Computer Science & Information Technology Database Management System

DPP: 1

- Q1 In a relation R(X,Y), we can say that a functional dependency $X \rightarrow Y$ holds if:
 - (A) All values of X are unique in R
 - (B) All values of Y are unique in R
 - (C) For every pair of tuples in R, if the X values are the same, then the Y values are also the same
 - (D) For every pair of tuples in R, if the Y values are the same, then the X values are also the same
- **Q2** Given $A \rightarrow B$, which of the following can be inferred using augmentation?
 - (A) $B \rightarrow A$
- (B) $AC \rightarrow BC$
- (C) $C \rightarrow A$
- (D) $AB \rightarrow C$
- Q3 Relation R(A, B, C, D) has FDs:
 - $A \rightarrow B$
 - $B \rightarrow C$

Which of the following is NOT derivable using Armstrong's axioms?

- $(A) A \rightarrow C$
- (B) $A \rightarrow BC$
- (C) $AD \rightarrow BCD$
- (D) $C \rightarrow A$
- **Q4** Consider relation R(A, B, C, D) with candidate keys: {AB}, {BCD}.

Which of the following best explains why {AB} is considered a minimal key?

- (A) Because {AB} has fewer attributes than {BCD}
- (B) Because no proper subset of {AB} is a key
- (C) Because {AB} determines all attributes of R
- (D) None of the above
- **Q5** Consider a relation R(A, B, C, D) with the following functional dependencies:
 - $A \rightarrow B$
 - $C \rightarrow D$
 - AC → BCD

Which of the following is a candidate key of R?

- (A) AC
- (B) A
- (C) C
- (D) ABC
- Q6 Consider a relation R(A, B, C, D, E, F) with the following set of functional dependencies:

 $F = \{A \rightarrow B, B \rightarrow C, CD \rightarrow E, E \rightarrow F, F \rightarrow A\}$

How many candidate keys does relation R have?

- Q7 Consider relation Employee(EmpID, Dept, Project, Manager, Location) with the following functional dependencies:
 - 1. EmpID → Dept, Project
 - 2. Project, Manager → Location
 - 3. Location → EmpID

Which of the following is/are prime attribute of Employee?

- (A) Project
- (B) Manager
- (C) Dept
- (D) Location
- **Q8** Consider a relation R(A, B, C, D, E) with the following set of functional dependencies F:
 - a. $A \rightarrow B$
 - b. $B \rightarrow C$
 - c. $CD \rightarrow E$

Which of the following functional dependencies is/are NOT the member of F^+ (i.e., implied by the given FD set)?

- (A) $A \rightarrow C$
- (B) $A \rightarrow E$
- (C) $AB \rightarrow C$
- (D) $BD \rightarrow E$
- **Q9** Consider relation R(A, B, C, D, E) with the following two FD sets:

F1:

- 1. A → BC
- 2. $B \rightarrow D$
- 3. $CD \rightarrow E$

F2:

- 1. A → B
- 2. $A \rightarrow C$
- 3. $B \rightarrow D$
- 4. $C \rightarrow E$,

5. D → E

Which of the following is correct?

- (A) F1 and F2 are equivalent (i.e., they imply the same set of FDs)
- (B) $F1 \subset F2$ (F1 is a subset of F2, but not equivalent)
- (C) $F2 \subset F1$ (F2 is a subset of F1, but not equivalent)\

(D) F1 and F2 are not comparable

Q10 Consider relation R(P, Q, R, S, T) with the following functional dependencies:

a. $P \rightarrow Q$

b. $R \rightarrow S$

c. $P \rightarrow T$

How many superkeys does R have?



Answer Key

O 1	(A,	~ \
Q1	(A.	C)



Hints & Solutions

Q1 Text Solution:

Because FD $X \rightarrow Y$ means: X uniquely determines Y. i.e., if two tuples agree on X, they must also agree on Y.

Q2 Text Solution:

Augmentation: If $X \rightarrow Y$, then $XZ \rightarrow YZ$.

Q3 Text Solution:

We can derive A \rightarrow C (transitivity), A \rightarrow BC (union), and AD \rightarrow BCD (adding D using augmentation).

But $C \rightarrow A$ cannot be derived.

Q4 Text Solution:

A candidate key is minimal if none of its proper subsets is a key.

Q5 Text Solution:

Check closure of AC:

AC⁺= {A, C} \cup {B (from A \rightarrow B), D (from C \rightarrow D)} = {A, B, C, D} = all attributes \rightarrow AC is a superkey. Now check minimality:

 $A^+ = \{A, B\}$ (not all attributes).

 $C^+ = \{C, D\}$ (not all attributes).

Neither A nor C alone is a key \rightarrow AC is minimal. Option D (ABC) is a superkey too, but not minimal since AC \subseteq ABC

Q6 Text Solution:

 $AD^{+} = \{A, B, C, D, E, F\}$

Along with AD, the following are also candidate keys:

FD

ED

CD

BD

Q7 Text Solution:

The candidate keys of the relation are {EmpID, Manager}, {Location, Manager}, and {Project,

Manager}.

Q8 Text Solution:

- (A) $A \rightarrow C$
 - A → B (given)
 - $B \rightarrow C$ (given)
 - So by transitivity: $A \rightarrow C$ (implied, in F^+)
- (B) $A \rightarrow E$
 - From A, we get B (via 1) and then C (via 2).
 - But to get E, we need CD → E. We don't

have D from A.

- So A \rightarrow E (NOT implied, NOT in F⁺)
- (C) $AB \rightarrow C$
 - AB contains A. From $A \rightarrow C$ (as shown in A).
 - So AB \rightarrow C (implied, in F⁺)
- (D) $BD \rightarrow E$
 - From BD, we already have B and D.
 - From $B \rightarrow C$, we get C.
 - Now we have C and D → E (via 3).
 - So BD \rightarrow E (implied, in F⁺)

Q9 Text Solution:

Step 1: Minimal Cover of F1

- A \rightarrow BC can be split: A \rightarrow B, A \rightarrow C
- So F1 becomes: $\{A \rightarrow B, A \rightarrow C, B \rightarrow D, CD \rightarrow E\}$ Step 2: Compare with F2
- F2 is: $\{A \rightarrow B, A \rightarrow C, B \rightarrow D, C \rightarrow E, D \rightarrow E\}$
- From CD → E in F1, we can't separately derive C
 → E or D → E (both needed together).
- So F2 allows more dependencies than F1. Conclusion: F1 is not equivalent to F2; instead, F2 is strictly stronger.

Q10 Text Solution:

The candidate key of this relation is PR. So, the remaining three attributes Q, S, and T have 2 choices each.

Hence, there are 8 superkeys possible