Effizientes Programmieren in C, C++ und Rust



Compilation Pipeline

Colin Bretl, Nikolai Maas, Peter Maucher | 24.10.2024

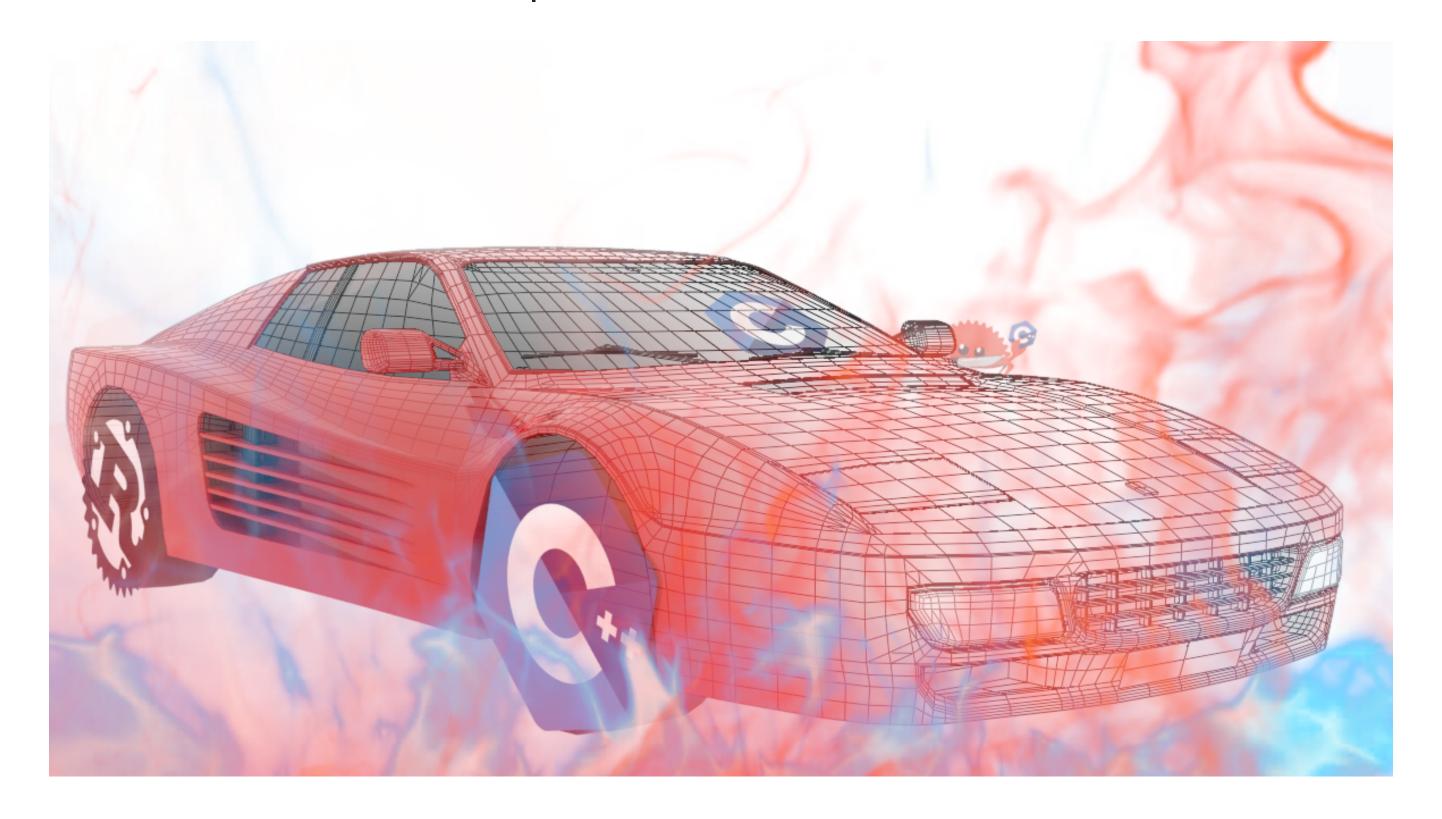


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CVSC++

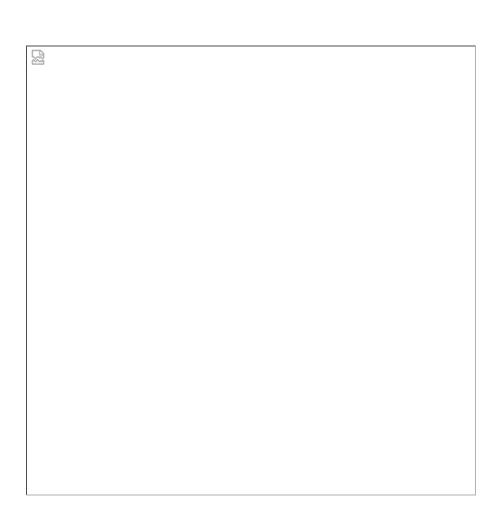


WHAT THE "PERFECT PROGRAM" SHOULD LOOK LIKE

- Knows exactly what to do (machine readable!)
- Knows exactly when to do it (in sequence, little to no branching) &
- Does only the things it needs to do (no overhead) it
- Does its job as fast as possible (processing efficiency) &
- Has everything it needs to do for its job at hand (memory efficiency) &



WHAT C++ PROGRAMS/ PROJECTS LOOK LIKE



- (Many) human-readable source file(s) (.cpp)
- (Many) human-readable header file(s) (.h)
 - #include "somelib"
 - #include <somestdlib>
- Source code, #pragma once, ...
- (Many) directories

src/, include/, lib/, tests/, doc/, .vscode/, build/

• ...





