IS593: Language-based Security

11. Modular Analysis

Kihong Heo



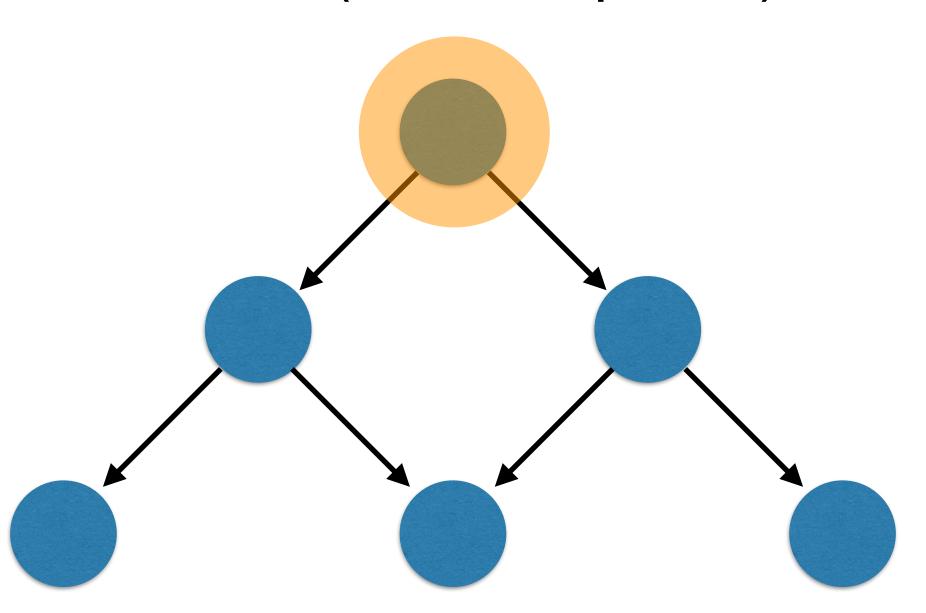
Modularity

- Key to build a scalable software system, in general
 - E.g., modular, incremental, and parallel compilation (make −j)
- How to make a static analysis modular?



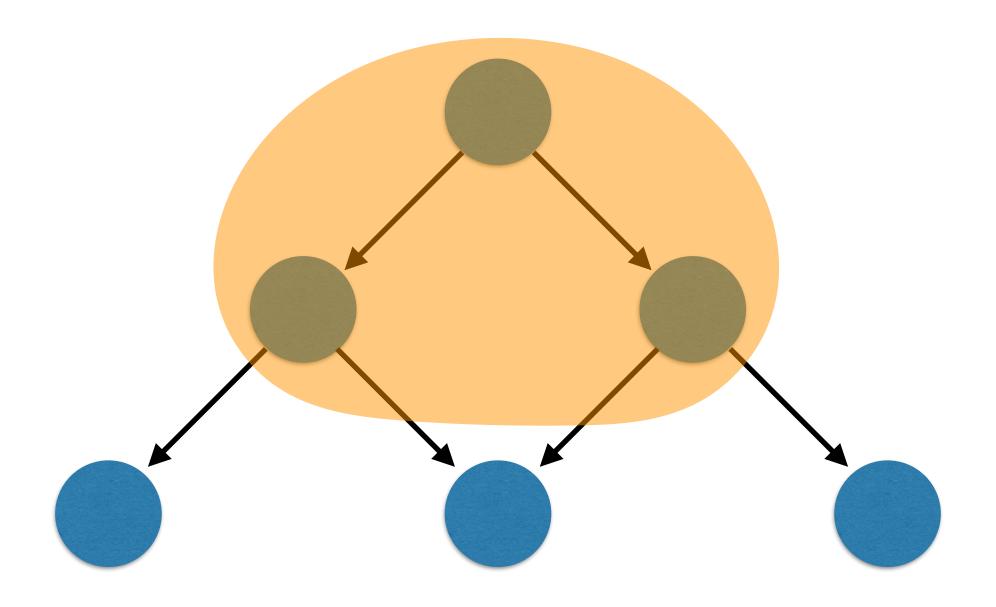
Global Analysis

- Analyze the whole program altogether
 - Starting from a root (e.g., main)
 - Similar to program execution (i.e., interpreter)



