

# End-to-End RSA Encryption with Lightweight Key Management in Edge Computing

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April 12, 2025

## ABSTRACT

Today, edge computing, also known as the edge, continues to grow in the market as technology continues to evolve to require higher efficiency, meaning faster computation and more power, and the edge not only allows devices to reduce their own workload, but also to offload from cloud computing and distribute it to multiple edge servers. This saves computing power and resources for both devices and the cloud, meaning longer battery life for devices and lower costs for cloud resources. However, despite the great potential of the edge for future technology, its implementation introduces security risks.

Due to the distributed nature of edge devices that continuously collect, process, and store large amounts of potentially sensitive and/or confidential data, the attack surface increases as more edge devices are in use. This paper aims to address a list of security concerns with edge computing and to demonstrate a solution using RSA encryption and digital signatures to protect and validate data. In future works, the solution can be further improved to more solidly resolve these concerns.

## 1. INTRODUCTION

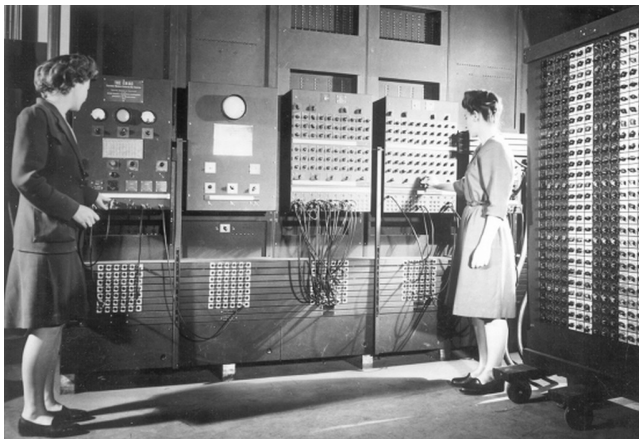


Figure 1: Operation of the ENIAC's control panel [6]

### 1.1 A Brief History of Computers

Computers have come a long way since the world's very first general-purpose electronic computer on the 14<sup>th</sup> of February 1946. The Electronic Numerical Integrator and Computer

(ENIAC) was a very large computer with a room-sized main-frame that was constructed at Moore's School of Electrical Engineering, now called Penn's School of Engineering and Applied Science, and it was one of the first technologies to electronically compute complex mathematical problems at what was considered very high speeds at the time.

The amount of time that it would take even the smartest people to solve complex math equations with the help of function tables would be exponentially reduced when using the ENIAC which could perform 5000 additions per second. The amazement of the increase in efficiency in its time inspired the revolution of computer science and electrical engineering that continues today [6]. Figure 1 shows an old picture of Jean Bartik on the left and Frances Spence operating the ENIAC's main control panel.

Computers today are much smaller and mobile, especially laptops and smartphones, and users can do so much more than perform mathematical calculations; people can communicate through emails, use social media, get the news, research, write documents and presentations, develop and/or advertise products, simulate, and program and operate machines, you name it. With innovation, technology has evolved to be able to efficiently run multiple processes and perform numerous tasks, and part of this is with the help of edge computing.

### 1.2 Evolution of Edge Computing

Just like computers, edge computing was also part of the technology evolution. Edge computing began in the 1990s with Akamai's launch of its content delivery network (CDN), which it allowed static media to be cached closer to end users, and in 1997, Brian D. Nobel along with numerous contributors wrote a paper demonstrating a prototype model that allows applications to transfer certain tasks to strong servers when operating on resource-scarce mobile devices. This ability to offload tasks to servers in a heterogeneous distributed manner has become a main principle for technologies from businesses such as Google, Apple, and Amazon, and edge computing continues to grow in the market to meet the demands for better efficiency in evolving technologies [8, 10, 11].

