

P - atom_codes (tiger, X)

$$X = [116, 105, 103, 101, 114]$$

Yes

? - atom_codes (tom, A)

$$A = [116, 111, 109]$$

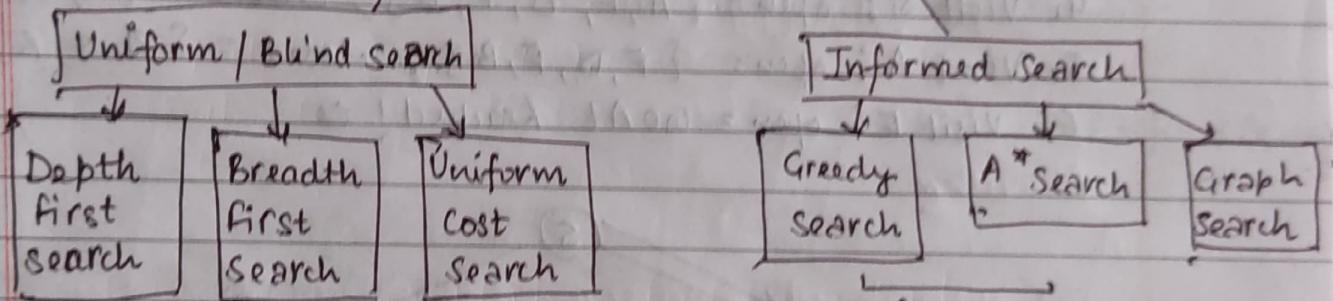
Yes

Ans:

write a program to find cube of a number.

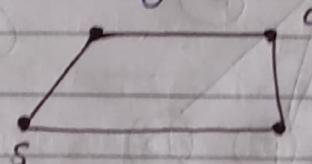
Types of search Algorithms

Friday
22/09/2022



Each of these algorithms will have :

10 A problem graph containing start node S & goal node G.



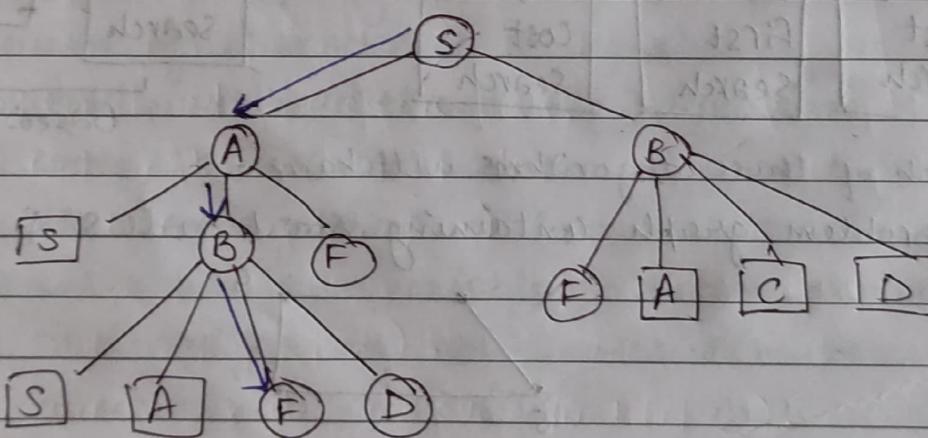
2. A strategy, describing the manner in which the graph will be traversed to get to G.
3. A fringe, while is a structure used to store all the possible states (nodes) that you can go from the current state.
4. A tree, that results, while traversing to the goal node.

5. A solution plan, while the sequence of nodes from S to G.

Depth First Search

$$\text{open} = \{S\}, \text{closed} = \{\}$$

0. Visit S; open = {A, B}, closed = {S}
1. Visit A; open = {S, B, F, B}, closed = {A, S}
2. Visit S; open = {B, F, B}, closed = {S, A, S}
3. Visit B; open = {S, A, B, F, D, F, B}, closed = {B, S, A, S}
4. Visit S; open = {A, F, D, F, B}, closed = {S, B, S, A, S}
5. Visit A; open = {F, D, F, B}, closed = {A, S, B, S, A, S}
6. Visit F; open = GOAL Reached!



