Best First Search

A* Algorithm

Notion of Heuristics

- Heuristics use domain specific knowledge to estimate the quality or potential of partial solutions
- Examples
 - Manhattan distance heuristic for 8 puzzle

Calculating Cost

$$f(n) = g(n) + h(n)$$

- g(n) Actual cost of traversing from initial state to state n
- h(n) Estimated cost of reaching to the goal from state n

A* Algorithm

- Given: [S, s, O, G, h] where
 - S is the (implicitly specified) set of states
 - s is the start state
 - O is the set of state transition operators each having some cost
 - G is the set of goal states
 - h() is a heuristic function estimating the distance to a goal
- To find:
 - Min. cost of sequence of transactions to the goal state

A* Algorithm

- 1. Initialize: Set OPEN = $\{s\}$, CLOSE = $\{\}$, Set f(s) = h(s), g(s)=0
- 1. Fail:
 - If OPEN = { }, Terminate with Failure
- 2. Select: Select the minimum cost state, n, form OPEN and Save n in CLOSE
- 3. Terminate:
 - If $n \in G$, terminate with SUCCESS

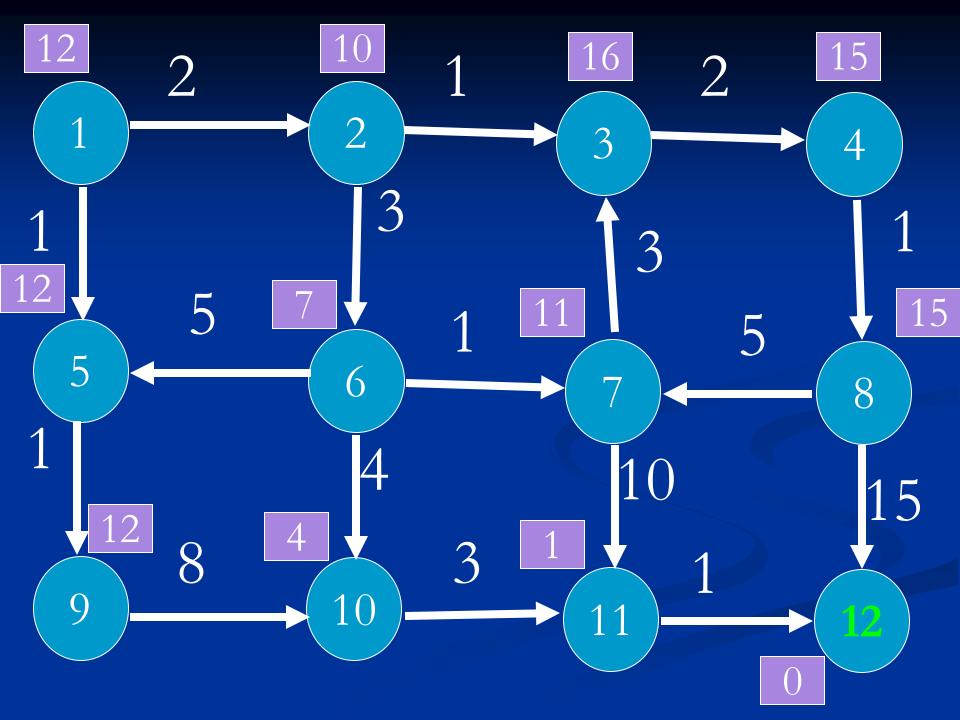
A* Algorithm

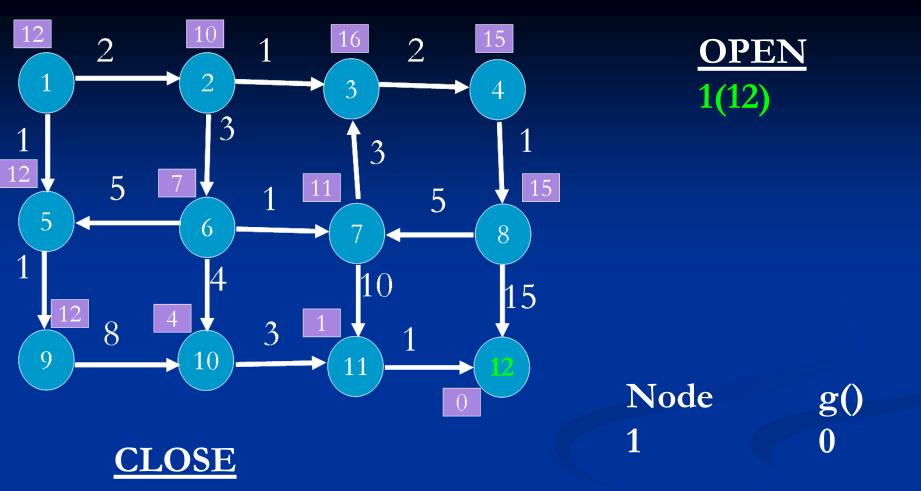
5. Expand:

