Parallel Particle-in-Cell Algorithm

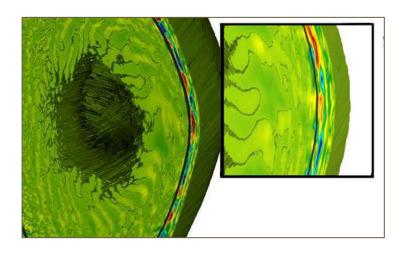
CS 205 - Final Presentation

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Motivation

Background

- Particle in Cell (PIC) is an attractive approach for simulating particle trajectories
 - N-body → O(N²) due to considering binary interactions between all particles
 - PIC → O(N) due to performing work on each particle
 only once
- Princeton Plasma Physics Laboratory develops XGC
 - PIC research code
 - o Informs design of new fusion systems
 - Planned for use on exascale computing architectures



A property heatmap inside a tokamak simulated by XGC Source: <u>insidehpc.com</u>

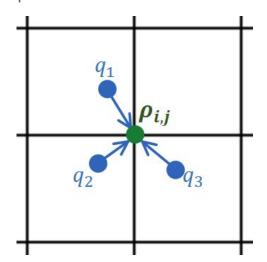
Scientific Goals

- Demonstrate capability of our PIC code to simulate plasma behavior at realistically large input values
 - Handle a large number of gridpoints → reduce spatial discretization error
 - Simulate at small timesteps → accurately resolve high frequency plasma behavior
- Serve as a proof of concept for accurately representing plasma phenomena
 - 10¹⁸ particles / m³ in real-world fusion problems

Algorithm

PIC Algorithm Overview

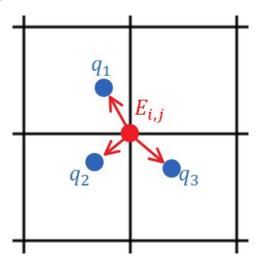
1. Interpolate from Particles to Mesh



2. Solve Discrete Poisson Equation on Mesh

$$\nabla^2 \phi = -\boldsymbol{\rho}, \qquad \nabla \phi = \boldsymbol{P}$$

3. Interpolate from Mesh back to Particles



4. Time-step Particle Locations

$$\frac{d\overline{v}}{dt} = q\mathbf{E}(\overline{x}), \qquad \frac{d\overline{x}}{dt} = \overline{\imath}$$

Data Structures

Particle: Not NUMA aware

```
Particle() {
      std :: vector < float > xx;
      std:: vector < float > xy;
     std:: vector < float > vx;
     std:: vector < float > vy;
      std:: vector < float > Ex;
     std:: vector < float > Ey;
                                                                                           \in \mathbb{R}^{N_p}
                                  P_1
                                                   P_3
                                                           P_{\Delta}
                                                                    Ps
```

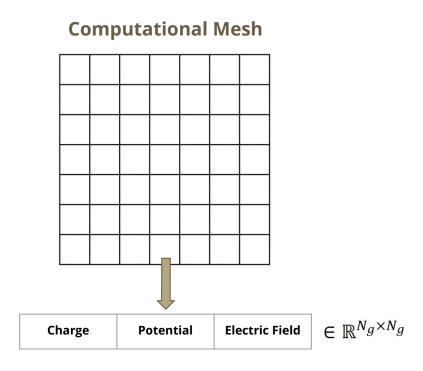
Particle* Particles = new Particle[N_p]();

Data Structures

```
Particle: Not NUMA aware
                                                                       Particle: NUMA aware
                                                               Particle(int N tin) {
Particle() {
                                                                    Nt = N tin;
     std:: vector < float > xx;
                                                                    xx = new float [Nt];
     std:: vector < float > xy;
                                                                    xy = new float [Nt];
     std:: vector < float > vx;
                                                                    vx = new float [Nt];
     std:: vector < float > vy;
                                                                    vy = new float [Nt];
     std:: vector < float > Ex;
                                                                    Ex = new float [Nt];
     std:: vector < float > Ey;
                                                                    Ey = new float [Nt];
                                                                                 \in \mathbb{R}^{N_p}
                              P_1
                                             P_3
                                                     P_{\Delta}
                                                            Ps
```

Particle* Particles = new Particle[N p]();

Data Structures

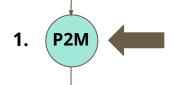


Eigen:: Matrix qGrid, pGrid, eGrid;

Parallel Implementation

Step 1: Particle-to-Mesh Interpolation

Algo. Steps



2. SOR

3. M2P

4. <u>Δt</u>

Parallel Implementation of P2M

- P2M reads from particles and writes to (shared) mesh
 - Implement thread-private grids and custom reduction
 - o Parallelize over number of particles