### 1.3 Security Issues

Edge computing has played a role in the evolution of many mobile devices, 5G technologies, and other intelligent terminals that are used today, such as autonomous cars, industrial IoT devices (e.g., oil and gas failure detection), and numerous smart devices such as security cameras, mobile phones, home appliances, and to name a few. With the ability to offload

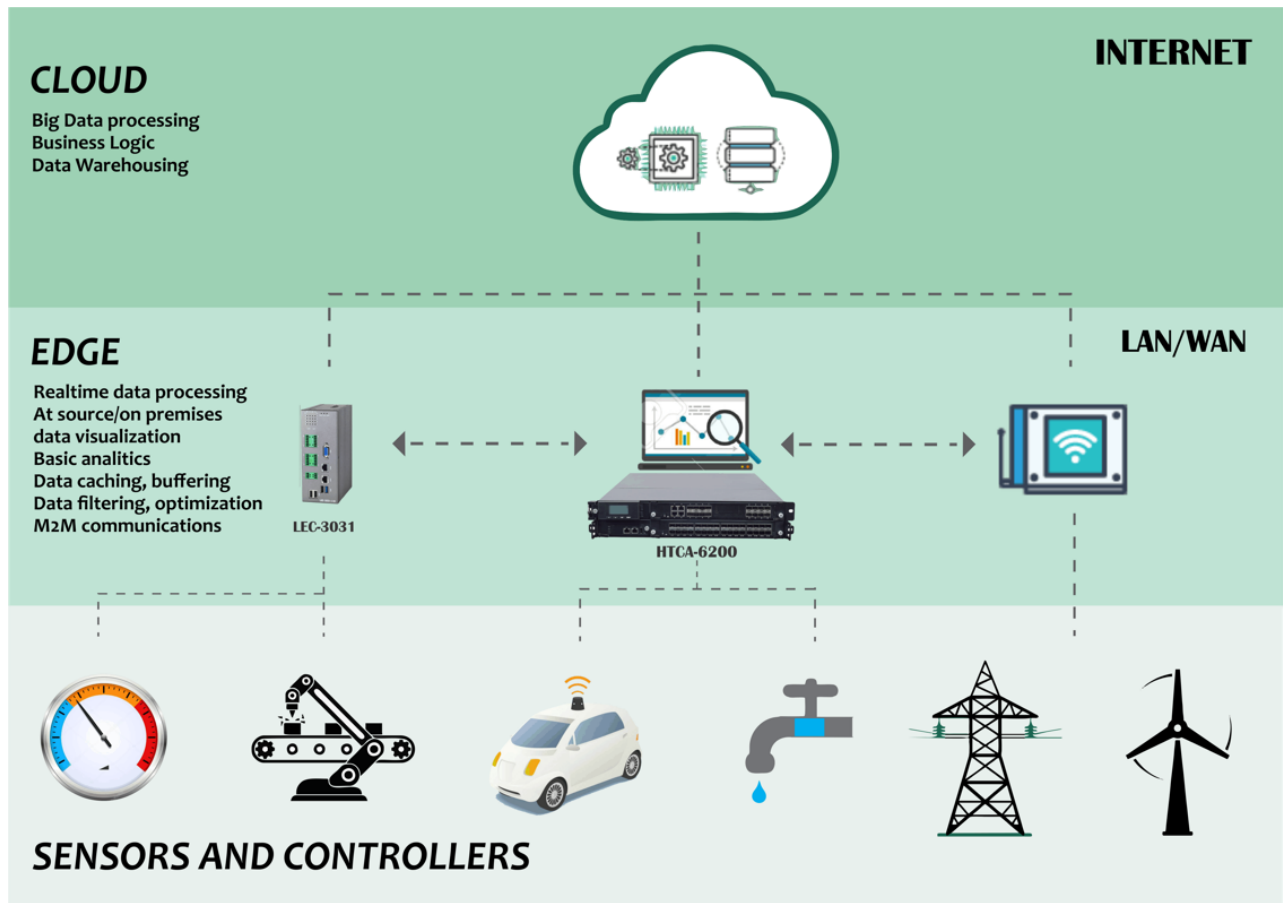


Figure 2: Edge Computing Diagram [3]

processes to edge servers instead of having to process directly on the user's end, the devices can perform faster and still save battery power and/or storage and memory space, allowing them to last longer.