- DFS: complexity

- Time complexity = $O(b^m)$ m is depth of sol.
- Space complexity = $O(bm)$

DFS & Algorithms

List open, closed, Successors = {};

Node root_node, current_node;

insert - first (root_node, open)

while not-empty (open);

current_node = remove-first (open);

insert - first (current_node, closed);

if (goal (current_node)) return current_node;
else

successors = successors of (current node);

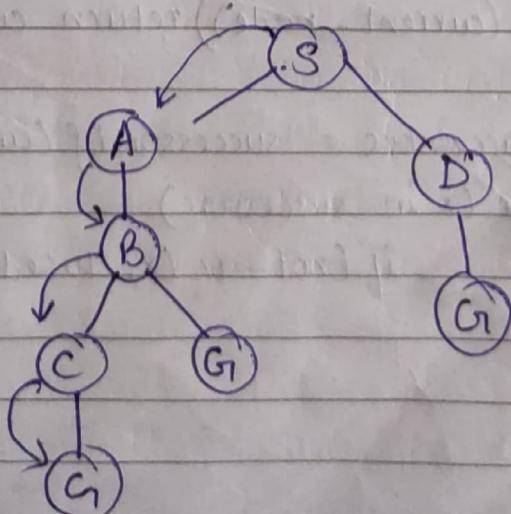
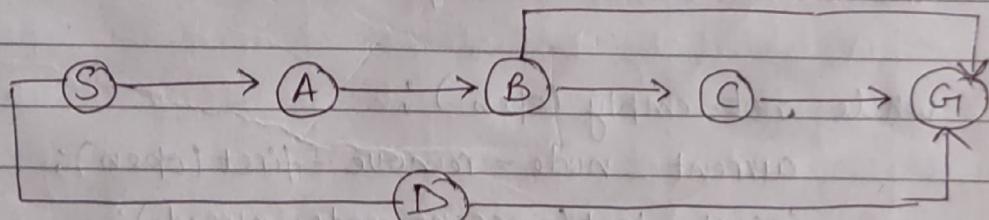
for (x in successors)

if (not-in (x, closed)) insert-fixed (x, open);

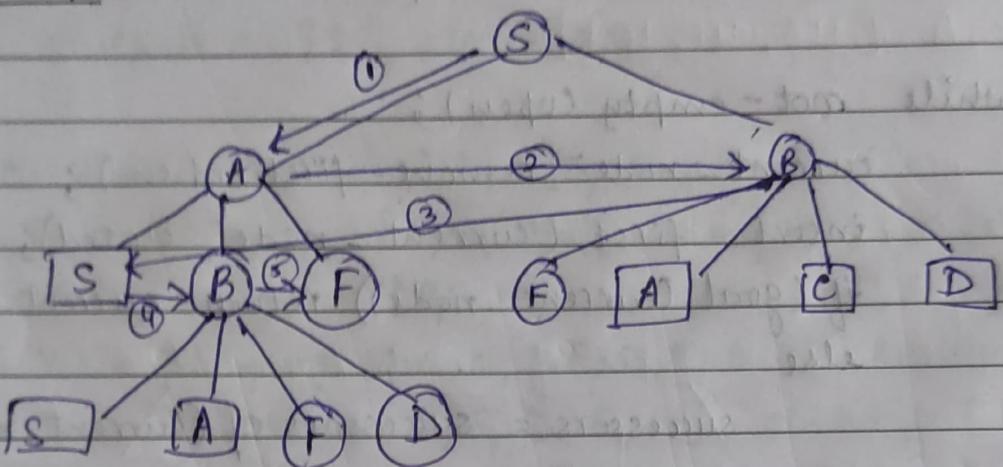
endif

endwhile

Page 27 DFS



Informed

Breadth first searchExample & AlgorithmExampleAlgorithm

list open, closed, successors = {};

Node root_node, current_node;

insert - last (root_node, open)

while not-empty (open);

current_node = remove-first (open);

insert - last (current_node, closed);

if (goal (current_node)) return current_node;

else

successors = successorsOf (current_node);

for (x in successors)

