

單元一 遞迴

□ Factorial 階乘

□ Greatest Common Divisor 最大公因數

□ Search in Array 搜尋

□ Fibonacci series 費氏數列

□ Combinatorial numbers 組合數

□ Towers of Hanoi 河內塔

divide and conquer 分而擊之

□ Four questions / steps

1. 遞迴定義

2. 問題簡化

3. 終止條件

4. 保證終止

心得：在一開始接觸遞迴時，常常會想到腦袋打結，會不太清楚使用遞迴的時機以及方式，經常上網看別人的方法再理解，但上過單元一之後，了解其大致的操作方式以及原理，雖然使用遞迴能讓程式更精簡，但自己還是更偏向使用迴圈。

單元二

A class

□ Attributes

• Typically

• Called

□ Behavior

• Typically

• Called

Three

□ Encap

• Object

• Hide

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單元二 資料抽象化

A class combines

□ Attributes (characteristics) of objects of a single type

- Typically data

屬性

- Called data members

□ Behaviors (operations)

- Typically operate on the data

運算

- Called methods or member functions

Three characteristics

□ Encapsulation 封裝

- Objects combine data and operations

- Hide inner details

□ Inheritance 繼承

- Classes can inherit properties from other classes

- Existing classes can be reused

□ Polymorphis 多型

- Object can determine appropriate operation at execution time.

□ Key Issues in Programming

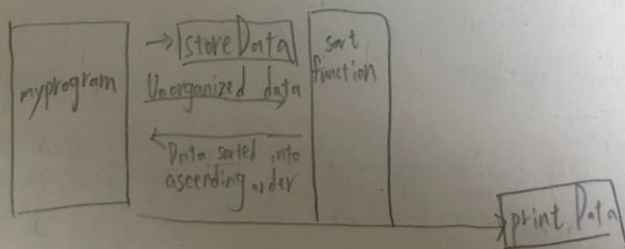
1. Modularity
2. Style
3. Modifiability
4. Ease of use
5. Fail-safe programming
6. Debugging
7. Testing

□ Cohesion - modules perform single well-defined tasks

- highly cohesive modules desired 高内聚

□ Coupling - measure of dependence among modules

- Loosely coupled modules desired 低耦合



Abstract

myprogram

Compose

□ ADT

- Create

- Destroy

- Detect

- Detect

- Insert

- Delete

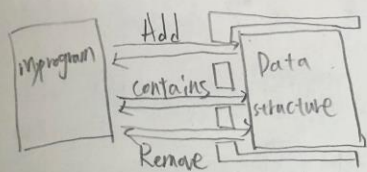
- Look

心个

有很

程立

Abstract Data Type (ADT)



Composed of $\begin{cases} \text{A collection of data} \\ \text{A set of operations on the data.} \end{cases}$

□ ADT List Operations

- Create an empty list 建構
 - Destroy a list 解構
 - Determine whether a list is empty 是否為空
 - Determine the numbers of items in a list 計算個數
 - Insert item in a given position in list 插入
 - Delete the item at a given position in list 刪除
 - Look at (retrieve) the item at a given position 檢索
- 心得: 相對於第一章, 較於理解, 且在寫程式時
有很多地方都可以使這章的內容實際使用後, 也能讓
程式看起來更整齊且容易理解。

單元3 鏈結串列

Pointers 指標 = 門牌 $\text{int}^* p$

$p = \text{new int};$ // 配置一個整數型態的空間

$*p = 1;$ // 存放值到此空間

Dynamic Allocation of Arrays 動態配置陣列

ex.

$\text{int array size} = 50;$

$\text{double}^* \text{anArray} = \text{new double} [\text{array size}];$

Struct Node

{ int item;

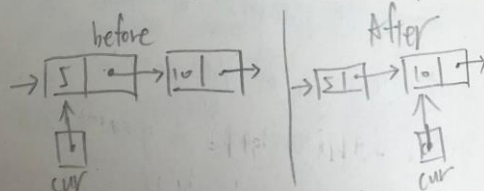
Node *next;

}; //

for (Node *cur = head; cur != NULL;

cur = cur->next)

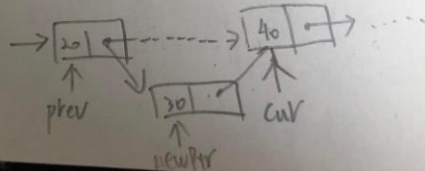
cur << cur->item << endl 走訪



新增 To insert a node between two nodes

$\text{newPtr} \rightarrow \text{next} = \text{cur};$

$\text{prev} \rightarrow \text{next} = \text{newPtr};$



A Poin

□ Public

- is Emp

- get Len

- insert

- remove

- retrieval

□ Private

- find

深層

head = 1

head →

ListNode

for (L

cur

{ new

new

new

newPtr

}

A Pointer-Based Implementation of the ADT List

□ Public Methods

- is Empty

- get Length

- insert

- remove

- retrieve

□ Private method

- find

□ Private data members

- head 串列首

- size 節數

□ Local variables to methods

- cur 現在節

- prev 前一節

深層複製

head = new ListNode;

head->item = aList.head->item;

