

Computer Networks & Internet Protocol

Data communication - Data → Information

↳ Exchange of Information between two devices via some form of transmission medium.

Computer network -

A network → set of devices connected by communication links.

interfaces across a distance → "peer to peer" interfaces.

local interfaces - 'services' interfaces.

Modules on each end are organized as sequence of function → layers.

Set of modules organized as layers → "Protocol stack".

Computer Networks (Based on transmission Mode)

Simplex

Communication is unidirectional.
Only one can transmit
the other can only receive

Half Duplex

Communication is Bidirectional,
Both station can send & receive but not at same time

Full Duplex

Both stations can transmit & receive simultaneously

Computer networks (Based on transmission time)

Synchronous

- Data is transmitted in continuous stream at constant rate.
- high speed communication
- clock signal is used
- sender & receiver must be synchronized.

Asynchronous

- Data is sent in individual characters or blocks, with start & stop bits.
- No need to synchronize the sender & receiver.

Based on authentication

Peer to peer connection

- ↳ no dedicated servers
- ↳ All computers are Equal so, termed as peer.
- ↳ function as both a client & server

Server Based Connection

- ↳ Dedicated server
- ↳ one computer serve as a server & can't be used as client.
- ↳ server control the clients

Based on Geographical location

LAN (Local Area Network)

- ↳ small high speed N/w.
- ↳ few no. of system interconnected
- ↳ distance ↑ speed ↓
- ↳ limited to few only
- ↳ used to link the devices in single office, building, campus.
- ↳ owned by private people

WAN (wide AREA Network)

- ↳ Network speed is smaller
- ↳ connects system indirectly
- ↳ spread over the world. over more than one city, country & continent.
- ↳ WAN is slower than LAN
- ↳ owned by single owner So, Enterprise network
- ↳ Combination of one more than one topology

MAN (metropolitan area network)

- ↳ Extension of local area network
- ↳ spread over city
- ↳ single network or a network of more than one local network & share the resources

Based on Reliability

Connection-oriented

- Reliable network service
- Establish a session connection before data can be sent.
- Guarantee that data will arrive in same order.

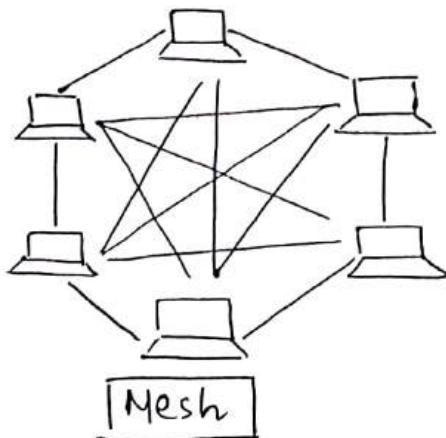
Connection less

- doesn't require a session connect
- provide minimal services.

Topology :-

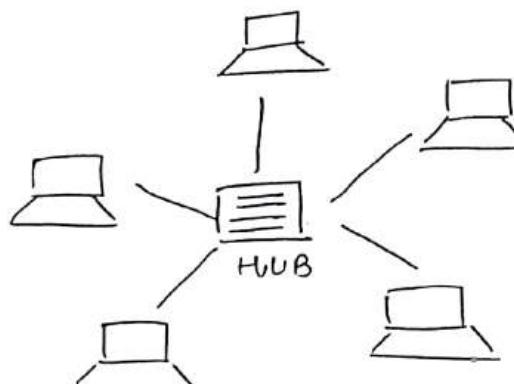
- Way in which a network is laid out physically
(दो या दो से ज्यादा devices connect hote hai link se to use topology bolte hai)
- Geometric Representation of Relationship of all the links & linking devices to one another

4 topologies



Each node is connected directly to Every other node in network

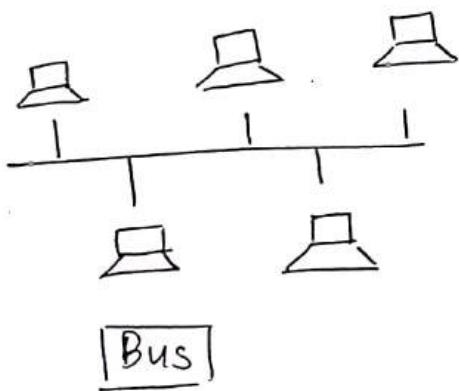
Point to point connection.



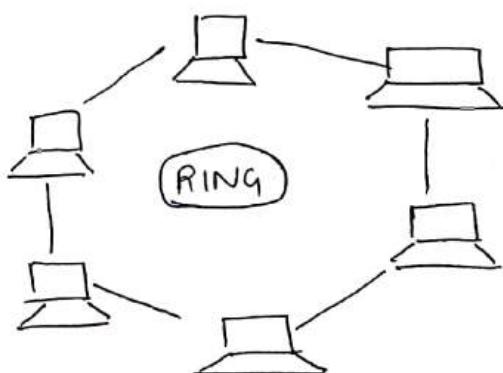
star topology

All nodes are connected to central points/device (hub)

indirectly connected to Every other node through the central device



- all nodes are connected to a single communication line.
- har node ke pass data jata hai bus process ho karta hai, jiske lie intended hota hai



Each node connected to exactly two other nodes.

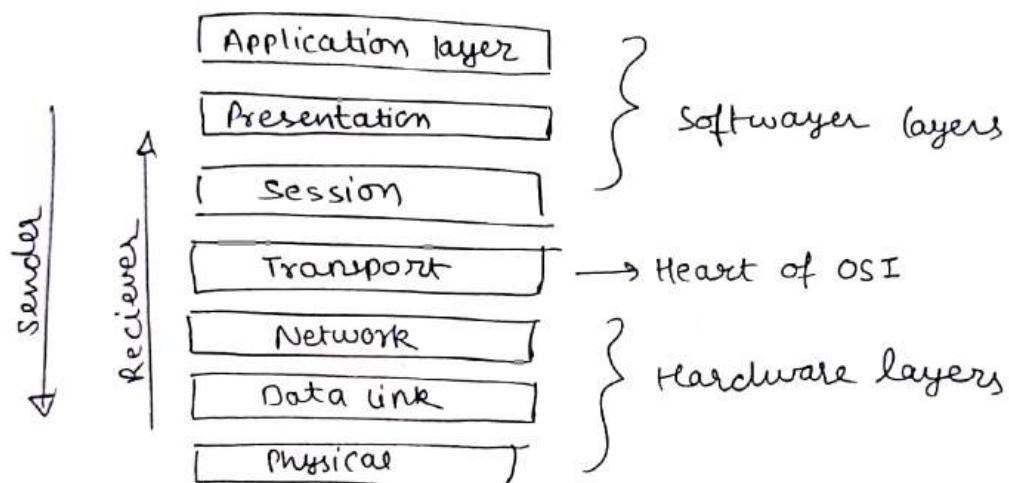
forms a close loop or Rig structure.

Data travels in one direction Pass through Each node until Reach its destination

Network Models

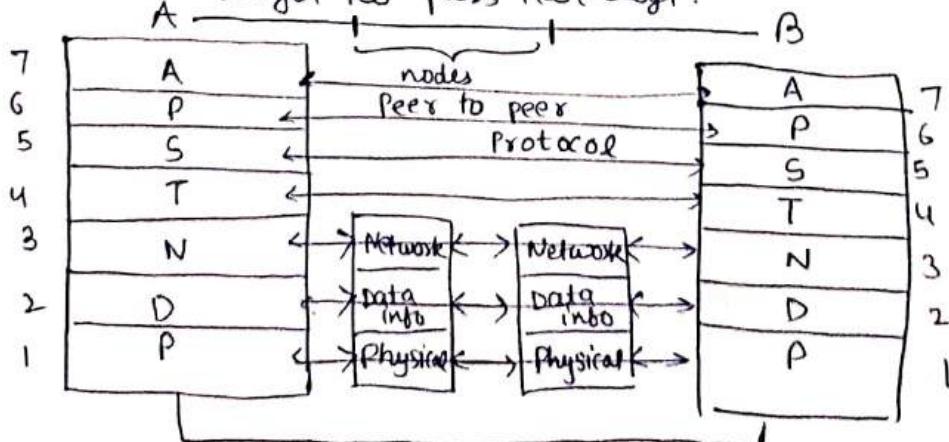
International organization standardization

- OSI model - 'Open system Interconnection)
 - ↳ ISO ne develop kia
 - ↳ facilitate communication between two different system bina change Kare logic ko hardware or software ke.
 - ↳ OSI model protocol nhi hai, model hai understanding & designing a network architecture jo flexible, robust & interoperable ho
 - ↳ 7 separate but related layer hoti hai, Har kisi ka apna kaam hota hai, poore process ke part ko define Karti hai.



Aaj Phir Se test nhi dena pada

- isme 2 device hong A & B.
- communication hoga move down through layers on Device A & Back up through B.
- Har layer sending device mein add Kargi information ko message mein, jo bhi upr wali layer se Recieve Karegi or apne se meeche wali layer ko pass Kar degi.

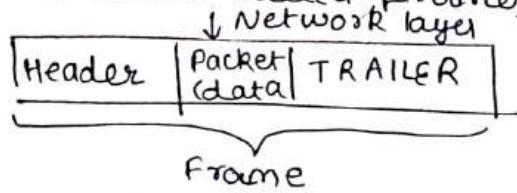


Physical layer

- defines the characteristics of the interface between devices & transmission medium.
- transmission rate (no. of bits sent each second) is also defined by Physical layer.
- Physical topology is handled.

Data link layer -

- Divides the stream of bits received from network into manageable data units called frames



↓
Physical layer

- agar frames distribute Karenge different System mein network ke. DLL adds a header & define sender & receiver
- agar Rate of Receiving Data less hoga production of data se, to yeh layer impose Karti hai flow control mechanism.

Network layer -

- ↳ responsible for source to delivery to destination of a packet.
- ↳ logical addressing
- ↳ agar 2 system connected hai same layer se to network layer ki tarot nahi hoti.

Transport layer -

- ↳ connection control
- ↳ error control
- ↳ divide karta hai data ko small parts mein / segments mein for transmission or assemble karta hai destination.
- ↳ flow control.

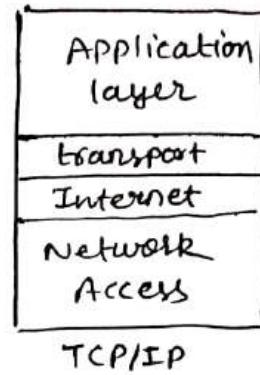
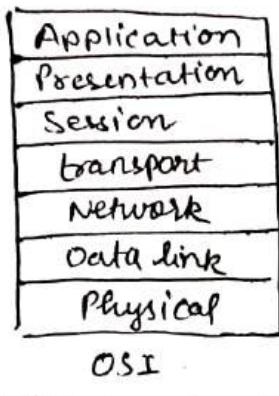
Session layer -

- ↳ Establish, maintain & synchronize Karti hai interaction among communicating Systems.
- ↳ Dialog control - communication b/w two process in half & full duplex.
- ↳ Synchronization - add check points to stream of Data

Presentation layer

- ↳ translation : converts Data formats
- ↳ Encryption : safer secures Data
- ↳ compression : compress data size

TCP/IP Protocol suit or Model



Process to Process
 Host to process
 Source to Destination
 Node to node

There are 4 layers in TCP/IP model.

Network access -

- ↳ Combination of Physical layer + Data link layer (of OSI)
- ↳ Define how data should be sent physically through the network
- ↳ mainly responsible for transmission of Data b/w two devices
- ↳ It encapsulate the IP datagram into frames transmitted by network & mapping of IP addresses into physical addresses

Internet

- ↳ Second layer
- ↳ main responsibility is to send the packets of any network & they arrive at destination irrespective of route

Protocols used by Internet layer -

IP Protocol

- ↳ IP addressing used by the Internet & higher layers to identify device
- ↳ Host to Host communication : Determine the path through which data is to be transmitted
- ↳ Data Encapsulation & Formattting Ensure the data is sent & received securely, encapsulate the data into message
- ↳ Routing : job source & destination don't have to, IP datagram indirectly the job has to.

Transport layer -

↳ Responsible for reliability, flow control & correction of Data
jo send hona hai network ko

Protocols

(1) UDP (User Datagram protocol) :

↳ Connectionless service & End-to-end delivery of transmission
↳ Discover error, & ICMP protocol reps the Error to Sender

(2) Transmission control protocol (TCP) :

↳ Full transport layer service to application

↳ virtual circuit between sender & receiver active during transmission

Application layer -

↳ top most layer

↳ Responsible for handling high-level protocols, issues of representation

↳ allow the user to interact with application.

↳ jab ek application layer duri application layer se communicate
Karna chahi hai to vo data ko transport layer pe forward kar
data hai.

Main protocols of application layer -

[HTTP] - allow to access data - over world wide web. transfer data in
form of plain text, audio video.

[SNMP] - Simple Network Management protocol
Provide framework used for managing the device

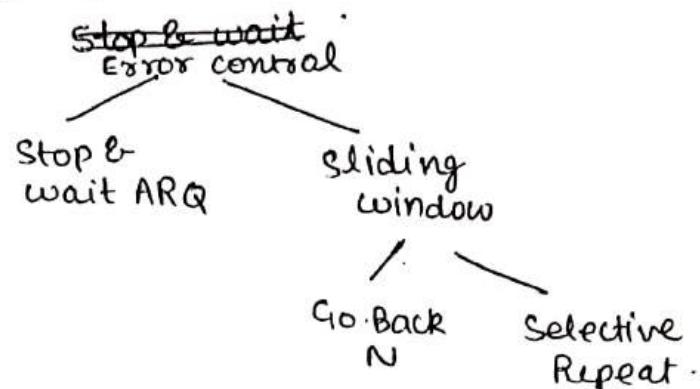
[FTP] - File transfer protocol (used for transmitting the files from
one computer to another)

Unit - II

Data link protocols

→ Stop & wait protocols

Error control mechanism - used to ensure data should be exactly same whatever sender has sent the data.



Stop & wait

- ↳ jo bhi sender data send kar raha hai, receiver ko. After sending Sender waits karta hai tab tk jab tk usko acknowledgement receive nahi ho jata receiver se.
- ↳ flow control protocol hai, service of data link layer
- ↳ used for transmitting the data over noiseless channels.
- ↳ unidirectional (ya to send hoega ya receive)

Primitives of Stop & wait protocol :-

Sender :

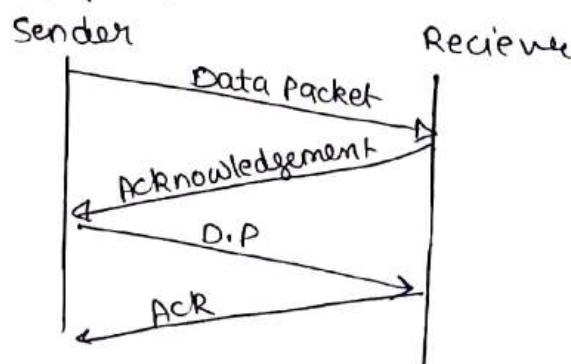
Rule 1 Sender sirf one data packet send kar skta hai ek baar m.

Rule 2 Sender next packet sirf tabhi send kar skta agar acknowledgement mil rhi hai previous packet ki toh.

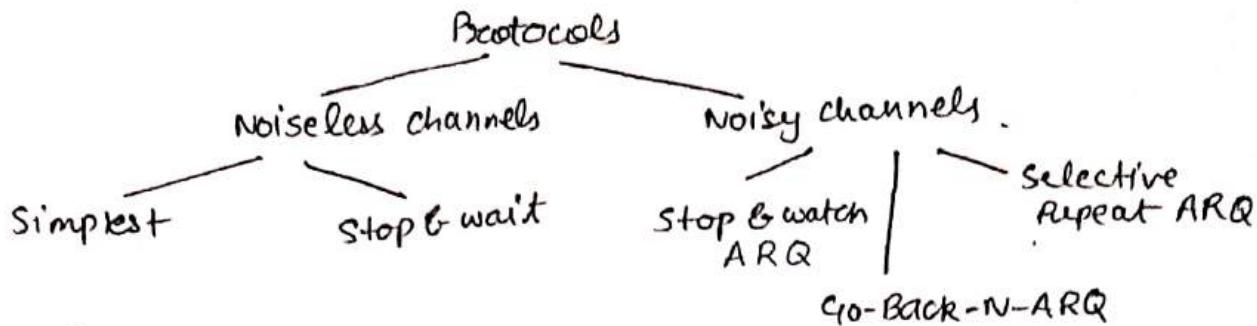
Receiver :

Rule 1 Receive & consume Data pack

Rule 2 jab consume ho jate ack send karo sender ko .