Godbolt

- Compiler Explorer https://godbolt.org/
- Developed by Matt Godbolt
- Incredibly useful to compare different compilers, compiler options, ...



Features

- different languages (C, C++, Rust, ...)
- preprocessor output
- different std versions
- different compilers
- target architectures
- libraries (e.g., boost)
- filtering options
- control flow graphs
- stack usage
- ctrl + s to save your code
- additional tools (e.g., Sonar)
- ...







The Compilation Pipeline

The Compilation Pipeline: IDE

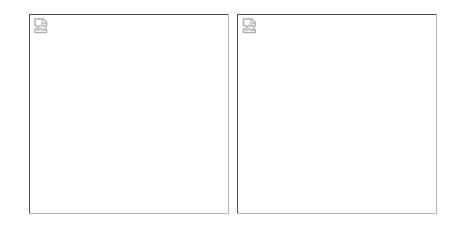


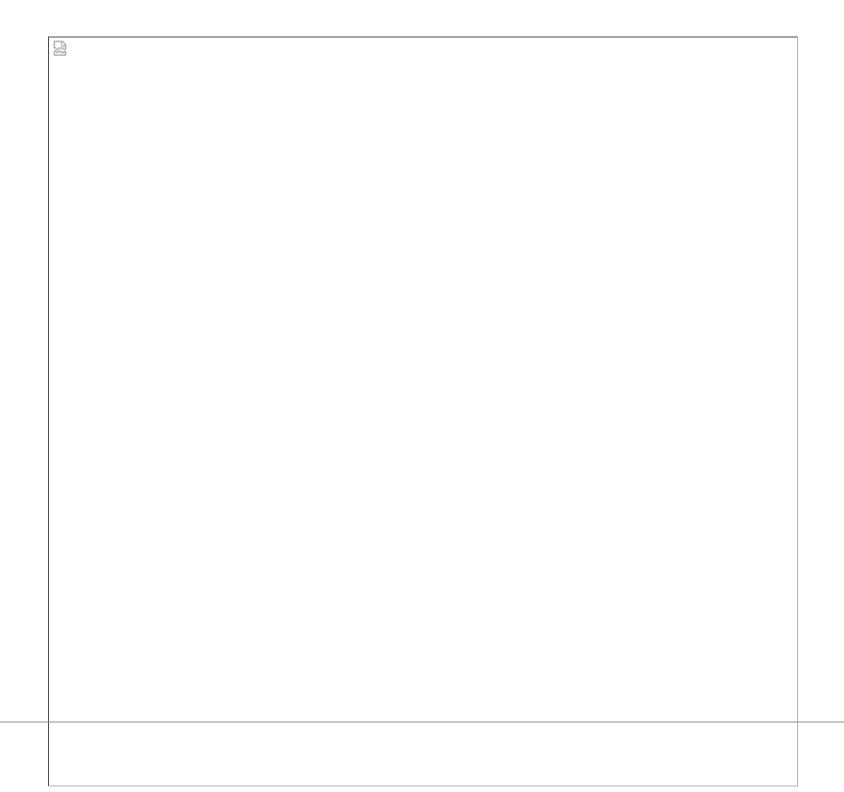
The Compilation Pipeline: IDE

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- Usually not considered part of the pipeline
- Performance aspect...?
 - Software-Engineerish reasoning
 - More important in larger projects
 - Must not slow work down!
- Examples
 - Show type of auto variables, show errors, autocompletion, fast definition lookup, tell you where you can use const
 - Pipeline orchestration (for large projects)
 - Tooling (tests, debugging, version control, coding guidelines, performance analyzer, ...)
 - E ...











- Text replacement (glorified 's/foo/bar/g')
- Driven by preprocessor directives (#)

```
#define PI 3.1415
#define circleArea(r) (PI *(r)*(r))
int main() {
    // first circle
    auto a1 = circleArea(5);
    // second circle
    auto a2 = circleArea(10);
    return 0;
}
```

```
# 0 "demo.cpp"
# 0 "<built-in>"
# 0 "<command-line>"
# 1 "/usr/include/stdc-predef.h" 1 3 4
# 0 "<command-line>" 2
# 1 "demo.cpp"

int main() {
    auto a1 = (3.1415 *(5)*(5));
    auto a2 = (3.1415 *(10)*(10));

    return 0;
}
```



Use Cases in C

- Symbolic constants: #define PI 3.1415
- (Pre-defined) function-like Macros
 #define circleArea(r) (PI *(r)*(r))
- (Recursive) header file inclusions: #include
- Include guards: #pragma once
- Conditional compilation: #if _WIN32 ... #else ... #endif
- Passing configuration options
- Compatibility macros

• ...