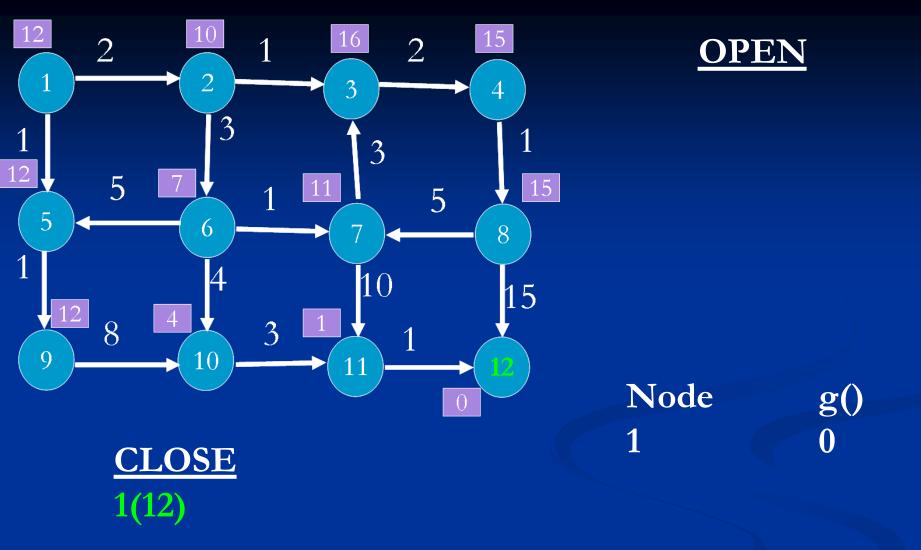
Generate the successors of n using O. For each successor, m, insert m in OPEN only if m ∉ [OPEN ∪ CLOSE] set g(m) = g(n) + C(n,m)set f(m) = g(m) + h(m)insert m in OPEN if $m \in [OPEN \cup CLOSE]$ Set $g(m) = min\{g(m), g(n)+C(m,n)\}$ set f(m) = g(m) + h(m)

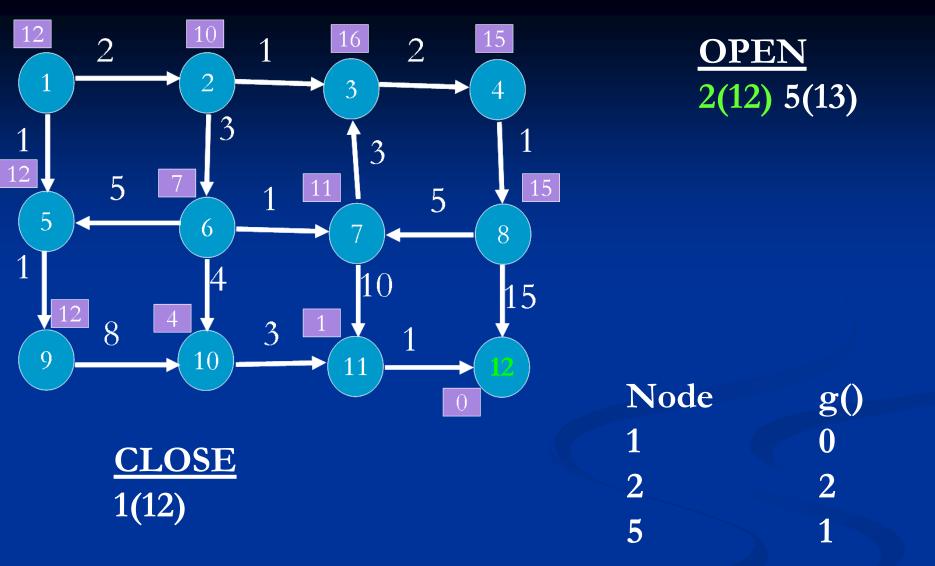
If f(m) has decreased and $m \in CLOSE$ move it to OPEN