ListNode *newPtr = head

for (ListNode *origPtr = aList.head->next;
origPtr != NULL; origPtr = origPtr->next)

{ newPtr->next = new ListNode

newPtr = newPtr->next;

newPtr->item = origPtr->item;

}
newPtr->next = NULL; }

Array-Based vs Pointer-Based

□ Size

- Increasing the size of a resizable array can waste storage and time
- Linked list grows and shrinks as necessary

□ Storage requirements

Array-Based < Pointer-Based for each item in the ADT.

□ Retrieval

檢索

The time to access the i^{th} item

- Array-Based: Constant (independent of i)
- Pointer-based: Depends on i

□ Insertion and deletion

新增/刪除

- Array-based: Requires shifting of data 搬移

- Pointer-based: Requires a traversal 走訪

心得
串列
會覺得
致不快
加速

心得：在大一計算機概論課程中，鏈結串列是讓我有印象深刻的之一，一開始學習到會覺得有些困難實做，常常會將節矣亂接而導致不知道程式碼哪裡有錯，但後期自己也越加理解運作，反而成為我其中之一擅長的單元。

單元 4 以遞迴解題

The Basics of Grammars

$\langle \text{number} \rangle = \langle \text{digit} \rangle \langle \text{number} \rangle \mid \langle \text{digit} \rangle$

$\langle \text{digit} \rangle = 0 \mid 1 \mid 2 \mid 3 \mid 4 \mid 5 \mid 6 \mid 7 \mid 8 \mid 9$

$\langle \text{addition} \rangle = \langle \text{digit} \rangle + \langle \text{addition} \rangle \mid \langle \text{digit} \rangle$

$\langle \text{digit} \rangle = 0 \mid 1 \mid 2 \mid 3 \mid 4 \mid 5 \mid 6 \mid 7 \mid 8 \mid 9$

$\langle \text{identifier} \rangle = \langle \text{letter} \rangle \mid \langle \text{identifier} \rangle \langle \text{letter} \rangle \mid \langle \text{identifier} \rangle \langle \text{digit} \rangle$

$\langle \text{letter} \rangle = a \mid b \mid \dots \mid z \mid A \mid B \mid \dots \mid Z$

□ Recognition algorithm 遞迴

`isPal (w: String): boolean`

if (w is empty or w is of length 1) 空字串 or 單一字
return true;

else if (w's first and last are same letter)
return isPal (w minus first and last)

else
return false;

Algebraic Expressions

□ Infix expressions 中序運算式

- An operator appears between its operands

□ Example: $a+b$

□ Prefix expressions 前序運算式

- An operator appears before its operands

□ Example: $+ab$

□ Postfix expressions 後序運算式

- An operator appears after its operands

□ Example: abt

$\langle \text{infix} \rangle = \langle \text{identifier} \rangle | \langle \text{infix} \rangle \langle \text{operator} \rangle \langle \text{infix} \rangle$ 中序運算式

$\langle \text{operator} \rangle = + | - | * | /$

□ Advantages of prefix and postfix expressions

- No precedence rules 優先權

- No parentheses 括弧

- No association rules 結合律

$\langle \text{pre fix} \rangle = \langle \text{identifier} \rangle \mid \langle \text{operator} \rangle \langle \text{pre fix} \rangle \mid \langle \text{pre fix} \rangle \langle \text{pre fix} \rangle$ 前序

□ Determine the end of a prefix expression

endPre (in first: integer): integer

last = strExp.length() - 1;

if (first < 0) or (last < first)

return -1;

ch = strExp[first];

if (ch is identifier)

return first

else if (ch is operator)

{ firstEnd = endPre (first+1)

if (firstEnd > -1)

return endPre (firstEnd+1);

else return -1;

else return -1;

前序

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ch = f

Delete

if (ch

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else

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前序轉成後序

convert (in: string, out: string)

ch = first character in pre

Delete first char in pre

if (ch is a lowercase letter)

post = post + ch

else

{ convert(pre, post)

convert(pre, post)

post = post + ch

}

心得：這是我第一次得知有前序和後序運算式，
所以對於這單元還不是那麼熟悉，但能夠
理解一些基礎以及如何運算，而像是將前序
轉為後序還是有些地方不懂，也只能多看多聽。