{1}{3}$$

$$\text{row 1: } 3x_1 - x_2 + 2x_3 = 1$$

$$\vec{x} = \left( \frac{1}{2}, \frac{7}{6}, \frac{1}{3} \right)$$

row 2:

$$\frac{10}{3} x_2 - \frac{11}{3} x_3 = \frac{8}{3} \quad \frac{11}{9} x_3 = \frac{8}{24} + \frac{11}{9} \cdot \frac{1}{30} \quad x_2 = \frac{7}{6}$$

$$x_1 = \frac{1}{2}$$

Outline:

$x_1 \quad x_2 \quad \dots$

$$A \begin{bmatrix} a_{11} & a_{12} & \dots & a_{1n} & b_1 \\ \boxed{a_{21}} & a_{22} & \dots & a_{2n} & b_2 \\ \vdots & \boxed{a_{n2}} & \dots & a_{nn} & b_n \end{bmatrix}$$

$$a_{21} \rightarrow 0$$

$$a_{21} - \frac{a_{21}}{a_{11}} \cdot a_{11} = 0$$

Step  $i \quad i = 1 \dots n-1$

Step 1: eliminate  $a_{21} \ a_{31} \dots a_{n1}$

for  $j = \cancel{i+1} \dots n$

$$t \leftarrow \frac{a_{j1}}{a_{\cancel{i+1}1}}$$

for  $k = j \dots n+1$

$$L \ a_{jk} \leftarrow a_{jk} - t \cdot a_{\cancel{i+1}k}$$

Step 2: eliminate  $a_{32} \ a_{42} \ a_{52} \dots a_{n2}$

for  $j = \cancel{i+1} \dots n$

$$t \leftarrow \frac{a_{j2}}{a_{\cancel{i+1}2}}$$

for  $k = j \dots n+1$

$$L \ a_{jk} \leftarrow a_{jk} - t \cdot a_{\cancel{i+1}k}$$

Step  $n-1$

Input: An  $n \times (n+1)$  augmented matrix,  $A$

Output: An upper triangular matrix  $A'$

$$A' \leftarrow A$$

for  $i = 1 \dots n-1$

    pivotrow  $\leftarrow i$

    for  $j = (i+1) \dots n$

        if  $(|A'[j,i]| > |A'[pivotrow,i]|)$

            pivotrow  $\leftarrow j$

    for  $k = i \dots (n+1)$

        swap  $A'[i,k], A'[pivotrow,k]$

    for  $j = (i+1) \dots n$

$t \leftarrow \frac{A'[j,i]}{A'[i,i]}$

        for  $k = i \dots (n+1)$

$A'[j,k] \leftarrow A'[j,k] - A'[i,k] \cdot t$

Find the "best" pivot row

$$\sum_{i=1}^{n-1}$$

Output  $A'$

$$\sum_{i=1}^{n-1} \sum_{j=i+1}^n \sum_{k=i}^{n+1} 1 = \sum_{i=1}^{n-1} \sum_{j=i+1}^n (n-i+2)$$

input size:  $n^2 = N$

$$= \sum_{i=1}^{n-1} \left[ \sum_{j=i+1}^n n - \sum_{j=i+1}^n i + \sum_{j=i+1}^n 2 \right]$$

$$= \sum_{i=1}^{n-1} \left[ n(n-i) - i(n-i) + 2(n-i) \right]$$

$$\begin{matrix} (n+1) - (i-1) \\ n+2-i \end{matrix}$$

$$= \sum_{i=1}^{n-1} \left[ n^2 - ni - ni + i^2 + 2n - 2i \right]$$

$$= (n-1)n^2 - 2n \sum_{i=1}^{n-1} i + \boxed{\sum_{i=1}^{n-1} i^2} + 2(n-1)n - 2 \sum_{i=1}^{n-1} i$$

$$= \textcircled{n^3} - n^2 \dots \dots \dots \downarrow$$

$$\in \Theta(n^3) \longrightarrow \Theta(N^{3/2})$$

~~Gauss's Formula~~

$$\sum_{i=1}^n i = \frac{n(n+1)}{2}$$

$$\left[ \begin{array}{cccccc} a_{11} & a_{12} & a_{13} & \dots & a_{1n} & b_1 \\ 0 & a_{22} & a_{23} & \dots & a_{2n} & b_2 \\ 0 & 0 & a_{33} & \dots & a_{3n} & b_3 \\ \vdots & \vdots & \vdots & \ddots & \vdots & \vdots \\ 0 & 0 & 0 & \dots & a_{nn} & b_n \end{array} \right]$$