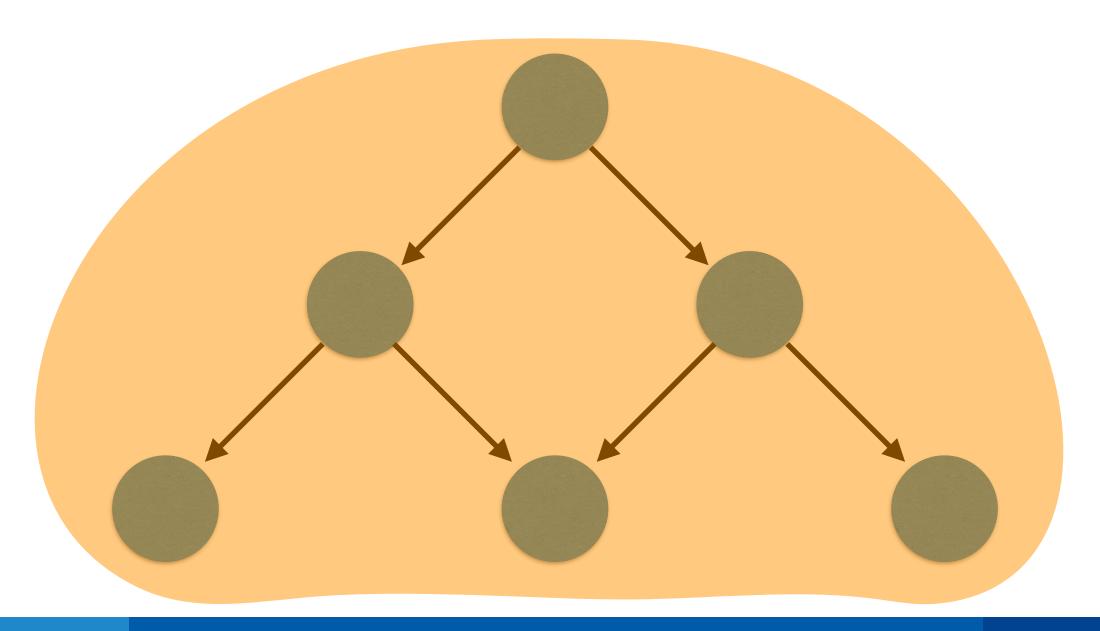
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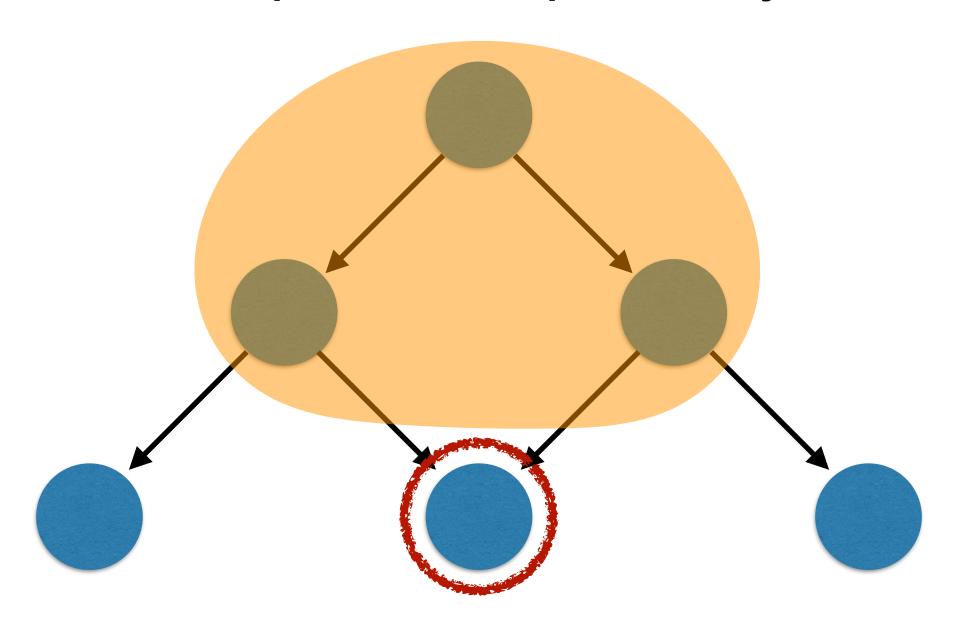
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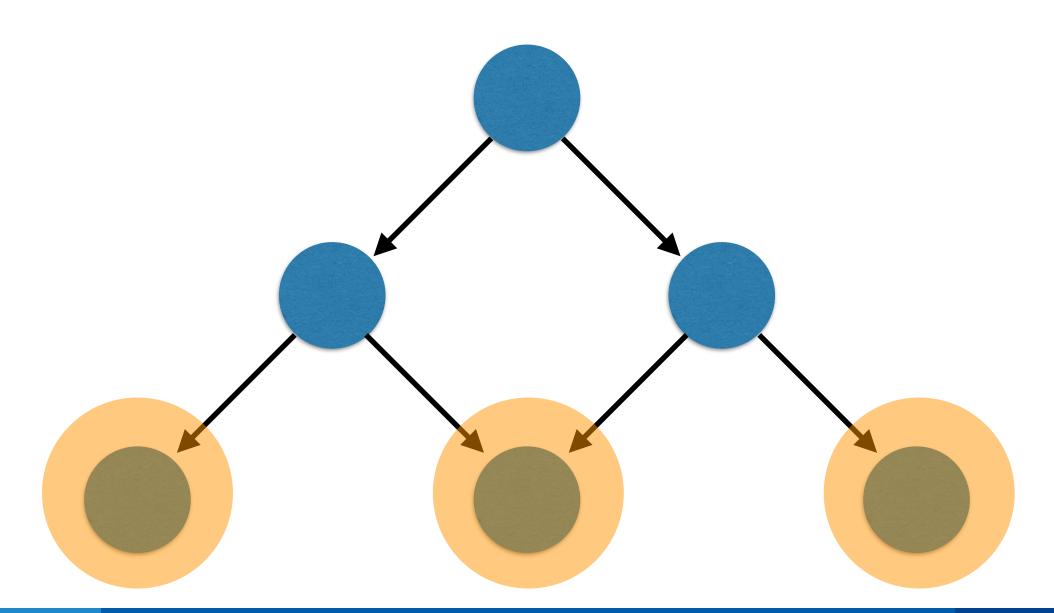
Pros and Cons of Global Analysis

