

# CS & IT ENGINEERING



## Database Management System

DBMS

Lecture No. 2

By- Ravindrababu Ravula Sir





# Recap of Previous Lecture

Topic

Introduction





# Topics to be Covered



## Topic

Integrity Constraints & ER Model (2 Marks)

## Topic

Normalization (2-4 Marks)

## Topic

Queries (Relational Algebra, SQL, Tuple  
Relational Calculus) (4 Marks)

## Topic

File Organization & Indexing(2-4 Marks)

## Topic

Transactions & Concurrency Control (2- 4 Marks)



## FOREIGN KEY (Referential key)

### Foreign Key:

- Used to relate tables. ✓
- Usually defined over two relations.
  - Referenced and Referencing Relation.





# FOREIGN KEY (Referential key)

- **Definition:** Set of one or more attributes referencing to a primary key or alternative key of the same relation or other relation.



- X & Y can be from same relation or from two different relations.

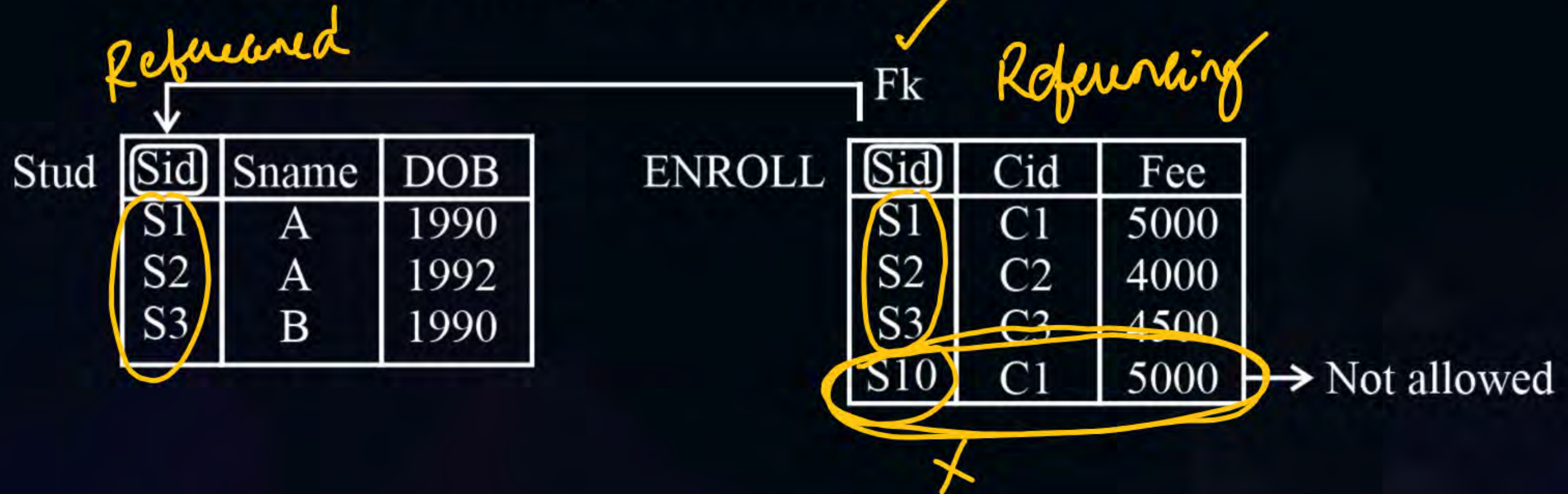




# FOREIGN KEY (Referential key)

**Example 1:** Relation 1: Stud (Sid, Sname, DOB)

Relation 2: ENROLL (Sid, Cid, Fee)