Back solve:

Determine values for

$$x_n, x_{n-1}, x_{n-2}, \dots, x_2, x_1$$

$$x_n = \frac{b_n}{a_{nn}}$$

$$x_{n-1} = \frac{b_{n-1} - a_{n-1,n} x_n}{a_{n-1,n-1}}$$

$$= \frac{b_{n-1} - a_{n-1,n} x_n}{a_{n-1,n-1}}$$

$$\begin{array}{ccc} x_{n-1} & x_n & \text{RHS} \\ a_{n-1,n-1} & a_{n-1,n} & b_{n-1} \end{array}$$

$$a_{n-1,n-1} x_{n-1} + a_{n-1,n} x_n = b_{n-1}$$

$$x_{n-2} = \frac{b_{n-2} - a_{n-2,n} x_n - a_{n-2,n-1} x_{n-1}}{a_{n-2,n-2}}$$

$$x_i = \frac{b_i - a_{i,i+1} x_{i+1} - a_{i,i+2} x_{i+2} - \dots - a_{i,n} x_n}{a_{i,i}}$$

Back Solve

Input: An  $n \times (n+1)$  upper triangular matrix

$n^2$

Output: A solution vector  $\vec{x} = (x_1, x_2, \dots, x_n)$

for  $i = n \dots 1$

$t \leftarrow A[i, n+1]$  // sets  $t$  to  $b_i$

for  $j = (i+1) \dots \text{~~(n+1)~~} n$

$t \leftarrow t - x_j \cdot A[i, j]$

$x_i \leftarrow t / A[i, i]$

Output  $(x_1, \dots, x_n)$

$$\sum_{i=1}^n \sum_{j=i+1}^n 1 = ?$$

$$\in \Theta(n^2)$$



Application:

• Solve Systems of linear equations

$$x \rightarrow x^{-1}$$

$$x \cdot \frac{1}{x} = 1$$

$$I = \begin{bmatrix} 1 & & & \\ & 1 & & \\ & & \ddots & \\ 0 & & & 1 \end{bmatrix}$$

Given  $A$ , what is  $A^{-1}$  (or does it exist?)

$$n \times \underbrace{\begin{bmatrix} A & | & I \end{bmatrix}}_{2n} \rightarrow \text{gaussian elimination} \left[ \begin{array}{c|c} \square & \square \\ \hline \oplus & \square \end{array} \right]$$

$\rightarrow$  again for the upper diagonal + divide by the diagonal

$$\left[ I \mid A^{-1} \right]$$

Detect

Indeterminate Systems (infinite solutions)

A row of all zeros  $\Rightarrow$  indeterminate system

Inconsistent systems (zero solutions)

A row of all zeros except the last column

$$\begin{array}{ccc} 1 & 1 & 2 \\ 2 & 2 & 5 \end{array} \rightarrow \begin{bmatrix} 1 & 1 & 2 \\ 0 & 0 & 1 \end{bmatrix}$$