Why did (do...) people use function-like macros?

• Inlining!



```
#define PI 3.1415
#define circleArea(r) (PI *(r)*(r))
int main() {
    // first circle
    auto a1 = circleArea(5);
    // second circle
    auto a2 = circleArea(10);
    return 0;
}
```

```
g++ -S demo.cpp
circleArea(double):
    [10 lines omitted]
    ret
main:
    [3 lines omitted]
   mov rax, QWORD PTR .LC1[rip]
   movq xmm0, rax
    call circleArea(double)
   movq rax, xmm0
   mov QWORD PTR [rbp-8], rax
   mov rax, QWORD PTR .LC2[rip]
   movq xmm0, rax
    call circleArea(double)
   movq rax, xmm0
   mov QWORD PTR [rbp-16], rax
    [2 lines omitted]
                               Potentially expensive
    ret
    .LCO: ... [2 lines omitted] function calls!
    .LC1: ... [2 lines omitted]
    .LC2: ... [2 lines omitted]
```



Calling a function (calling convention) requires...

- ...pushing the return address on the stack
- ...pushing arguments onto the stack/ registers
- ...jumping to the function body
- ...executing the function
- ...executing a return instruction when finished

```
g++ -S demo.cpp
circleArea(double):
    [10 lines omitted]
    ret
main:
    [3 lines omitted]
   mov rax, QWORD PTR .LC1[rip]
   movq xmm0, rax
    call circleArea(double)
   movq rax, xmm0
   mov QWORD PTR [rbp-8], rax
   mov rax, QWORD PTR .LC2[rip]
   movq xmm0, rax
    call circleArea(double)
   movq rax, xmm0
   mov QWORD PTR [rbp-16], rax
    [2 lines omitted]
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```



```
#define PI 3.1415
#define circleArea(r) (PI *(r)*(r))
int main() {
    // first circle
    auto a1 = circleArea(5);
    // second circle
    auto a2 = circleArea(10);
    return 0;
}
```

```
main:
    push rbp
    mov rbp, rsp
    movsd xmm0, QWORD PTR .LC0[rip]
    movsd QWORD PTR [rbp-8], xmm0
    movsd xmm0, QWORD PTR .LC0[rip]
    movsd QWORD PTR [rbp-16], xmm0
    mov eax, 0
    pop rbp
    ret
.LC0: ... [2 lines omitted]
.LC1: ... [2 lines omitted]
```

- For 3 multiplications, function call overhead seems to be disproportionate
- Macro forces area caluation "in line" (inlining)
- Macros have certain disadvantages \
- Side quest: what type do a1 and a2 have?



Use Cases in C

- Symbolic constants: #define PI 3.1415
- (Pre-defined) function-like Macros & #define circleArea(r) (PI *(r)*(r))
- (Recursive) header file nclusions: #include
- Include guards: #pragma once
- Conditional compilation:
 #if _WIN32 ... #else ... #endif
- Passing configuration options
- Compatibility macros

• ...

SYMBOLIC CONSTANTS

- constexpr double PI = 3.1415;
- Declares that the evaluation of the variable (or function) is possible at compile time
- Wherever compile time constant expressions are allowed
- Implies constness

HEADER FILE MANAGEMENT

See C++20 modules ↑

CONDITIONAL COMPILATION

- if constexpr (DEBUG_MODE) { ... }
- Discards branches of an if statement at compile time
- Since C++17



Use Cases in C

- Symbolic constants: #define PI 3.1415
- (Pre-defined) function-like Macros & #define circleArea(r) (PI *(r)*(r))
- (Recursive) header file nclusions: #include
- Include guards: #pragma once
- Conditional compilation ### ### _WIN32 ... #else ... #endif
- Passing configuration options
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INLINING VIA inline...?

- "The original intent of the inline keyword was to serve as an indicator to the optimizer that inline substitution of a function is preferred over a function call [...]"
- So, let's inline double circleArea(double r) { ... }



```
demo.cpp

#define PI 3.1415

inline double circleArea(double r){Tell the compiler that you return PI * r * r; think inlining is beneficial

int main() {
    // first circle
    auto a1 = circleArea(5);
    // second circle
    auto a2 = circleArea(10);
    return 0;
}
```

```
g++ -S demo.cpp
circleArea(double):
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   movq rax, xmm0
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   mov rax, QWORD PTR .LC2[rip]
   movq xmm0, rax
   call circleArea(double)
   movq rax, xmm0
   mov QWORD PTR [rbp-16], rax
    [2 lines omitted]
                               Still function calls!
    ret
    .LC0: ... [2 lines omitted]
    .LC1: ... [2 lines omitted]
    .LC2: ... [2 lines omitted]
```



```
demo.cpp

#define PI 3.1415

double circleArea(double r){
    return PI * r * r;
}

int main() {
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    return 0;
}
```

```
main:
    push rbp
    mov rbp, rsp
    movsd xmm0, QWORD PTR .LC0[rip]
    movsd QWORD PTR [rbp-8], xmm0
    movsd xmm0, QWORD PTR .LC0[rip] Compiler inlines!
    movsd QWORD PTR [rbp-16], xmm0
    mov eax, 0
    pop rbp
    ret
.LC0: ... [2 lines omitted]
.LC1: ... [2 lines omitted]
```

No actual compiler output. Used for demonstration purpose only and to build up a meme. Your compiler is even smarter.'