Protocols

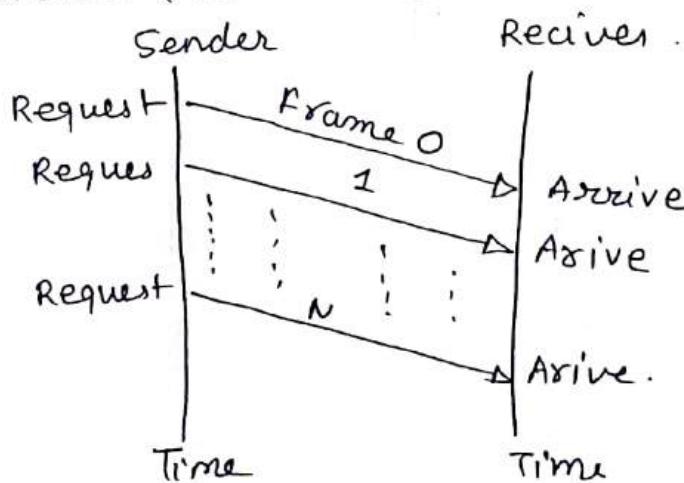


Protocols -

- ↳ defined set of Guidelines & Regulations to control karte hai communication b/w diff. devices in a network.
- ↳ yeh guideline specify karte hai ek way jisme data format, timed, sequenced & check for error during transmission
- ↳ ensure accurate, efficient, clear communication between ideas.

Simplest Protocol -

- ↳ involve direct transfer of data from the source to destination.
- ↳ without any intermediate processing.
- ↳ noisless protocol hai (metlab jaisa data jaayega waise hi receive hoga)
- ↳ Doesn't have any mechanism for controlling the flow of data or detecting & correcting
- ↳ one way protocol (allow data to flow from sender to receiver)
- ↳ unidirectional (no need for acknowledgement)



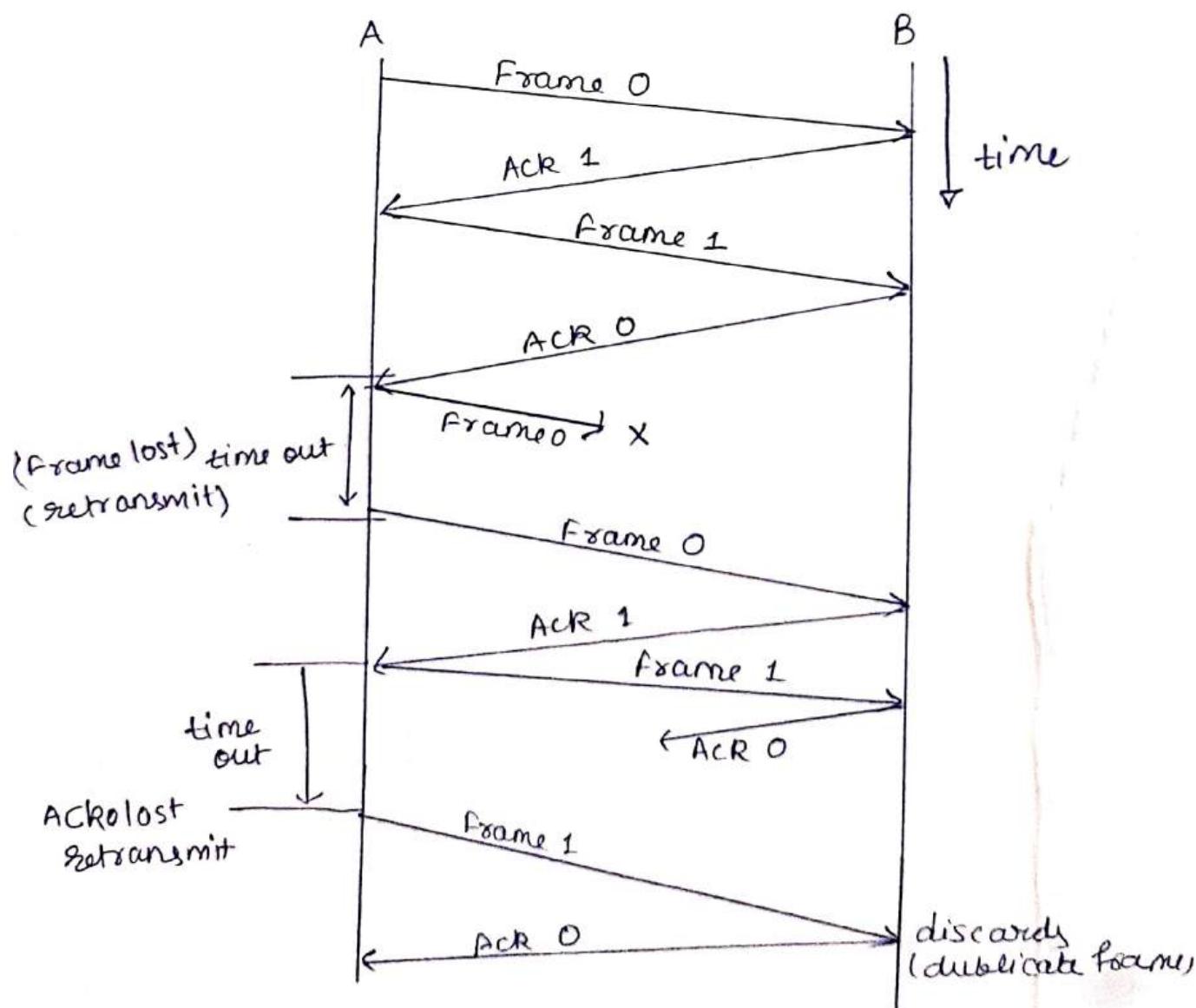
Stop & wait

↳ (Peeche ho gya) 

Noisy channel protocol -

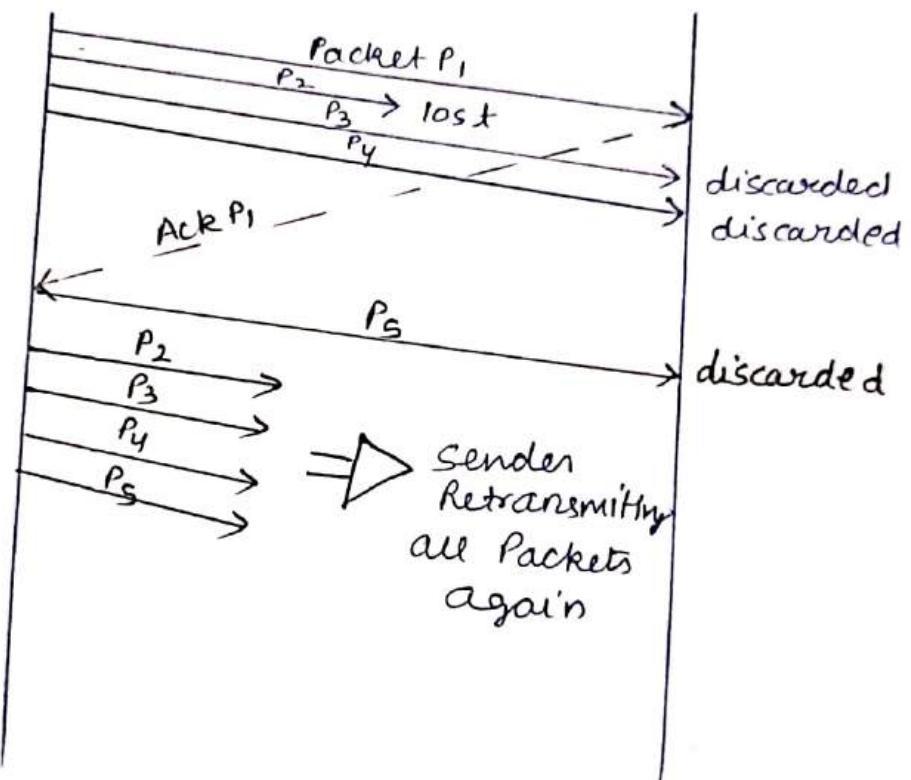
Stop & wait ARQ -

- ↳ used for reliable data transmission over a noisy channel
- ↳ Sender sirf ek frame send karta hai & wait karta hai.
Ack ke lie before sending next frame.
- ↳ yeh help karta hai janne mein ki receiver ke pass
data pauch gya (kunki yeh noisy channel hai)
- ↳ sender continuously monitors the channel for errors,
agar error detect hota hai to next Ack ka wait
karta hai, frame ko dubara bhejne se pehle.



Go-Back N ARQ

- ↳ Error control protocols used in data communication to ensure reliable delivery of data over noisy channel
- ↳ type of sliding window protocol
 - ↳ Sender ER window of packets bhejta hai receiver ko.
 - OR Receiver ACK send kardhi hai
 - ↳ in case agr sender ko shi time pe ACK receive nahi hوتا, to sender entire window ko dubara send karta hai



Sender side

- Packet send karega Receiver Ko jisme (i+N-1) frame honge.
- Sender set karega tim for each packet of window.
- Sender wait karega ACK ka

Receiver

- receiver packet receive karega & check karega error.
- correct Recive hua hai to ACK send karega with seq.no jisme Receiver hua the
- agr error ke sath receive hua hogya to discard karaya Nak bhejya.

Sender

(if No ack Rec)

- agr jo time hai ack Recieve
Kone ka unme ack Recieve
nhi hui to sender vhi window
Retransmit karega.
- time Reset Karega
- wait kaega ACK from Receiver

Receiver (In case of Nak)

- agr NAK Recieve hua
to vhi packets dubara
send karega jisme
Error aya or Recieve
nhi hue Receiver ko
dheng se
- Reset timer for each
packet
- wait for ACK.

Selective Repeat ARQ

- transmit a window of packets to Receiver, Receiver
ack send karta to sender.
- agr error detect hota hai to Nak send karta hai Reg. for
retransmission.
- sender maintains a timer for each packet, agr ack nahi
recieve hui to retransmid hoega.
- agr packet correctly Recieve hua to Receiver Ack send
karega seg. no. ke sath, agr error ke sath Recieve
hua to discard karega Nak send karega & Retransmit
ke lie Reg. karega.
- buffer maintain rehta hai, throughout the process.

Sender

- (i+N-1) Packets send Karenge Ek
window mein.
- time set Karenge
- Ack ka wait Karenge

Sender (if ^{no} ack Recive)

1. time out se pehle ack nahi
mila to sender Retransmit karega
2. time set karega.
3. wait karega ack ka.

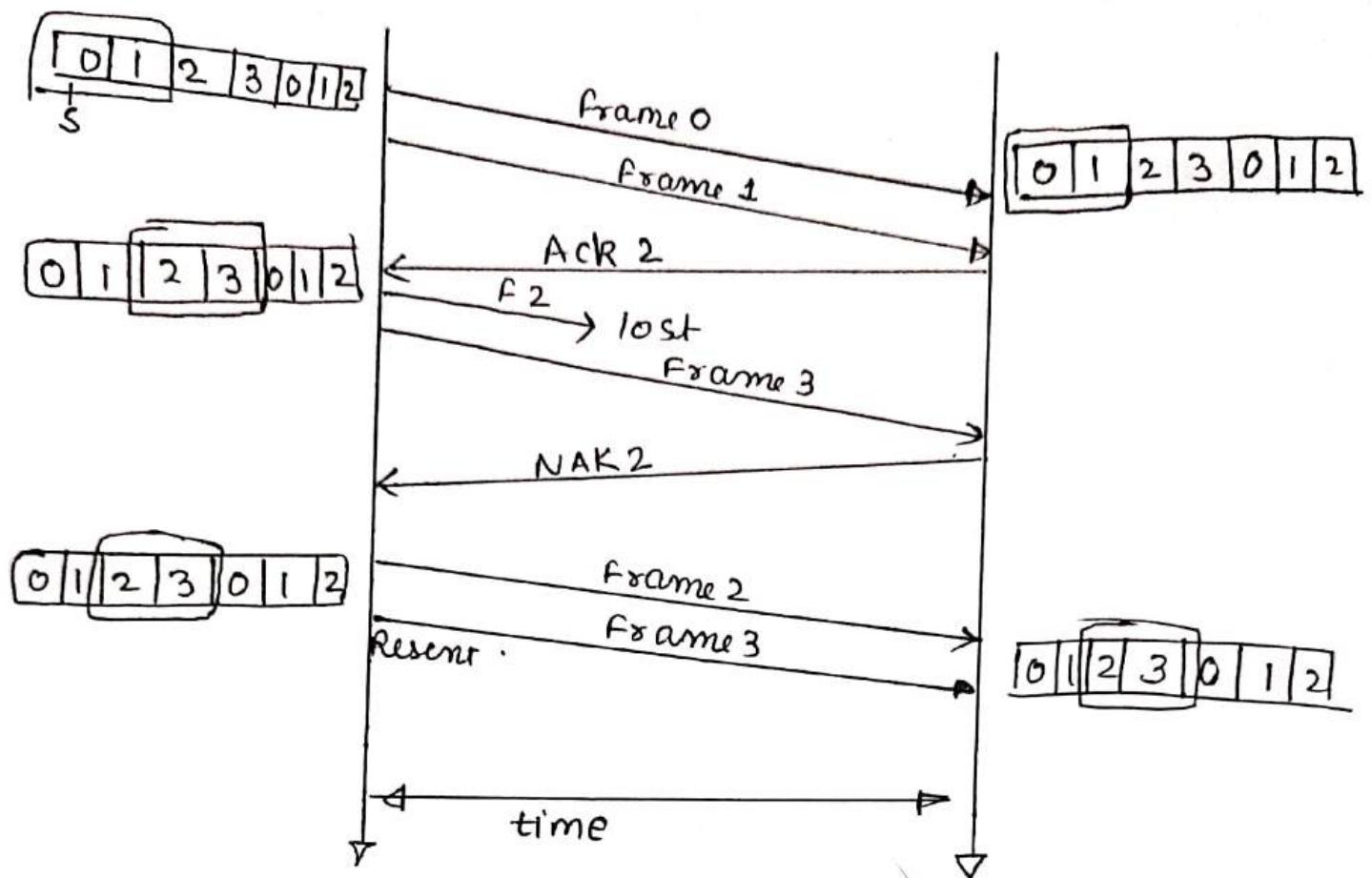
Receiver

1. check for errors.
2. Recieve correctly - then ack
the packet
3. otherwise Nak & Reg. for
Retransmit

(if Nak Recieve)

1. agr NAK Recieve hua
sender sirf vo Packets transmit
karega jo correct Recieve nahi hue
2. timer set karega
3. ack no wit karega.

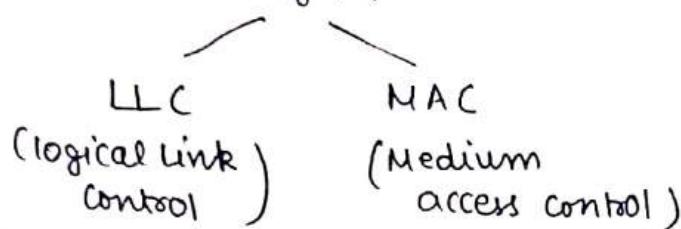
M'



MAC layer

OSI Model layered networking framework hai jo communication mein help Karta hai 2 heterogeneous systems mein.

Data link layer (2nd lowest layer)



Functions of MAC -

- ↳ Determine Karta hai Kaise device access gain Karta hai transmission medium mein, data ko transmit krne ke lie.
- ↳ used for addressing & identifying devices.
- ↳ Adds necessary header & trailer information to data to facilitate proper transmission & reception
- ↳ Error detection & Ensure data integrity.
- ↳ Regulate the Rate of transmission between Devices

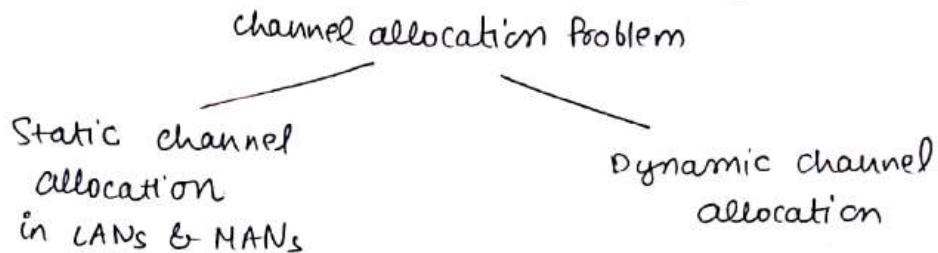
MAC address -

- Unique identifier allotted to network interface controller (NIC) of a device.
- used as a network address for data transmission.
- assigned to network adaptor at time of manufacturing.
- MAC address comprises of 6 groups. (hexadecimal digits) Separated by hyphens, colons or no separators.

Ex 00: 6A: 89: 5B: F0: 11.

Channel allocation problem -

- ↳ signal channel ko divide Karte hai & allocate Karte hai to multiple user taki vo different task perform kar sake.
- ↳ Multiplexing is used
- ↳ channel allocation problem can be solved by two schemes



Static channel allocation in LANs & MANs.

- classical or traditional approach of allocating a single channel among multiple competing users using Frequency Division Multiplexing (FDM)
- if there are N users, the frequency channel is divided into N equal sized portions
- not suitable for large no. of users.
- not efficient to divide into fixed no. of chunks.

$$T = 1/(U * C - L)$$

$$T(FDM) = N * T(1/U(C/N) - L_N)$$

T = mean time delay

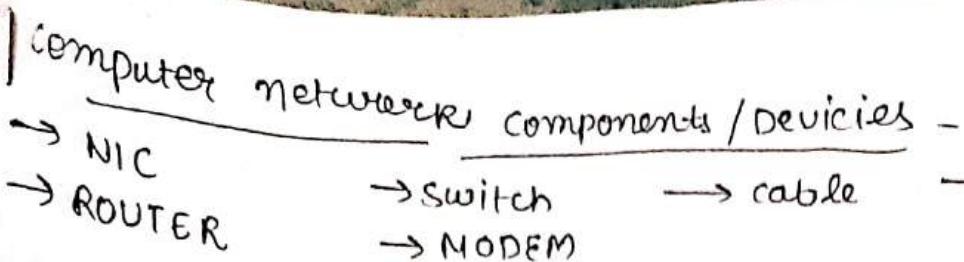
C = capacity of channel

L = arrival rate of frames.