Despite the ability to make data processing and storage more efficient, implementation of the edge also introduces some major security risks considering that devices nowadays also process sensitive and confidential data that need to be kept safe from unauthorized access. In addition to using the cloud and the increasing number of edge devices, this puts data privacy and integrity even more at risk [8]. Such security issues and risks include, but are not limited to:

- Password attacks
- Man-in-the-Middle (MitM) attacks
- Packet and Malware injection
- Compromised data from the cloud
- Manipulated data
- Account theft

#### 1.4 Proposed Solution

For this paper, the main focus is to protect sensitive data over an edge network and maintain its integrity, and the

proposed solution is to implement an edge network model that enables end-to-end protection and validation of a digital communication between an end user (client) and an edge server by using Rivest Shamir Adleman (RSA) encryption and digital signatures.

The RSA algorithm allows users to create two keys, a public and private key, that are mathematically linked to each other. The receiver can share their public key to anyone that needs to send protected data to them, and then the receiver can decrypt the data with their private key. Similarly, the sender can encrypt with their private key to digitally sign a message, and then the receiver decrypts with the shared corresponding public key to validate the signature [13]. For this experiment, a client and an edge server will generate their own key pair, and they will exchange each other's public keys to enable end-to-end encryption and validation. Then the client will be the first to send a unique, digitally signed encrypted message to the edge server, and once received, validated and decrypted, the edge server will return a "processed" digitally signed encrypted message back to the client.

To check for successful end-to-end encryption, a packet sniffer will be used to ensure the messages remain "hidden" during transit, and this packet sniffing will also serve as a passive man-in-the-middle attack demonstration. Upon arrival to the recipient, the decrypted message should be verified and match the original plaintext to check for successful decryption.

tion. As stated before, the keys are mathematically linked, so if both the public and private key get compromised, then it risks the privacy and integrity of data to being compromised as well for both communicating parties. The owner(s) of the keys must, at minimum, practice lightweight key management to minimize the risk of data security and integrity, such as practicing periodic key rotation and keeping the private key securely stored.

## 2. BACKGROUND

This section goes through more background information and vocabulary in detail related to edge computing, cybersecurity, and the RSA algorithm to help readers better understand the concepts in this paper.

### 2.1 Edge Computing

Table 1 gives a list of terms and definitions relevant to understanding edge computing.

Term	Definition
Cloud computing	Also known as the cloud, public cloud, or commercial cloud; A centralized computing model including massive amounts of resources that service end users by computing and storing large loads of data for them.
Edge computing	Also known as the edge, fog computing, edge cloud, or local computing; A more distributed computing model where storage and processing happens closer to the data source.
Edge device	Any device located at the edge that can generate and process data locally.
Edge server	A more powerful computing resource compared to an edge device; Located close to edge devices for more efficient data processing and storage.
Offload	To unload something from some kind of source.

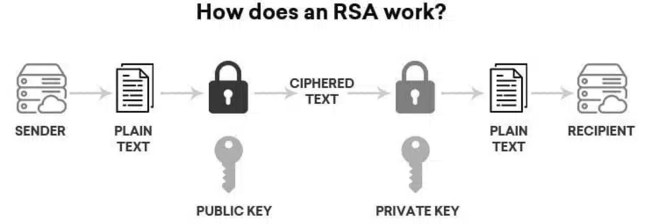
**Table 1: Edge computing terminology**

Edge computing is widely used in the Internet of Things (IoT) considering that most of what we do today relies on the internet. With how evolved technology has become, the data volume for edge devices has also increased, meaning that to process larger volumes of data at reasonable speeds, especially in applications where fast, real-time processing is mandatory or dire, they require a lot more power to meet that demand. The same goes for resources in cloud computing, which is more centralized and remote compared to edge computing; by itself, the cloud would have to process large loads of data from millions of end-users almost at once, and would not be very efficient and doable for this centralized kind of fashion, so by implementing the edge, it decreases the amount of data that needs to be sent to the cloud by decentralizing computation to edge servers for data to be processed more locally, and this saves money and resources for

the cloud [2, 3, 8]. Figure 2 shows a diagram of the general network with the edge in place.