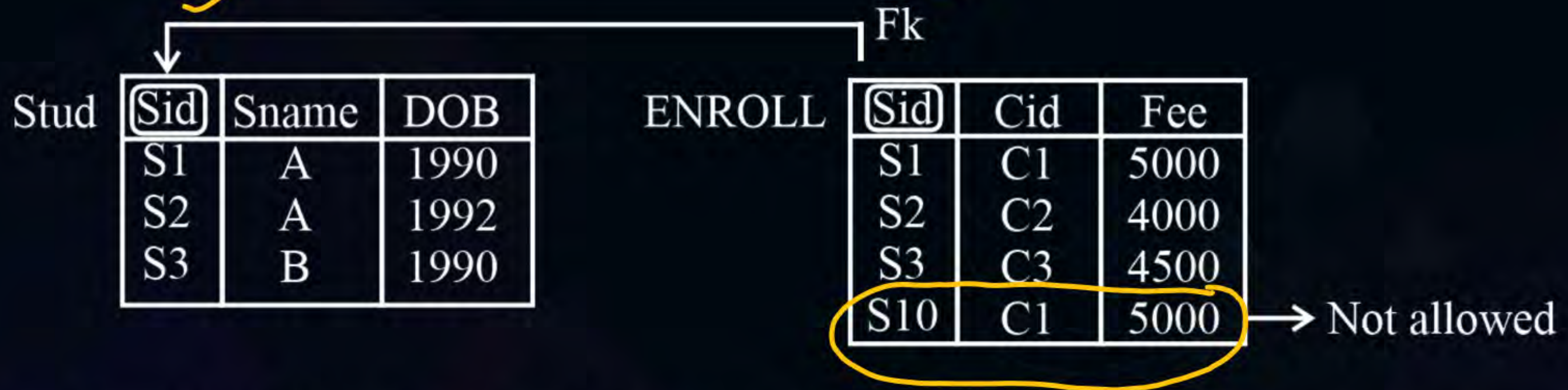




## FOREIGN KEY (Referential key)

**Example 1:** Relation 1: Stud (Sid, Sname, DOB)

Relation 2: ENROLL (Sid, Cid, Fee)



Here, Stud Relation is called Referenced Relation and ENROLL Relation is called Referencing Relation.



## FOREIGN KEY (Referential key)



**Example 2:** Self Referencing





# FOREIGN KEY (Referential key)

## Example 2: Self Referencing

EMP

eid	ename	SupID
<u>E1</u>	N1	<u>E2</u>
<u>E2</u>	N2	<u>E3</u>
<u>E3</u>	N2	NULL
<u>E4</u>	N3	<u>E10</u>

Fk ✓

No Supervisor, highest Authority

doesn't exist, so not allowed



# Inspiring Stories : Malati Mem

**Background:** From tea-garden community in Assam.



**Struggles:** British pushed opium addiction to control labor.

**Achievements:** Fought against it, taught villages about its danger—but got shot by police.

**Impact:** Showed how everyday people can stand up for health and freedom.





## FOREIGN KEY (Referential key)

- FK attributes allows NULL values.
- Referencing Related Records whose Foreign Key value is NULL, is not related to any referenced related record.