```
demo.cpp

#define PI 3.1415

inline double circleArea(double r) {Let the compiler happily return PI * r * r; ignore your humble opinion :-)

int main() {
    // first circle
    auto a1 = circleArea(5);
    // second circle
    auto a2 = circleArea(10);
    return 0;
}
```

- The original intent of the inline keyword was to serve as an indicator to the optimizer that inline substitution of a function is preferred over function call, [...].
 - Compiler w/o optimization does not inline (debugging)
 - Optimizing compilers decide on their own whether to inline a function call or not (both calls are good candidates here)
 - inline is almost certainly ignored
 - Takeaway: Inlining is good. Inline is not inlining.
 - inline __attribute__((alwys_inline))
 serves as proper hint to the compiler (widely
 supported)
- Let's run our example with an actual optimizing compiler now...



```
#define PI 3.1415

inline double circleArea(double r){
    return PI * r * r;
}

int main() {
    // first circle
    auto a1 = circleArea(5);
    // second circle
    auto a2 = circleArea(10);
    return 0;
}
```

```
g++ -S -03 demo.cpp

main:
    xor eax, eax
    ret
```



INTERPRETED

JUST-IN-TIME COMPILATION







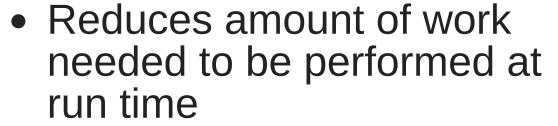
 Compilation happens before execution

- Executes instructions written in a another programming or scripting language
- Continuosly analyses running code, (re-)compiles to gain speedup

immediately executes it

Translates bytecode to

machine code and



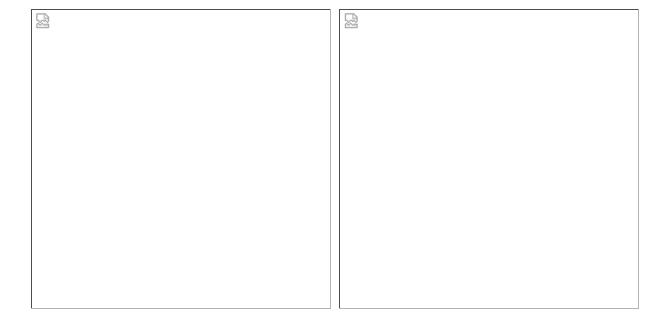


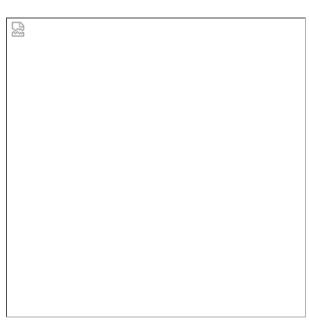






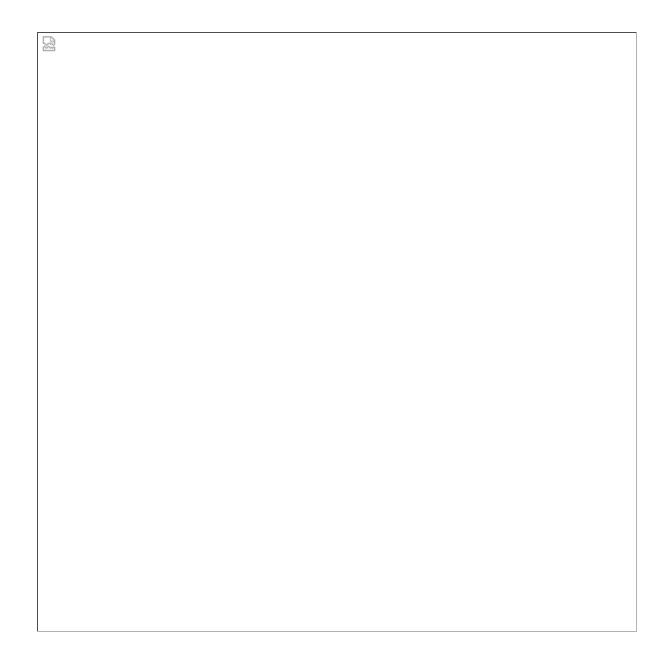
- 1. What ProPa teaches you, i.e.,