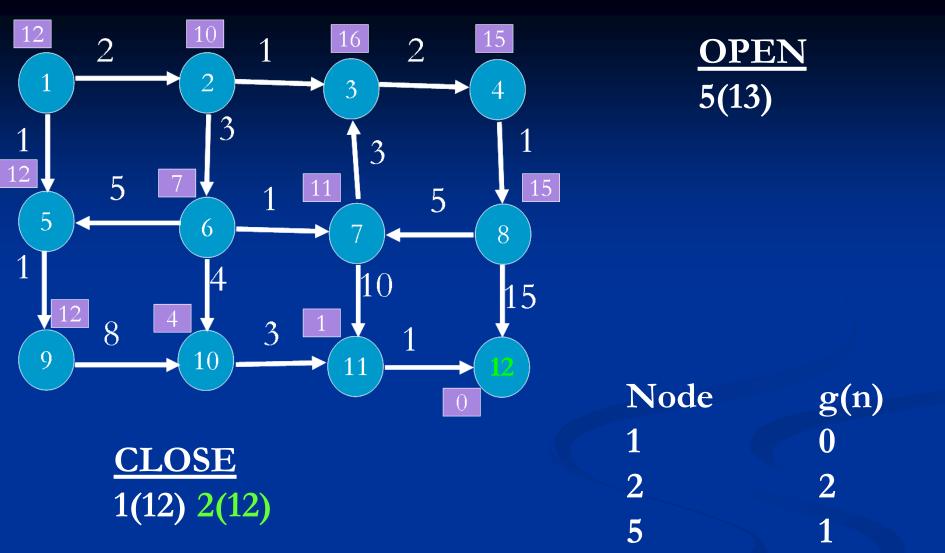
6. Loop: Goto step 2

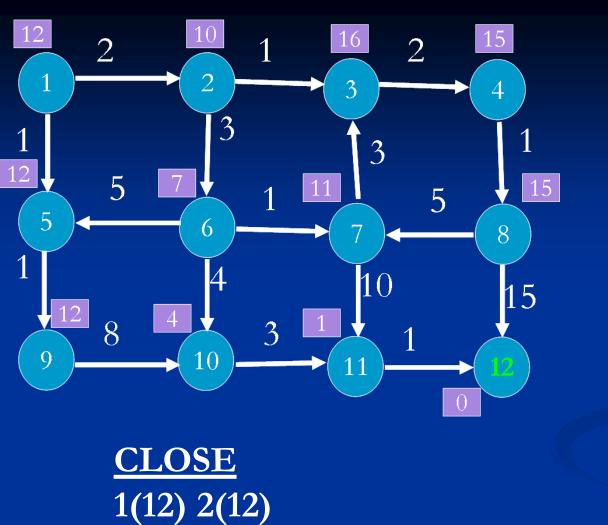






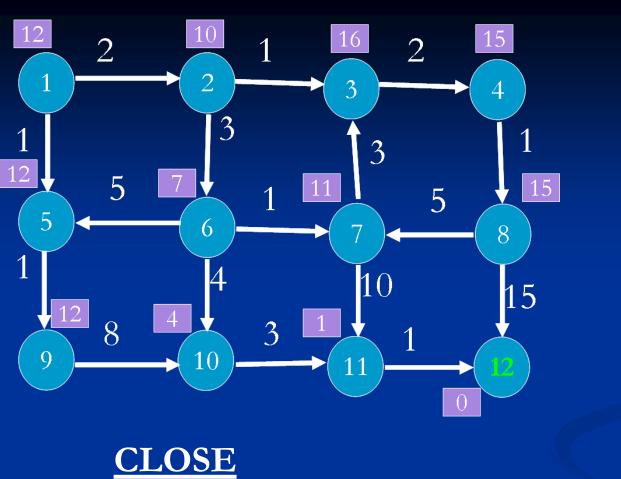






OPEN 5(13) 3(19) 6(12)

Node	g (
1	0
2	2
5	1
3	3
6	E



1(12) 2(12) 6(12)

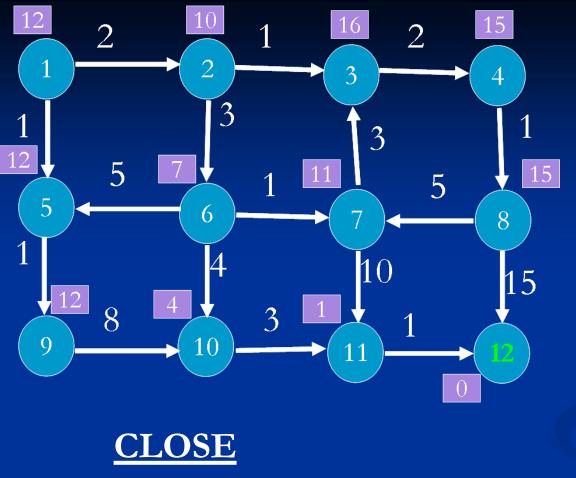
OPEN 5(13) 3(19)

g() 0 2 5 1 3 3

5

Node

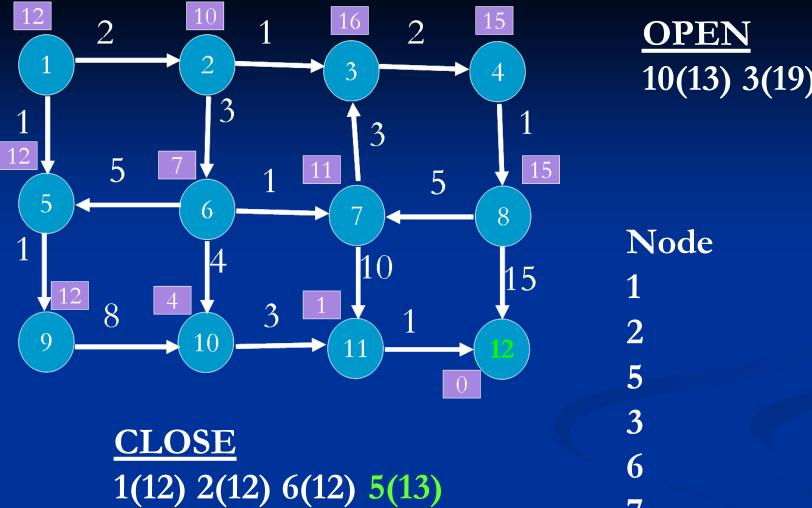
6



1(12) 2(12) 6(12)