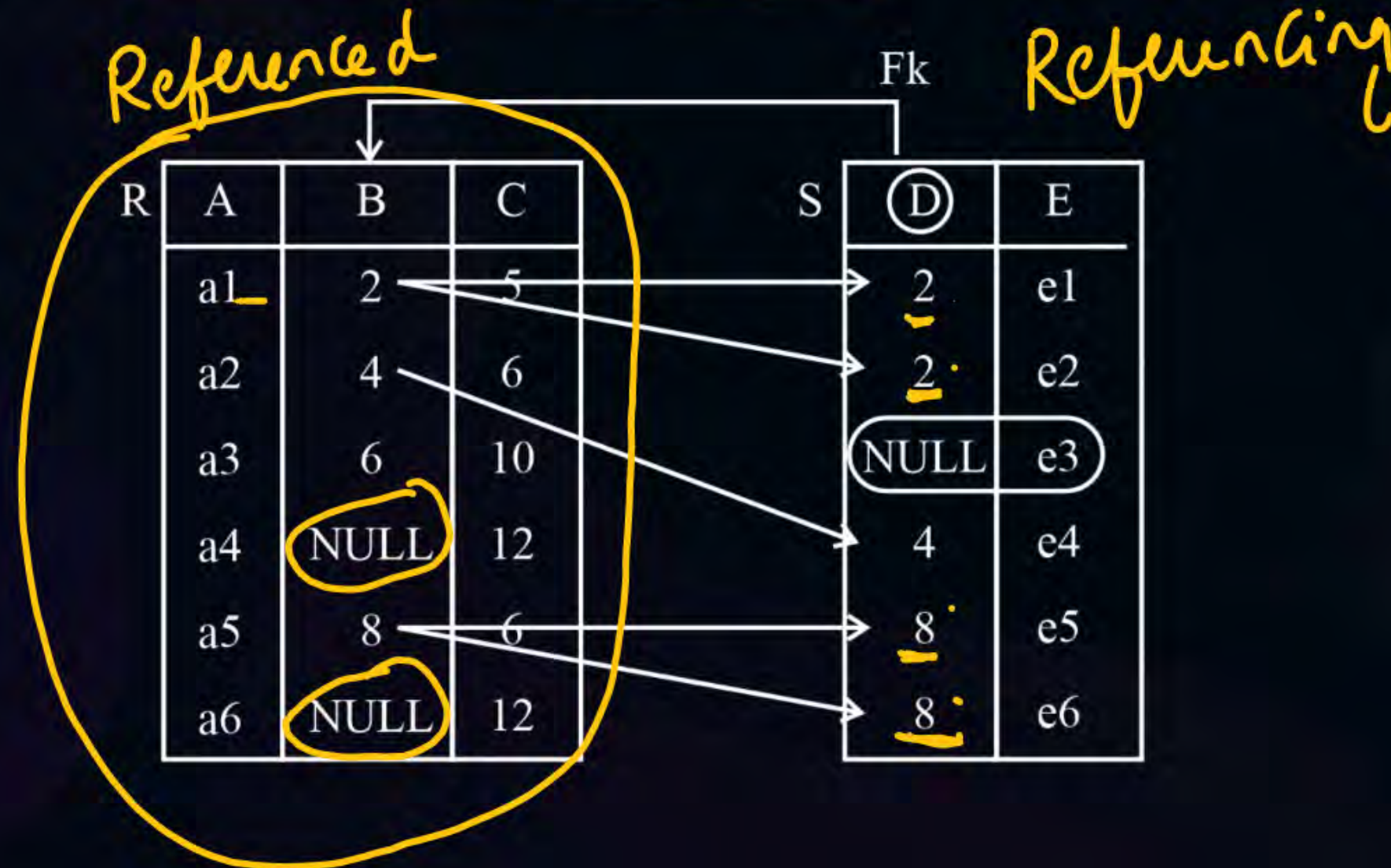


# FOREIGN KEY (Referential key)

- Referenced and Referencing relation have 1 : Many Relationship.

Referenced → Referencing (Many)

Referencing → Referenced (At Most one)







# Referential Integrity Constraints

## Foreign key Integrity Constraints with Referenced Relation

### (a) Insertion:

No violation so No action is needed.

A B C			D E F		
✓	1		3	1	
X	2		2		
	3				

### (b) Deletion:

May cause violation, these are handled in 3 ways.

#### (i) ON DELETE NO ACTION (default)

- Deletion restricted if Foreign Key violation occurs.





# Referential Integrity Constraints

## (ii) ON DELETE CASCADE

- Forced to delete related referencing records.

						Fk
		↓				
R	A	B	C	S	D	E
a1	2	5		2	e1	
a2	4	6		2	e2	
a3	6	10		NULL	e3	
a4	NULL	12		4	e4	
a5	8	6		8	e5	
a6	NULL	12		8	e6	

after  
removing a row  
(a5, 8, 6)  
in R

R	A	B	C
a1	2	5	
a2	4	6	
a3	6	10	
a4	NULL	12	
a6	NULL	12	

S	D	E
	2	e1
	2	e2
	NULL	e3
	4	e4





# Referential Integrity Constraints

## (iii) ON DELETE SET NULL ✗

- Set Null in the rows that are deleted

						Fk
		↓				
R	A	B	C	S	D	E
	a1	2	5		2	e1
	a2	4	6		2	e2
	a3	6	10		NULL	e3
	a4	NULL	12		4	e4
	a5	8	6		8	e5
	a6	NULL	12		8	e6

after  
removing a row  
(a5, 8, 6)  
in R

R	A	B	C
a1	2	5	
a2	4	6	
a3	6	10	
a4	NULL	12	
a6	NULL	12	

S	D	E
	2	e1
	2	e2
	NULL	e3
	4	e4
	NULL	e5
	NULL	e6





# Referential Integrity Constraints

## (c) Updation:

- May cause violation, this is also handled in ways.

### (i) ON UPDATE NO ACTION (default)

Updation is restricted if there are foreign keys dependent on the row deleted.