 - 1. Lexical analysis (turn sequence of symbols into tokens and maintain identifiers)
 - 2. **Syntactical analysis** (tell if sequence of tokens is a valid word of the context-free language implied by C++)
 - 3. Semantic analysis (context-sensitive analysis of names introduced by declarations, types and language rules)
- 2. Produce intermediate code, report meaningful errors





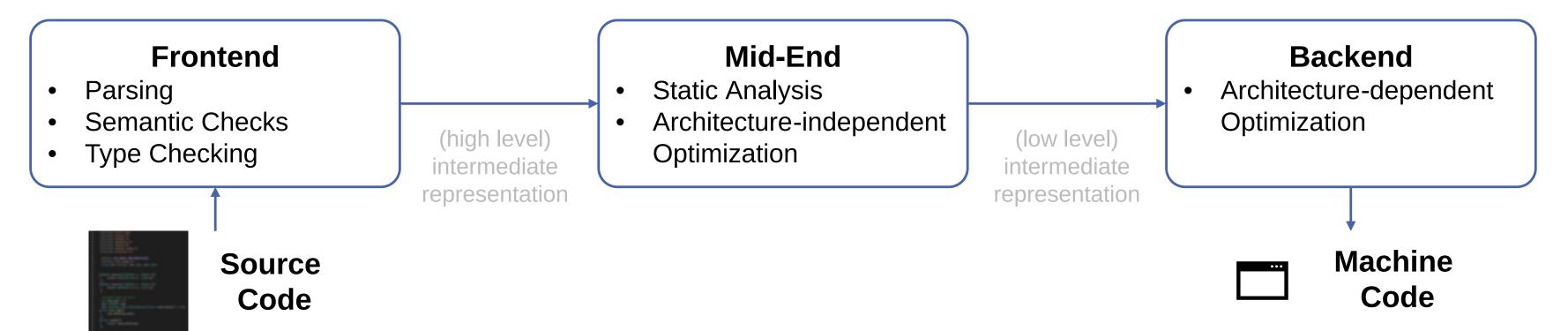


- 1. What ProPa teaches you, i.e.,
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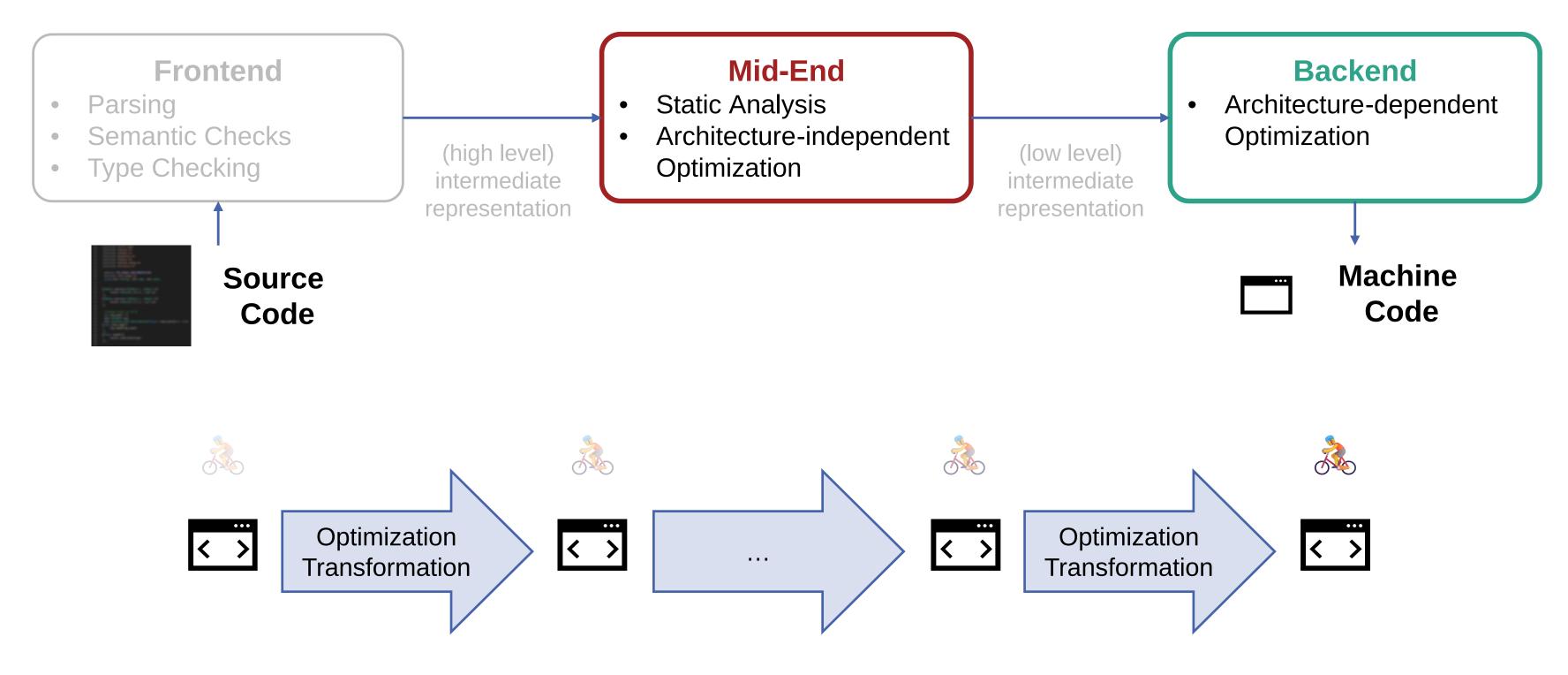
Hier schon :-)





- Mostly irrelevant for optimization
- But you can still have fun with this if that's your cup of tea!
- C++: Monomorphization /







- Output must be semantically equivalent to the input (as-if rule)
- C++: Compiler is permitted to perform any changes to the program as long as the "observable behaviour" does not change
 - E.g., regrouping operators (reductions)
 - auto dummy = (b + c); auto assoc = (a + b) + c
 - Ok if a, b and c are int, forbidden if they are float
- No Elder Wand! (but a little bit like a magic wand occasionally)



- Output must be semantically equivalent to the input (as-if rule)
- Exceptions to the as-if rule
 - Undefined behaviour /
 - Copy elision
 - Return value optimization



- Output must be semantically equivalent to the input (as-if rule)
- Exceptions to the as-if rule
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- These are **not** semantically equivalent if creation/ destruction of BLOB has sideeffects (custom types!)
 - Explicitly allowed by the standard
 - To disable with gcc: -fno-elide-constructors



Local vs. global