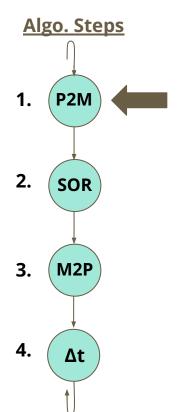
$$\phi_{M4}(x) = \begin{cases} 1 - \frac{5}{2}|x|^2 + \frac{3}{2}|x|^2 & \text{if } |x| \le 1\\ \frac{1}{2}(2 - |x|)^2(1 - |x|) & \text{if } 1 < |x| \le 2\\ 0 & 2 \le |x| \end{cases}$$

M4 Interpolation Kernel

$$\rho_d = q_p \times \phi_{M4}(x_{p,d}) \times \phi_{M4}(y_{p,d})$$

Rule to update grid point "d" with signature from particle "p"

Step 1: Particle-to-Mesh Interpolation



Parallel Implementation of P2M

- P2M reads from particles and writes to (shared) mesh
 - Implement thread-private grids and custom reduction
 - o Parallelize over number of particles

```
// perform the custom reduction to add Eigen::Matrix objects
#pragma omp parallel for reduction(MatrixPlus: chargeGrid)
   for (unsigned int p = 0; p < N_p; p++) {
        P2M_M4_kernel(particles[p], chargeGrid);
   }
}</pre>
```

Custom reduction definition and usage

Step 2: Successive Over-Relaxation

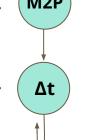
Algo. Steps



2. (SOR)

3. (M2P

4.



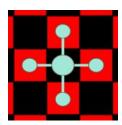
Parallel Implementation of SOR

- SOR corresponds to stencil:
 - Read from Self and Neighbors, Write to Self
- Parallelization over number of grid points in Red-Black sweeps

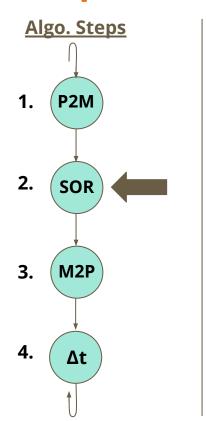
$$V_{i,j}^{k+1} = (1-\omega^k)V_{i,j}^k + \omega^k \frac{V_{i-1,j}^k + V_{i+1,j}^k + V_{i,j-1}^k + V_{i,j+1}^k + h^2\rho_{i,j}}{4}$$

SOR update rule for grid point (i, j) in iteration "k"

Example stencil for update during black sweep

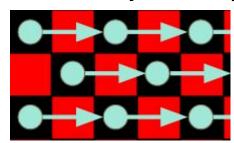


Step 2: Successive Over-Relaxation



Parallel Implementation of SOR

Parallelization of black sweep



Visualization of parallel black sweep + importance of memory layout!

Step 3: Mesh-to-Particle Interpolation

Algo. Steps P₂M **SOR** M₂P

Δt

Parallel Design of M2P

- M2P reads from shared mesh and writes to particles
- Parallelize over the number of particles

$$\vec{E}_p = \sum_{d \in \mathbb{D}_{M4}} \vec{E}_d \times \phi_{M4}(x_{p,d}) \times \phi_{M4}(y_{p,d})$$

Rule to update particle "p" with signature from surrounding grid points $d{\in}\mathbb{D}_{M4}$

Parallelization of M2P interpolation

Step 4: Time-stepping

Algo. Steps



- 2. SOR
- 3. (M2P
- 4. <u>\Data t</u>

Parallel Design of Time-stepping

- Time-step using Leapfrog Method
 - Read from and write to same particle
 - Independent per particle →parallelize over number of particles

$$\vec{v}_{t+1/2} = \vec{v}_{t-1/2} + \Delta t \frac{q}{m} \vec{E}(\vec{x}_t)$$
$$\vec{x}_{t+1} = \vec{x}_t + \Delta t \vec{v}_{t+1/2}$$

Leapfrog Method rule to update a given particle's kinematic data

Step 4: Time-stepping

Algo. Steps **P2M SOR** 3. M₂P Δt

Parallel Design of Time-stepping

- Possible load imbalance when enforcing periodic wraparounds
- Address via dynamic scheduling with large chunk size

```
#pragma omp parallel for schedule(dynamic, 500)
    for (unsigned int p = 0; p < N p; p++) {
        // timestep the velocity
        particles[p].vx[ti + 1] = particles[p].vx[ti] + constantFactorV * particles[p].Ex[ti];
        particles[p].vy[ti + 1] = particles[p].vy[ti] + constantFactorV * particles[p].Ey[ti];
        // timestep the position
        float xNew = particles[p].x[ti] + constantFactorX * particles[p].vx[ti + 1];
        float yNew = particles[p].y[ti] + constantFactorX * particles[p].vy[ti + 1];
        // enforce periodic boundary condition
        while (particleOutsideGrid) {
            periodicWraparound_x(xNew, particleOutsideGrid);
            periodicWraparound y(yNew, particleOutsideGrid);
        // write the new particle position
        particles[p].x[ti + 1] = xNew;
        particles[p].v[ti + 1] = vNew;
```