Y_U = bits / frame.

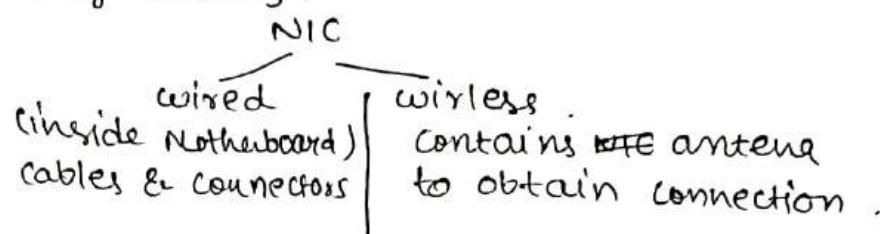
N = No. of sub channel

$T(FDM)$ = Frequency Division Multiplexing time



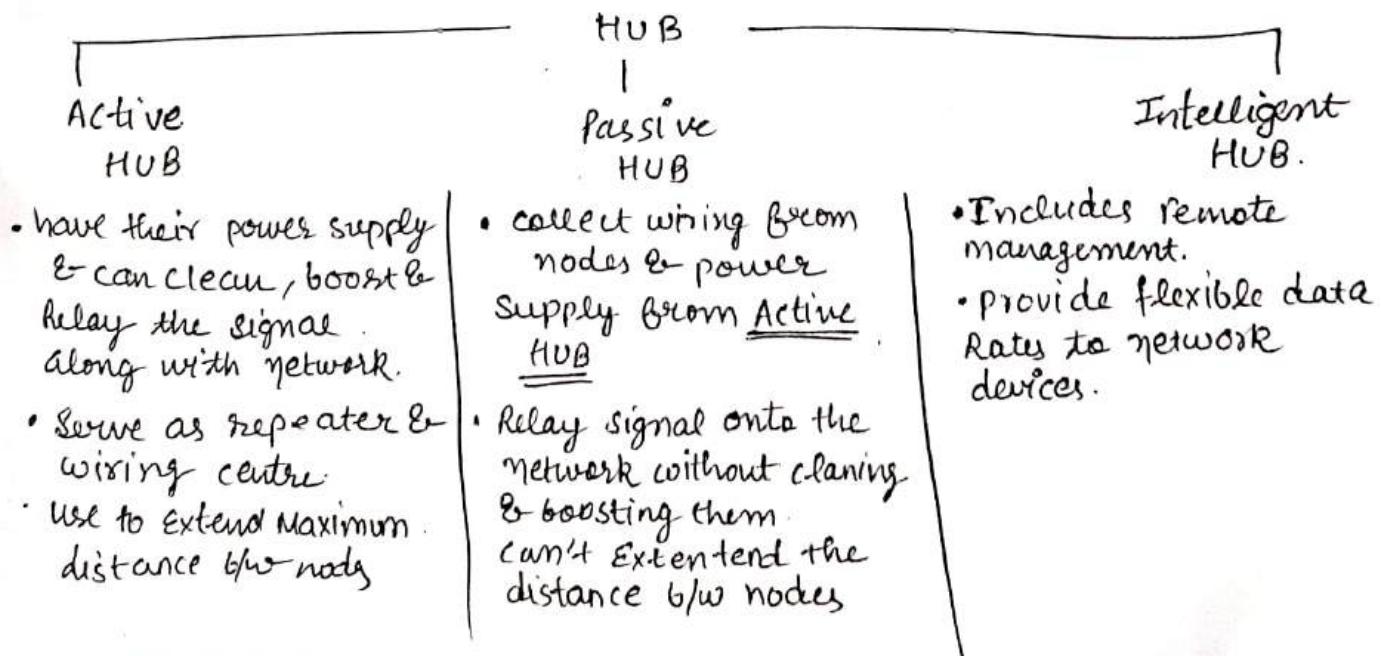
NIC - Network Interface card -

- Used to connect computer with another component onto Network.
- Support a transfer Rate 10, 100 to 1000 Mb/s
- NIC ke MAC address / Physical address Encoded hota hai jo ki assigned hota hai, IEEE as jo ki uniquely define Karta Network Card ko.
MAC address store hota hai "PROM me (Programmable Read only memory)



HUB - hardware device, divides Network connection among multiple device.

- Jab computer request Karta hai Kisi information ke lie. to vo Request Karta hai Sbse pehle Hub se cable ke through HUB parne broadcast ho vo Req Send Karta hai Sb device check Karte hai ki Request unse Related hai ya nhi, agr nhi to drop the Req.
- HUB limits the amount of communication.



Switch -

- multiport bridge with buffer + a design can boost its efficiency & performance.
- Data link layer Device.
- can perform error checking before forwarding data.
- switch / divides the collision domain of hosts, but broadcast domain remains same.

Types

1. unmanaged switches - home networks & small business
 - as they plug in & start work instantly start doing their work.
 - require small cable connection only.
 - allow device on a network to connect with each other.
 - least expensive.
2. Managed switches - highest security, precision control & full management.
 - used in organization containing a large network.
 - customize to enhance the functionality of a certain network.
 - most costly, but ideal.

Router - hardware device, used to connect a LAN with Internet connection.

- used to receive, analyze & forward the incoming packets from another packet.
- Network layer mein kaam karta hai, router ki.
- forward karta hai packet based on info available in routing device.
- determine best path from available for transmission.

Advantages

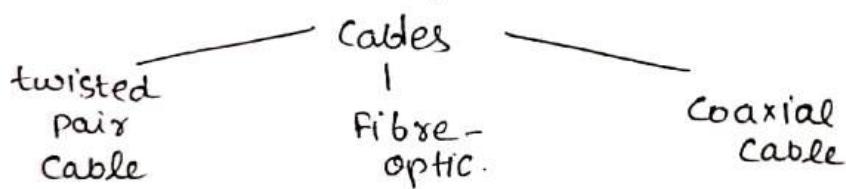
- Security
- Reliability
- Performance

Modem

- ↳ hardware device allow computer to connect to internet over the existing telephone line.
 - ↳ not integrated with motherboard, installed in the PCI slot.
 - ↳ stands for modulator. converts digital data to analog signal
- modem
- (Based on speed & transmission speed)
- | | | |
|--|-------------------|----------------|
| Standard PC modem or Dial up modem | cellular modem | cable modem |
|--|-------------------|----------------|

Cable & connectors -

- ↳ transmission media, used for transmitting a signal.



Repeater - . operates at physical layer .

- regenerate the signal over the same networks isse pehle ki vo weak / corrupted ho.
- jab signal weak ho jata hai to vo ~~execute the gate~~ copy karte hai bit by bit & Regenerate karte hai ,
- It is 2-port device.

Bridge - operates at data link layer, repeater hota hai & type ke functionality → filter content by Reading the MAC address of source & destination.

- ↳ Transparent
- Stations are unaware of Bridge Existence.
- reconfiguration of stations is unnecessary

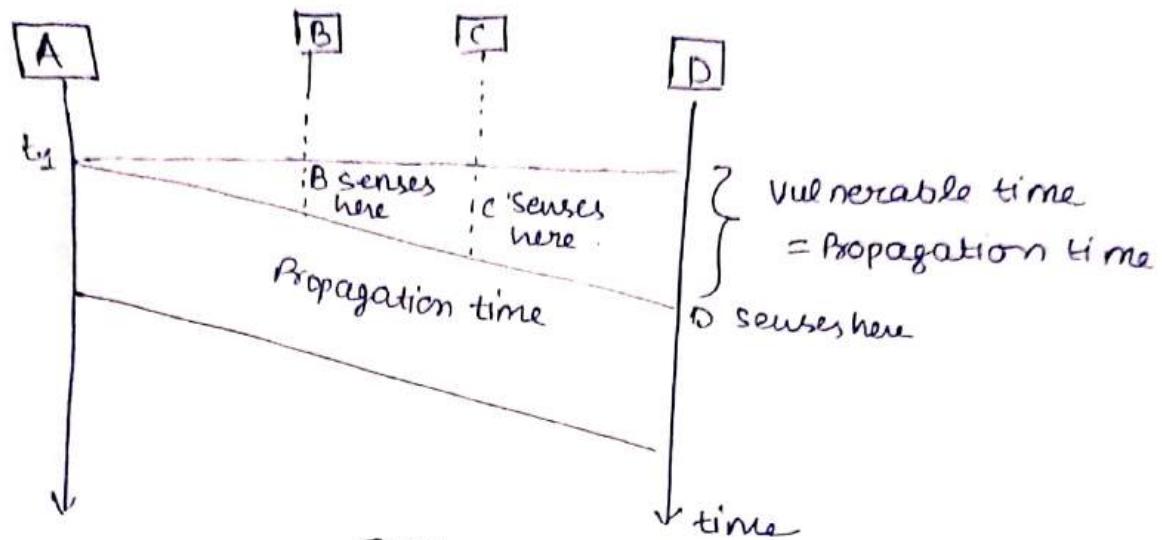
Bridge

Source-Routing

- performed by source station station & frame specifies which route to follow

SMA (Carrier Sense Multiple Access)

Developed to decrease the chances of collision job 2 or 2 se jyada stations signal sent karne hai data link layer ko.



Type CSMA access mode

1-Persistent

senses - shared channel
First & deliver
the data right
away if channel
is idle.
agr nahi hai to
continuously track
for channel jb tk
idle nahi ho jata
or the broadcast
Karega frame ko
without any
condition as soon as it
does.

→ aggressive transmission
algorithm.

Non persistent

access the channel
before transmitting
data, agr channel
idle hai to node
transmit Karega data
agr nahi hai to
station wait Karega
ER arbitrary amount
of time ke lie
(not continuously)
or jab dekhega ki
ab Empty/Idle hai
to frame send
Karega.

P Persistent

(1 persist + Non persist)
Har node observe
Karegi channel ko
1 persistent main
ek channel idle
hai to send
Kareg frame
P probability
ke sath, agr
data transfer
nhi hua to,
frame restart
Karega. Kisi
bhi time slot
ke sath for
q (1 - P)
Random
period

0-persistent
supervisory
node order
deti hai ha
node ko
transmission
Ke lie. Node
wait karti
hai apne time
period ke acc.
jo allotted
nota hai
on demand

Advantages of CSMA

- (1) Efficiency ↑ → ensure karta hai 1 time pr 1 device communicate karne, jisse collision ram noga or efficiency accuracy badhgi.
- (2) Simplicity - easy to implement & doesn't require complex hardware or software.
- (3) Flexibility - wide range of network environments use ho skte h.
- (4) Low Cost - jyada complex hardware or software ki zaroorat nahi hoti.

Disadvantages -

- (1) Limited Scalability - inefficient no. of devices hote hain network pe.
- (2) Delay - wait karna Medium ko sense karne ke lie or availability of channel ke lie delay kar skta hai.
- (3) Limited Reliability - affected by interference.

Collision Detection in CSMA/CD.

→ measure for effective transmission.

→ This protocol decides which station will transmit when so, that Data Reaches the destination without Corruption.

Working -

1) Check sender Ready hai packet serve karne ke lie ya nhi.

2) Check if transmission link is idle.

Sender check Karta Rehta hai ke transmission link/medium idle hai, iske lie vo continuously sense Karता other transmission nodes ko. Sender dummy data send karega link pe or agr koi signal Recieve nhi hua collision ko, istka matlab idle hai us moment pe, over again agar hua to Refrain Kar

3) transmit the Data, check for collision.

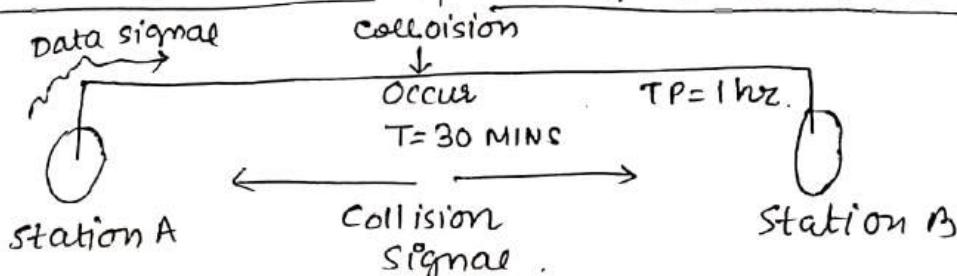
Sender transmit Karega data ko link pe. CSMA/CD doesn't use an "Ack" system. Yeh check Karta hai successful & unsuccessful transmission through collision signal.

Agr transmission ke samay collision signal Recieve hota hai node ko to transmission Stop ho jayega. pher station transmit Karega jam signal link pe and wait Karega Random time interval ke lie. before it Resend the frame.

Pher Kuch Random time ke Baad vo again attempt Karega transfer Karne ke lie Data. and Repeat Karega above process.

4) agr no. collision was detected in propagation, Sender complete its frame transmission & reset the counters.

How does a station know if its Data collide -



Collision free protocols in computer Network

Almost sbhi collisions avoid kie ja skte hai, CSMA/CD mein phir sbhi vo ho jate hai contention period mein. Or yeh adversely affects Karte hai system performance ko.

Yeh tb hota hai jab cable long hoti hai or length of packet short. Or yeh problem bht serious bn jati hai jb fibre optic use karte hai.

Protocols banne hai unko resolve karne ke lie.

↳ Bit map protocol

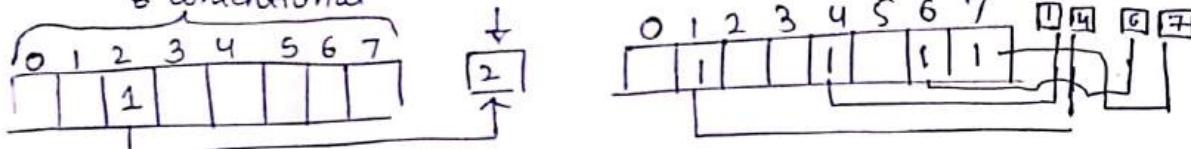
↳ Limited contention

↳ Binary countdown

↳ Adaptive tree walk protocol.

1. Bit map protocol

- Each contention period Exactly N slots ka hota hai.
- Agar kisi station ko frame send karna hai, to usko i bit corresponding slot ko transmit karni hogi.



Iski limitation yeh thi ke aage, hume 1000 frame send karne hai to, 1000 bits chahiye honge. Matlab for large no. of frame we need same no. of bits.

Binary countdown - Binary station address use hota hai.

Jo station chalta hai send karna usk address length, or vo apni ^{successive} MSB compare karega, or karta rehega till end. Aage MSB 0 hogi to ~~next~~ wo give up kar dega otherwise, aage wali bit pe jayega compare karega

For Exp.

| Bit time | 0 0 0 1 | 1 0 0 1 | 1 1 0 0 | 1 0 1 1 |
|-----------|---------------|-------------|---------|-------------|
| check MSB | 0 0 (give up) | 1 | 1 | 1 |
| | 1 — — | 0 (give up) | 1 | 0 (give up) |

1100 will transmit data/frame.

Limited contention protocol -

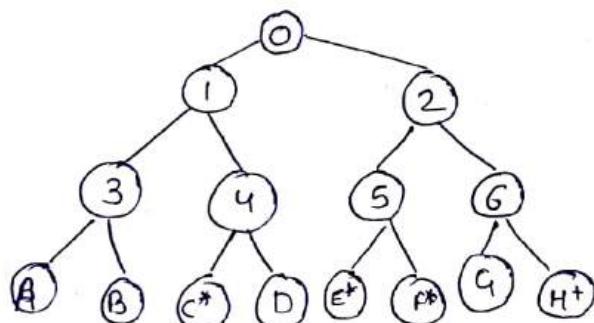
- collision Based protocols (Pure & slotted ALOHA, CSMA/CD) are good when network load low
- Collision free protocol (Bitmap, Binary countdown) are good when load high.

Adaptive tree walk protocol

- Partition the group of station & limit contention of each slot
- under light load \rightarrow each slot like aloha
- under high load \rightarrow only a group can try for each slot

How?

- treat every stations as leaf of binary tree
- First slot (after successful transmission), 8 stations try karen ga pna slot lena ka (under root)
- If no conflict, well & good.
- else, only nodes under sub tree get to try for next one (DFS)



Slot 0 : C*, E*, F*, H* (all nodes 0 ke under hai) conflict .
8b send karne ka try karenji .

Slot 1 : C* (all nodes under 1 can try) , C sends .

Slot 2 : E*, F*, H* (all nodes under 2) all try) , conflict .

Slot 3 : E*, F* (all nodes under 5 all try) , conflict .