- In general, context-aware but unscalable
 - Pros: aware of calling contexts (e.g., parameters, global variables, etc)
 - Cons: reanalyze the same portion repeatedly



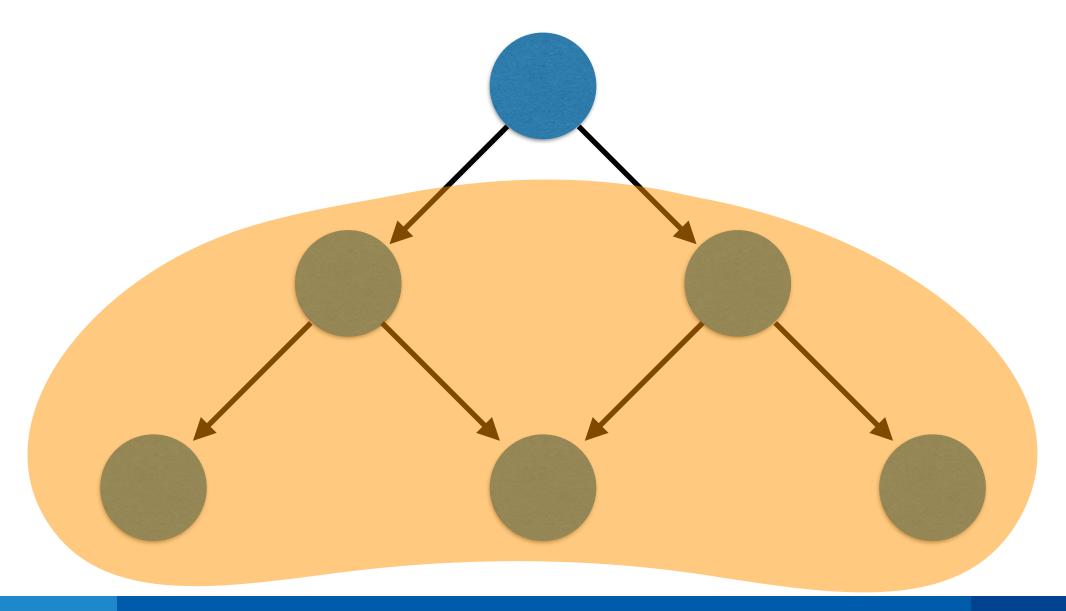
Modular Analysis