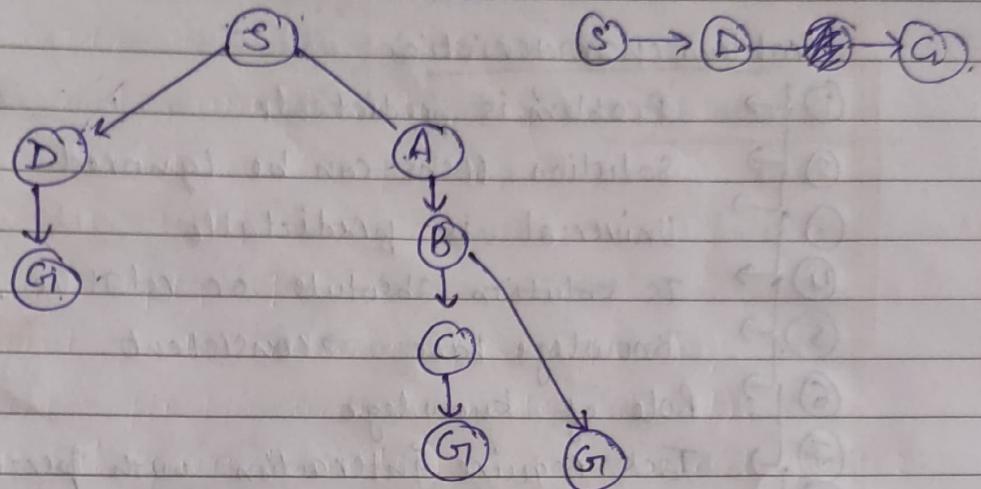
if (not-in (x, closed)) insert - last (x, open);

endif

endwhile

Pg 36

BFS

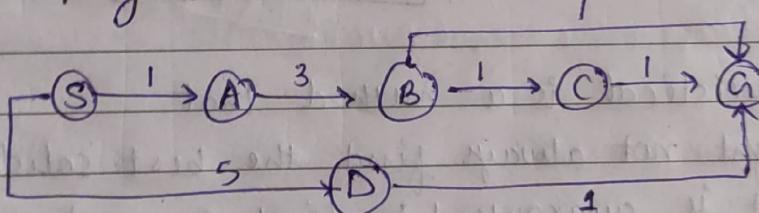


$$S \rightarrow D \rightarrow \cancel{A} \rightarrow G_1$$

Pg 39

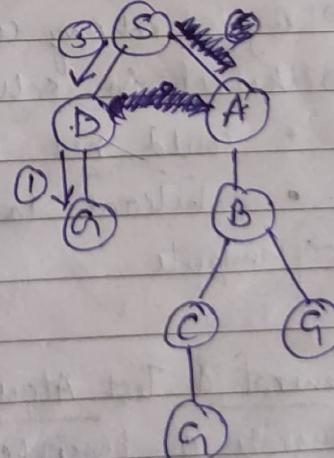
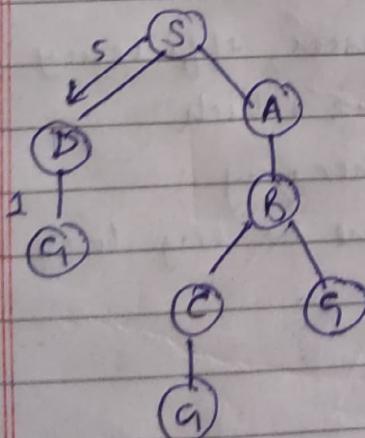
Uniform Cost Search

- USS is different from BFS & DFS because here the costs come into play



DFS (cost = 6)

BFS (cost = 6)



Informed Search Algorithms

Problem characteristics

- ① → Problem is predictable
- ② → Solution steps can be ignored
- ③ → Universal is predictable
- ④ → Is solution absolute or relative.
- ⑤ → Knowledge base → consistent
- ⑥ → Role of knowledge
- ⑦ → Task require interaction with person
- ⑧ → Problem classification

Tuesday

Pg 44

03/10/2023

Heuristic search

- A heuristic is a method that
 - might not always find the best solution
 - but is guaranteed to find a good solution in reasonable time.
 - By searching completeness it increases efficiency.
 - Useful in solving tough problems which
 - could not be solved any other way.
 - solutions take an infinite time or very long time.
 - compute

General & Test Algorithm Pg 45

- Generate a possible solution which can either be a point in the problem space or a path from the initial state.
- Test to see if this possible solution is a real solution by comparing the state reached with set of goal states.
- If it is a real solution, return. otherwise repeat from 1.

Pg 46

Hill climbing

→ Drawback → It is not complete unless we introduce backtracking.