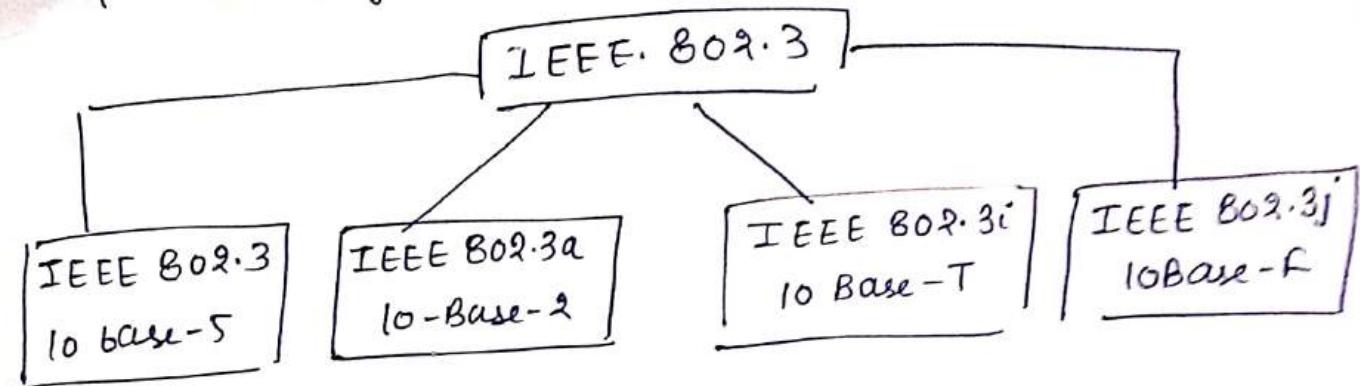
Slot 4 : E* (all nodes under E) all try E sends .

Slot 5 : F* (all nodes under F) all try F sends .

Slot 6 : H* (all nodes under H) all try H sends

IEEE 802.3

- set of standards & protocol that define Ethernet-based networks.
- use LANs, iske use se yeh MANs mein bhi & WANs mein bhi use ho sakte hai.
- defines Physical layer & medium access control (MAC) of Data link layer



Frame format -

- preamble - 7 bytes. starting field provide alert & timing pulse for transmission.
- start of frame Delimiter : 1 byte field.
 - ↳ contains an alternating pattern of ones & zero ending with two ones.
- Destination address : 6 byte field.
 - ↳ physical address of destination.
- Source address - 6 byte field.
 - ↳ physical address of sending station.
- length : 2 Bytes
 - ↳ store the no. of bytes in Data field.
- Data : variable sized carry data from upper layers.
(Max = 1500)
- Padding : 46 bytes → minimum data ke saath combination mein nota hai.
- CRC - cyclic redundancy check . error detection (4 bytes)

| | | | | | |
|---------------------|---------------|---------------|---------------|--|----------------|
| Preamble 7 Bytes | SOF 1 Byte | DA 6 bytes | SA 6 bytes | DATA+Padding Min=46 Max=1500 bytes | CRC 4 bytes |
|---------------------|---------------|---------------|---------------|--|----------------|

IEEE Standard

802.11

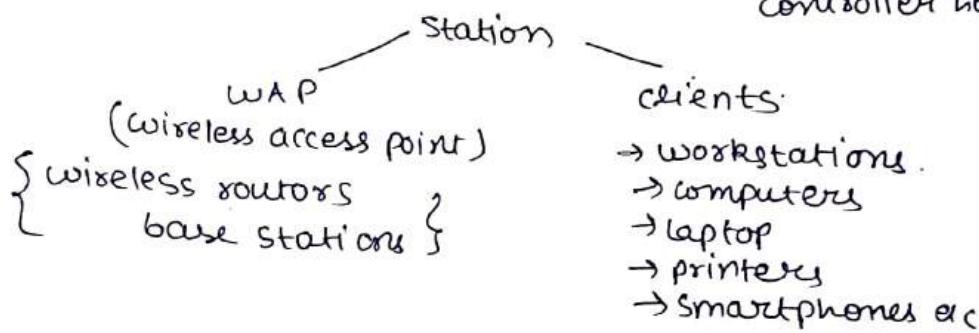
IEEE 802.11 standard, popularly known as wifi.
architecture & specifications of wireless LAN's (WLANs),

↓
use high-frequency Radio waves
instead of cables for connecting
device in LAN.

Architecture

Components

→ Stations (STA) - comprises of all devices & equipment that are connected to wireless LAN (wireless network interface controller hotsai)

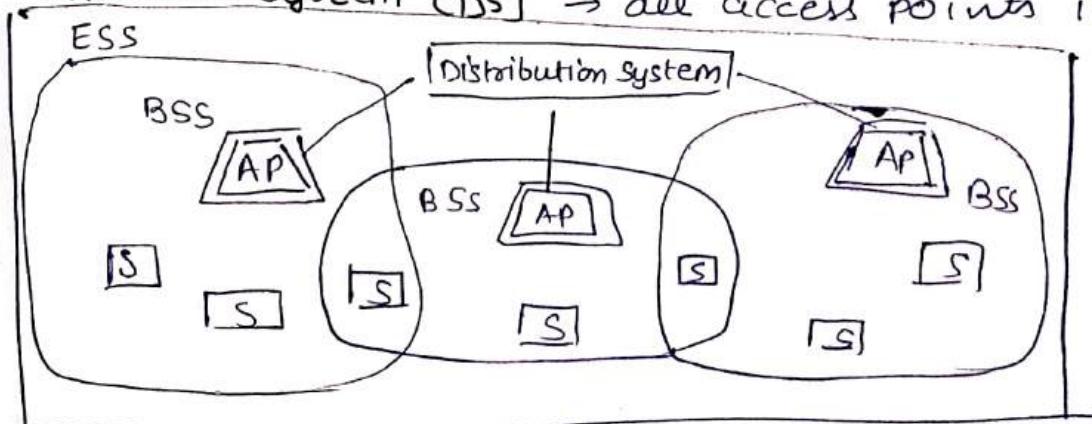


→ Basic Service Set (BSS): Group of station communicating at physical layer level.



→ ESS (Extended service set) → all connected BSS

→ Distributed system (DS) → all access points in ESS



in field of frame.

frame control - 2 byte starting field composed of 11 subfields

↳ control information

• duration - 2 byte field

↳ specify how long time period for which the frame & acknowledge occupy how much channel to

• Address fields - 6 byte (three)

↳ contain add. of source, immediate destination & final end point

• Sequence - 2 byte 3-field that

↳ stores frame no.

• Data - variable sized

↳ carry data from upper layers. (Max = 2312)

• check sequence - 4-byte field

↳ contain error detection information.

| | | | | | | | |
|----------------------------|-------------------------------|-------------------|-------------------|-------------------|---------------------|----------------------|-----------------------|
| FRAME CONTROL 2 Byte | DURATION 2 BYTE 2 Bytes | ADDRESS 1 ⑥ | ADDRESS 2 ⑥ | ADDRESS 3 ⑥ | sequence 2 bytes | Data 0-2312 bytes | check seq. 4 bytes |
|----------------------------|-------------------------------|-------------------|-------------------|-------------------|---------------------|----------------------|-----------------------|

FDDI - Fiber distributed data interface

- ↳ work on OSI model's Physical & Data link layer.
- ↳ use optical fibre to transfer data much quicker
- ↳ Dual ring system, transmit data in opposite direction.
- ↳ suitable for LANs, MANs, CANs but not WAN

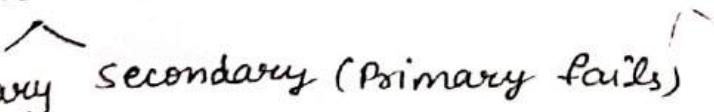
Working

- ↳ works using token ring protocol, used timed tokens
i.e., if koi host message send karne chaha hai, to vo transmitted token capture karega.
- ↳ host can transmit data for fixed time (Kuch der baad token further pass kar dia jayega)
- ↳ sb host access kar skte hai transmitted data ko & check kar skte hai "MAC address"
- ↳ Network host accept karega data ko, agar destination address match hota hai.
- ↳ Host transmit & receive kar skta hai data as frames

Frame format

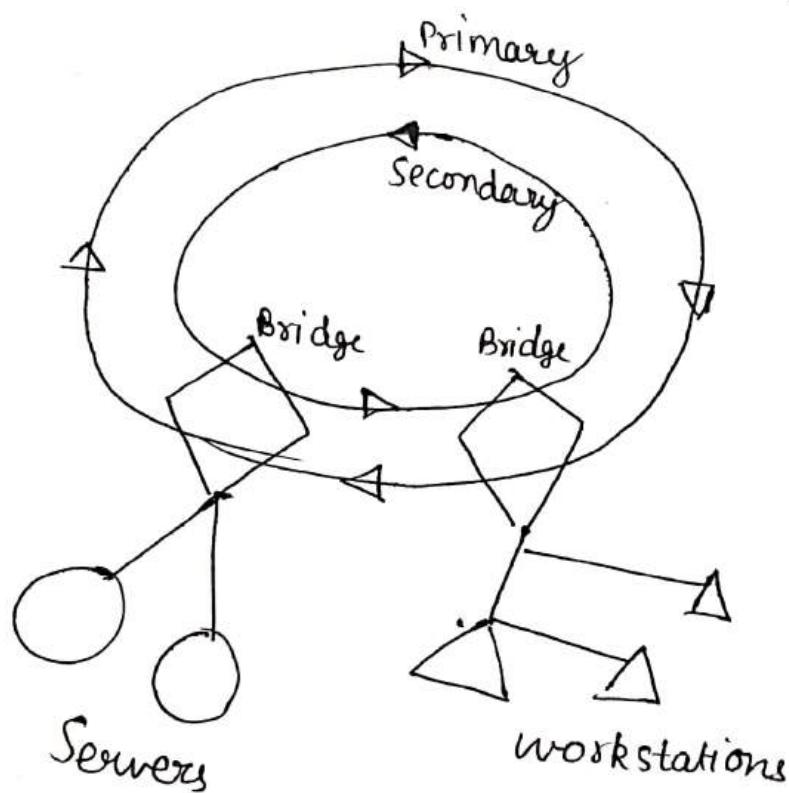
| frame header | start delimiter | frame control | destination address | source address | Payload | check sum | end delimiter |
|--------------|-----------------|---------------|---------------------|----------------|-----------------|-----------|---------------|
| 1 byte | 1 byte | 1 byte | 2-6 bytes | 2-6 bytes | Variable length | 4 bytes | 1 byte |

Topology & Design -

- ↳ has two token passing ring, transmit in opp. direction


Ring topology

- station connection
- single attached Dual attached
- ↳ Stations connected to a Primary or a Secondary Ring, lack fault tolerance (Kisi ek se)
 - ↳ can use them to connect Server to network backbone
 - ↳ connect to both Rings of FDDI Fibre networks.
 - ↳ Fault tolerance most ^{hoti hai}
 - ↳ Excellent for network backbones ~~of fault~~



Network Layer: Design Issues

1. Addressing

IP addressing - Each device on network must have unique address

Hierarchical - manage address efficiently & Reduce size of Routing table.

2. Routing -

Path Determination - best path through the network

Scalability - protocols must be able to scale as network grows

3. packet Forwarding -

Forwarding technique - how packets are forwarded

4. Error handling & Diagnostics -

Error Detection & correction - Detect & correct errors in transmitted Packet

5. Congestion control -

Avoid Congestion - techniques to prevent network congestion

Handling - packet dropping.

6. QoS (Quality of Service) -

Traffic prioritization - high priority traffic

Resource Reservation - along with network path

7. Security -

Secure Routing - protection from being tampered with

Network layer

- Third layer of OSI
- handles services from transport layer & further forwards the service requests to Data link layer.
- translates logical address to physical address
- determines the route (source to destination) manage traffic problems - switching, Routing & Controls the congestion.
- **Role** - Network layer is to move packets from sending host to receiving host.

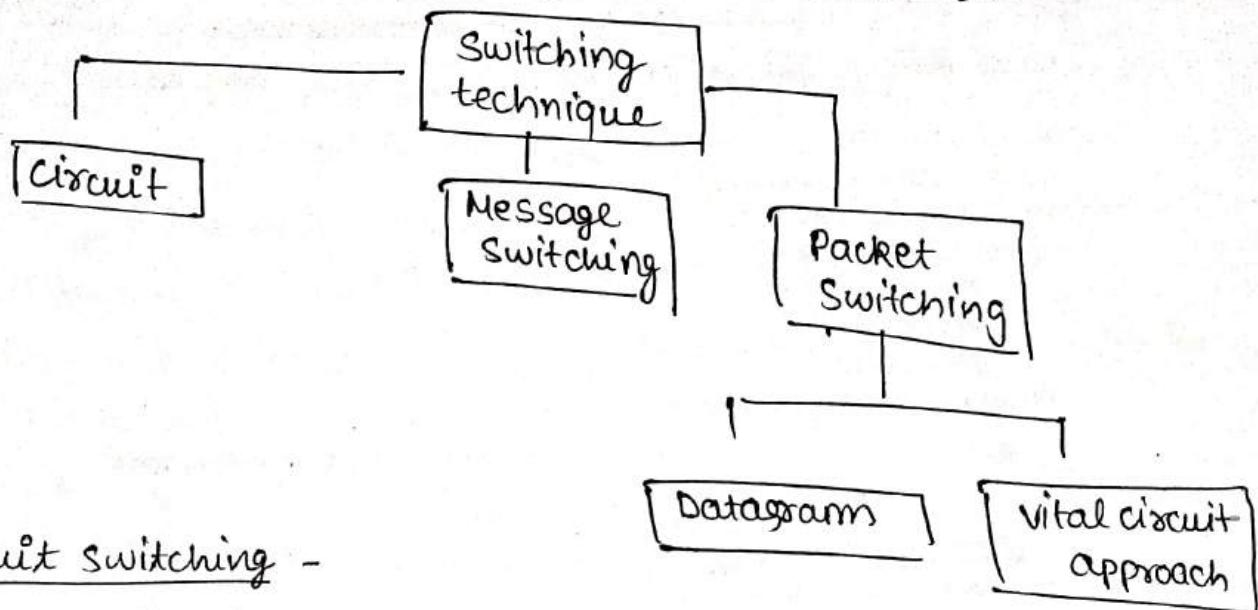
Main functions -

1. Routing - jab packet router input link pe data hai to Router usko output link pe bhej data hai.
2. Logical Addressing - implements physical addressing & network layer implements the logical addressing. used to distinguish b/w source & destination system.
It adds header to packet include logical address of sender & receiver
3. Fragmentation - breaking of packets into smallest individual data units

PACKET SWITCHING

Switching → used to connect the system for making one to one communication.

Decide best route for data transmission.



Circuit switching -

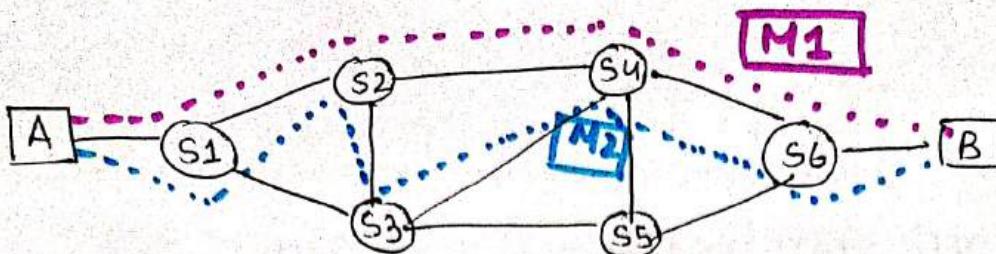
- ↳ established a dedicated path between Sender & Receiver
- ↳ Ek baar jab connection establish ho jata hai, jo dedicate Path hota hai it will remain Exist until connection is terminated.
- ↳ Before Data transfer, connection will be established first.

Phases

- 1. Connection Establishment.
 2. Data transfer.
 3. Connection Disconnection.

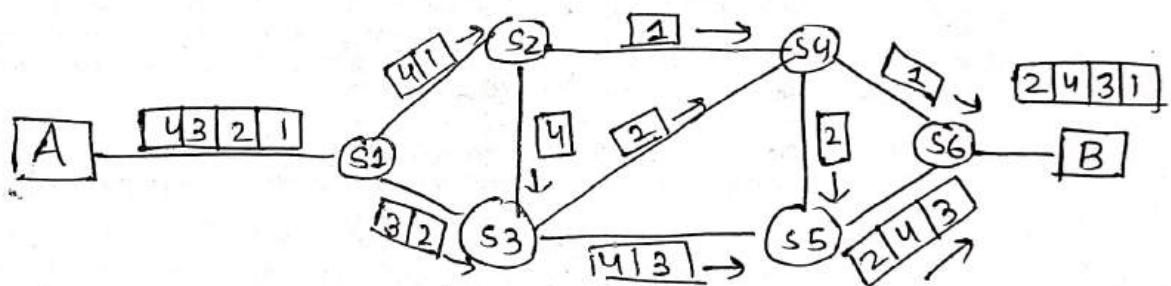
Message switching -

- ↳ store & forward mechanism.
- ↳ Message is transferred as a complete unit & routed through immediate nodes at which it is stored & forward
- ↳ no establishment of dedicated path.
- ↳ Programmed in such a way so that they can provide the most efficient routes.



