佇列 Queue

- First in First out

EX: Convert a sequence of digits into the decimal value 247.53

```
do { aQueue.dequeue(ch)          // 移除空格
} while (ch is blank)
n=0
done=false
while (!done && ch is digit){
   n=n*10 + integer of ch
   if (aQueue.isEmpty())         // 判斷是否為空
     done=true
   else aQueue.dequeue(ch)       // 移除佇列最前面ch
}// while
if (!done && ch=='.'){
 aQueue.dequeue(ch)              // 移除小數點
 p=0                             // p判斷小數後有幾位
 while (!done && ch is digit){
   n=n*10 + integer of ch
   p++
   if (aQueue.isEmpty())
     done=true
  else  aQueue.dequeue(ch)
 }// while
 n=n*(0.1)^p
}// if
```

判斷迴文：比較 stack 頂端和 queue 前端

EX: Recognizing Palindromes.

```
isPal
aQueue.createQueue()
aStack.createStack()
for(the next character ch in str){
  aQueue.enqueue(ch)
  aStack.push(ch)
}// for
charEqual = true
while(!aQueue.isEmpty() && charEqual){
```

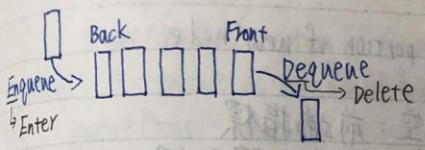
<法1>

```
  aQueue.getFront(front)
  aStack.getTop(top)
  if(front == top){
    aQueue.dequeue()
    aStack.pop()
  }// if
  else charEqual=False
}// while
```

<法2>

```
aQueue.dequeue(front)
aStack.pop(top)
if(front != top)
  charEqual=False
```

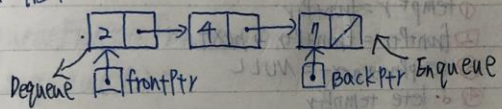
Implementations of the ADT Queue



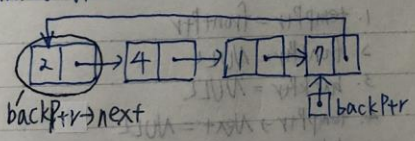
- A circular linked list with one external reference
  - Only a reference to the back (環狀：只有後端)

優：有一個變數

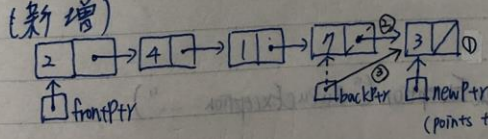
linked list



Circular linked list

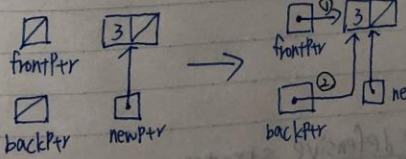


Enqueue (新增)



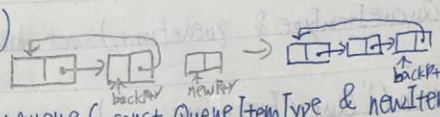
① newPtr→next = NULL
② backPtr→next = newPtr
③ backPtr = newPtr

若一開始為空



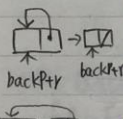
① frontPtr = newPtr
② backPtr = newPtr

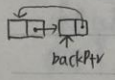
Circular(環狀)

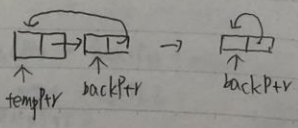



```
void Queue :: enqueue ( const QueueItemType & newItem ) { // 新增
    QueueNode * newPtr = new QueueNode;
    newPtr -> item = newItem;
    if ( isEmpty() )                          // 0 -> 1 node
      newPtr -> next = newPtr;                // point to itself
    else {                                    // k -> k+1 nodes, k>0
      newPtr -> next = backPtr -> next;       // point to the front
      backPtr -> next = newPtr;               // put behind the back
    } // else
    backPtr = newPtr;          // new node at the back
} // enqueue

void Queue :: dequeue() throw ( QueueException ) { // 移除
    if ( isEmpty() )
      throw ... ;
    else {
      QueueNode * tempPtr = backPtr -> next;   // tempPtr 指向 front
      if ( backPtr == backPtr -> next )        // 只有一個節點
        backPtr = NULL;                        // one node -> empty
      else
        backPtr -> next = tempPtr -> next;     // the next front
      tempPtr -> next = NULL;                  // defensive strategy
      delete tempPtr;                          // release space
    } // else
} // dequeue
```