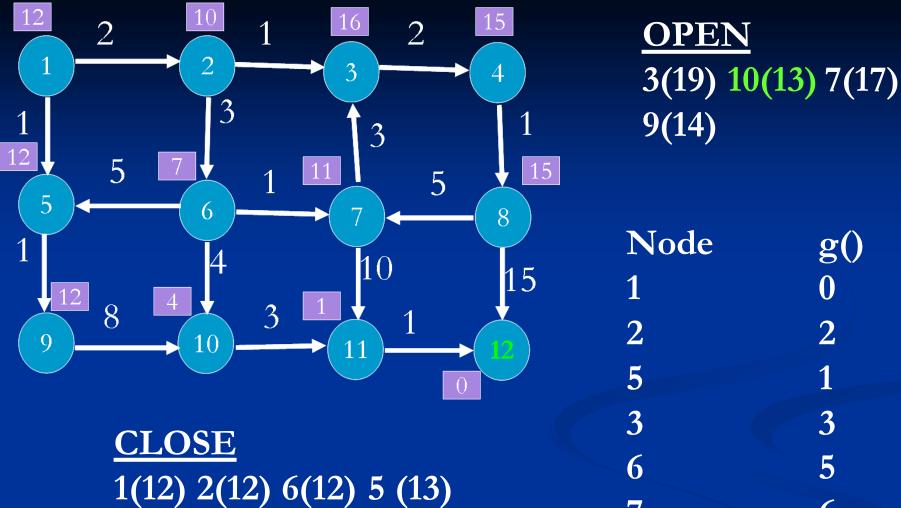
<u>OPEN</u> 10(13) 3(19) 5 (13) 7(17)

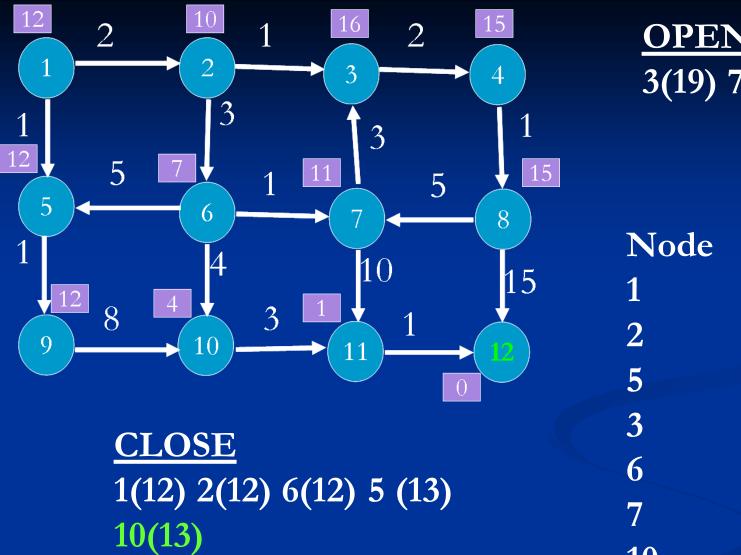
Node	$\mathbf{g}($
1	0
2	2
5	1
3	3
6	5
7	6
10	9



10(13) 3(19) 7(17)

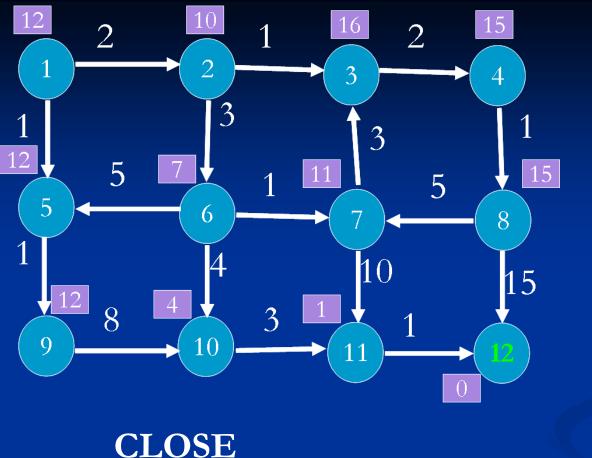
Node	g
1	0
2	2
5	1
3	3
6	5
7	6
10	9





OPEN 3(19) 7(17) 9(14)

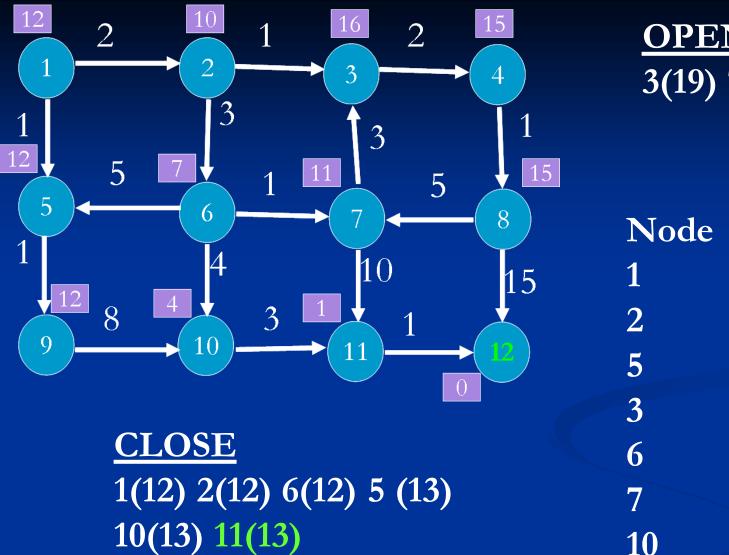
Node	g()
1	0
2	2
5	1
3	3
6	5
7	6
10	9
9	2