Packet switching

- Message is sent in one go, but divided into small pieces
- Message splits into smaller pieces as packets & given a unique number to identify order at receiving end.
- Every packet contains some info in header as source address, destination add & Seq. no.
- Packet travel across network, take shortest path as possible.
- all packet reassemble at receiver end.
- If any packet is corrupted or missing then message will be sent to Resend the message.
- If correct then ack.

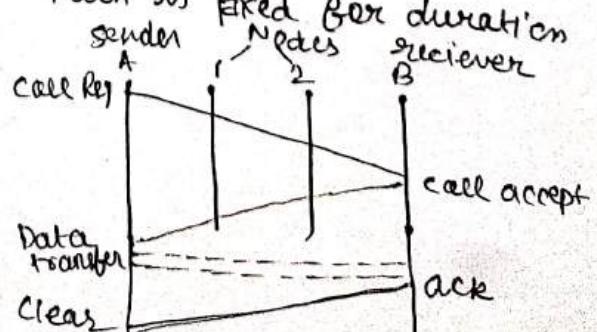


Datagram

- Packet Switching technology
 - ↳ known as Datagram
 - considered as an independent entity
- Packet contains info about destination & switch uses this info to forward the packet to correct destination.
- Packets are reassembled at Receiving end
- the path is not fixed.
- Connection less switching.

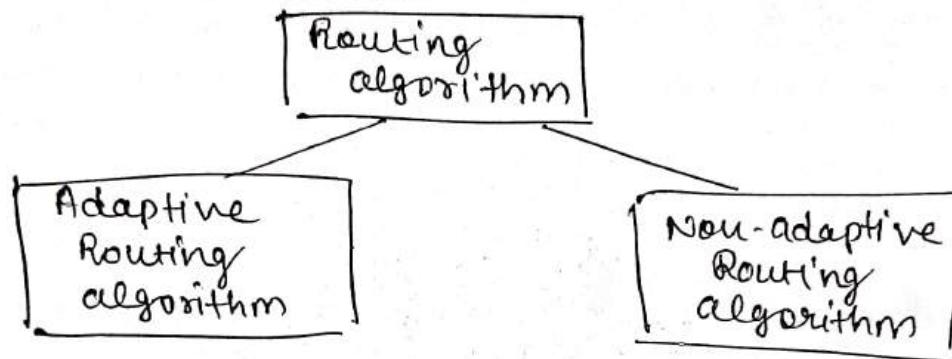
vital

- connection oriented.
- preplanned route is established before the message are sent
- call Request & accept - packets are used to establish the connection b/w sender & receiver
- Path is fixed for duration

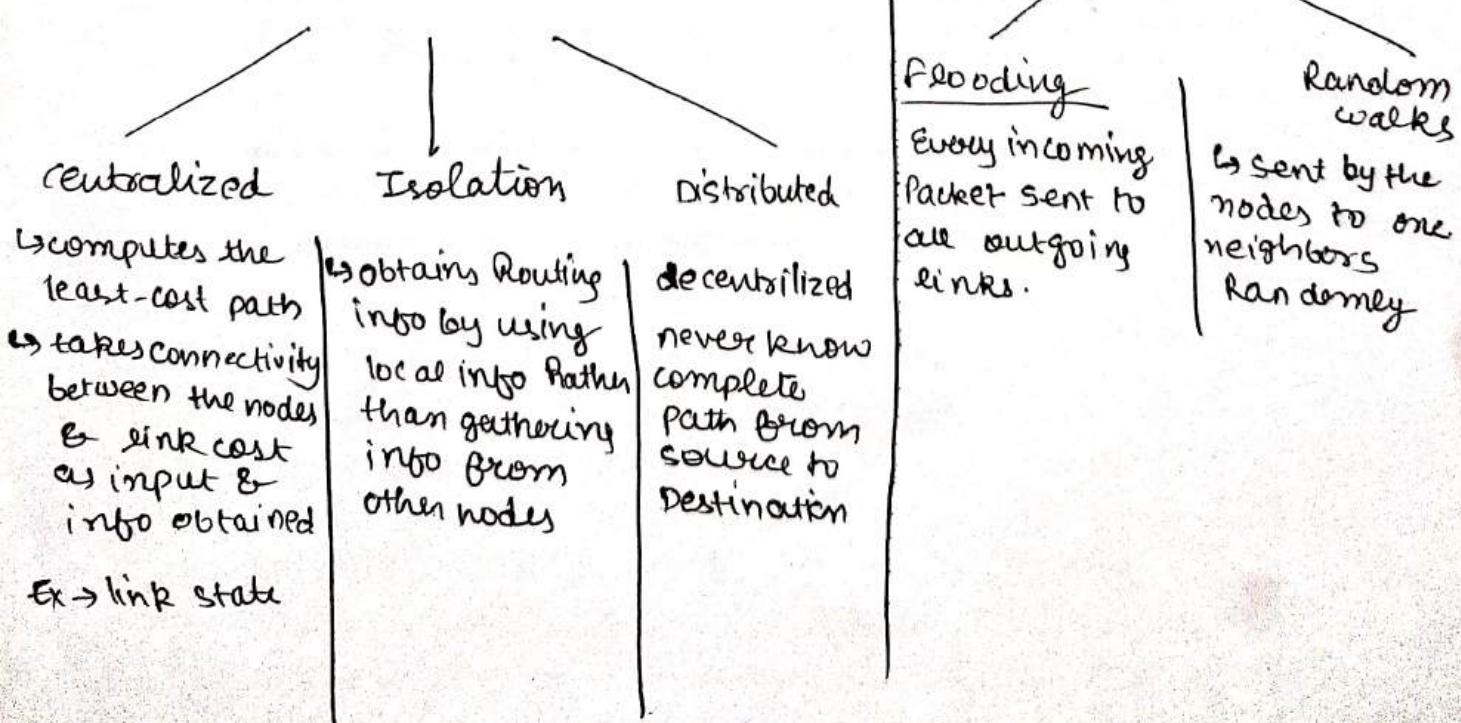


ROUTING ALGORITHM

- determines best route from Source to Destination through which packets can be transmitted.
- Routing protocol at network layer provide this job.
Provides the best path → "least cost path"
- Routing - process of forwarding the packets from source to Destination



- ↳ Dynamic Routing algorithm
- ↳ makes the Routing ~~algo~~ decision based on topology & network traffic.
- ↳ Parameters - hop count, distance, estimated transit time.



(dynamic algo) Distance vector Routing Algorithm -

Iterative
its process continues until no more info is available to be exchanged between neighbours

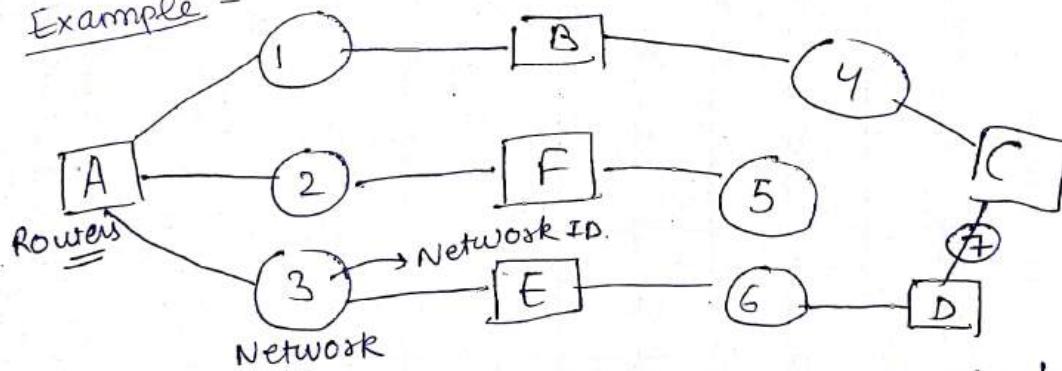
Asynchronous
doesn't require all of its nodes operate in lock step with each other.

Distributed
(each node receive Kargi info from one or more directly attached neighbours . calculation perform Kargi or result back to Neighbours)

- Each router maintains a distance table → vector algorithm -

- (1) A router transmits its distance vector to each neighbors in packet .
- (2) Router receives & saves the most recently received distance vector from each of its neighbour .
- (3) Router recalculate distance vector when
 - (1) it receives a distance vector from neighbour contain diff. info .
 - (2) discover that link to a neighbour has gone down .

Example -



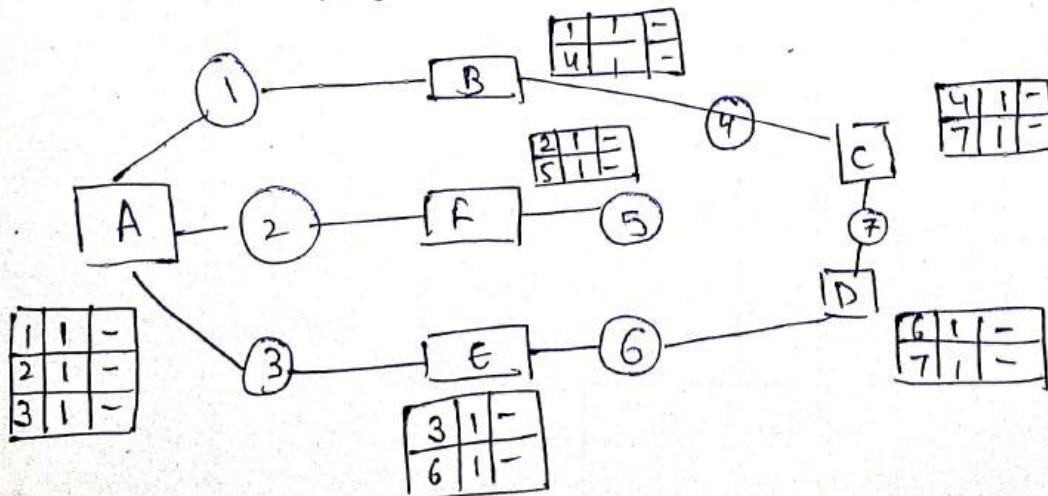
A → BFE
B → AC

C → BD
F → A
E → AD

D → EC

Routing table

| ID | Cost | Next Hop |
|----|------|----------|
| | | |



Updating table

| | | |
|---|---|--|
| 1 | 1 | |
| 4 | 1 | |

+

Tone loop

=

| | | |
|---|---|---|
| 1 | 2 | B |
| 4 | 2 | B |

| | | |
|---|---|---|
| 1 | 1 | - |
| 2 | 1 | - |
| 3 | 1 | - |

| | | |
|---|---|---|
| 1 | 2 | B |
| 4 | 2 | B |

→

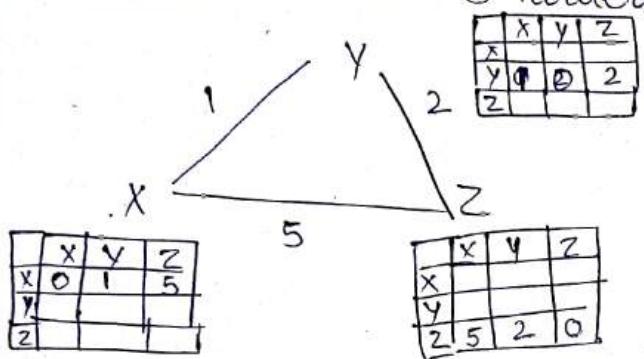
| | | |
|---|---|---|
| 1 | 1 | - |
| 1 | 2 | B |
| 2 | 1 | - |
| 3 | 1 | - |
| 4 | 2 | B |

duplicate data
{sirf whi Rakhega jiski cost kam hogi}

| | | |
|---|---|---|
| 1 | 1 | - |
| 2 | 1 | - |
| 3 | 1 | - |
| 4 | 2 | B |

For EXP

consider 3 Routers



Exchange Routing Table

X will get Z & Y Routing table

So, new Routing table will be like this -

| X | X | Y | Z |
|---|----------|---|----------|
| X | 0 | 1 | infinity |
| Y | 1 | 0 | 2 |
| Z | infinity | 2 | 0 |

now as we can see when

Y is intermediate node b/w

X & Z it have low cost ($3 < 5$)

| | X | Y | Z |
|---|---|---|---|
| X | 0 | 1 | 3 |
| Y | 1 | 0 | 2 |
| Z | 3 | 2 | 0 |

Same steps for Y & Z.

| | | | |
|-----|---|---|---|
| Y = | X | Y | Z |
| | 0 | 1 | 3 |
| | 1 | 0 | 2 |
| | 3 | 2 | 0 |

| | | | |
|-----|---|---|---|
| Z = | X | Y | Z |
| | 0 | 1 | 3 |
| | 1 | 0 | 2 |
| | 3 | 2 | 0 |

circuit
 QO
 Conn
 Ad
 Cef
 R
 F

Link state

Technique in which each router shares the knowledge of neighbourhood with every other router in internetwork.

Phases

Reliable flooding

Initial State

(Each node knows cost of its neighbours)

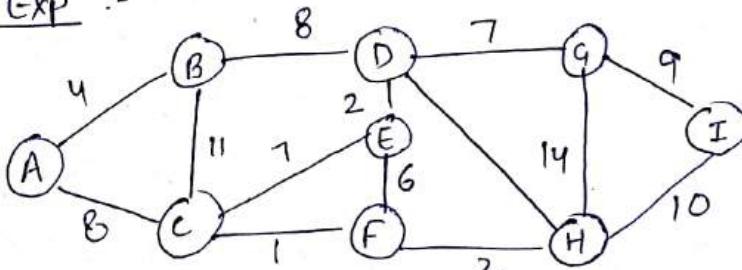
Final state

(Knows entire graph)

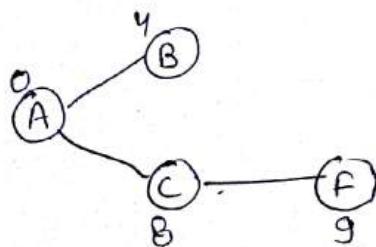
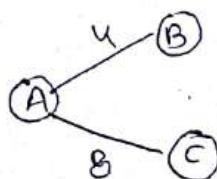
Route calc.

↳ Dijkstra's algo
 (used to find shortest path from one to every other node)
 ↳ iterative property hoti hai after kth iteration of algorithm the least cost path are well known for K destination nodes.

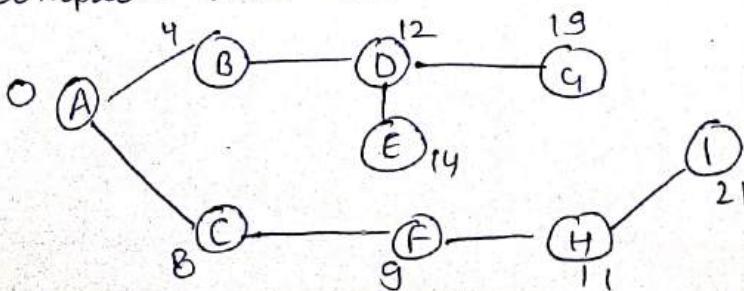
For Exp :-



Let's start with A vertex,



Complete will be



Circuit Flooding -

QC
con

communication method where data packets are sent to all connected devices. Even if they are not intended recipient of data.

Why flooding

1. Address Resolution - translate IP to actual MAC.
2. Routing protocol - Link state Routing
3. Multicast tree Establishment - Create Multicast Distribution tree
4. Emergency Broadcast - regardless of location make sure that message reach all network device.

Features

1. Broadcasting - send to every device connected to network
2. No prior knowledge - precise address of device optional
3. Simple & Robust - straightforward, simple doesn't rely on intricate routing algorithms
4. Loop prevention - use Time to live counters to stop data packets
5. Security consideration - expose private info to recipient who weren't intended

Broadcasting -

- allow message to be received by all nodes of network
- transmission of signals from Radio or television.

Broadcasting

unicast

ek source info send
karta hai ek hi destination
Ro.

Ex → Phone call

multicast

multiple destination ko
hi ek hi time

Ex → Radio



circ

QO

con

Ad

EF

R

T

..

Flooding

Broadcasting

- Packets ko har node tak bhejna, Bina kisi specific Routing ke.

