Calculation of signatures

CRISM workshop: Statistics for Differential Equations Driven by Rough Paths

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Table of Contents

- Signatures
- 2 Log signatures
- 3 Timings
- 4 (Autodifferentiation)

Signatures

The signature of a path is a set of iterated integrals. Consider a path in \mathbb{R}^3 parameterised by the variable t ranging from 0 to 1, given by

$$t \mapsto \gamma(t) = (\gamma_1(t), \gamma_2(t), \gamma_3(t))$$

Then, for example, the element 2,3 of the signature is

$$\int_0^1 \left[\int_0^t \gamma_2'(s) \, ds \right] \, \gamma_3'(t) \, dt = \int_0^1 \int_0^t d\gamma_2(s) \, d\gamma_3(t)$$

and element 2,1,2 of the signature is

$$\int_0^1 \int_0^t \int_0^s d\gamma_2(r) d\gamma_1(s) d\gamma_2(t).$$



Signatures

The mth level of the signature of a path in \mathbb{R}^d given as a function from [a,b] is the d^m values of the elements with m integrated integrals. It is denoted $X_{a,b}^m$ and takes values in $(\mathbb{R}^d)^{\otimes m}$. Given a piecewise linear path, we can compute the first m levels of its signature.

For a straight path with displacement x, the signature is

$$\left(1, x, \frac{x \otimes x}{2!}, \frac{x \otimes x \otimes x}{3!}, \dots\right) \qquad \Omega(d^m)$$

• Chen's identity for the signature of the concatenation of paths, $a \le b \le c$

$$X_{a,c}^m = \sum_{k=0}^m X_{a,b}^k \otimes X_{b,c}^{m-k} \qquad \qquad \Omega(md^m)$$



Log-Signature demonstration

There is redundancy in the signature. For example, in \mathbb{R}^2 , the first four levels of the signature look like this

$$1 + (\cdot \cdot) + \begin{pmatrix} (\cdot \cdot) \\ (\cdot \cdot) \end{pmatrix} + \begin{pmatrix} \begin{pmatrix} (\cdot \cdot) \\ (\cdot \cdot) \end{pmatrix} \begin{pmatrix} (\cdot \cdot) \\ (\cdot \cdot) \end{pmatrix} \end{pmatrix} + \begin{pmatrix} \begin{pmatrix} \begin{pmatrix} (\cdot \cdot) \\ (\cdot \cdot) \end{pmatrix} \begin{pmatrix} (\cdot \cdot) \\ (\cdot \cdot) \end{pmatrix} \begin{pmatrix} (\cdot \cdot) \\ (\cdot \cdot) \end{pmatrix} \end{pmatrix} \end{pmatrix}$$

- that is 2+4+8+16=30 numbers while the log signature is only 2+1+2+3=8 numbers.

Log-Signature domain

The log signature lives in the free Lie algebra, a subspace of the tensor space. We can pick any basis to express it, e.g. in the Lyndon basis.

$$a\mathbf{1} + b\mathbf{2} + c\mathbf{12} + d\mathbf{112} + e\mathbf{122}$$

 $a\binom{1}{0} + b\binom{0}{1} + c\binom{0}{-1}\binom{1}{0} + d\binom{0}{0}\binom{1}{0}\binom{0}{0} + e\binom{0}{0}\binom{0}{0}\binom{1}{0}}{\binom{0}{1}\binom{0}{0}} + e\binom{0}{0}\binom{0}{0}\binom{0}{1}$

is the log of

$$1 + \left(\begin{smallmatrix} a \\ b \end{smallmatrix} \right) + \tfrac{1}{2} \left(\begin{smallmatrix} a^2 & ab + 2c \\ ab - 2c & b^2 \end{smallmatrix} \right) + \tfrac{1}{6} \left(\begin{smallmatrix} \left(\begin{smallmatrix} a^3 & a^2b + 6d + 3ac \\ a^2b - 12d & ab^2 + 6e + 3bc \end{smallmatrix} \right) \\ \left(\begin{smallmatrix} a^2b + 6d - 3ac & ab^2 - 12e \\ ab^2 + 6e - 3bc & b^3 \end{smallmatrix} \right) \right) + \cdots$$