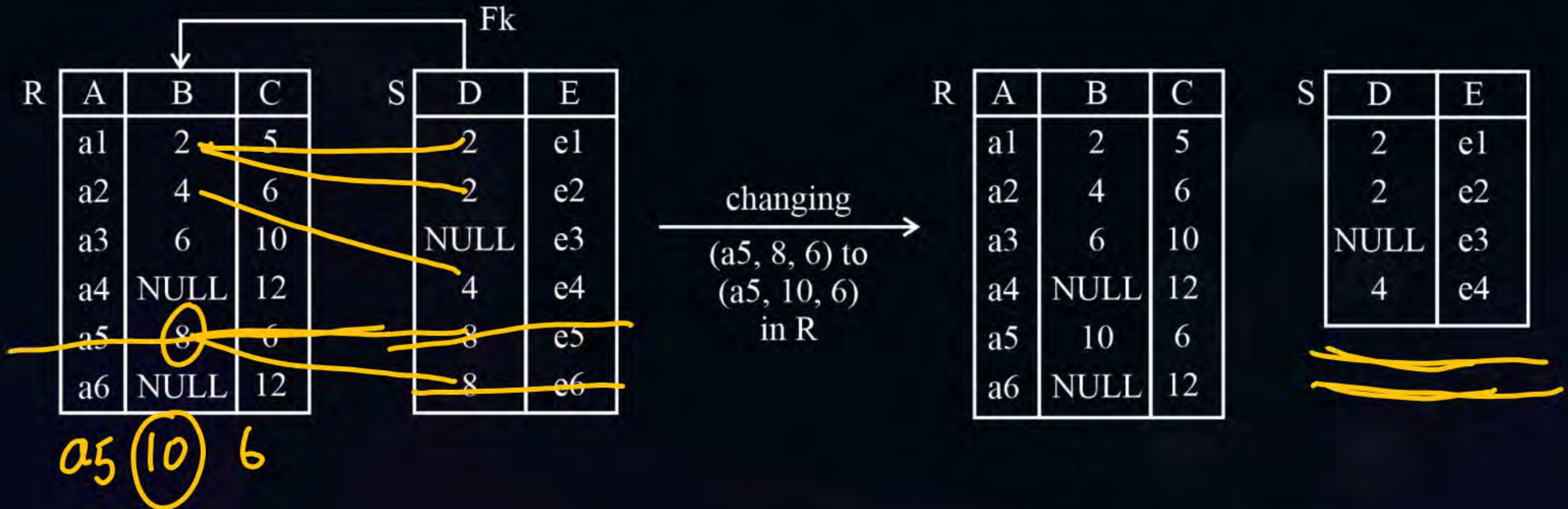




# Referential Integrity Constraints

## (ii) ON UPDATE CASCADE

Forced to delete the records referencing the row delete updated.