It is not optimal. → Solution find is a local optimum.

Pg 47

56 - Pseudo code

Hill climbing Algorithm | Greedy search Method | Greedy local Search
→ used for optimizing mathematical problems.

Hill climbing node has 2 components

statevalue

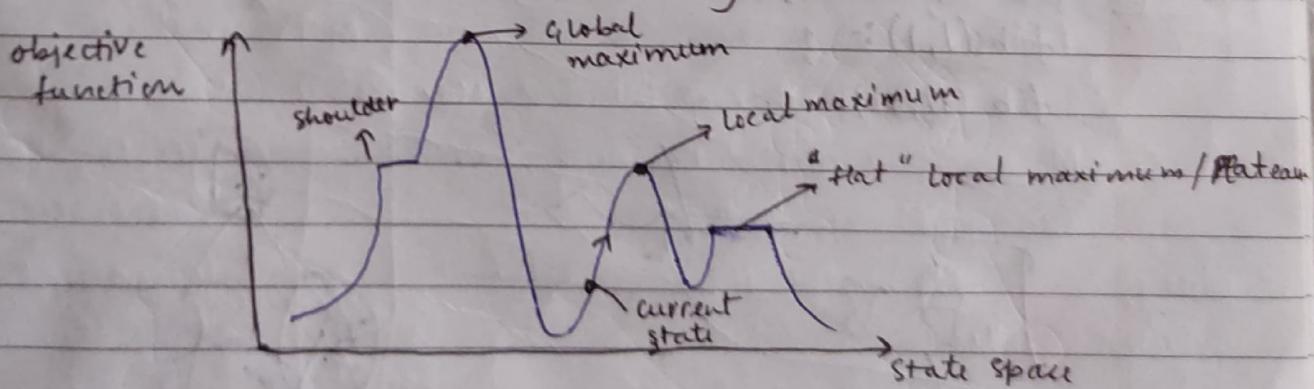
Pg 48

Features of Hill climbing

1. Generate & Test variant → produce feedback which helps to decide which direction to move in.
2. Greedy approach → Hill climbing search moves in direction the search space which optimizes the cost.
3. No backtracking.

Pg 49

State space Diagram for Hill climbing



2. Local Maximum: State better than its neighbouring state, but there is also another state which is higher (Global Max.)

2. Global Max :- Best possible state of state space landscape.
→ Highest value of objective function.

3. Current state :- State where agent is currently present.

4. Flat Local Maximum :- All neighbouring state have same value.

Good Write

Types of Hill Climbing Algorithm

(SAHC)

1. simple Hill Climbing (SHC)

→ less optimal solⁿ & solⁿ is
not guaranteed.

→ less time consuming.

→ simplest way to implement
a hill climbing algorithm.

Evaluates neighbour node at a
time & selects the 1st one which
optimizes current cost & set it as a
current state.

2. steepest-Ascent Hill-Climbing

→ examines all the neighbouring
nodes of the current state and
selects one neighbour node which
is closest to goal node

→ more time consuming as
it searches for multiple
neighbours.