```
double r = 5;

double a = circleArea(r);

int a2 = static_cast<int>(a)* 2;
return a2;
}
```

- All that happens here is transforming (compile-time) constants
- Same as return 157; (not refactoring friendly)
- But we cannot necessarily make this assumption if we don't know what circleArea does!



- Local vs. global
 - Inlining to the rescue!
 - Enables further compiler optimization!
 - Decided per call
- And much, much more! ///
- Modern compilers excel at this!

```
double r = 5;

double a = circleArea(r);

int a2 = static_cast<int>(a)* 2;
return a2;
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```



Taming the Beast

- Controlling optimization (gcc)
- -O0: default, no optimization, faster compilation time
- -O1: turn on optimization (predefined set of optimization flags)
- -O2: even more optimization (more flags)
- -O3: even more optimization (more flags)
- -Os: optimize for binary size
- Ofast: disregards strict standards compliance (e.g., non-associativity of float). You have been warned!
- Fine-tuning with specific options
 - -fno-inline
 - -flop-parallelize-all
 - -fwhole-program

•

What could possibly go wrong?



- Takes preprocessor output as input
- Calling the compiler produces an object file (.0)
- Objects (functions, global variables) have symbolic addresses
- Object files can refer to symbols, that are not defined (declaration w/o definition)
- Different "blocks of code" (separate .cpp files) are compiled separately
- Blocks of code = Translation Units
 - Single source file (.cpp)
 - + any header (in-/)directly included (.h)
 - lines ignored using conditional preprocessing





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Taming the Beast

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Header Guards

Explicitly

```
#ifndef MY_UNIQUE_NAME
#define MY_UNIQUE_NAME

// code goes here

...
// endif
```

Most compilers support #pragma once



Taming the Beast

Remaining Problem

- Takes preprocessor output as input
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Remaining Problem





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- Compiler converts source code (.c, .cpp) into (machine-readable) object code (.o)
- Object files are not executables
- Example on the right:
 - main.o is from valid C++
 - main.o knows that there is some void log() (and how to call it)
 - main.o does not know, where to find what void log() does (missing binding)
- Linker takes object files and combines them into a single executable (or library file or another object file)
- Maps symbolic addresses to (proper) memory addresses (can involve relocating code)



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Static Linking (.a, .lib)

 All objects/ library functions (that are used) are copied into the executable



Dynamic Linking (.so, .dll)

- Executable contains...
 - ...undefined symbols
 - …list of all objects/ libraries that provide definitions
- Resolving undefined symbols is deferred until runtime





Static Linking (.a, .lib)

- All objects/ library functions (that are used) are copied into the executable
- Advantages
 - Prevents "DLL hell"
 - Not the entire library must be installed if only a few functions are actually needed
 - Portability
- Disadvantages
 - ((Requires more memory))
- Static linking allows link-time optimization

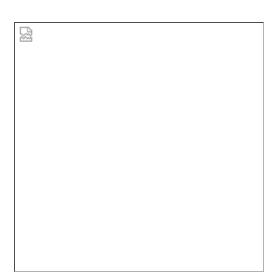
Dynamic Linking (.so, .dll)

- Executable contains...
 - ...undefined symbols
 - …list of all objects/ libraries that provide definitions
- Resolving undefined symbols is deferred until runtime
- Advantages
 - Often-used libraries (e.g., standard system libraries) are not duplicated
 - No re-linking after, e.g., a bug fix necessary
- Disadvantages
 - Incompatible update libraries can break executables (DLL hell)

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Link-time Optimization

- Compiler requires holistic view of involved files to do certain optimizations
 - Inlining across files
 - Removal of unused code across files
- At best: same performance/ optimizations as without linking
- Additionally: information on final object layout can allow micro optimizations (e.g., for jump instructions)