CLOSE 1(12) 2(12) 6(12) 5 (13) 10(13)

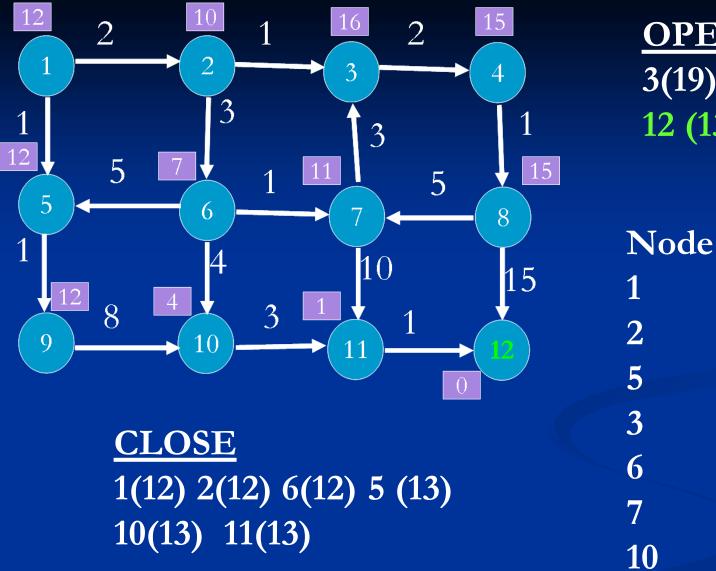
OPEN 3(19) 7(17) 9(14) 11(13)

Node	g()
1	0
2	2
5	1
3	3
6	5
7	6
10	9
9	2
11	12



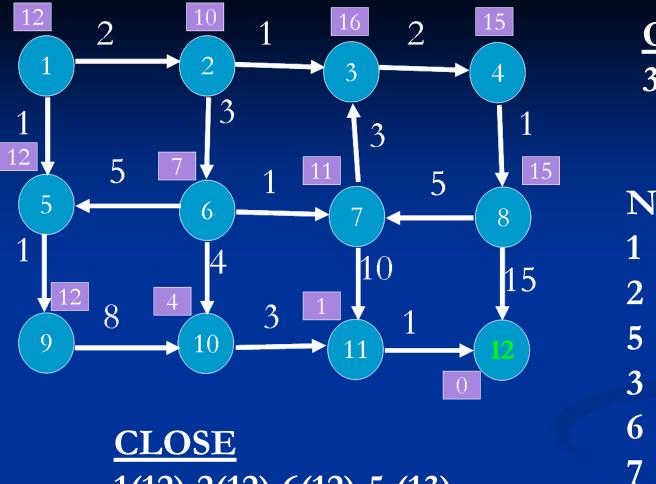
<u>OPEN</u> 3(19) 7(17) 9(14)

Node	g()
1	0
2	2
5	1
3	3
6	5
7	6
10	9
9	2
11	12



OPEN 3(19) 7(17) 9 (14) 12 (13)

Node	g()
1	0
2	2
5	1
3	3
6	5
7	6
10	9
9	2
11	12
12	13



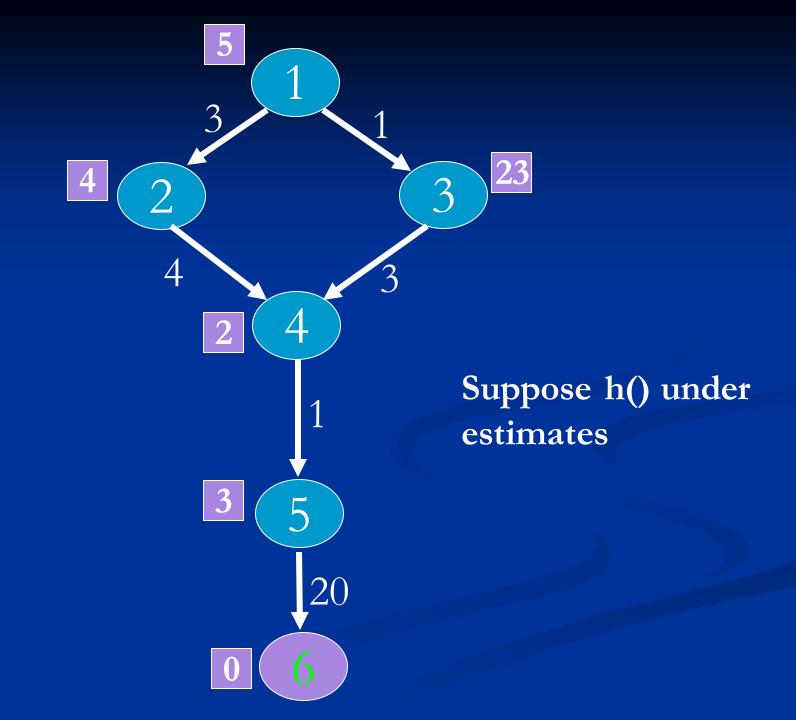
<u>CLOSE</u> 1(12) 2(12) 6(12) 5 (13) 10(13) 11(13) 12 (13)