Detour on Lyndon words

The Lyndon basis gets an element from each Lyndon-word using a procedure which splits a Lyndon word into its longest Lyndon suffix and the rest, e.g.

$$\begin{array}{l} \underline{1122} \rightarrow [1,\underline{122}] \rightarrow [1,[\underline{12},2]] \rightarrow [[1,[1,2]],2] \\ = [[1,12-21],2] \\ = [1(12-21)-(12-21)1,2] \\ = [(112-121)-(121-211),2] \\ = ((112-121)-(121-211))2 \\ -2((112-121)-(121-211)) \\ = 22121-21212+1122-2211 \end{array}$$

Log-Signature domain

The log signature of a path with displacement x is just x. BCH formula expresses $A \bullet B := \log(\exp A \otimes \exp B)$ as

$$A + B + [A, B]/2 + ([A, [A, B]] + [B, [B, A]])/12 + \cdots$$

$$(a1 + b2 + c12 + d112 + e122 + \cdots) \bullet (a'1 + b'2)$$

$$= (a1 + b2 + c12 + \cdots) + (a'1 + b'2)$$

$$+ (ab'12 - a'b12 - a'c112 + cb'122 + \cdots)/2 + \cdots$$

$$= (a + a')1 + (b + b')2 + (c + \frac{ab' - a'b}{2})12 + \cdots$$

Anagrams

Code example

```
void joinSegmentToSignatureInPlace(
 FLogSignature<2,3>& a, const FSegment<2>& b) {
  a[4] +=0.5*a[2]*b[1]
        +0.08333333333333*a[1]*a[1]*b[0]
        +0.08333333333333*a[0]*b[1]*b[1];
  +0.08333333333333*a[1]*b[0]*b[0]
        +0.08333333333333*a[0]*a[0]*b[1]
        -0.5*a[2]*b[0]
        a[2] +=0.5*a[0]*b[1]
        -0.5*a[1]*b[0];
  a[1] += b[1];
  a[0] += b[0];
```

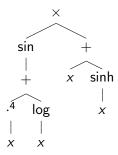
Some comparative timings

(D,M)	(2,6)	(2,10)	(3,10)	(5,5)	(10,4)
Time (s) for 100 (log) signatures of length-100 paths:					
Signature	0.021	0.201	8.005	0.212	0.404
Compiled	0.005	0.491	276.174	0.325	0.578
Projection	0.028	0.389	16.491	0.389	0.829
Time (s) for 100 (log) signatures of length- 1000 paths:					
Signature	0.017	2.427	100.572	2.400	4.935
Compiled	0.042	4.931	2857.072	3.341	5.845
Projection	0.211	3.285	123.143	3.297	7.701
Time (s) for a single call of the preparation function:					
Compiled	0.107	0.805	439.38	0.372	1.466
Projection	0.001	0.037	4.386	0.103	0.156

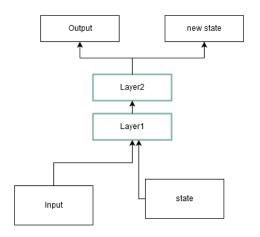
Autodifferentiation

Consider evaluating, for some fixed value of x,

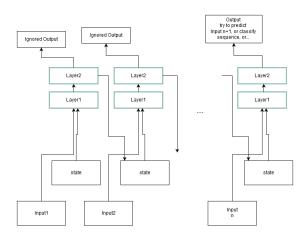
$$\frac{d}{dx}\left[\left(\sin(x^4+\log x)\right)(x+\sinh x)\right]$$



Recurrent Neural Networks



Recurrent Neural Networks



Thanks!









Dr Ben Graham