- Why even bother?
- Before execution, a program must be loaded into main memory
- If need be, dynamic linking must be performed
- Binary loaders require absolute binary form of a program (always located in the same memory area)
- Relocating loaders translate addresses
 - Set a relocation bit to each machine language instruction
 - Group all information in a relocation table
- Run program by jumping to the first instruction (main)

The Compilation Pipeline: Summary



-	





- API (Application Programming Interface)
 - Source code level
 - Types, structures, functions, ... used to access functionality of external components (libraries, OS functionality, ...)
 - Which functions are part of the library
 - E.g., order of arguments to a function
- ABI (Application Binary Interface)
 - How compiled types, structures, function, ...
 (binary data) are accessed
 - How code is stored inside the library
 - E.g., how arguments to a function are passed (registers used, stack, ...)
- Whenever ABI changes, all programs using the affected library must be recompiled



- ABI specifies how C++ compilers compile
 - Example 1: data layout

```
class StdLibraryObj {
private:
    char c;
    int* p;
public:
    void foo(); // API
};
int bar(StdLibraryObj lo) { ... }
int main(){ // App
    return bar(...);
```



- ABI specifies how C++ compilers compile
 - Example 1: data layout

```
class StdLibraryObj {
private:
    char c;
    int* p;
    float f; // new private member
public:
    void foo(); // API does not change!
};
int bar(StdLibraryObj lo) { ... }
int main(){ // App
    return bar(...);
```



- ABI specifies how C++ compilers compile
 - Example 1: data layout
 - Example 2: calling conventions

- Low-level rules for function calls
- Calling conventions define...
 - ...where arguments and return values are stored (stack/ registers)
 - ...what is the ordering of arguments
 - ...who is responsible for cleaning up after return
- Performance consideration
 - return value (or pointer to it) in register, arguments on stack
 - vs. return value (or pointer to it) in register and (first few) arguments on stack
- Prefer register-passed arguments!





- ABI specifies how C++ compilers compile
 - Example 1: data layout
 - Example 2: calling conventions
 - Example 3: std::unique_ptr (register-passed arguments)

Consider

```
int deref(std::unique_ptr<int> ptr){
   *ptr = 7; // assign 7 to pointee
}
```

 What dereferencing a std::unique_ptr should look like:

```
movDWORD PTR [rdi], 7Keep pointee (int) in register (fast 3%)
```

- What dereferencing a int looks like:
 mov rax, QWORD PTR [rdi], 7
 mov QWORD PTR [rdi], 7
 - Keep reference to stack (pointee) in register
 - Required for "non-trivial" objects in C++ by standard
- Must not mistake pointee and reference (otherwise: random access of stack address!)
- Smart pointers are still preferred for many other reasons

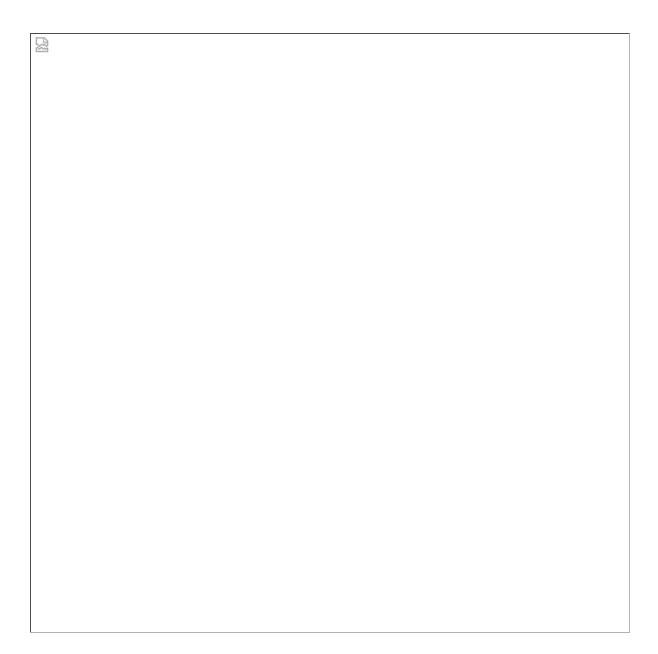


- ABI specifies how C++ compilers compile
 - Example 1: data layout
 - Example 2: calling conventions
 - Example 3: std::unique_ptr (register-passed arguments)
 - Stack layout and register usage
 - Name mangling
 - Managing template instances
 - Passing of hidden function parameters (this)
 - Throwing and catching exceptions
 - Other standard library details
 - E ...

- In hindsight, some standard library features that are ABI dependent cannot be changed
- see std::unique_ptr example
- Maybe there is an ABI break one day...