ER source se multiple recipients tk info bhejna

- redundancy avoid ke lie

info delivery ke lie

- limited to network infrastru.

wide Range

- minimal control, Packets har node tak pauchte hai

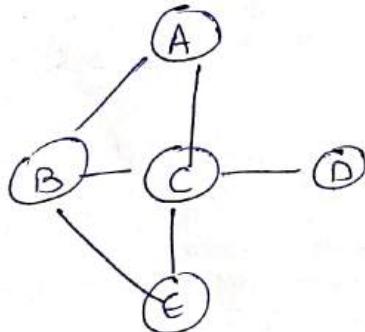
controlled, specific to chosen medium/audience

- network congestion & collision

efficiently reach large audience

Flow Based Routing -

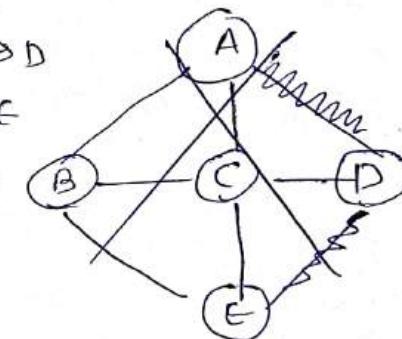
- ER gaetwork routing technique hai jisme network traffic ko different flows mei catagorize karke route karte hai.
- har individual packet ko aay se route kia jata hai.



$$(A \rightarrow D) = A \rightarrow C \rightarrow D$$

$$(B \rightarrow E) = B \rightarrow C \rightarrow E$$

$$(A \rightarrow E) = A \rightarrow B \rightarrow F$$



Multicast Routing -

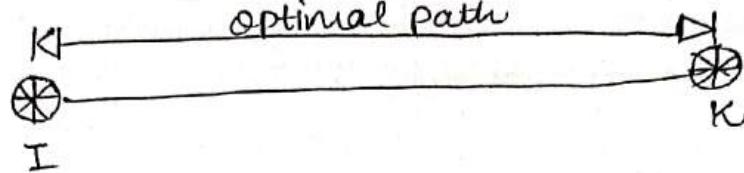
↳ ER single source se multiple destinations tk data packets efficiently bheje jate hai.

↳ sirf unko bhjti hai jo actively participate kar sake hai.

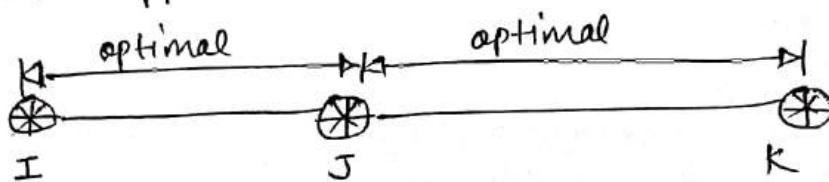
↳ Bandwidth converse karti hai.

Optimality principle

Statement - if the router J is on the optimal path from I to Router K, then optimal path from J to K also falls along same route.



Now suppose J lies between I & K



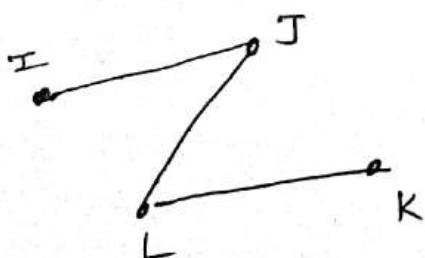
Importance -

- * Efficiency → avoid wasting resources by exploring unnecessary paths.
- * Loop prevention - Routing algorithm can prevent creating loops in network.

Limitations -

- * Dynamic networks - assumes all routers have complete & up-to-date info. about network.
- * cost metrics - optimal path kisi ke lie kuch bhi ho skta hai
least cost, fastest, shortest or combination of factors.

For EXP . Let IJK, L are the routers & connection between them.



Let suppose $I \rightarrow J \rightarrow L \rightarrow K$ is optimal path from I to K then $J \rightarrow L \rightarrow K$ is the optimal path from J to K.

Flow control -

- ↳ technique used to regulate the flow of Data b/w Different nodes in a network
- ↳ Ensure Karta hai Sender ER dm se overwhelm na Karo, (Matlab bolt Sara Data ER sath Send na Karo)
- ↳ goal : to prevent buffer overflow , lead to dropped packets & poor network performance.

Congestion control -

(GTS - GTS)

- ↳ technique to prevent congestion in a network
- ↳ yeh occur hota hai jab bolt Sara Data ER sath bheja jata hai. network pe.
- ↳ lead to ~~pro~~ dropped packets & poor network performance.

Similarities

- (1) Regulate the flow of Data.
- (2) prevent packet loss.
- (3) improve network Efficiency.

Flow

- (1) Traffic from Sender to Receiver is controlled
- (2) used in DDL , transport
- (3) Keeps the data from overwhelming
- (4) Sender sends data slowly to Receiver

congestion

- Sender has to control/modulate own rate
- used in network & transport
- keep networks from being congested
- Slowly transmitted into network by transport layer .

| | Datagram | virtual circuit network |
|---------------------------|---|---|
| circuit step | Not needed | Required |
| QoS | Difficult | Easy if enough resources |
| congestion control | Difficult | Easy if enough resources |
| Address | Each pack contain the full source & Destination address | contains a short vc number |
| Effect of Router failures | None except for Packets lost during crash | VC's passed through failed network are terminated |
| Routing | Independent | choose vc, all packets follow it |

Internetworking = inter + networking = connections b/w two completely
 → isme switches, gateways jaise devices use hone hai.
 → jo data ko ek network se dusre network ke pauchate hai.

① Internet - network ko ek satth jodta hai jaise private, edu. organization etc.

② Intranet - private network ek organization ke internal use ke lie hota hai. Ex - employees, Department

③ Extranet - extension of intranet (external parties se connect karta hai) limited access

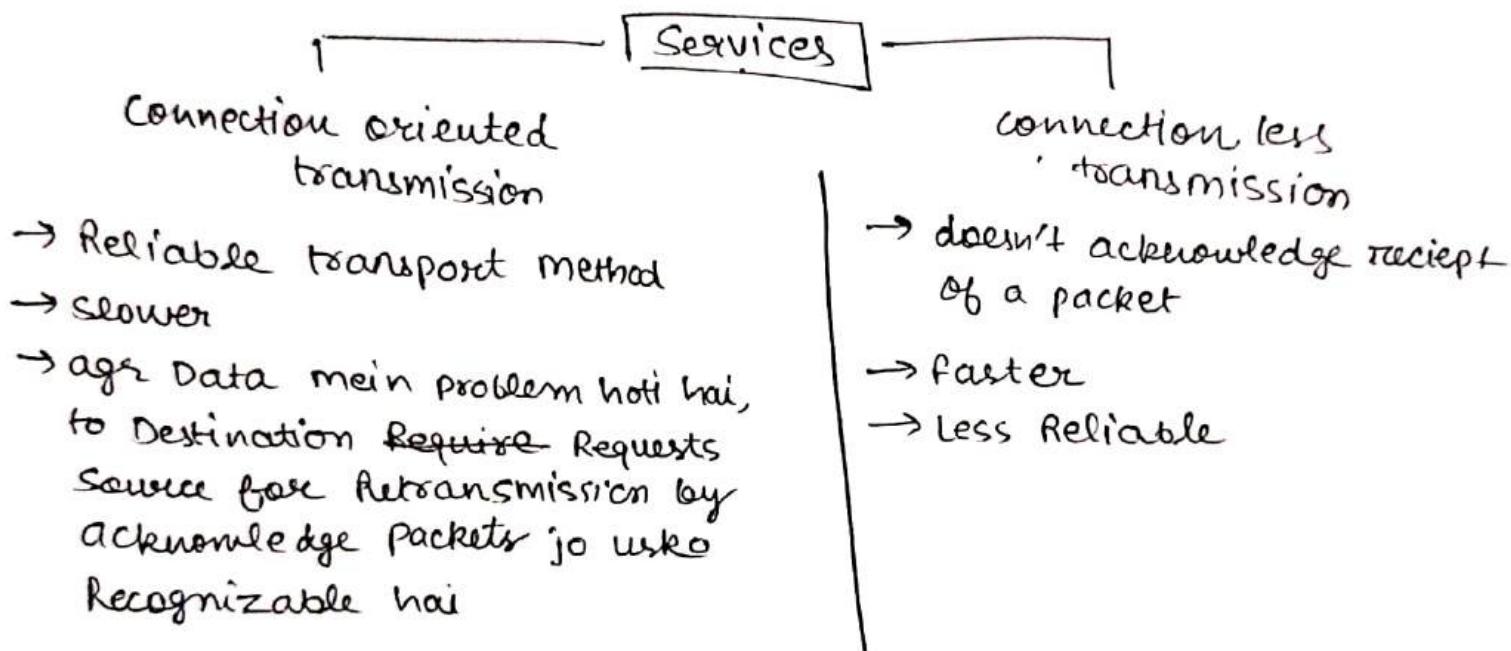
④ VPN → privacy & security

⑤ P2P → file aur resource sharing

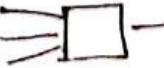
Unit-4

Transport layer - (End to End layer)

- ↳ Ensure jo packets hai arrive hao order mei.
- ↳ Provide acknowledgement of successful data transmission
 - ↳ " " ob retransmits data agr error paya jata h.
- ↳ Ensure Karta hai jo data hoga vo Error free hoga with no losses or Duplications.
- ↳ Service provide → to application & take from network
- ↳ Divides message Received from upper layer into packets → source or destination Resemble Karti hai again into message



functions of Transport Layer -

1. segmentation of Messages into packet and Resembly of Packets into Message.
2. Message acknowledgement - Reliable End to End Message delivery
3. Session multiplexing -  - ER mei convert kar deti hai data packets to
4. protocols - TCP, ATP, NWLINK
5. flow control - Ensure the Rate they of communication they both can handle.

Transport Layer Design issues -

1. Reliability - Data transmit between host is delivered reliable. TCP Provides but UDP doesn't provide
2. Flow control - Process of Regulating the flow b/w two network nodes. It Results in better network utilization by avoiding packet loss
3. Congestion - Preventing network from becoming ~~congested~~^{traffic}. TCP uses effective congestion control to prevent packet loss.
4. Multiplexing & Demultiplexing -


(DLD wala)
5. Connection establishment & termination - Before data transfer, ~~ER~~ connection establish hata hai jo after completion terminate ho jata hai.
TCP uses - 3 way handshake to establish
4 way handshake to terminate
6. Quality of service - Ensures to provide an acceptable level of QoS for traffic.
minimum Bandwidth, max delay

Quality of Service (QoS)

↳ critical application require Reliability & timely data delivery.

↳ something a flow seeks to attain.

Flow characteristics

Reliability

Means loosing a packet or acknowledgement entails retransmission.

Delay

Applications can tolerate delays in different degrees.

Jitter

Variation in delay for packets belonging to same flow
High → delay large
Low - small

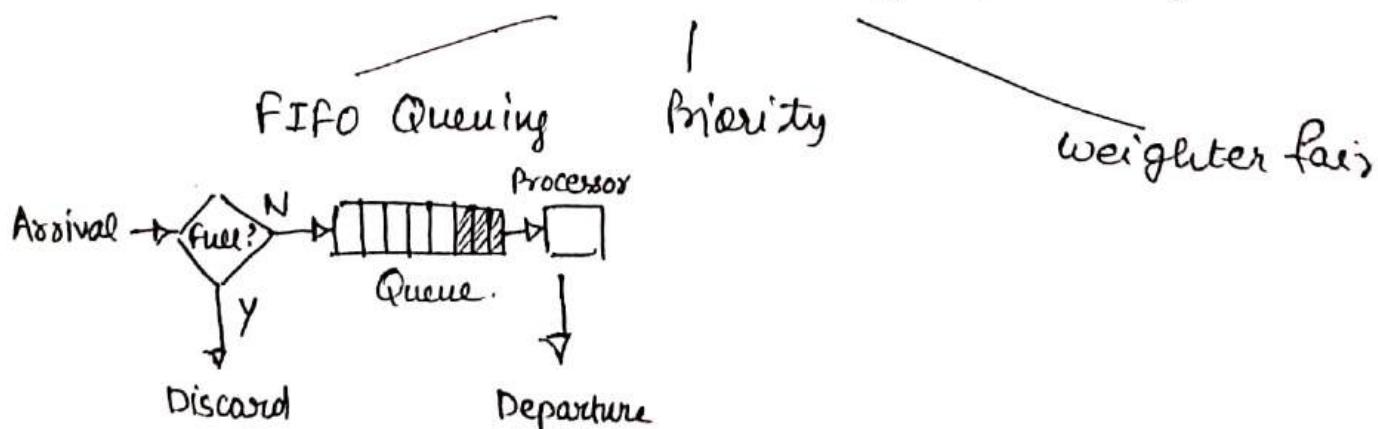
Bandwidth

Different application
Different bandwidth.

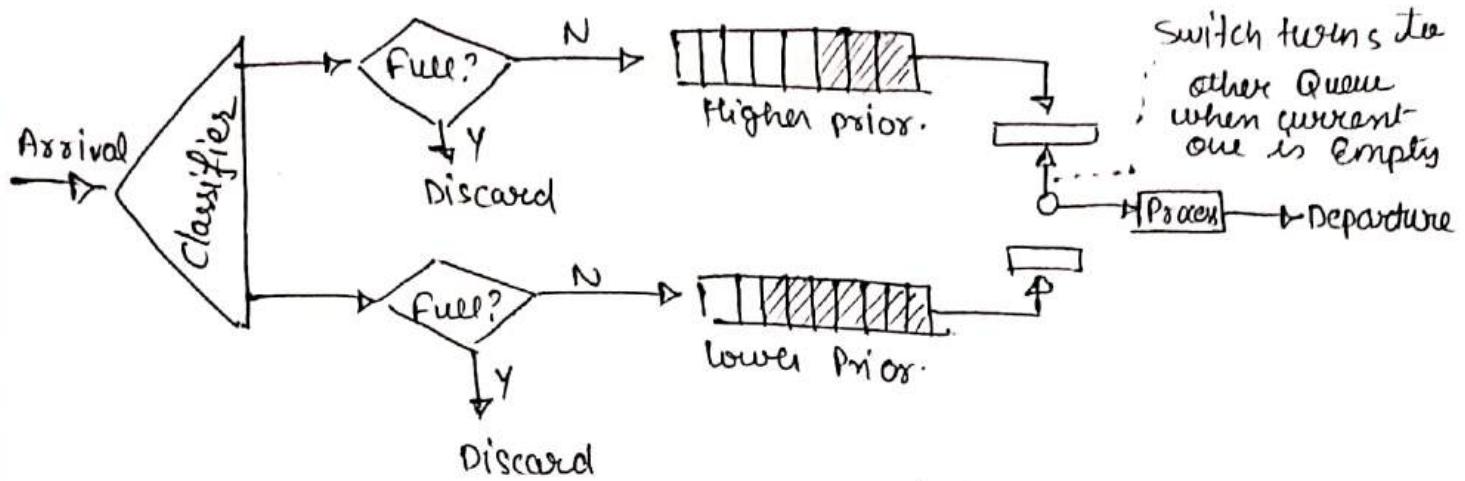
Techniques to Improve QoS -

1. Scheduling
2. Traffic shaping
3. Resource Reservation
4. Admission control

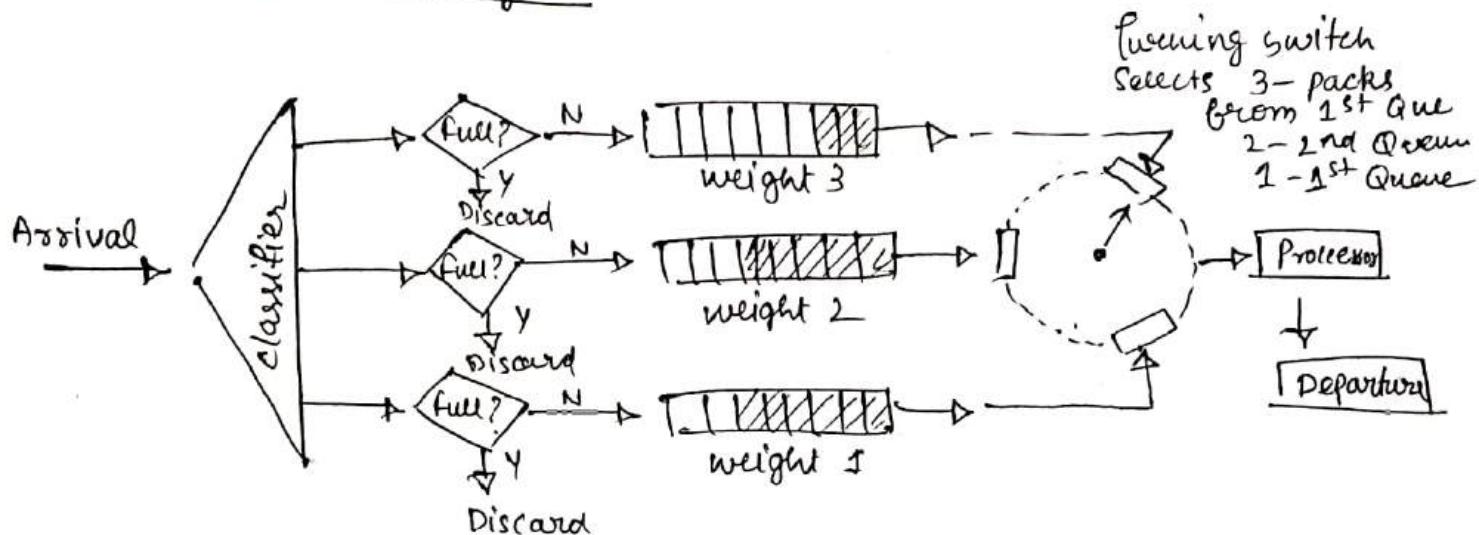
1. Scheduling - Packets from different flows arrive at a switch or router for processing



Priority Queue



Weighted Fair Queuing



2. Traffic Shaping - Traffic is shaped before it enters network.
Controls rate at which packets are sent.

algorithms

Leaky Bucket

- ↳ constant output data rate
- ↳ If buffer overflows then Packets Discarded
- ↳ Results in uniform flow of Packets.
- ↳ when Packets → same size one Packet Per tick is okay.
- ↳ for variable, allow fixed no. of bytes/tick

Token Bucket

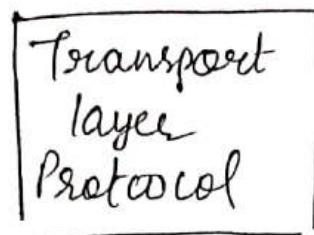
- ↳ allows output rate to vary
- ↳ bucket holds token
- ↳ for one packet host must capture & destroy one token
- ↳ token generated at rate of one token every At sec

Resource Reservation

↳ flow data needs resources such as buffer, bandwidth, CPU etc.