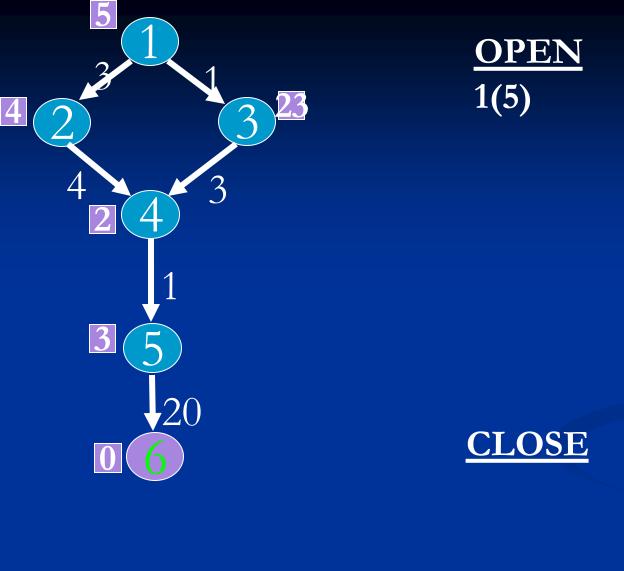
<u>OPEN</u> 3(19) 7(17) 9 (14)

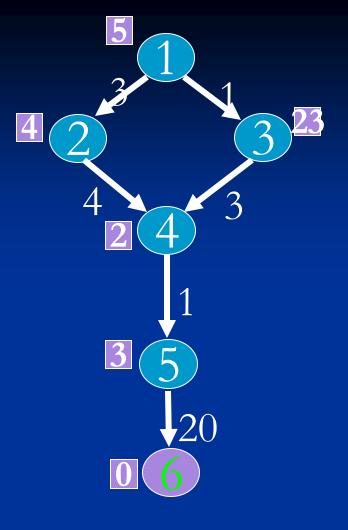
Node	g()
1	0
2	2
5	1
3	3
6	5
7	6
10	9
9	2
11	12
12	13

Comparing with OR Graph Search

- Instead of 11 nodes (OR) we expanded only 7 nodes in A*
- Inference: Nodes which looked promising initially were found to be not so good later on and were ignored/left off

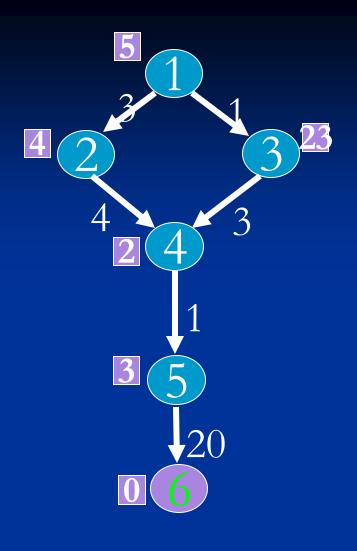






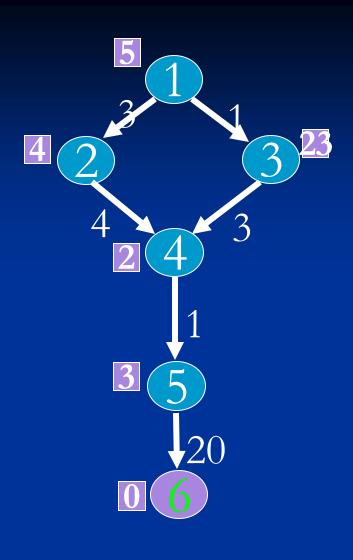
<u>OPEN</u>

CLOSE 1(5)



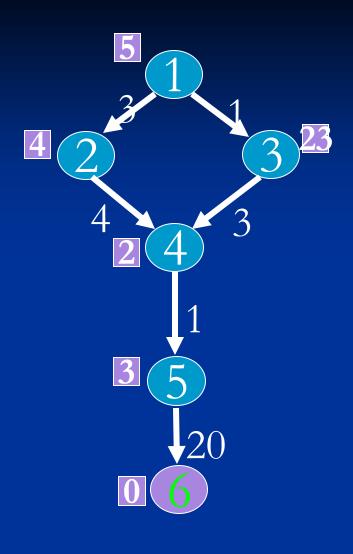
OPEN 2(7) 3(24)

CLOSE 1(5)



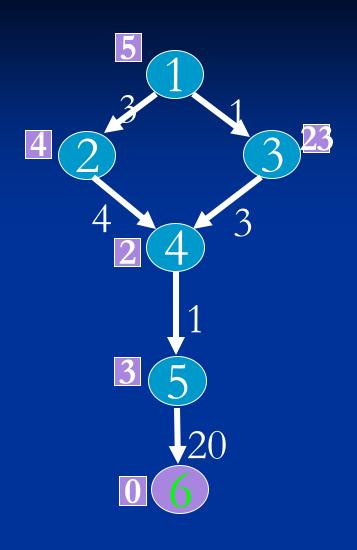
OPEN 3(24)

CLOSE 1(5) 2(7)