Parallel Results

Problem Size

Small Case

 10^6 particles (N_p) 400 grid points (N_g) 100 timesteps (N_t)

Large Case

 3×10^7 particles (N_p) 1500 grid points (N_g) 100 timesteps (N_t)

Memory Requirement

$$Mem(N_p, N_g, N_t) = 4 \times (N_p \times N_t \times (2 + 2 + 2) + N_g^2 \times (1 + 1 + 2)) \times 10^{-9}$$

Note: All performance benchmarks and scaling analysis run on 1 Intel Broadwell node (~126 GB memory) with the large problem size

Problem Size

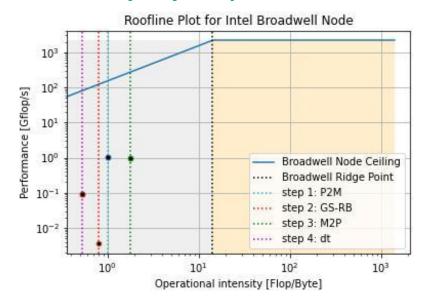
Problem Size	"Small" Case	"Large" Case
Wall clock time [hours]	0.03	1.5
Memory per node [GB]	2.4	72.1

Our "large" problem size requires 72.1 GB which is less than the Intel Broadwell node capacity, and is still computationally reasonable

Roofline Analysis: Sequential

- All four kernels are memory bound
- Gauss-Seidel largest bottleneck with performance < 0.01 Gflop/s
- Focus on analyzing M2P, because it has highest ceiling for improvement

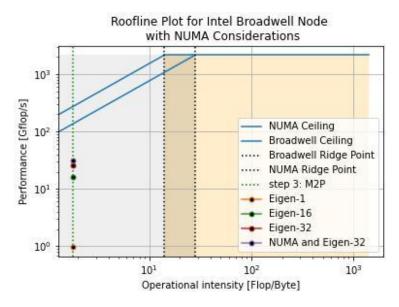
Roofline for Sequential Code



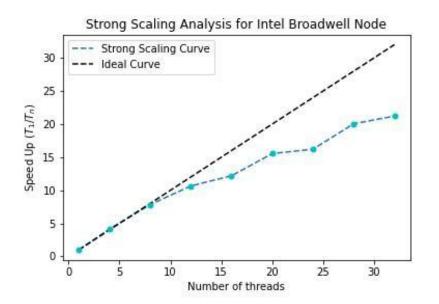
Roofline Analysis: M2P

- Improves from 1.0 Gflop/s (sequentially) to
 32.7 Gflop/s (32 threads, NUMA)
- Difference between NUMA aware and unaware is small
 - M2P infrequently interacts with the
 Particle data structure
- Other kernels (e.g. Timestepping) with more
 Particle memory accesses benefit
 significantly from NUMA

Roofline for Parallel M2P Code



Strong Scaling Analysis



Threads	$Time\ (T_n)$	Speed Up (T_1/T_n)
1	5214.2856 s.	1.0
4	1260.3978 s.	4.1
8	665.9816 s.	7.8
12	489.5262 s.	10.7
16	427.8973 s.	12.2
20	$335.5052~\mathrm{s}.$	15.5
24	321.7918 s.	16.2
28	260.1357 s.	20.0
32	245.8093 s.	21.2

Implementing shared memory parallelism with OpenMP, leveraging the Eigen library, and including NUMA first touch considerations produced a speed-up of \sim 21.2X with 32 threads

Conclusions

Key Takeaways

- Implement and optimize high-performance code using:
 - Tools: OpenMP, Eigen, Pybind11
 - Techniques: Data layout/access patterns, scheduling, effective data structures
 - Software-Hardware Interplay: Roofline Analysis, NUMA first-touch policy
- Kevin will conduct his master's thesis implementing a new style of PIC algorithm with a group at the Princeton Plasma Physics Laboratory

Future Work: Holistic Improvements

- Implement efficient I/O pipeline
 - Increase usability and efficiency
 - Explore capabilities available with HDF5 library
- Switch to SoA data structures to enable SIMD vectorization
- Migrate to MPI + OpenMP Hybrid memory model
 - Enables much larger problem sizes

Future Work: Kernel-Specific Improvements

• Particle to Mesh Interpolation:

- Leverage partially shared grids or cell list data structures (reduce memory footprint)
- Implement a higher order interpolation kernel (increase accuracy and operational intensity)

• Discrete Poisson Equation:

• Use a better algorithm such as multigrid or FFT (*increase accuracy and efficiency*)

Mesh to Particle Interpolation:

Implement a higher order interpolation kernel (increase accuracy and operational intensity)

Timestepping:

Utilize a higher order ODE Solver (increase accuracy and operational intensity)

Thank You