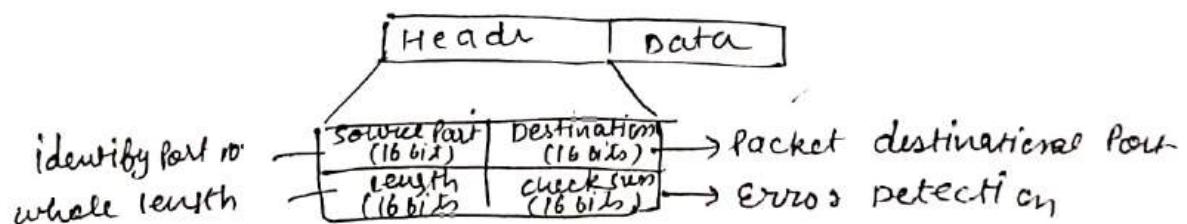
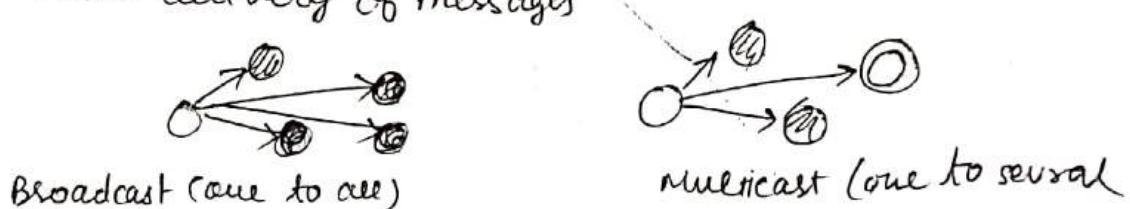
Admission control -

↳ a mechanism used by Router, switch, accept, reject a flow based on predefined parameters.



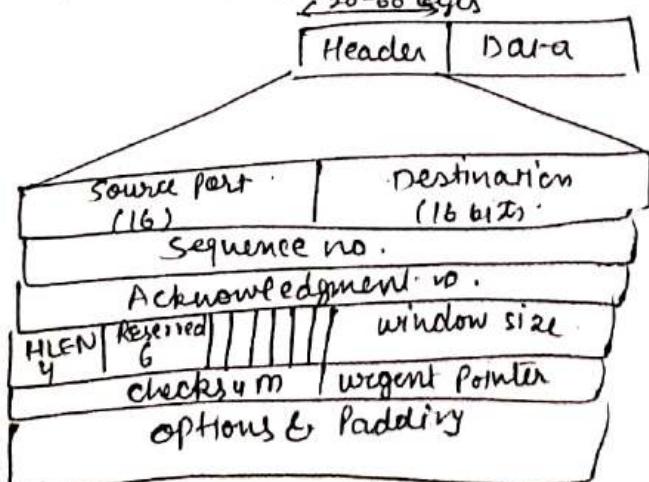
1. UDP - User Datagram protocol.

- non sequential transmission
- connectionless (speed & size is imp)
- adds checksum, error control
- faster delivery of messages



2. TCP - Transmission control protocol.

connection oriented
virtual circuits connect the sender & receiver
20-60 bytes



SCTP

- ↳ Stream control transmission protocol.
- ↳ connection oriented.
- ↳ full duplex mode
- ↳ Reliable & Fairer

IPv4 address

- 32 bit address uniquely & universally defines connection of device (a computer/router) to internet
- unique one address defines only one connection to internet.
- Two devices can never have same address at same time.
- ER address thode time ke lie ER device ko detek or thode time Baad usse vapis leke kisi or node skte hai
- age device operate kar rha hai network layer pe "n connections ke sath, to use pass n addresses hone chahiye"
- address space - total no. of addresses used by protocol.
- age protocol N bits use kar rha hai to define address. then address space = 2^N because each bit can two different values (0, 1) & N bits can have 2^N values

IPv4 notations

1. Binary

32 bit/4 byte
Octet \rightarrow byte

Ex: 0110101 10010101

00011101 00000010

2. Dotted Decimal

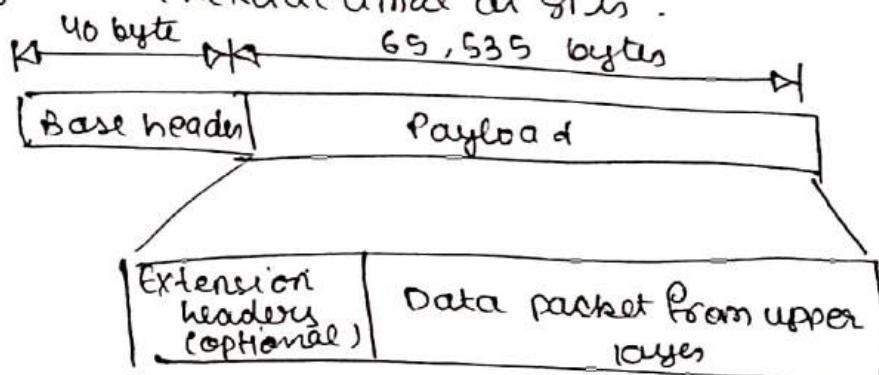
Ex: 117.149.29.2



- IPv4 host to host communication karata tha
- well designed
- deficiencies → unsuitable in fast growing
 - Subnetting, class addressing & NAT address depletion long term problem in internet
 - Internet must accommodate Real time audio & video transmission which requires minimum delays, Strategies & Reservation of resources (jo ki IPv4 mei nahi tha)
 - No encryption & authentication.

in Kamyon ko poora karne ke lie IPv6 proposed kia gya.

- Extensively modified to accomodate unforeseen growth of internet.
- address consists of 16 bytes (128 bits long)
- To make address readable, specifies hexadecimal colon
- 128 bits divid into 8 Sections, 2 byte Each.
- 2 bytes = 4 hexadecimal digits.



IPv4

- # Provides 32 bit address
- # Security X
- # No Protocol Enhancement
- # divided into 5 classes
- # Can be converted into IPv6.
- # header can be from 20-60 bytes
- # Checksum field
- # Representation in Decimal

IPv6

- 128 bits -
- authentication, integrity & confidentiality ✓
- Features Hierarchical addressing
- doesn't have any
- can't be to IPv4
- 40 bytes
- not .
- hexadecimal.

SESSION LAYER

| | |
|------|---------------------------------|
| Data | session Interhost communication |
|------|---------------------------------|

- Provides Reliable & Secure communication b/w two devices by Establishing managing & terminating sessions.
- Regulates Data flow, defines format of data sent to connections.
- manage - Kaise data send kar skta hai, in a certain amount of time & for how long.
- Reconnect Session if disconnects.
- Protocols - NetBIOS, Mail slots, Names pipes & RPC

functions -

1. Session Establishment - Establish connection b/w two devices before transmission begins.
2. Session management - Keep track off session throughout its duration. In case of error session term.
3. Session termination - After completion.
4. Session Security - Encryption, authentication & authorization.
5. Session Recovery - if connection lost or interrupt. Keep tracks.
6. Dialog management - jab device connect. hota hai to session layer responsible hoti hai for determining konsa device communication mei part le skta hai as well as control the amount of data that can be transmitted.
Types of dialog control
 - Simplex
 - half duplex
 - full duplex.

F. Synchronization - Handles synchronization between incoming & outgoing data stream by adding synchronization points → checkpoints
Jiski help se session layer Retransmitting mei help Kar pati hai easily & fastly

Quality of Service (QoS)

↳ Ability of network to prioritize & deliver data based on importance & ensure certain level of performance.

Achieved through -

1. Traffic prioritization - Ensures high priority data send first. Achieved by setting different priorities for each session Based on importance of Data.
2. Bandwidth management - monitors amount of data being transmitted & allocating bandwidth on priority.
3. Congestion control - manage rate & amount of data being transmitted which prevents Network overload, reduce packet loss.
4. Error handling & Recovery - Detects & correct Errors If error paya jata hai to SL Action keli hai.

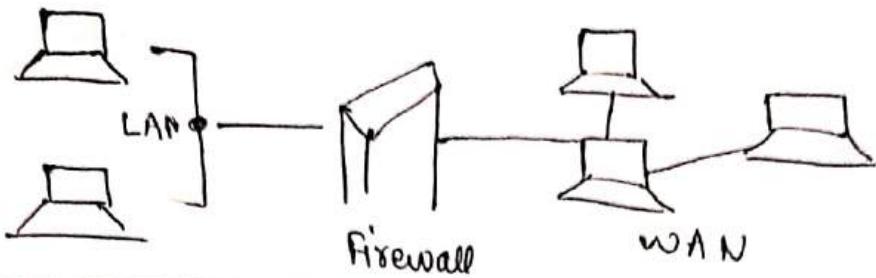
Firewalls

- Security devices monitor & control network traffic based on predetermined security policies.
- Helps to prevent unauthorized access & malicious attacks and ensure only authorized sessions allowed.

Accept → allow the traffic

Reject → Block "unreachable error"

Drop → Block traffic with no Reply



Separate private from open internet

firewall

Packet

- Operate inline at junction points where devices → Routers, switches do their work.
- Do not route the packet rather compare every received packet with established criteria
Ex → Port no, ip address etc.
- Flagged packets are dropped

- circuit
- Operate to monitor & control individual session
 - Establish circuit between ~~two~~ communicating devices, to monitor control of flow
 - Session ID, source or destination ke basis pe filter karta hai.

Advantages of firewall

- 1) Security
- 2) Prevention
- 3) Control of network access
- 4) Regulation compliance
- 5) Monitoring of network activity

Disadvantages

- 1) Complexity
- 2) Limited visibility
- 3) Cost
- 4) Limited VPN support

APPLICATION LAYER

- ↳ topmost layer of OSI model & enable communication b/w application on different hosts
- ↳ Designed to communicate with specific application
- ↳ Establish, manage, terminate communication sessions & enable data exchange for specific application.

Protocols -

1. HTTP (Hyper text transfer protocol) -

Designed for world wide web.

Used to transfer hypertext document & other data b/w web & clients.

2. FTP (File transfer protocol) -

Used to transfer files over network.

Enable sharing & copying of files b/w computers located on different networks.

3. SMTP (simple mail transfer protocol)

Used to send email message from one server to another.

4. SNMP (simple network management protocol)

Used for network management & monitoring

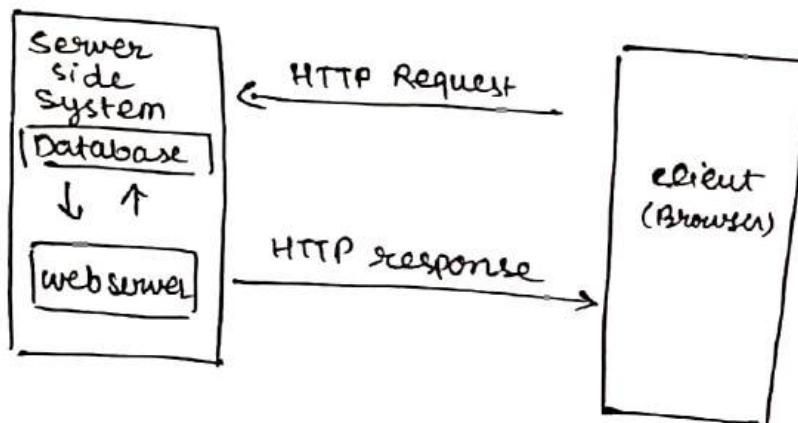
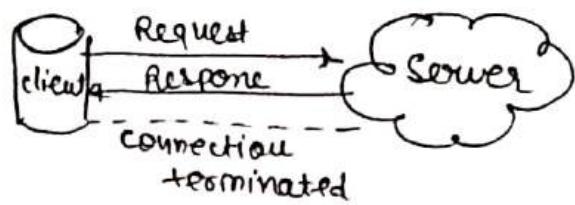
Enables network devices to be monitored, managed & controlled remotely.

5. DNS (Domain Name system)

Maps domain names to IP address & help translate human readable domain names to computer readable IP address.

HTTP

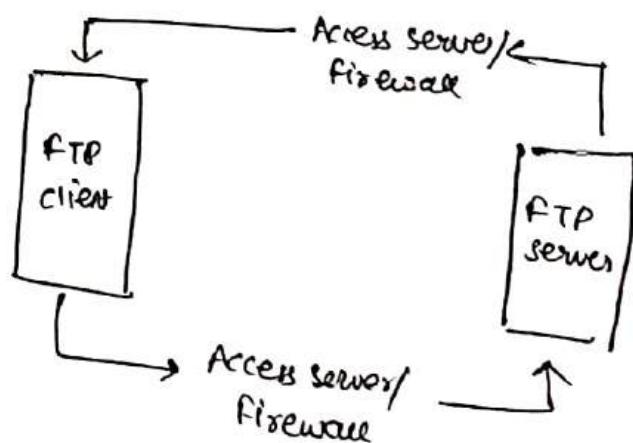
- 1) The URL sent to DNS
 ↓
Check records for URL in Database
 ↓
Return IP address
 ↓
Browser send Req to server
 ↓
Server send data,
connection closed



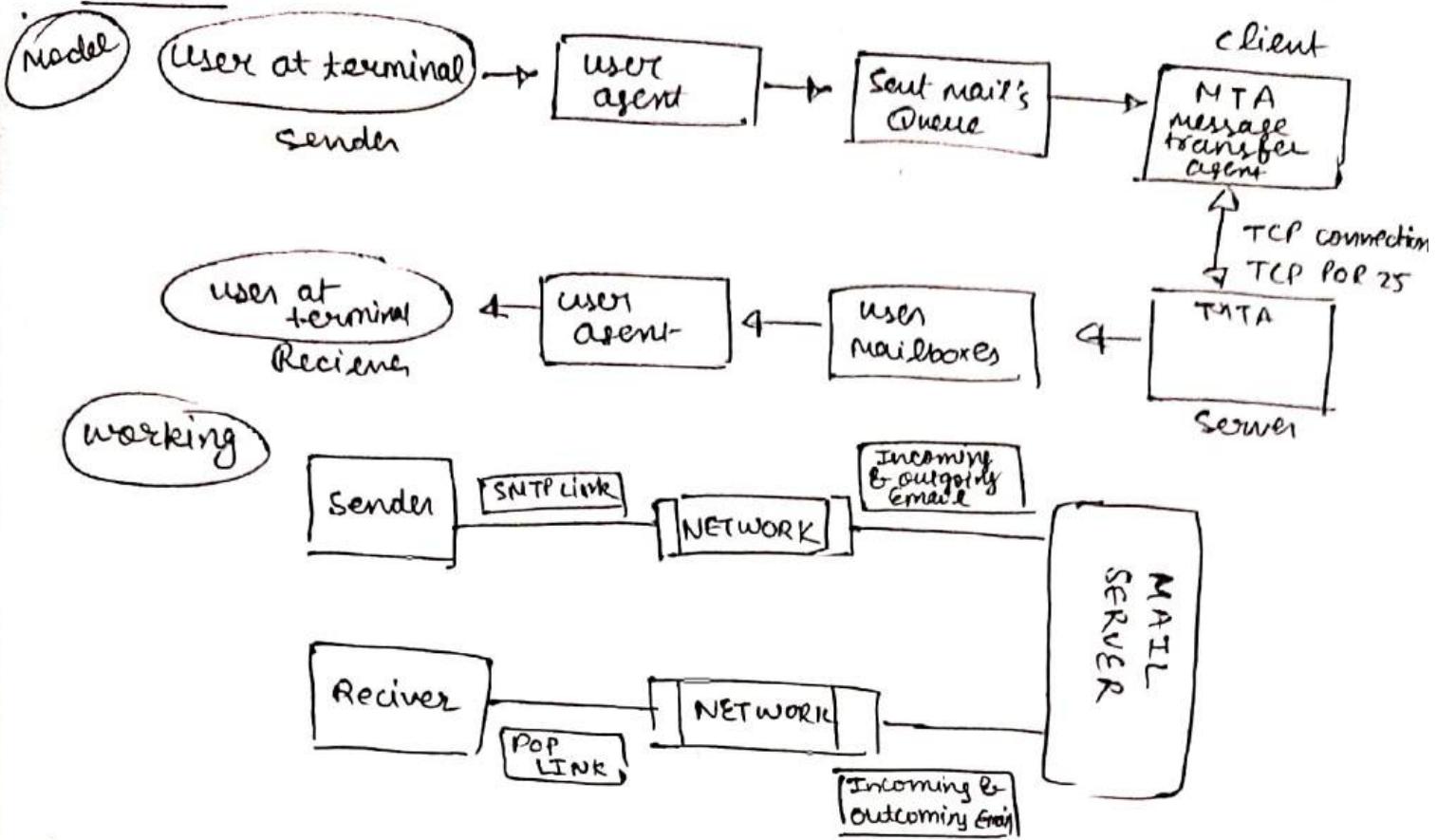
FTP (file transfer protocol)