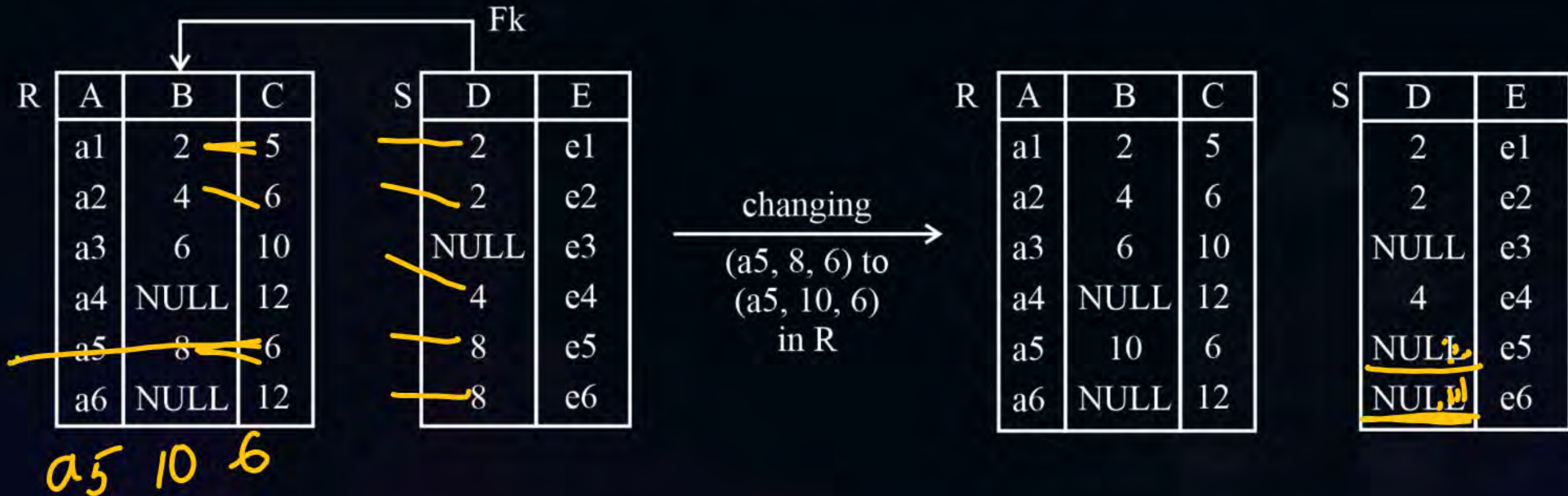




# Referential Integrity Constraints

## (iii) ON UPDATE SET NULL

Sets NULL in the rows that are referenced to the updated value.







# Referential Integrity Constraints

## Foreign key Integrity Constraints with Referencing Relation

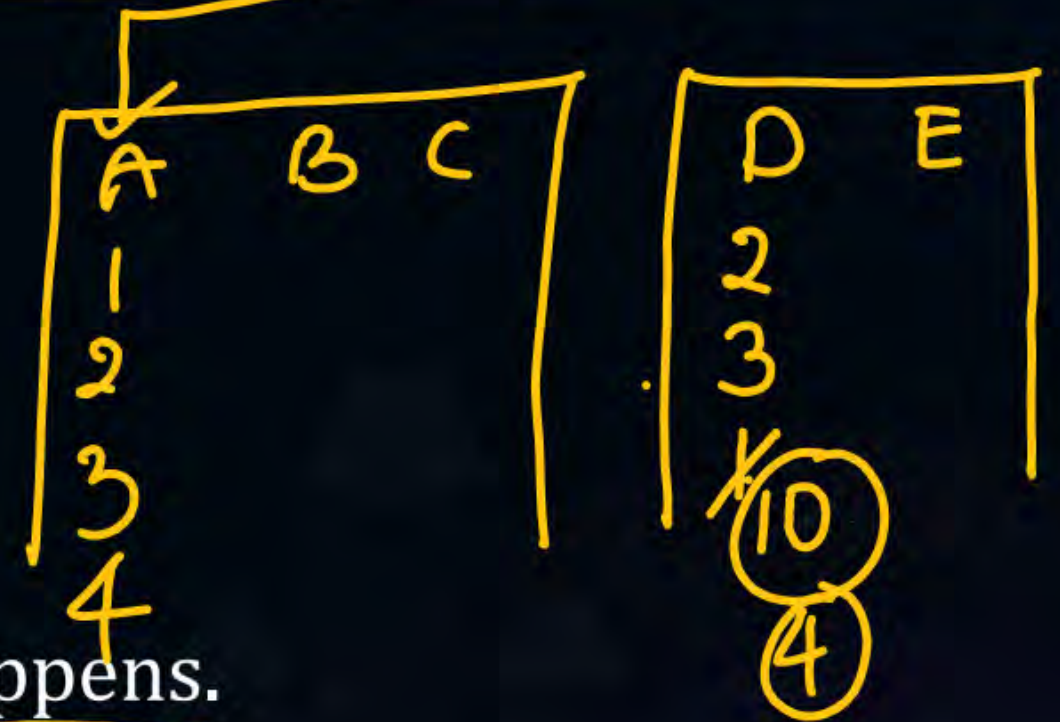
**(a) Insertion:** May cause foreign key violation

Sol: Restricted the Insertion if violation occurs.

**(b) Deletion:** No Violation occurs ✓

**(c) Updation:** May cause Foreign key violation ✓

Sol: Updation is restricted if Foreign key violation happens.