- Analyze each subcomponent (e.g., function) independently and compose
 - Starting from leaf elements
 - Similar to compilers



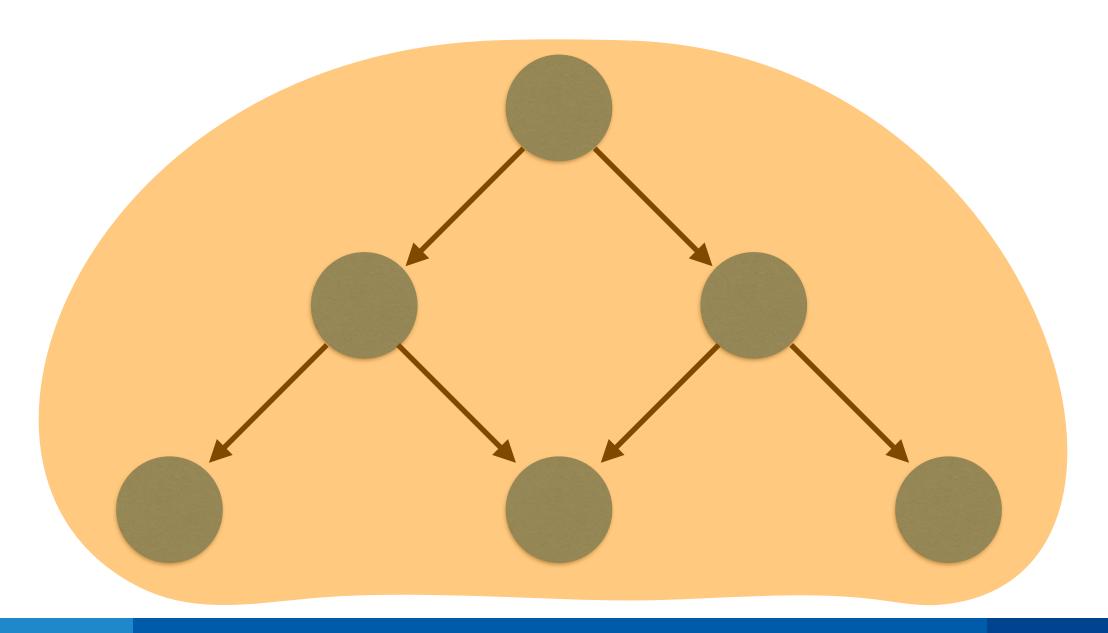
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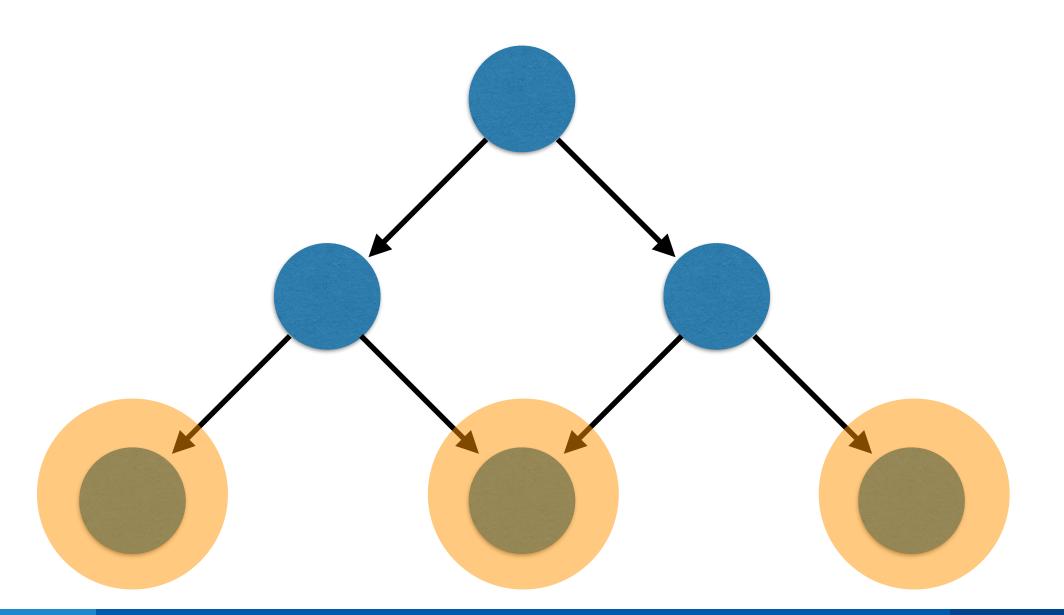
Modular Analysis