- Encourages direct use of remote computers.
- Promotes sharing of files or other types of data

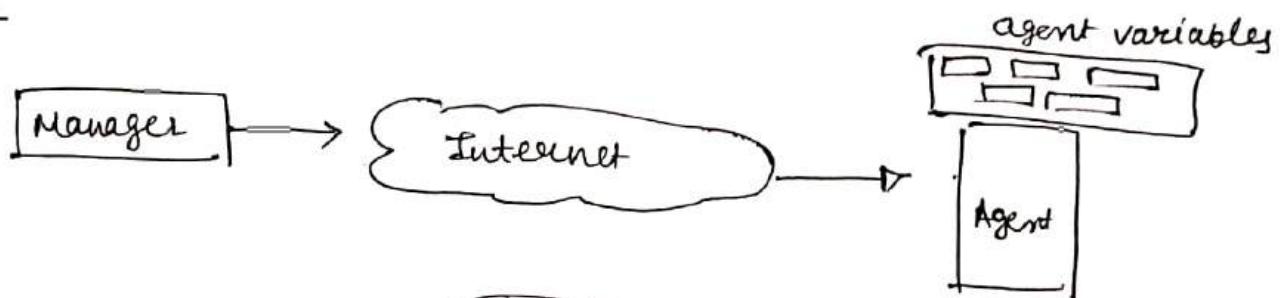
- Client contacts FTP server
 ↓
Obtain authorization
 ↓
Browse Remote directory
 ↓
Server Receives a command
 ↓
after transferring sever close connection



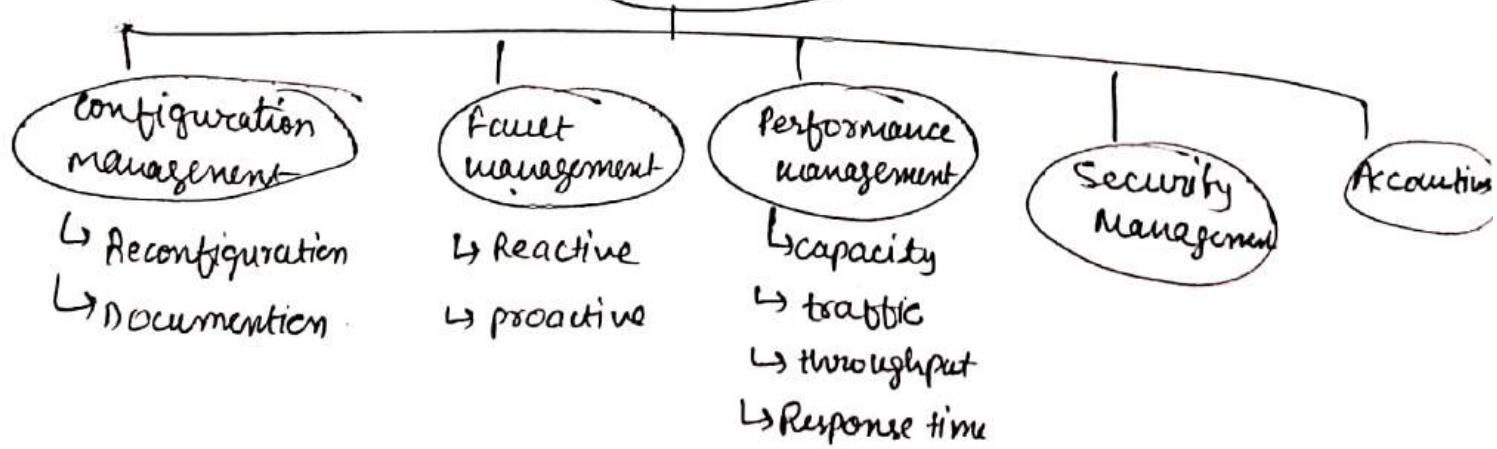
SMTP



SNMP



functions



CNIP

IPv4 ?



Internet protocol version 4 → primary version
brought into action for production ARPANET in 1983

32-bit addresses are 32-bit integers, Expressed in Decimal
notation.

Parts of IPv4 .

1) Network part: distinctive variety appointed to network
conjointly identifies category of network & assigned .

2) Host part - uniquely identifies the machine on network .
assigned to every host

3) Subnet number - non-obligatory part of IPv4

local networks that have massive no. of host are
divided into subnets and subnet no's are appointed to it .

Characteristic -

- 1) 32-Bit IP address
- 2) numeric, & Separated by a dot .
- 3) header fields → twelve ~~no. of header~~
no. of header field = 20
- 4) unicast, broadcast & multicast style of address .
- 5) support virtual length subnet mask

Advantage

- permits encryption .
- Network allocation significant .
- Easy to attach to multiple devices
- Quality service
- redefined & permit flawless coding .

Limitations

- Each network unique IP address .
- complex .

- IPV6 - most common version of Internet protocol.
 - ↳ Internet protocol version 6.
 - ↳ Designed by Internet Engineering task force (IETF)

Types of IPV6 address.

Unicast

only one interface is specified by unicast.

Packet moves from one host to destination host

Multicast

represent a group of IP devices & can only be used as destination of datagram

Anycast

same as multicast
anycast address varies from other addresses
it can deliver the same IP address to several servers or devices.

Advantages

- Routing efficiency
- Reliability
- allocates address on its own.
- Security .
- faster
- stronger
- Simple aggregation

Disadvantages

- conversion (long period .
- communication .
- cross protocol

IPV6

128 bit address
hexadecimal
header → 40 bytes
checksum field not available
doesn't support VLSM

IPV4

32-bit address
Decimal → Add. representation
header → 20-60 Bytes .
available
support VLSM.

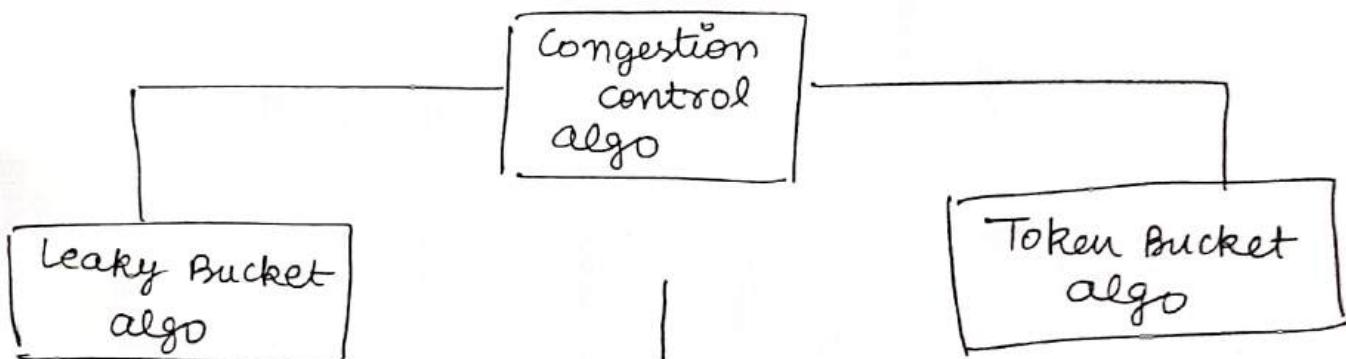
Congestion control

Effects

- ↳ happen because of availability of many packets in network.
- ↳ decreases network performance.
- ↳ Packet delivery to receiver is delayed
(Packet loss)

Congestion control algo -

- ↳ mechanism which control the entry of data packets into network, enable a better use of shared network infrastructure & avoid congestive collapse.



- used for shaping traffic
- used to control the rate at which traffic is sent to network, shape the burst traffic

Algorithm -

Each network interface ke Pass leaky Bucket hoti hai

1. jab host packet send karta hai, to packet throw karte hai Bucket mei
2. Bucket leak karta hai constant rate pe, meaning → network transmit packets at constant rate
3. Bursty traffic → uniform.
4. Bucket → finite que bn jayega jo output dega finite rate pe.

leaky Bucket algo Kabhi-kabhi Disadvantage De deti hai, because agr Kabhi boht jyada value mei Packets aae in the Network tb bhi Ek particular value mein flow jayega jiski vjh se problem hogi.

is algo mei bucket ~~se~~ tokens contain Karta hai. Each token define Karta hai packet jo ki Ek predetermined size ka hota hai. Token delete kie jate hai buckets mein se arc to ability to share Packer

algo steps -

1. Regular interval mei throw Karenge packet.
2. Bucket ki Max Capacity hogi.
3. agr packet Ready to sent hai, to token Remove Karenge Bucket se or Packet Send Karenge.
4. agr Bucket mei koi ~~packet~~ token nahi hai to ~~token~~ 'packet' send nahi hoga.



Flow Control -

- ↳ A technique that gives permission to two stations (Sender & Receiver) that are working to just communicate to one another.
- ↳ It restricts & coordinates number of frames or amount of Data sender can send. just before it waits for an acknowledgement.
- ↳ Set of procedures. jo Explain Karta hai, Sender Kitna the frames of Data transfer or transmit Kar skte before Data overwhelms receiver.

Approaches to flow control

Feedback - Based

Sender transmit Data or information to receiver, Receiver transmits data back to Sender. Sender & also allow Karta hai. Sender ko transmit more Data

→ Sender transmit Karta hai Data or Receiver acknowledge Karta hai on Receiving Data

(easy & simple)

Stop & wait

Message / Data broken down into multiple frames. Receiver indicates readiness to receive frame of data. Ack receive holi hai tabhi, sender next frame send Karta hai. Yeh process continue raha hai until Sender transmit (EOT) End of transmission.

Rate - Based

Sender sends or transfer Karta hai Data at faster speed or Receiver able nahi hota Receive Karne ke Data ko us speed pe to mechanism known as built in mechanism in protocol transmit or restrict Karta hai overall Rate ko. jispe information transfer or transmit hoti hai by Sender

Techniques to control Flow in DLL

(efficient)

Sliding window

Reliable in order delivery of packets. None of other entity tries to communicate until current data or frame gets transferred. Sender send various packets before receiving acknowledgement. Sender & Receiver both agree Karta hai, total no. of data frames transfer.isme Sender to bolte se send Kar skta pr Receiver Ek hi receive karaga.

Firewall -

- ↳ Network Security Device
- ↳ Monitors both incoming & outgoing traffic using Predefined
↳ Protective Barrier. set of Security
[Set of security Rules]

- Accept → allow the traffic
- Reject → Block the traffic but reply with an "unreachable error"
- Drop → Block the traffic with no reply.

[working] -

1. Packet Inspection → when Data packets travel b/w networks, Firewall Examine Karta hai source. & other attributes to determine Security criteria meet kar she hai ya nhi
2. Filtering - compares the packet info with set of Rules agr match ho jata hai, allow to Pass otherwise Reject.
3. Access control - Block at unauthorized access attempt
4. Network access translation - hides internal IP address of devices, add extra layer of security by making it difficult for attackers to directly target.
5. Logging & monitoring - Keeps a record of all Network traffic & events. Generate logs & alerts to notify administrators about potential security threats.

HTTP

- ↳ used for communication b/w web browser & web server.
 - ↳ allows transfer of hypertext documents, such as webpage.
- (working) → type a website url, it send an Http request
server process Karta hai Request ko & send
back an Http response jo ki contain Karta hai
requested web, or wo display ho jata hai

SMTP → Simple mail transfer protocol

- o Protocol used for sending & Receiving Email messages.
 - o allow the transfer emails b/w mail server & internet.
- (working)
 - send an Email, client used Smtp to send message to recipient's mail server.
or use Baad Smpt use nota hai receiver ke dvara to deliver email to recipient mail box.

FTP → File transfer protocol

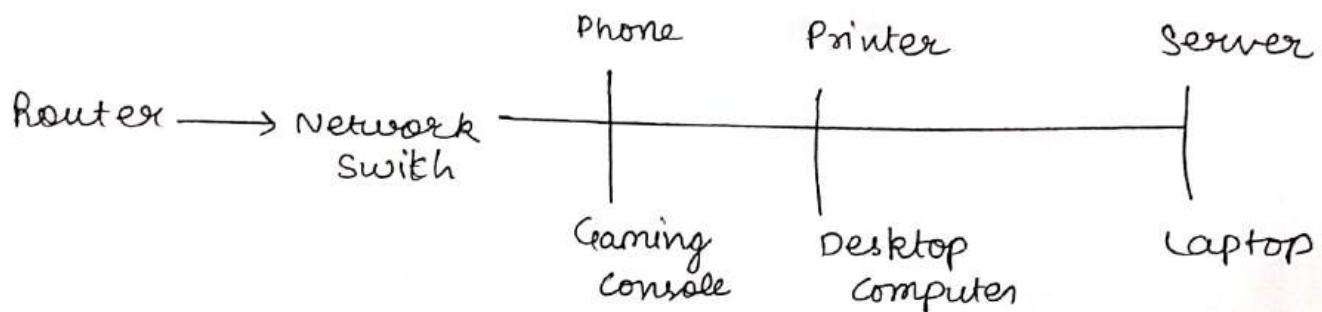
- ↳ used for transferring files.
 - ↳ Enable uploading, downloading, managing files on Remote server.
- (working)
 - connect to Remote Server, authenticate karo khud ko, browse directories → transfer file

Switching -

- ↳ Process of transferring data packets from one device to another using specific devices → Switches.
- ↳ take place at Data link layer of OSI Model.

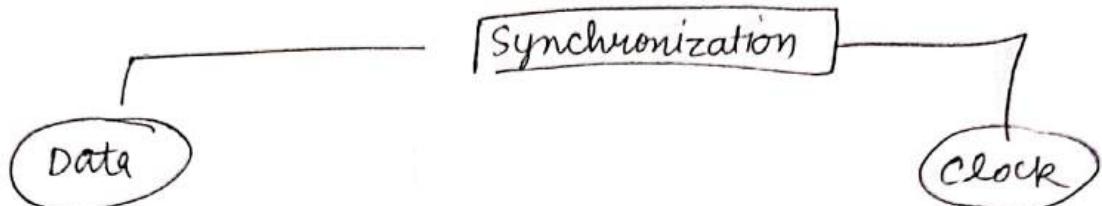
Network switching -

- ↳ Dedicated piece of computer hardware (incoming Data packets & transferring them to destination)



Synchronization

- ↳ coordination of data exchange b/w different devices or systems.
- ↳ ensure data is transmitted & received correctly without errors or collisions.



Keep multiple copies of data consistent across different location.

File
↳ Replication, mirroring & versioning technique

Database
↳ update, insert, deletions
Replication.

Data transfer relies on precise timing. network device use karta hai synchronization clocks to ensure ke data packets arrive ho correct order mein or expected time pe.

↳ Global clock - accurate event logging, transaction sequencing.

↳ Network time protocol distribute time info to clients.

↳ Precision time protocol high accuracy application

Routing algorithms -

- ↳ to determine Best path for data travel from Source to destination
- ↳ Find (shortest) & most efficient paths for data transmission.

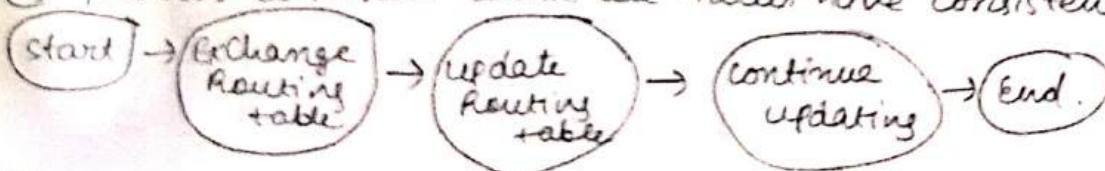
① Shortest path -

1. Start at source node
2. Initialize shortest Distance to all other nodes as infinity.
3. Explore neighboring nodes & update shortest distance if shorter path found
4. Continue Exploring & updating distance until all nodes have been visited
5. Select shortest path



② Distance vector

1. Each node maintains its own routing table which stores shortest distance to other nodes.
2. Nodes exchange their routing tables with neighbouring nodes.
3. If Routing table receives info from another node, each node updates its own Routing table by considering shortest path.
4. Process continues until all nodes have consistent Routing tables.



③ Link state Routing algo -

1. Each node gathers info about its directly connected links, include state & cost.
2. Node info exchanges with other nodes in the network.
3. Receive the link state information from all nodes & construct Network map.
4. Each node calculates shortest path to others using Dijkstra's algorithm.
5. Update Routing table calculate shortest paths.



Design issue →

- ① Reliability → Data is delivered accurately in correct order even if there are errors.
- ② Flow control → Rate of data transmission.
- ③ Congestion control → Peeche jake padha
- ④ Multiplexing → Allow both same application or services to use same network connection sat mein
- ⑤ Error Detection & correction →
- ⑥ Performance → minimize delay, latency & overhead
- ⑦ Connection Establishment & termination -
Managing the setup & teardown of connection b/w sender & receiver.