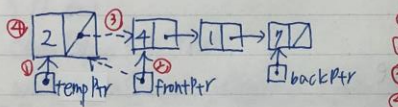




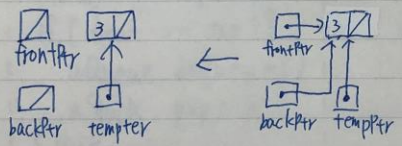


```
void Queue :: enqueue (const QueueItemType & newItem){ // 擷取後移除
    QueueNode * newPtr = new QueueNode;
    newPtr -> item = newItem;    // set data portion of new node
    newPtr -> next = NULL;       新節點指標
    if (isEmpty()) frontPtr = newPtr;   空：前端指標
    else backPtr -> next = newPtr;      非空：後端指標下一個
    backPtr = newPtr;            // new node is at the back
}// enqueue
```

Dequeue (移除)



① tempPtr = frontPtr
② frontPtr = frontPtr → next
③ tempPtr → next = NULL
④ delete tempPtr



1. tempPtr = frontPtr
2. frontPtr = NULL
3. backPtr = NULL
4. tempPtr → Next = NULL
5. delete tempPtr

```
void Queue :: dequeue() {
  if (isEmpty()) throw QueueException("QueueException: ...")
  else {
    QueueNode * tempPtr = frontPtr;
    if (frontPtr == backPtr) {   // one node only
      frontPtr = NULL;
      backPtr = NULL;
    }// if
    else frontPtr = frontPtr -> next;
    tempPtr -> next = NULL;      // defensive strategy
    delete tempPtr;              釋放空間
  }// else
}// dequeue
```

陣列實作佇列ADT (使用環狀陣列)

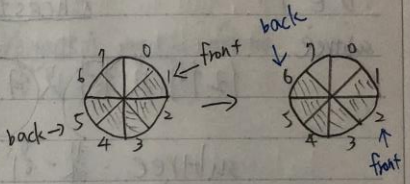
To initialize the queue:

```
front = 0
back = Max_Queue - 1
count = 0
```

Insert:

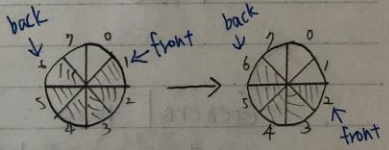
```
back = (back+1) % Max_QUEUE;
items[back] = newItem;
++count;
```

back初始值 = Max_Queue - 1



Deleting:

```
front = (front+1) % Max_QUEUE;
--count;
```

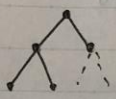


全空條件: count = 0

全滿條件: count = Max_QUEUE

Binary Tree

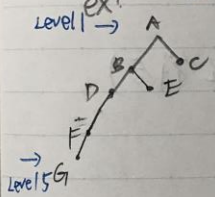
最多 2個 node



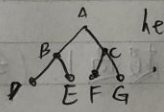
< left subtree
  right subtree

* Height of a tree (樹高)
  → number of nodes along the longest path from root to a leaf
  → 影響效率" (越小越好)

ex:



height 5



height 3

* Level of a node n in a tree T (階層)
  * 最大階層 = 樹高
  height(T) = 1 + max{ height(T_L), height(T_R) }

* 樹裡無 cycle

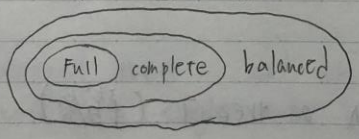
Full Binary Tree (完全樹)

height h is full
- nodes at levels <h have two children each

complete Binary Tree (完整樹)

① full to level-1 ex: level:5, level 4是full binary tree
② Level h is filled from left to right
③ at levels <=h-2 have two children each

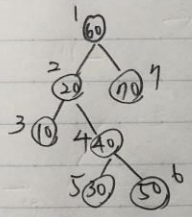
balanced
→左樹高，右樹高差距不超過1. (每個點都要符合)
→效率穩定，省空間



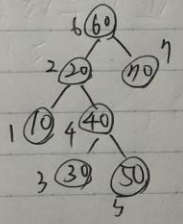
Traversal

- preorder ⇒ visit root before visiting its subtrees.
- Inorder ⇒ visit root between visiting its subtrees.
- Postorder ⇒ visit root after visiting its subtrees

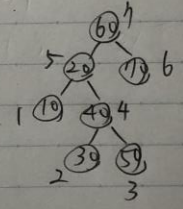
Preorder



Inorder



Postorder



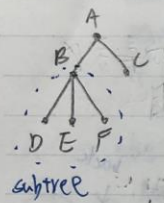
左邊先

Inorder = sort

Trees

目的：紀錄資料間關係（二維）

內容導向

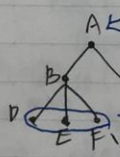


Trees are hierarchical
- Parent - child 親子 relationship between two nodes
- Ancestor - descendant 祖孫 relationship among nodes

任何資料只有一個直屬長官適用於 trees

subtree：某一部份資料形成小樹
子樹

General tree：one or more nodes
root：a single node



Parent of node B：A
-Leaf 葉節點：A node with no children
siblings 兄弟點：Node with a common parent
（父節點為同一人）

Ancestor of node B 祖先節點 A D
Descendant of node B 子孫節點 D A