

Leet Code Weekly 262

Grey Wang

1. Two Out of Three

Description

Given three integer arrays, return a **distinct** array containing all the values that are present in **at least two** out of three arrays. You may return the values in **any order**.

Input

nums1 = [1, 1, 3, 2]

nums2 = [2, 3]

nums3 = [3]

Output

[3, 2]

Constraints

- $1 \leq \text{nums1.length}, \text{nums2.length}, \text{nums3.length} \leq 100$
- $1 \leq \text{nums}[i] \leq 100$

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如何解一道题？

1. 读题
2. 理解数据
3. 数据规模？
4. 思路？
5. 复杂度？
6. Coding

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思路1 朴素枚举 $O(n^2)$

has(nums, i):

for num in nums:

if num == i:

return 1;

return 0;

for i := 1..100:

if has(nums1, i) + has(nums2, i) + has(nums3, i) >= 2:

i in result

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思路2 预处理(哈希) + 枚举 $O(n)$

make_hash(nums):

h = [0 for i in 1..100]

for i in nums:

h[i] = 1;

return h;

h1 := make_hash(nums1)

h2 := make_hash(nums2)

h3 := make_hash(nums3)

for i := 1..100:

if h1[i] + h2[i] + h3[i] >= 2:

i in result

2. Minimum Operations to Make a Uni-Value Grid

Description

Given a 2D integer grid of $m * n$ and an integer x . In an operation, you can **add** x to or **subtract** x from any element in the grid.

A **uni-value grid** is a grid where all the elements of it are equal.

Return the **minimum** number of operations to make the grid **uni-value**. If it is not possible, return -1.

Input

grid = [[2, 4], [6, 8]]

x = 2

Output

4

Constraints

- $m == \text{grid.length}$
- $n == \text{grid}[i].\text{length}$
- $1 \leq m, n \leq 10^5$
- $1 \leq m * n \leq 10^5$
- $1 \leq x, \text{grid}[i][j] \leq 10^4$

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解题的常见角度？

1. 题目类型？ 筛选常见策略

最优解问题：搜索，动态规划，枚举答案，贪心。。。

2. 数据规模？ 排除一些策略。

10^5 个元素，搜索不太可能。。

10^4 的值域 + 10^5 个元素，状态空间太大，DP不太可能。。

枚举 10^4 的值域 * 10^5 个元素的验证，枚举答案不太可行。。

要不试试贪心？

3. 问题简化？

二维跟一维有区别吗？ 没有。

4. 边界条件？ 由简单再到一般化。

$[a] \Rightarrow 0$

$[a, b] \Rightarrow |b - a| / x$

5. 构造贪心策略？

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思路 中位数 $O(N^{\log N})$

```
nums := sort(flatten(grid))
```

```
uni_val := nums[nums.size() / 2]
```

```
ans := 0
```

```
for t in nums:
```

```
    if abs(uni_val - t) % x != 0:
```

```
        impossible!
```

```
    else
```

```
        ans += abs(uni_val - t) / x
```


3. Stock Price Fluctuation

Description

Given a stream of **records** about a particular stock. Each record contains a **timestamp** and the corresponding **price** at that timestamp. Some records may be incorrect, another record may appear later in the stream **correcting** the price of the previous wrong record.

Design the **StockPrice** class that:

- ***update(int timestamp, int price);***
- ***int current();** // Returns the **latest price** of the stock*
- ***int maximum();** // Returns the **maximum price***
- ***int minimum();** // Returns the **minimum price***

Constraints

- **$1 \leq \text{timestamp}, \text{price} \leq 10^9$**
- **At most 10^5 calls made in total.**
- current, maximum, and minimum will be called **only after** update has been called at least once.

1. 题目类型?

纯数据结构题。

2. timestamp => price 映射关系, Map

3. maximum / minimum 优先队列?

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Constraints

- **1 <= timestamp, price <= 10⁹**
- **At most 10⁵ calls made in total.**
- current, maximum, and minimum will be called **only after** update has been called at least once.

