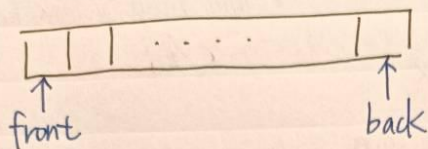


• Queue: 排隊, 先進先出



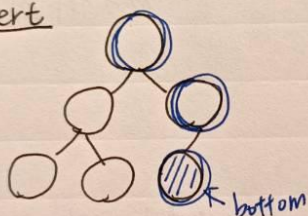
• Heap

- complete binary tree

- root 為優先級最高的元素

{ min heap: root 為最小值  
max heap: root 為最大值

Insert

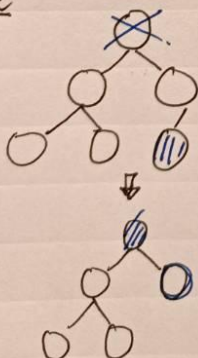


1. 將資料 insert 於 bottom 位置.

2. 往上比較, 若優先級較大則往上交換.

3. Efficiency:  $O(\log n)$

Delete



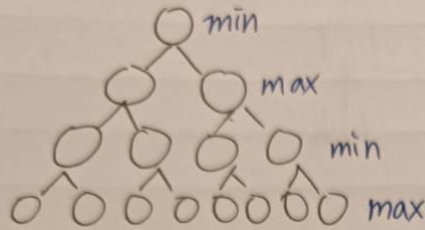
1. 刪除 root

2. bottom 移到 root 位置

3. 往下比較/交換

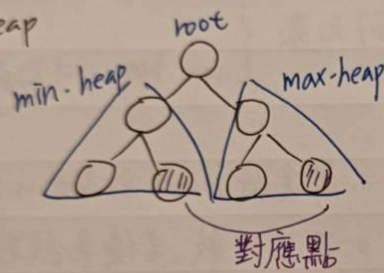
4. Efficiency:  $O(\log n)$

min-max heap



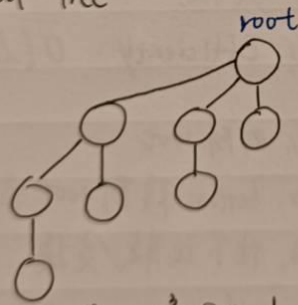
- min heap 和 max heap 交錯

Heap



- root 不存放資料
- root 左子樹為 min heap  
右子樹為 max heap

Binomial Tree

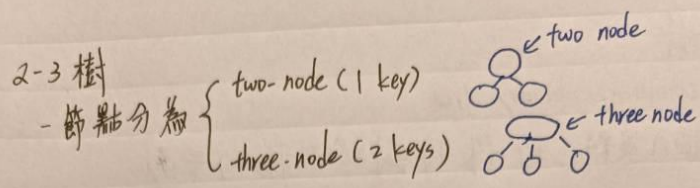


- root 有  $k$  個孩子  
⇒ order  $k$  ( $B_k$ )

$$B_3 \Rightarrow 2^3 = 8 (\text{nodes})$$

## 2-3 樹

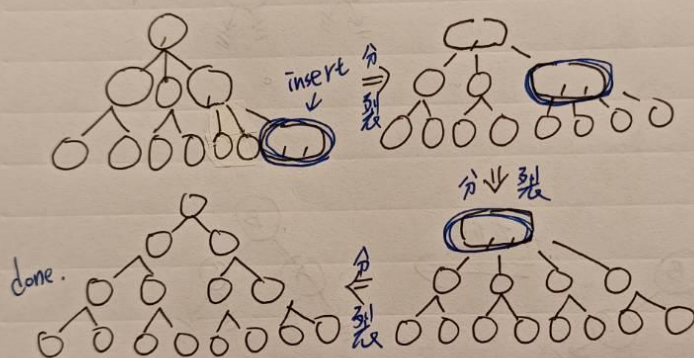
- 節點分為



- full tree

- 插入資料: insert 於樹葉節點

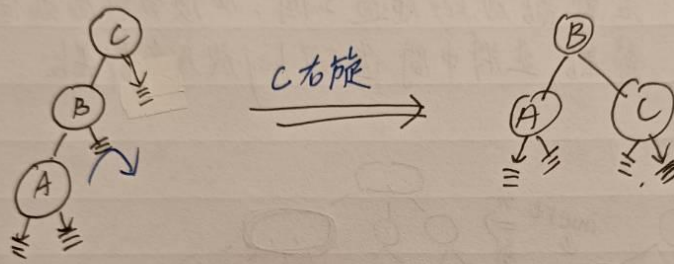
- 分裂: 若節點內 key 超過 2 個, 必須分裂為兩個節點, 並將中間值的 key 放入父節點。



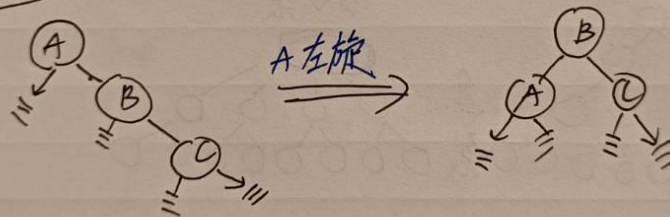
## AVL 樹

- complete binary tree
- 透過旋轉保持樹狀結構的平衡

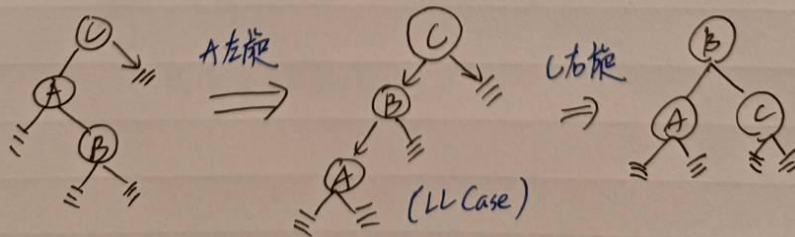
### LL Case.



### RR Case

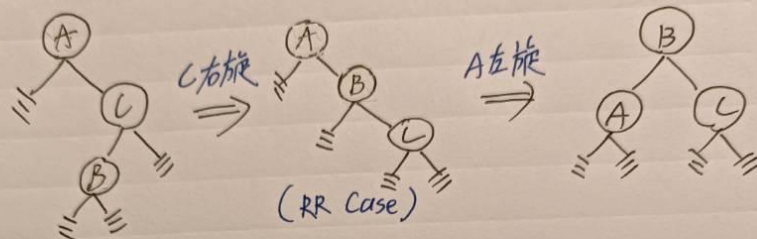


### LR Case



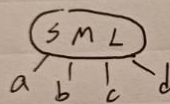


## RL Case

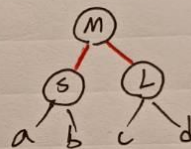


## 2-3-4 樹

4-node

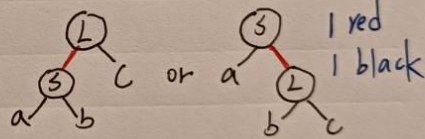


## 紅黑樹

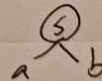
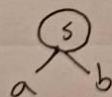


2 red pointers

3-node



2-node



2 black pointers

- full tree
- 節點內最多可有 3 keys

- complete binary tree
- root 到 leaf 的路徑上不可有連續的紅 pointer
- 每條路徑會有相同數量的黑 pointer