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Pros and Cons of Modular Analysis

- In general, scalable but context-unaware
 - Pros: each component is analyzed only once
 - Cons: a mechanism for handling pre-state is needed



Challenges

- How to design generic yet accurate modular analysis?
 - Effectively reusable analysis results of subcomponents
- Traditionally, simple or inductively defined properties
 - E.g., nullness, tree, list, etc

```
void foo(void* p){
  *p = 0;
}
```

Safe if p is not null

How about numerical properties?
(e.g., buffer-overrun)



```
void bar(list* p){
  while(!p)
  p = p->next;
}
```

Safely terminate if p is a well-formed singly linked list

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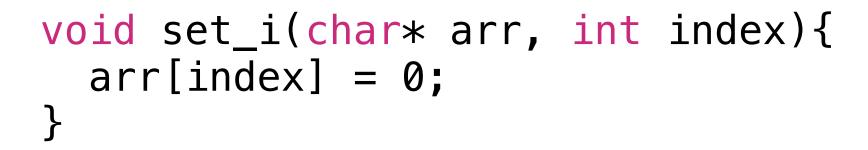
Example: Inferbo

- Facebook's Infer-based buffer overrun analyzer
- Function-level modular: analyze each function and then link them
- Parameterization: symbolic parameter for unknown calling context
- Summary-based: derive safety conditions for each function
- Scalability: enabled by modular and incremental analysis
- Availability: provided as a checker of Facebook Infer (https://github.com/facebook/infer)

Example

```
char* malloc_wrapper(int n) {
  return malloc(n);
}
```

Var	Val
n	[s0, s1]
ret	(offset: [0, 0], size: [s0, s1])