$$0=1$$

$$x = z - y$$

$$x + y = 2$$

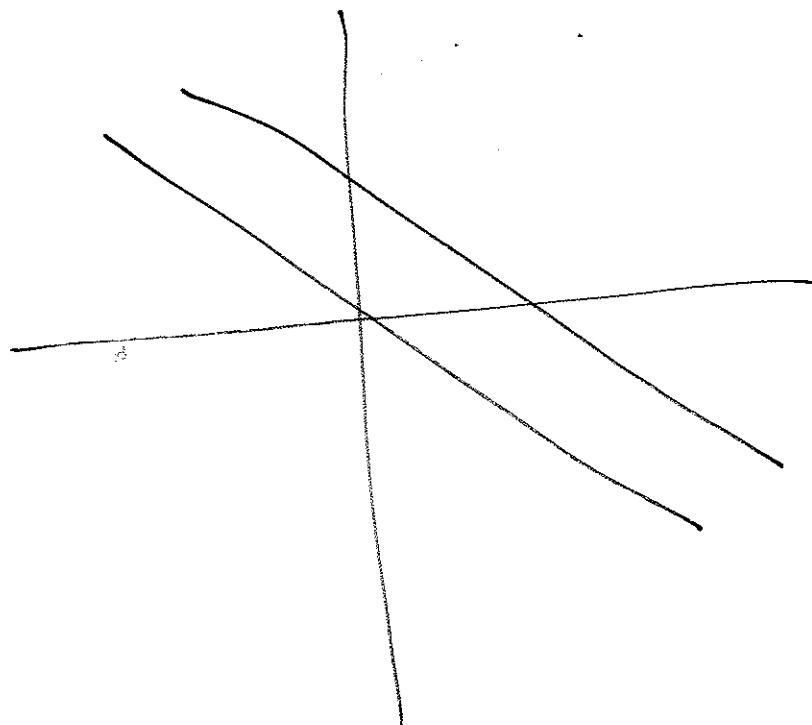
$$2x + 2y = 5$$

$$2(z - y) + 2y = 5$$

$$4 - 2y + 2y = 5$$

$$4 = 5$$

$$1 =$$



0 solution,

1 1 2

0 0 4

$$y = -x + 2$$

$$x + y = 2$$

$$x = 2 - y$$

$$2x + 2y = 4$$

$$2y = -2x - 4$$

$$y = -x - 2$$

$$2(2 - y) + 2y = 4$$

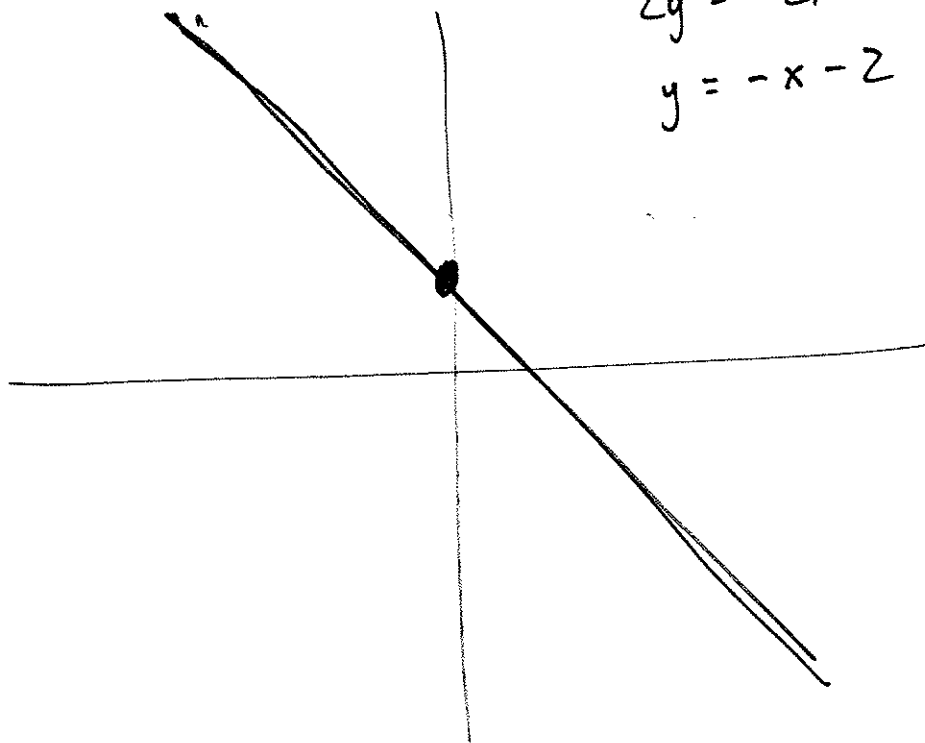
$$4 - 2y + 2y = 4$$

$$4 = 4$$

$$x = ?$$

$$y = ?$$

$\infty$  sol.



# Linear Programming

Given a collection of linear constraints

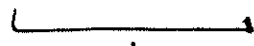
$$x \geq 0$$

$$y \geq 0$$

$$x \leq 9$$

$$2y + x \leq 12$$

$$2y - x \leq 8$$



linear constraints

objective function:

$$\text{maximize } 4x + 3y$$



'payoff' function

- you can "invest" resources into 2 (or generally  $n$ ) categories
- maximize profit
- minimize cost