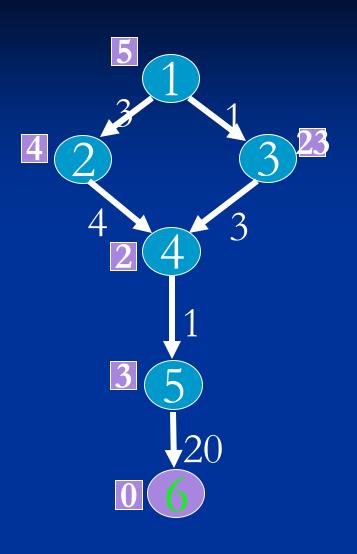
OPEN 3(24) 4(9)

CLOSE 1(5) 2(7)



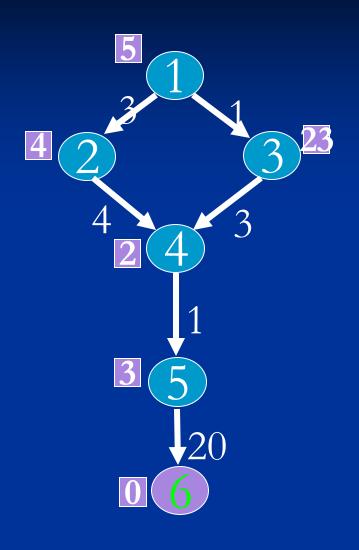
OPEN 3(24)

CLOSE 1(5) 2(7) 4(9)



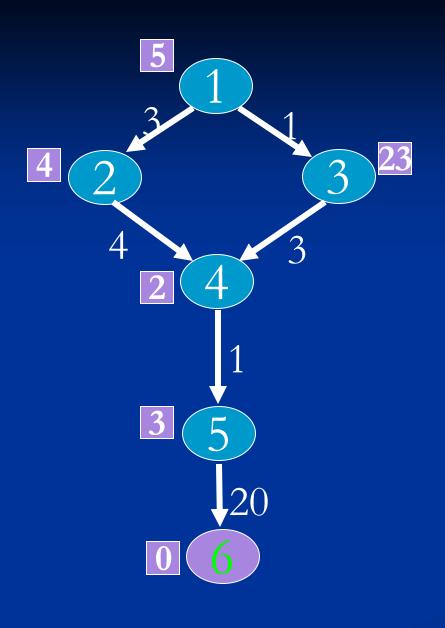
OPEN 3(24) 5(11)

CLOSE 1(5) 2(7) 4(9)



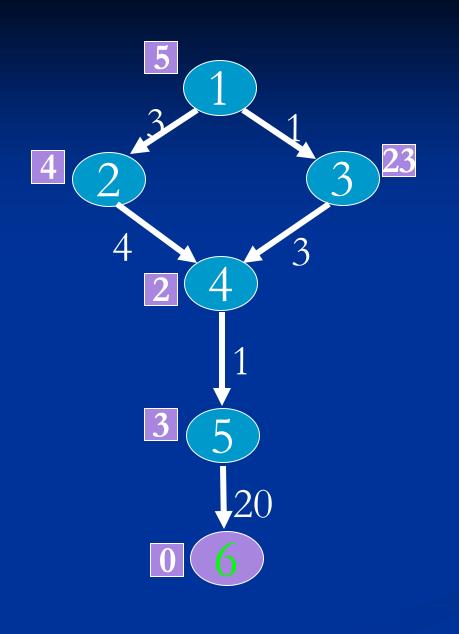
OPEN 3(24)

CLOSE 1(5) 2(7) 4(9) 5(11)



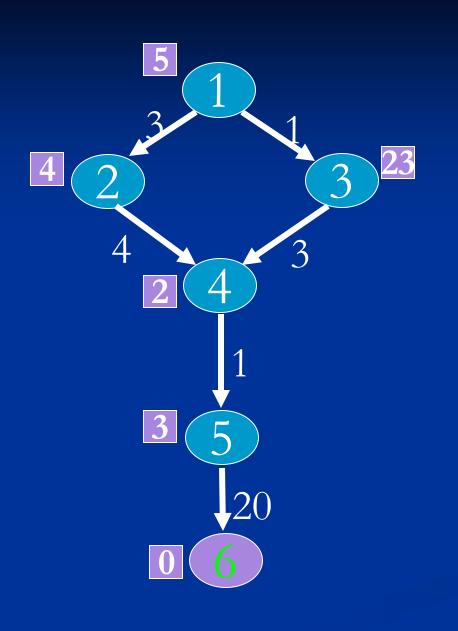
OPEN 3(24) 6(28)

<u>CLOSE</u> 1(5) 2(7) 4(9) 5(11)



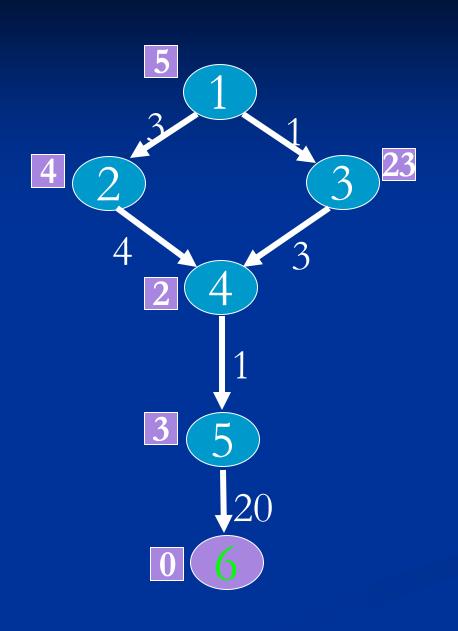
OPEN 6(28)