# GATE PYQ on Integrity Constraints

True ✓

<u>Referred Table</u>	<u>Insertion</u>	<u>No Violation</u>
	<u>Deletion</u>	<u>No Action, Cascade or Set NULL</u>
	<u>Updation</u>	<u>No Action, Cascade or Set NULL</u>
<u>Referencing Table</u>	<u>Insertion</u>	<u>May cause violation, Restricted If yes</u>
	<u>Deletion</u>	<u>No Violation</u>
	<u>Updation</u>	<u>May cause violation, Restricted If yes</u>



# Inspiring Stories : Helen Lepcha

**Background:** From Lepcha tribe in Sikkim.



**Struggles:** British disrespected tribal ways; movement was suppressed.

**Achievements:** Followed Gandhi, helped Subhas Chandra Bose escape, wrote for tribal people.

**Impact:** Kept hope alive in mountains; a tribal heroine. ✓





## GATE PYQ on Integrity Constraints

### Q1: ISRO CSE 2016

Let  $R(a,b,c)$  and  $S(d,e,f)$  be two relations in which  $d$  is the foreign key of  $S$  that refers to the primary key of  $R$ . Consider the following four operations on  $R$  and  $S$

I. Insert into  $R$  ✓

III. Delete from  $R$

II. Insert into  $S$  ✓

IV. Delete from  $S$

Which of the following can cause violation of the referential integrity constraint above?

(A) Both I and IV

(B) Both II and III

(C) All of these

(D) None of these





## GATE PYQ on Integrity Constraints



### Q1: ISRO CSE 2016

Let  $R(a,b,c)$  and  $S(d,e,f)$  be two relations in which  $d$  is the foreign key of  $S$  that refers to the primary key of  $R$ . Consider the following four operations on  $R$  and  $S$

I. Insert into  $R$

II. Insert into  $S$

III. Delete from  $R$

IV. Delete from  $S$

Which of the following can cause violation of the referential integrity constraint above?

(A) Both I and IV

(C) All of these

☒ (B) Both II and III

(D) None of these







## GATE PYQ on Integrity Constraints



### Q2: GATE CSE 2017 SET-2

Consider the following tables T1 and T2.

In table T1, P is the primary key and Q is the foreign key referencing R in table T2 with ondelete cascade and on-update cascade. In table T2, R is the primary key and S is the foreign key referencing P in table T1 on-delete set NULL and on-update cascade.

In order to delete record (3,8) from table T1, the number of additional records that need to be deleted from table T1 is \_\_\_\_\_.

a) 0   b) 1   c) 2   d) 3

T1		T2	
P	Q	R	S
2	2	2	2
3	8	8	3
7	3	3	2
5	8	9	7
6	9	5	7
8	5	7	2
9	8		





## GATE PYQ on Integrity Constraints

ON DELETE CASCADE  
ON UPDATE CASCADE



### Q2: GATE CSE 2017 SET-2

- Since S(FK) refers to P(PK), removing (3, 8) row in T1, makes (8, 3) row in T2 to (8, NULL) due to the ON DELETE SET NULL constraint
- Since Q(FK) refers to R(PK), deletion in a foreign key doesn't affect any table.

So zero entries are removed. ✓

Answer: 0 ✓

T1		T2	
P	Q	R	S
2	2	2	2
3	8	8	3
7	3	3	2
5	8	9	7
6	9	5	7
8	5	7	2
9	8		

ON DELETE SET NULL  
ON UPDATE CASCADE





## GATE PYQ on Integrity Constraints

### Q3: GATE CSE 2005

The following table has two attributes A and C where A is the primary key and C is the foreign key referencing A with on-delete cascade.

The set of all tuples that must be additionally deleted to preserve referential integrity when the tuple (2,4) is deleted is:

A	C
2	4
3	4
4	3
5	2
7	2
9	5
6	4

Handwritten annotations: A yellow arrow points from the header 'A' to the first row. Yellow lines connect the value '2' in column A to the values '4', '2', and '5' in column C. Checkmarks are placed next to the rows (5,2), (7,2), and (9,5).