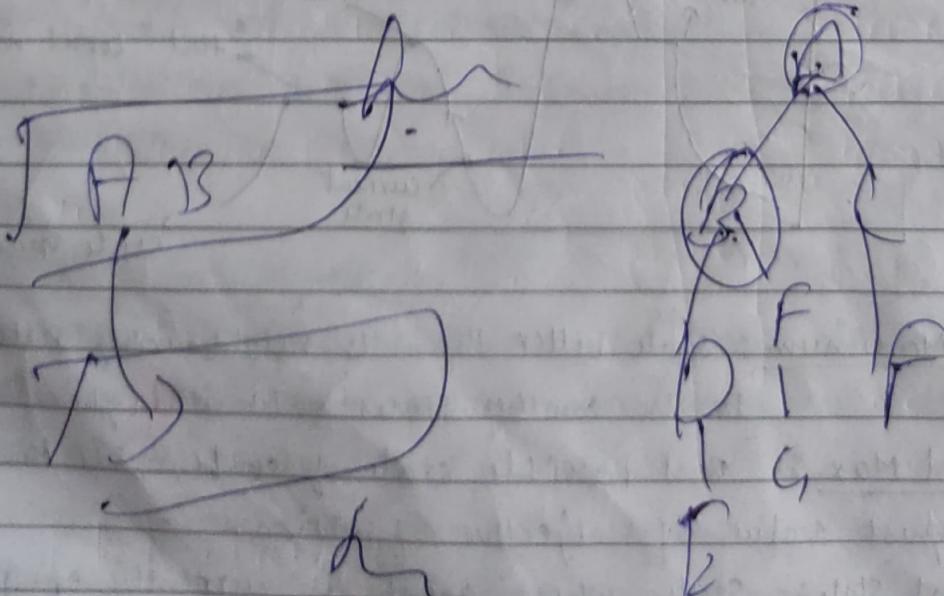
Algorithm for SHC , (SAHC) Pg 54

Fibonacci Series

fib(0,1):- !.

fib(1,1):- !.

fib(N,F)



SECTION - A

DATE: _____
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Ex 1

Fibonacci series

$\text{fib}(0, \emptyset) :- !.$

$\text{fib}(1, 1) :- !.$

$\text{fib}(N, F) :-$

$N \geq 1,$

$N_1 \text{ is } N-1,$

$N_2 \text{ is } N-2,$

$\text{fib}(N_1, F_1),$

$\text{fib}(N_2, F_2),$

$F \text{ is } F_1 + F_2.$

* Calculate factorial

factorial [0, 1].

factorial [N, M] :-

$\square N \geq 0,$

prev is $N-1,$

factorial (prev, M1),

$M \text{ is } M1 * N.$

Ex 2

local variables & conditional statements

E

$\text{eats}(\text{plant}, \text{grasshopper}).$

output
numbers.

$\text{eats}(\text{grasshopper}, \text{frog}).$

$\text{eats}(\text{frog}, \text{snake}).$

$\text{eats}(\text{snake}, \text{hawk}).$

$\text{food-web}(X, Y) :-$

$\text{eats}(X, Y);$

$\text{format('A ~ W eats W', [X, Y])},$

numbers :-

$\text{write('Enter A:')} =$

$\text{read(A)},$

$\text{nl},$

$\text{write('Enter B:')} ,$

$\text{read(B)},$

$\text{nl},$

$\text{condition}(A, B).$

$A > B, \text{ write('A value is greater than B value')};$

$A := B, \text{ write('A value is equal to B value')};$

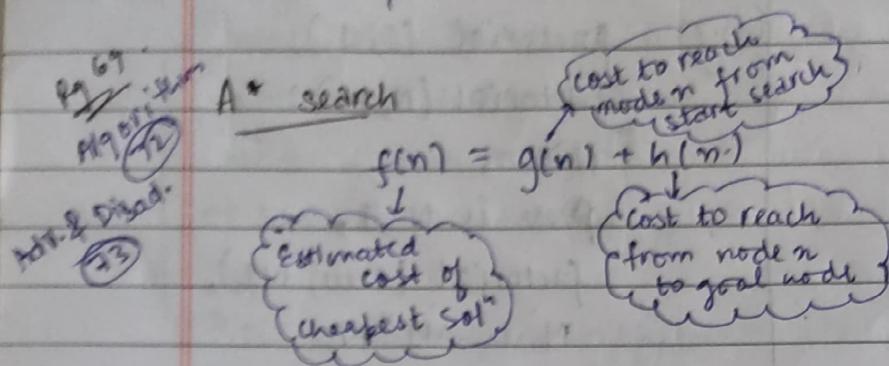
Good Write

$A < B, \text{ write('A value is less than B value')};$

Pg No. 20

Best First Search (Greedy Search)

↳ complete but not optimal.



- find shortest path.
- expand less search tree.
- back point to check the cover distance.

* 40* Algorithm → Solves AND, OR problems

(lab)

(SECTION - B)

Ex - 2
3facts for statement
°C to °F.

11/10/2023 Wednesday

Unit - II• Knowledge RepresentationExample 1: Family relation

Pg 5

Facts:

- Sarah is the mother of Tal and Mor
- Mosche is married to Sarah
- Fanny is the mother of Gal

Query: Is Mosche the father of Tal?Deduction? Yes

2): Circuit Diagnosis Pg 6

What is knowledge?