$$\max_{x,y} 4x + 3y \rightarrow y = -\frac{4}{3}x + b$$

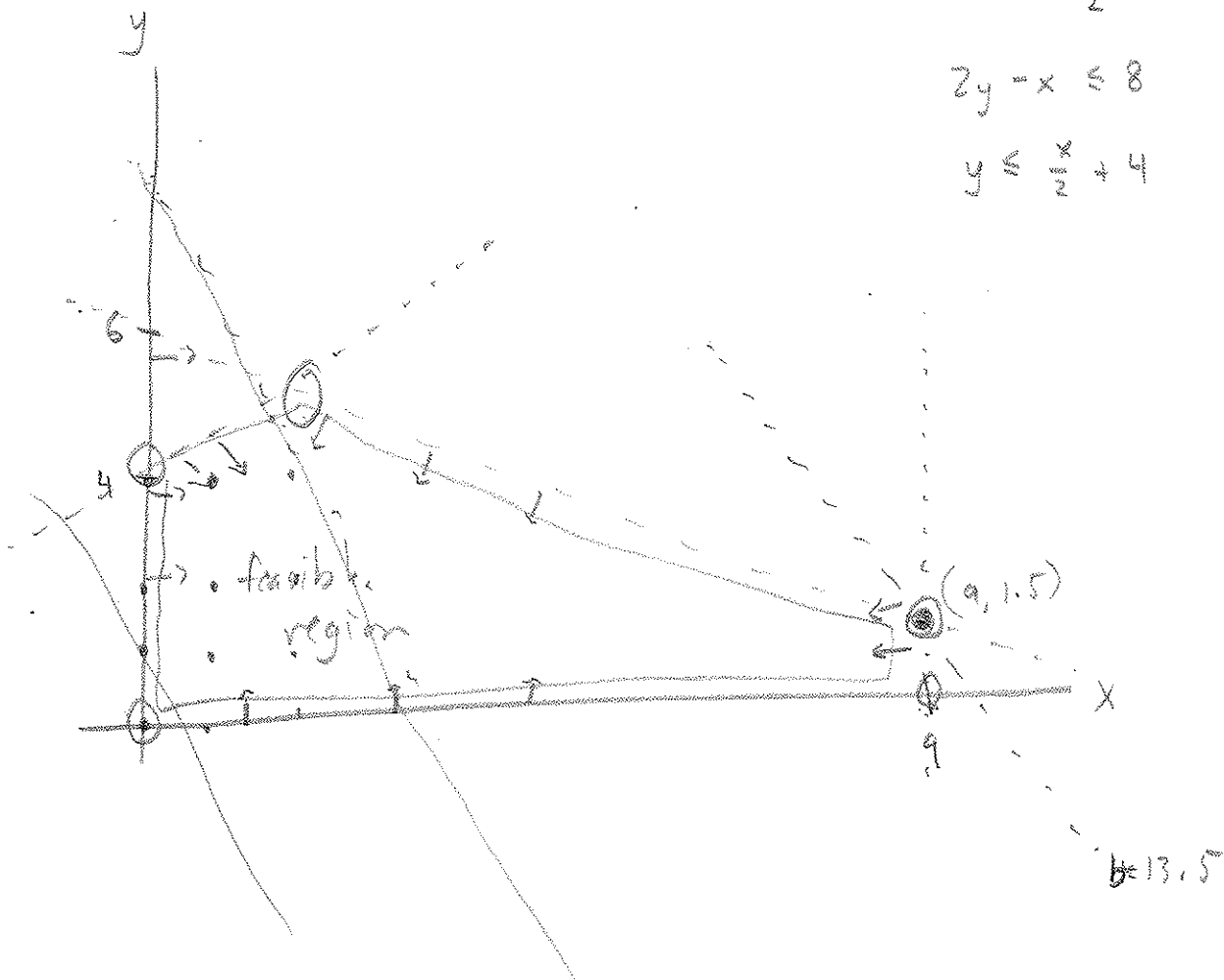
$$2y + x \leq 12$$

$$y \leq -\frac{x}{2} + 6$$

$$2y - x \leq 8$$

$$y \leq \frac{x}{2} + 4$$

point	value
(0,0)	0
(9,0)	36
(9,1.5)	40.5
(2,5)	23
(0,4)	12



(fractional)  
0-1 Knapsack Formulation

maximize your value subject to  $W$

$n$  items :  $a_1 \dots a_n$

$n$  values :  $v_1 \dots v_n$

weights  $w_1 \dots w_n$

capacity  $W$

$$x_i = \begin{cases} 1 & \text{if item } i \text{ is taken} \\ 0 & \text{if item } i \text{ is omitted} \end{cases} \quad \leftarrow \text{indicator variables}$$

relaxation :

$$0 \leq x_i \leq 1$$

$$\text{maximize } \sum_{i=1}^n v_i \cdot x_i$$

$$\text{subject to } \sum_{i=1}^n w_i x_i \leq W$$

$$\underline{x_i \in \{0, 1\}}, \quad 1 \leq i \leq n$$

ILP =  
integer linear program