## GATE PYQ on Integrity Constraints

### Q3: GATE CSE 2005

- On removing (2,4), the rows (5,2) and (7,2) are deleted.
- On removing (5,2), (9,5) is also deleted
- Ans: (5,2), (7,2) (9,5) are deleted

A	C
2	4
3	4
4	3
5	2
7	2
9	5
6	4



ON DELETE CASCADE





## Topic : NORMALIZATION ✓



1. Introduction to DB Anomalies ✓
2. Functional Dependency ✓
3. Attribute closure & Membership test ✓
4. Finding Candidate Keys \*\*\* ✓
5. Lossless join and Dependency Preserving Decomposition \*\* ✓





## Topic : NORMALIZATION



### 6. Normal Forms

- 1 NF, 2 NF, 3 NF, BCNF → Theory Question ✓
- Finding highest NF & Relational Schema \*\*
- Decomposition into higher Normal Form
- Multivalued Dependency and 4 NF ✗

80%

### 7. Canonical Cover of FD set (Minimal Cover)



# Inspiring Stories : Pasaltha Khuangchera



**Background:** Mizo warrior from the hills.

**Struggles:** British troops marched into his land.

**Achievements:** Fought them in 1890 to protect villages; died in battle but did not run.

**Impact:** First known Mizo hero to fall fighting the British; a legend in Mizoram.





## Topic : NORMALIZATION



### NORMALIZATION

- Process used to eliminate or reduce redundancy in Database tables.
- When two or more independent relations are stored in a single relation, there is a possibility of Redundancy.

#### Example:

For a Relation Student lets say the dependency is, Sid → Sname, DOB

For Course Table, Cid → Cname, Instructor

For Enroll Table, Sid, Cid → Fee





## Topic : NORMALIZATION

Stud, Course, enroll

Now lets say a table is created on these three relations, then the Relationship

formed is: R

Sid	Sname	DOB	Cid	Cname	Instructor	Fee
S1	A	1990	C1	DB	Korth	—
S2	A	1990	C1	DB	Korth	—
S3	B	1998	C1	DB	Korth	—
S3	B	1998	C2	Algo	Coremon	—
S3	B	1998	C3	OS	Galwin	—

Sid, Cid → Candidate key

- There is redundant data. ✓





## Topic : NORMALIZATION

### Problems because of Redundancy : [DB Anamolies] ✓

#### (i) Insertion Anamoly:

S<sub>4</sub> ✓      C<sub>4</sub>

- The insertion of student data without a course or viceversa is not possible.

R	Sid	Sname	DOB	Cid	Cname	Instructor	Fee
	S1	A	1990	C1	DB	Korth	—
	S2	A	1990	C1	DB	Korth	—
	S3	B	1998	C1	DB	Korth	—
	S3	B	1998	C2	Algo	Coremon	—
	S3	B	1998	C3	OS	Galwin	—





## Topic : NORMALIZATION

### Problems because of Redundancy : [DB Anamolies]

#### (ii) Deletion Anamoly:

- If someone wants to delete S3 data, then it deletes the C1 & C2 course as well.

R	Sid	Sname	DOB	Cid	Cname	Instructor	Fee
	S1	A	1990	C1	DB	Korth	—
	S2	A	1990	C1	DB	Korth	—
	S3	B	1998	C1	DB	Korth	—
	S3	B	1998	C2 ✓	Algo	Coremon	—
	S3	B	1998	C3 ✓	OS	Galwin	—





## Topic : NORMALIZATION



### (iii) Updation Anamoly:

- If the row (S3, B, 1998, C1, DB, Korth) has got updated to (S3, B, 1989, C1, DB, korth), then the data of S3 is inconsistent, like S3 has 2 different DOB's 1988 & 1989.

R	Sid	Sname	DOB	Cid	Cname	Instructor	Fee
	S1	A	1990	C1	DB	Korth	—
	S2	A	1990	C1	DB	Korth	—
	S3	B	1998	C1	DB	Korth	—
	S3	B	1998	C2	Algo	Coremon	—
	S3	B	1998	C3	OS	Galwin	—





## Topic : NORMALIZATION

- These Anamolies indirectly causes Inconsistency.
- The Goal of Normalization is to remove the Redundancy and Inconsistency.
- If there is 0% redundancy in DB table then there are No Database Anamolies present.



# Inspiring Stories : Pa Togan

**Background:** Young Garo leader.

**Struggles:** British raided the hills and taxed hard.

**Achievements:** Led a bold night attack on a British camp (1872);  
fell while fighting.

**Impact:** Garo people still honor him every year as a symbol of  
courage.







Telegram channel







**THANK - YOU**