思路1 平衡树 + 优先队列 $O(n * \log(n))$

last_timestamp := -INF

prices := map<int, int>()

q_max = max_priority_queue<record>()

q_min = min_priority_queue<record>()

update(timestamp, price):

 last_timestamp := max(last_timestamp, timestamp)

 prices[timestamp] = price;

 q_max.insert(record(timestamp, price))

 q_min.insert(record(timestamp, price))

current():

 prices[last_timestamp]

maximum():

 while record = q_max.pop()

 if prices[record.timestamp] == record.price:

 return record.price

minimum():

 while record = q_min.pop()

 if prices[record.timestamp] == record.price:

4. Partition Array Into Two Arrays to Minimize Sum Difference

Description

Given an integer array `nums` of $2 * n$ integers. You need to partition `nums` into two arrays of length `n` to **minimize the absolute difference** of the **sums** of the arrays. To partition `nums`, put each element of `nums` into one of the two arrays.

Return the **minimum** possible absolute difference.

Input

[3, 9, 7, 3]

Output

$|(3 + 9) - (7 + 3)| = 2$

Constraints

- $1 \leq n \leq 15$
- `nums.length == 2 * n`
- $-10^7 \leq \text{nums}[i] \leq 10^7$

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2. 数据规模？ 排除一些策略。

$15 * 2$ 个元素，大概率可以搜索。。

$2 * 10^7$ 的值域 * 2^{30} 的集合状态，DP不太可能。。

枚举答案？ 不确定。。

贪心？ 不确定。。

3. 问题简化？

选定 n 个元素的和为 `now`, 那么另一半元素为 `sum - now`, 对应当前解为: $|\text{sum} - 2 * \text{now}|$

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- $1 \leq n \leq 15$
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思路1. 朴素搜索 $O(2^{2n}) \rightarrow 2^{30} = 2 * 10^9$ 超时

```
n := nums.size() / 2
```

```
sum := nums.sum()
```

```
ans := +INF
```

```
dfs(now, cnt, i):
```

```
    if cnt == n:
```

```
        ans := min(ans, abs(sum - now * 2))
```

```
    else if i < 2 * n:
```

```
        if cnt < n:
```

```
            dfs(now + nums[i], cnt + 1, i + 1)
```

```
            dfs(now, cnt, i + 1)
```

```
dfs(0, 0, 0)
```

Tip: 算法比赛里, 10^8 一般可视为1秒

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思路2. 双向搜索 + 二分查找 $O(2^n * \log(2^n)) \rightarrow O(n * 2^n)$

Tip: 降低搜索深度

```
n := nums.size() / 2
sum := nums.sum()

dfs(now, cnt, i, nums, &values):
    values[cnt].add(now)
    if i < n:
        if cnt < n:
            dfs(now + nums[i], cnt + 1, i + 1, nums, &values)
        dfs(now, cnt, i + 1, nums, &values)
```

```
l_values = [{} for i in 0..n]
r_values = [{} for i in 0..n]
dfs(0, 0, 0, nums, l_values)
dfs(0, 0, 0, nums + n, r_values)

ans := min(abs(sum - 2 * l_values[0][0]), abs(sum - 2 * r_values[0][0]))
for i in 1..(n - 1):
    for l_value in l_values[i]:
        find nearest r_value to (sum / 2 - l_value) in r_values[n - i]
        ans := min(abs(sum - l_value - r_value))
```

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思路3. 状态压缩 + 双向搜索 + 二分查找 $O(n * 2^n)$

```
n := nums.size() / 2
sum := nums.sum()

l_values = [{} for i in 0..n]
r_values = [{} for i in 0..n]
for mask in 0..(1 << n):
    cnt := l_value := r_value := 0
    for i in 0..n:
        if mask & (1 << i):
            cnt += 1
            l_value += nums[i]
            r_value += nums[i + n]
    l_values[cnt].add(l_value)
    r_values[cnt].add(r_value)

ans := min(abs(sum - 2 * l_values[0][0]), abs(sum - 2 * r_values[0][0]))
for i in 1..(n - 1):
    for l_value in l_values[i]:
        find nearest r_value to (sum / 2 - l_value)
```

Tip: 状态压缩，将集合状态表示成比特位的形式

例：{3, 9, 7, 3} 选取 {3, 9}
可表示为：1100 = 3