Pg 7 Types of Knowledge Pg 8

Universal
→ Priori Knowledge truenot relia
→ Posteriori knowledge

→ Procedural knowledge knowing how to do

Difference
Pg 8
→ Declarative knowledge is something True / False

→ Tacit knowledge

knowledge not easily expressed by language.

Part I Representation

↳ symbols standing for in the world or symbolic encoding of propositions.



→ First Aid



→ women

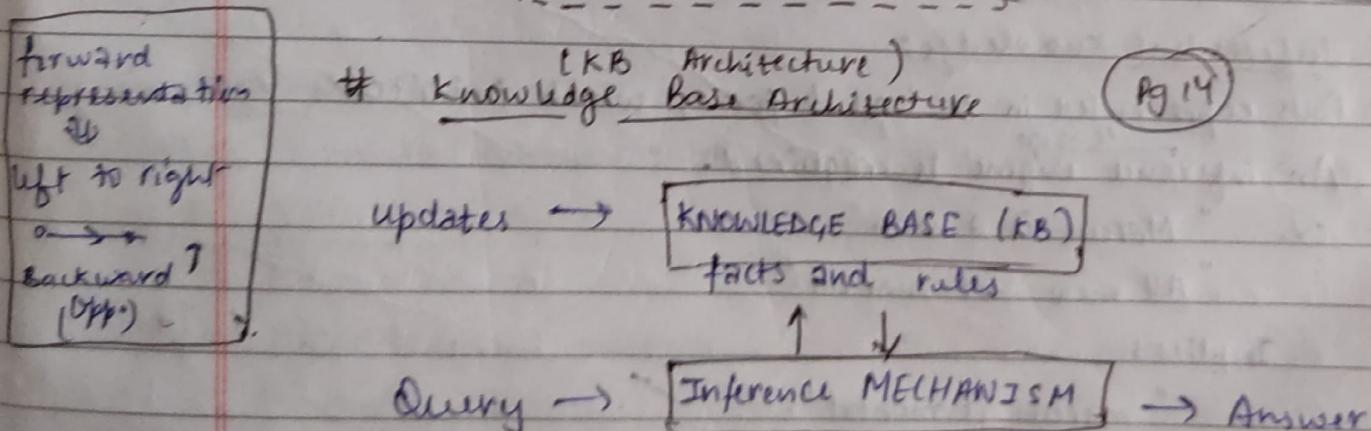
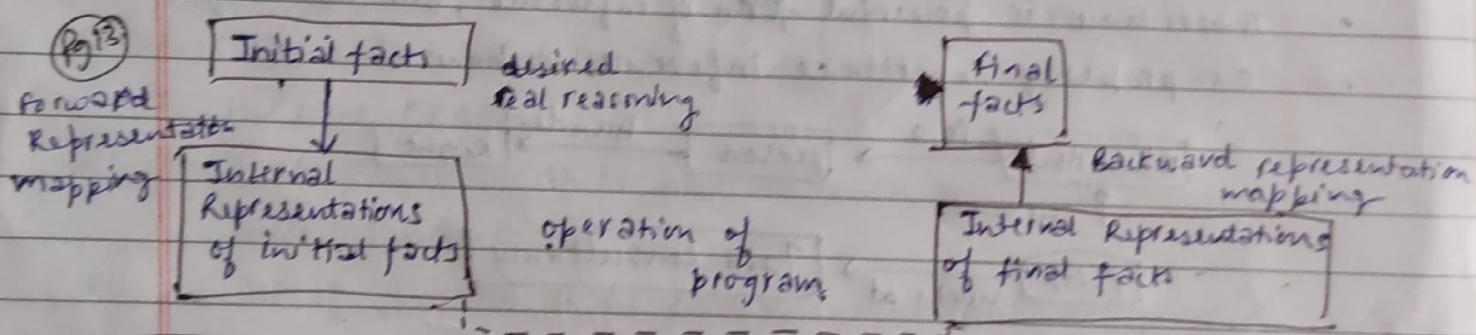
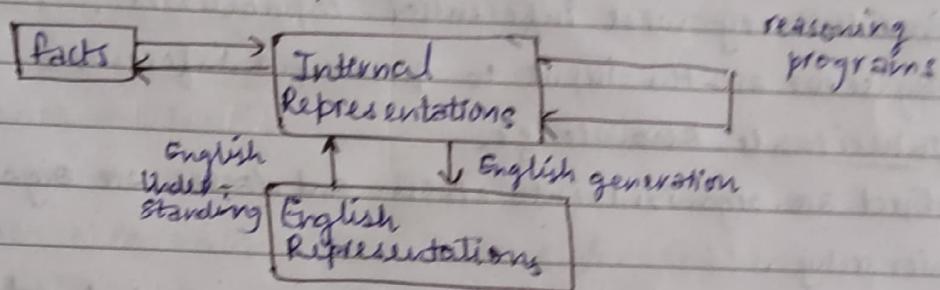
Good Write