Var	Val
arr	(offset: [s4, s5], size: [s6, s7])
index	[s8, s9]

Safety Condition[s4 + s8, s5 + s9] < [s6, s7]

Example (Cont'd)

Summary of malloc_wrapper

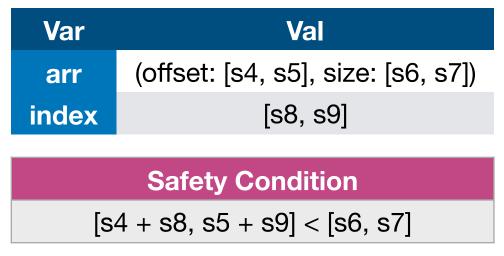
Var	Val
n	[s0, s1]
ret	(offset: [0, 0], size: [s0, s1])

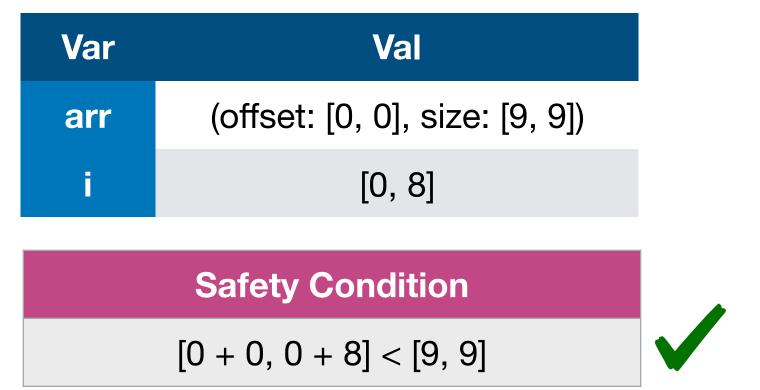
 Var
 Val

 arr
 (offset: [0, 0], size: [9, 9])

Example (Cont'd)

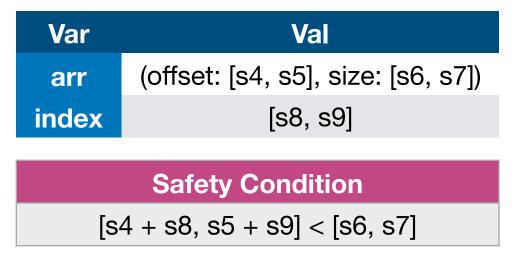
Summary of set_i

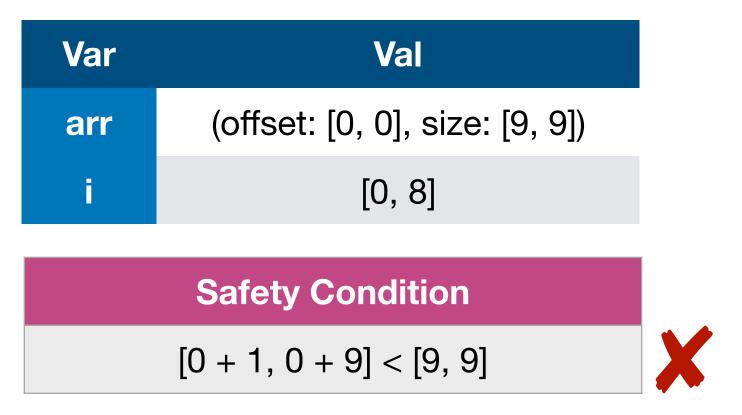




Example (Cont'd)

Summary of set_i





Unsound Design Choices

- Unsoundness is a necessary evil for better scalability and accuracy
- Inferbo is designed to be unsound for the following parts:
 - Aliasing of parameters
 - Global variables
 - Recursive calls

Summary

- Modular analysis: separately analyze each subcomponent and then link
- Key point: design of generic and accurate summary
- Inferbo: a function-level modular analysis with symbolic interval domain
- In practice, unsound design choices may be needed for better performance