CLOSE 1(5) 2(7) 4(9) 5(11) 3(24)



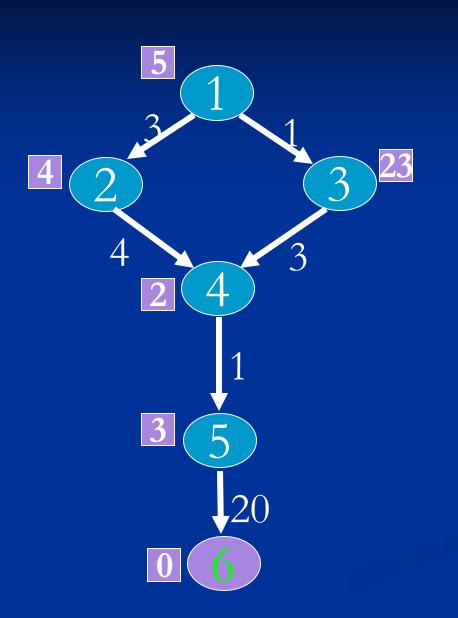
OPEN 6(28) 4(6)

CLOSE 1(5) 2(7) 4(9) 5(11) 3(24)



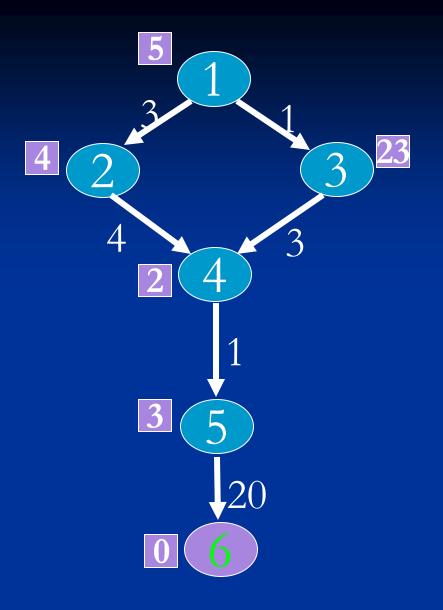
OPEN 6(28)

CLOSE 1(5) 2(7) 4(9) 5(11) 3(24) 4(6)



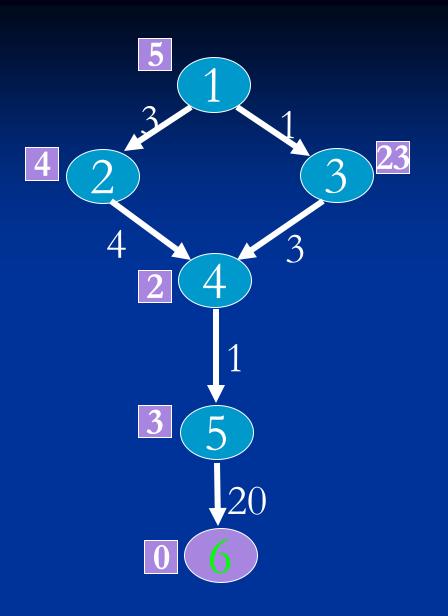
OPEN 6(28) 5(8)

CLOSE 1(5) 2(7) 4(9) 5(11) 3(24) 4(6)



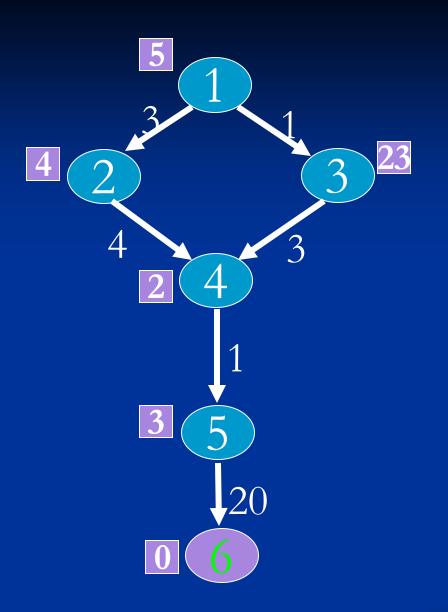
OPEN 6(28)

CLOSE 1(5) 2(7) 4(9) 5(11) 3(24) 4(6) 5(8)



OPEN 6(28) 6(25)

CLOSE 1(5) 2(7) 4(9) 5(11) 3(24) 4(6) 5(8)



OPEN

CLOSE 1(5) 2(7) 4(9) 5(11) 3(24) 4(6) 5(8) 6(25)

Results

- A heuristic is called admissible if it always under estimates, that is, we always have $h(n) \le f^*(n)$, where $f^*(n)$ denotes the minimum distance to a goal state from state n
- For finite state spaces, A* always terminates
- Algorithm A* is admissible, that is, if there is a path from start state to a goal state, A* terminates by finding an optimal path