- There are changes (quality of life, performance, ...) that could be unlocked by an ABI break
 - std::regex
 - container layouts
 - fit std::unique_ptr in register
 - · ...
- There are problems that could arise from an ABI break
 - Source code lost, only object files from (now bankrupt third party) exist





C ABI

- No "de-jure" standard but well-defined on a specific OS and architecture (de-facto standard, e.g., System V)
- Target triplet: architecture, OS, compiler
- "Lingua franca" of programming languages

C++

- Faces more "ABI design decisions" (name mangling, ...)
- Guarantees C ABI for subset of C++ compatible with C
- extern "C" int f(...) {...} to make function f explicitly have C linkage (for plain old datatypes, PODs)

Other Languages (e.g., Python)

- Internal (not necessarily stable) ABI for intralanguage calls
- External ABI for inter-language calls (C ABI)
- For a deeper dive from a Python POV: https:// azhpushkin.me/posts/python-c-under-the-hood





- Configuring the build pipeline for large projects is tedious (debug version(s), release version(s), different target machines, multiple libraries, many source files, testing, corresponding header files, linking everything together,...)
- Common Issues
 - Hard-coded paths
 - No continuous integration
 - No cross-platform support
 - Incomprehensible project structure
 - · ...



- Build systems to the rescue!
 - GNU Make on Linux (Makefiles)
 - Visual Studio Generator on Windows (VS projects)
 - Platform/ compiler-dependent
- Meta build systems to the rescue from the rescue!
 - CMake is de-facto standard for C/C++
 - Supported by most/ all (relevant) IDEs (even VS)
 - Modern CMake (3.5+, 3.29+, any version that came out after your compiler)



Sample Project Structure



Sample Project Structure

```
project/
  CMakeLists.txt
  src/
    engine.cpp
    utils.cpp
  include/
                    // -> /usr/include
    project/
                    // library
      engine.h
      utils.h
  apps/
                    // int main(){...}
    app1.cpp
  .gitignore
                    // /build*
  REAMDE.md
  LICENSE.md
                    // helper modules
  cmake/
                    // documentation
  docs/
  tests/
                    // tests
  scripts/
                    // helper.py
                    // git submodules
  extern/
                    // python bindings
  python/
```

project/CMakeLists.txt

```
#this is a comment
cmake_minimum_required(VERSION 3.29)
#project name, useful settings
project(
  ExampleProject
  VERSION 0.1
  DESCRIPTION "An example project"
  LANGUAGES CXX)
#for main project: set C++ standard
if(CMAKE_PROJECT_NAME STREQUAL PROJECT_NAME)
  set(CMAKE_CXX_STANDARD C++17)
endif()
#add Boost::boost
#cmake -help-module-list shows all modules
find_package(Boost) #find_package(Boost REQUIRED)
if (NOT Boost_FOUND)
 # show output
  message(WARNING "Boost not found")
endif()
add_subdirectory(src)
add_subdirectory(apps)
```



Sample Project Structure

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                    // python bindings
  python/
```

project/src/CMakeLists.txt

```
#find header files by globbing expression
#globbing expressions are simplified regexes
file(GLOB HEADER_LIST CONFIGURE_DEPENDS
"${ExampleProject_SOURCE_DIR}/include/project/*.h")

#build library (dynamic/ static based on user #setting)
as a first target
add_library(my_lib engine.cpp utils.cpp ${HEADER_LIST}

#include
target_include_directories(my_lib PUBLIC ../include)

#depend on other (header-only) library
target_link_libraries(my_lib PRIVATE Boost::boost)
```



Sample Project Structure

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 CMakeLists.txt
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                  // -> /usr/include
   project/
             // library
     engine.h
     utils.h
  apps/
             // int main(){...}
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   CMakeLists.txt
                   // /build*
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  REAMDE.md
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                   // helper modules
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                   // tests
 scripts/
                   // helper.py
                   // git submodules
  extern/
                   // python bindings
  python/
```

project/apps/CMakeLists.txt

```
add_executable(my_app app1.cpp)
#optional pre-processor definitions
option(USE_EXPERIMENTAL "Enable experimental
stuff" OFF)
#(platform-dependent) compile options
if(CMAKE_SYSTEM_NAME STREQUAL "Windows")
  target_compile_options(my_app PRIVATE /W4)
elseif(CMAKE_SYSTEM_NAME STREQUAL "Linux")
  target_compile_options(my_app PRIVATE -Wall -
                         Wextra -Wpedantic)
endif()
link with previously targeted library
#target_link_options(my_app ...)
target_link_libraries(app1 PRIVATE my_lib
  fmt::fmt)
```



Sample Project Structure

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  extern/
                    // python bindings
  python/
```

Build Project

- Run cmake from CMakeLists.txt in root directory and build in current directory
 - cmake -B build && cmake --build build
- cmake -DUSE_EXPERIMENTAL=ON ...
- cmake -DCMAKE_BUILD_TYPE=Debug ...
- cmake -DCMAKE_BUILD_TYPE=Release ...
- CMakeCache.txt
 - CMake caches user configurable options after first build configuration
 - cmake --target clean
- CMake also features testing, ...



Takeaways

Takeaways



- Compiler turns human-readable code (.cpp files, .h files) into machine-readable code
- Must obey as-(almost)-if rule
- Compilation pipeline consists of
 - (IDE)
 - Preprocessor
 - Compiler: Frontend, Mid-End, Backend
 - Linker
 - (Loader)
- Compiler and Linker are capable of compiletime and link-time optimizations

- CMake as meta-build-system to automate the build process
- ABI stability limits optimization potential (but an ABI break is not trivial)



Thank You!