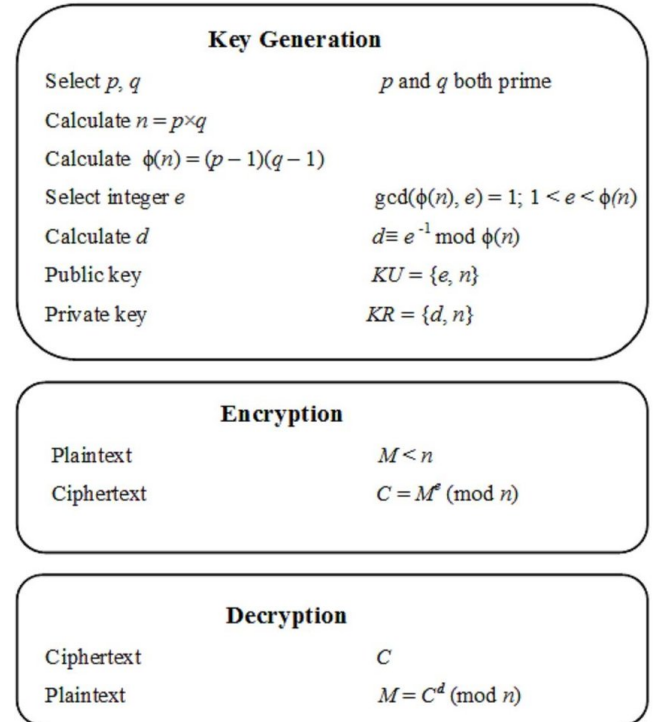
The edge is built so that storage and computations are closer to the data source, allowing for lower latency, higher bandwidth, and better reliability for data storage and processing than when relying on cloud resources alone [9]; in cases where certain tasks are too voluminous to handle, require further analysis beyond the edge server(s)'s capabilities, or simply do not require real-time processing, edge servers can selectively offload those tasks to the cloud to maintain efficiency [7].

### 2.2 RSA Algorithm for Cybersecurity



**Figure 3: RSA Encryption Decryption [4]**

Table 2 contains a list of terms relevant to understanding man-in-the-middle attacks and the use of RSA encryption for data protection in general networking.



**Figure 4: RSA Algorithm [12]**

The RSA algorithm begins with the creation of the key pairs, the public and private keys. To generate the keys, first select two large prime numbers  $p$  and  $q$  (preferably of similar-

ish size), then use  $p$  and  $q$  to calculate  $n$  and  $\phi(n)$ . Next, select an integer  $e$  that is greater than 1 and less than  $\phi(n)$ , and it must be co-prime with  $\phi(n)$  (meaning that the greatest common denominator between  $e$  and  $\phi(n)$  is 1), then use  $e$  to calculate for  $d$ . The result is a public key of  $(e, n)$  and a private key of  $(d, n)$ .

To perform RSA encryption of a plain text message  $M$ , the message must be less than  $n$ , and you use the public key values  $e$  and  $n$  to calculate the output ciphertext  $C$ .

To read what is behind the ciphertext, RSA decryption needs to be performed on the ciphertext  $C$  by using the private key values  $d$  and  $n$  to calculate the output plaintext  $M$ . Refer to Figure 3 for a general understanding of RSA encryption and Figure 4 for the specific formulas used in each part of the RSA algorithm.

For digital signatures, the message to be sent must first go through a hashing algorithm (e.g., SHA-1, SHA-256 in Figure 5, etc.) and then encrypted using their own private key, then the sender attaches the signature to the original message and shares their public key for the receiver to use to decrypt the hash. The receiver can validate the message by putting the original message through the same hashing algorithm and check if theirs matches the decrypted hash value, and if they match, the message has not been altered and has maintained integrity.

The recommended minimum key size for strong enough network security is 2048 bits, and along with proper implementation of RSA encryption and verification in the edge and using simple key management practices, sensitive data can be sent securely through the edge network, and end users can maintain integrity of the data they receive.

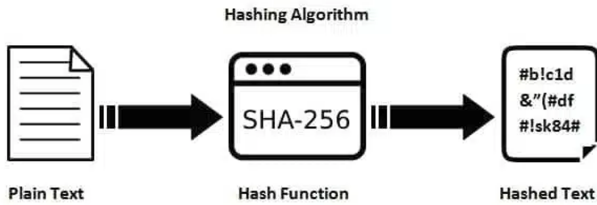


Figure 5: Hashing with SHA-256 [5]

### 3. IMPLEMENTATION

This portion of the paper will cover all the materials that were used to implement the experiment for end-to-end encryption and validation over an edge network. The section will first go over the general edge computing model built for the simulation, and then we will explain how the simulation runs for the model step-by-step as well as how the data is checked to verify proper transit. Note that this was all done on an Ubuntu 20.04 virtual machine in Oracle VirtualBox v6.1.16.