Reasoning

↳ Manipulation of symbols encoding propositions to produce representations of new propositions.

$$\begin{array}{ccc} \text{Binary} & \rightarrow & 1101 + 10 = 1101 \\ & & (11) \quad (2) \quad (13) \end{array}$$

(Pg 12) # What is Representation And Mapping?



Representation & Mapping

- Spot is a dog.
- Every dog has a tail.
spot has a tail.

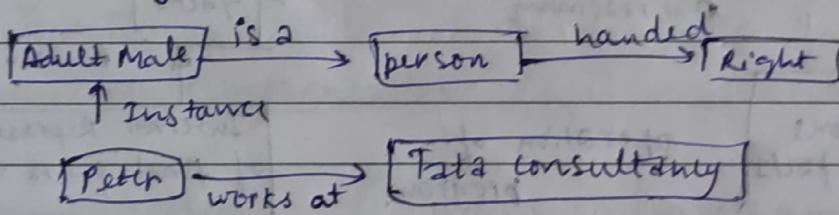
dog(spot)
forall x : dog(x) → hasTail(x)
hasTail(spot)
hasTail(spot)

Knowledge base issues

- ① Representation language: expressive is it?
Should be uniform universal to avoid it symbol.
- ② Inference procedure:-
- + is it sound? conclusion rationally follows from the facts & rules.
 - + is it complete? " " " from the KB, can I deduce it?
 - + is it efficient? Does it take time polynomial in the no. of facts & rules.

Approach to KB

- ① Simple factual relational knowledge
- Provides very weak inferential capabilities.
 - May serve as the input to powerful inference.
- ② Inheritable knowledge
- Objects are organized into classes and classes are organized in a generalization hierarchy.
 - Inheritance is powerful form of inference, but not adequate.
 - ex - property inheritance inference mechanism.



③ Inferential knowledge:

- Facts represented in a logical form, which facilitates reasoning
 - An inference engine is required.
- ex. "Marcus is a man"
"All men are mortal."

Implies:

"Marcus is mortal".

④ Procedural knowledge

- how to make it rather than what it is.

Issues in KB

Related to Path

1. Important attributes: Isa and Instance attributes.
2. Relationships among attributes: inverses, existence in a Isa hierarchy.
Good Write Single-valued attributes, techniques for reasoning about values.

First order logic = Predicates verb (sub, obj)
Axioms: Relationship b/w statements

DATE: _____
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3. Choosing the granularity

Ex - "John spotted sue".

spotted (John, sue)

4. Representing set of objects:

Finding the right structure as needed.

word "fly" have multiply meanings

"John flew to new york." 1. "John flew in a rage." 2. "John flew a kite."

Domain Model "ontology"

Representation language

KB Inference procedure

* Implementation

2. Domain theory

12/10/2023
Thursday

A sentence is valid if it is true for any truth assignment.
 $(A \vee \neg A)$.

Satisfiable if there exists a truth assignment that make it true $(A \wedge B)$.

unsatisfiable if there exist no truth assignment that makes it true $(A \wedge \neg A)$

model of sentence is an interpretation that satisfies the sentence.

Proposition • sentence which is true or false.

Proposition symbols/variables : P, Q, S

Propositional logic $\wedge, \vee, \Rightarrow, \Leftrightarrow, \circ / \neg$

Properties of sentence

→ satisfiability → contradiction

Ex - Japan is capital of India.

→ Validity
Delhi is the capital of India.

