

# UE21CS351B - Cloud Computing

#### **Mini Project Report**

### Building a Task Management Application with Raft Consensus Algorithm and MySQL

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6<sup>th</sup> Semester D Section

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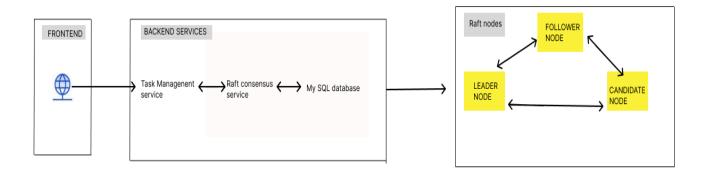
# DEPARTMENT OF COMPUTER SCIENCE AND ENGINEERING FACULTY OF ENGINEERING PES UNIVERSITY

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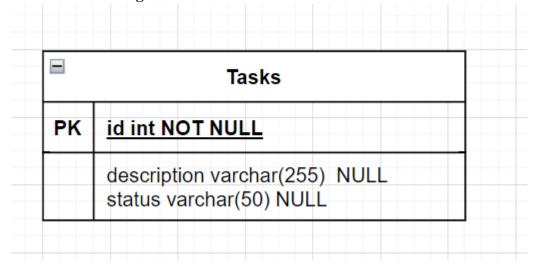
# Weekly Deliverables:

#### Week1:

- 1. Designing the Application Architecture
- Architecture



• Schema diagram



#### WEEK-2

## 1.Implementing Raft Consensus Algorithm:

In the Raft consensus algorithm, each node plays a specific role in the process of achieving distributed consensus. The roles include Leader, Follower, and Candidate. Here's a brief explanation of each role:

#### Leader:

- The Leader is responsible for coordinating the operations of the Raft cluster.
- It receives client requests, appends them to the log, and replicates the log entries to all Followers in the cluster.
- The Leader is the only node that can accept client updates to the system.
- In the event of a split vote during an election, Raft will require a new election to select a new Leader.

#### Follower:

- Followers are passive nodes that replicate log entries from the Leader.
- They respond to RPC (Remote Procedure Call) requests from Candidates and Leaders, granting votes during leader election rounds.
- Followers do not initiate elections; they follow the leadership decisions made by the Leader or Candidate nodes.

#### Candidate:

- When a node wishes to become the Leader, it transitions to the Candidate state and initiates a leader election.
- During the election process, a Candidate requests votes from other nodes in the cluster.
- If a Candidate receives votes from a majority of the nodes, it becomes the new Leader.
- If no node wins the election, a new election is triggered after a randomised timeout period.

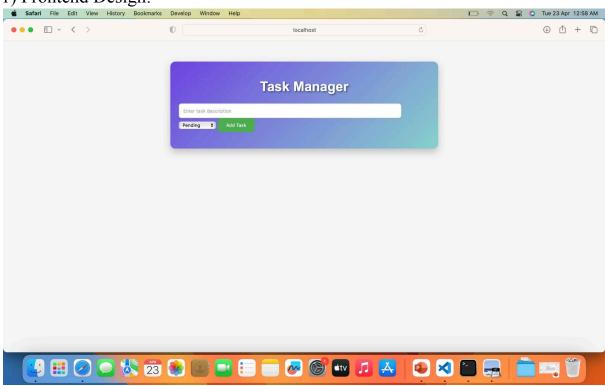
```
[Running] go run "d:\6th sem\CLOUD_COMPUTING\cc_project\raft_sql.go"
Connected to MySQL database!
Initiating leader election from node 1
Node 1 sending vote request to node 2
Node 1 received 1 votes, status: alive
Initiating leader election from node 2
Node 2 sending vote request to node 1
Node 2 received 1 votes, status: alive
Server is running on port 8501
```

In summary, the Raft consensus algorithm maintains a single Leader node responsible for managing the cluster's state and ensuring consistency among all nodes. Followers replicate the Leader's log and await instructions, while Candidates compete for leadership during leader election rounds. This division of roles ensures that the cluster remains stable, even in the presence of failures or network partitions.

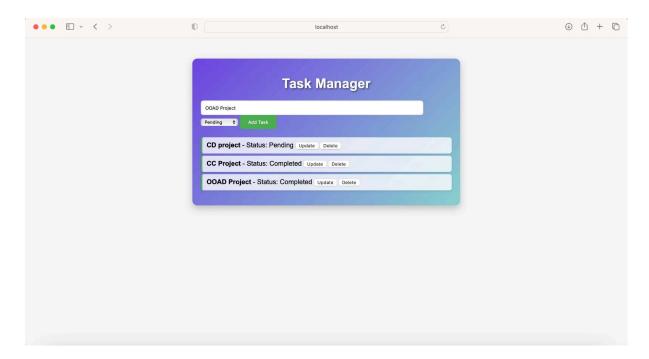
- 2. Communication protocols implemented:
- Heartbeat Mechanism: Implemented via periodic AJAX calls from the leader to followers, ensuring continuous monitoring and leader authority.
- Log Replication Protocol: Utilises HTTP POST requests for the leader to replicate logs to followers, with JSON payloads containing log entries.
- Election Process: Involves a combination of HTTP GET requests for voting and POST requests for announcing election results with timeouts.
- 3. Developed Backend:

#### **WEEK-3 PROGRESS:**

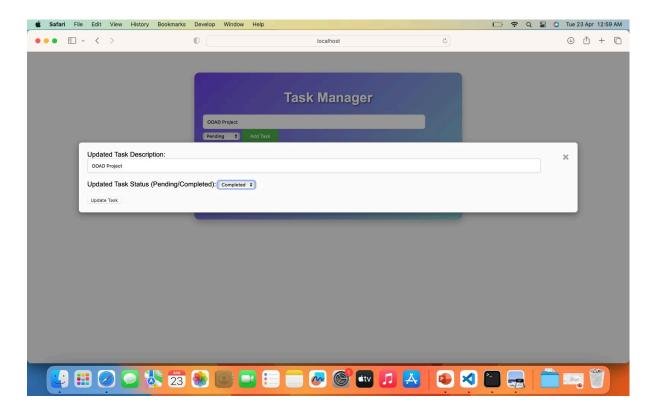
1) Frontend Design:



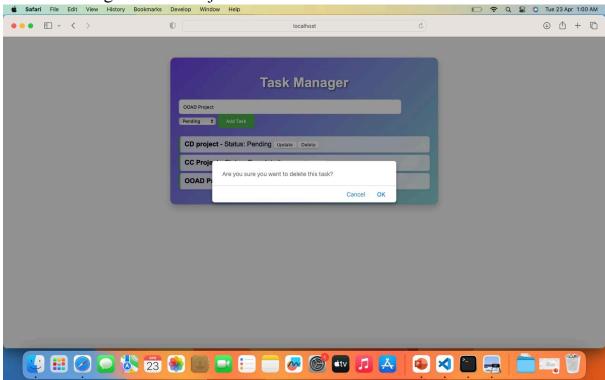
Added task OOAD project



• Updated task OOAD project from Completed to Pending



• Deleting task CD Project



• Table Structure



• Updated table after performing CRUD operations