Q.P

laws of Algebra of Propositions

• Idempotent

$$P \vee P \equiv P$$

$$P \wedge P \equiv P$$

• Commutative :

$$P \vee q \equiv q \vee P \quad P \wedge q \equiv q \wedge P$$

• Complement

$$P \vee \neg P \equiv T$$

$$P \wedge \neg P \equiv F$$

• Double Negation

$$\neg(\neg P) \equiv P$$

Good Write

- Associative

$$p \vee (q \vee r) \equiv (p \vee q) \vee r$$

$$p \wedge (q \wedge r) \equiv (p \wedge q) \wedge r$$

- Distributive

$$p \vee (q \wedge r) \equiv (p \vee q) \wedge (p \vee r)$$

$$p \wedge (q \vee r) \equiv (p \wedge q) \vee (p \wedge r)$$

- Absorption

$$p \vee (p \wedge q) \equiv p$$

- Identity

$$p \vee T \equiv T, p \vee F \equiv p, p \wedge T \equiv p, p \wedge F \equiv F$$

- De Morgan's

$$\sim(p \vee q) \equiv \sim p \wedge \sim q$$

$$\sim(p \wedge q) \equiv \sim p \vee \sim q$$

- Equivalence of contrapositive

$$p \rightarrow q \equiv \sim p \vee q$$

$$p \leftrightarrow q \equiv (p \rightarrow q) \wedge (q \rightarrow p)$$

Inference rules :- logical inference \rightarrow used to create new sentence that logically follow from given set of predicate calculus sentence (KB)

Sound rules of inference - A rule is sound if its conclusion is true whenever the premise is true.

User provides

Constant Symbols
green, 3, s, mary

Function Symbols
father_of(mary) = John

Predicate Symbols
- greater(s,3)
- green(4, grass)

Tuesday

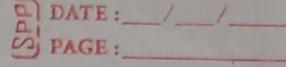
17/10/2023

(Pg 60) For illustrated models for FOL

to Pg 72

- Crown has a King. & Crown is a person. $\text{Crown}(\text{King}) \wedge \text{Crown}(\text{Person})$
- Richard and John both are siblings. $\text{siblings}(\text{John}) \leftrightarrow \text{siblings}(\text{John})$
- Richard and John have left leg. $\text{left_leg}(\text{Richard}, \text{John})$
- All students have a roll no. $\forall x : \text{student}(x) \rightarrow \text{roll_no}(x)$
- None of the students like Icecream. $\neg \exists x : (\text{student}(x)) \rightarrow \text{likes}(x, \text{icecream})$
- Every student likes healthy food. $\forall x : (\text{student}(x)) \rightarrow \text{likes}(x, \text{healthy})$
- Some students like ice cream. $\exists x : (\text{student}(x)) \wedge \text{likes}(x, \text{icecream})$

(Pg 82) X8/13u
Laws of Algebra & De-morgan's law
Good Write

<u>Universal instantiation</u> $\vdash \forall x P(x) :: P(A)$ instance of $\vdash \neg P(A) :: \exists x P(x)$	$\neg P(A) \wedge P(B) \dots :: \vdash \forall x P(x)$ [Universal generalisation]	[Existential instantiation] $\exists x P(x) :: P(F)$
		 DATE: _____ / _____ / _____ PAGE: _____

Unification with example → Similar sentences that makes them looks the same
 $\text{UNIFY}(\text{knows}(\text{John}, x), \text{knows}(\text{John}, \text{Jane})) = \{x / \text{Jane}\}$

and computes the substitut

Generalized Modus Ponens (GMP)

A_1' is King(John)
 only John is King

$$\frac{A \Rightarrow B, A}{B}$$

$$\frac{A_1 \wedge A_2 \wedge A_3 \dots A_n \Rightarrow B, A'_1, A'_2, \dots \text{Sub}(\sigma, B)}{B}$$

forward chaining

$$\text{Subst} \quad \text{SUBST}(\sigma, A_i) = \text{SUBST}(\sigma, A'_i)$$

Inferences unification resolution.

slide no 4
solution

Algo → Putting Axioms into clausal form

1. biconditional / Implication $\Leftrightarrow \rightarrow / \Rightarrow / \Leftarrow$ replace with $\vee + \neg$
2. Move \neg inward. $\neg A - \neg (\neg \text{Animal}(y) \vee \text{loves}(x, y)) \Rightarrow \neg \neg \text{Animal}(y) \neg \text{loves}(x, y)$

Constraint Satisfaction

Cryptarithmetic

Wednesday
25/10/2023