#### 3.1 Edge Model

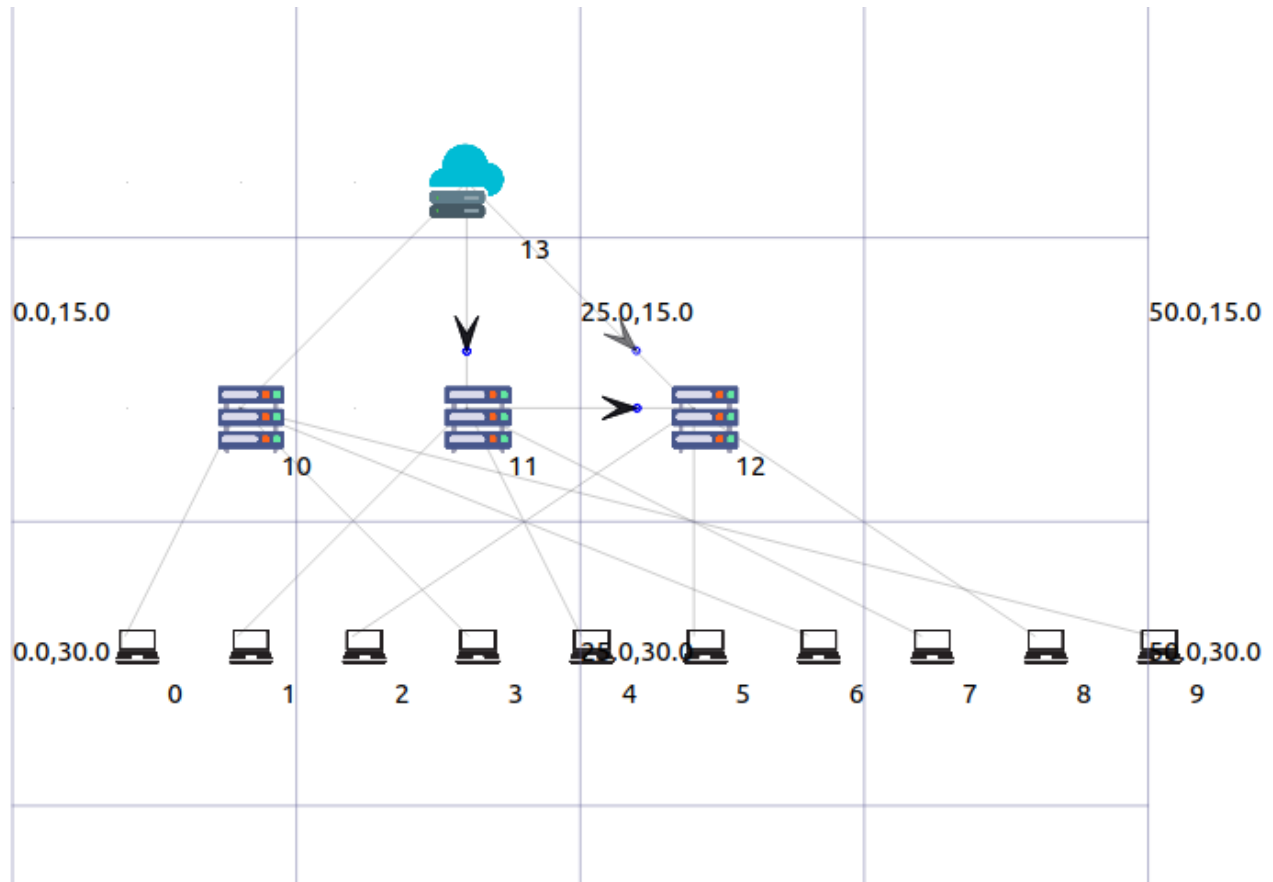
The overall edge computing model was developed in C++ with the NS-3 library v.42 can be found [here](#). The code includes the creation of 10 client/end-user nodes, 3 edge server nodes, and 1 cloud node, and each edge node is connected

Term	Definition
Rivest Shamir Adleman (RSA)	A public-key cryptosystem that enforces a mathematical relationship between a public and private key for encrypting and decrypting data, and it can be used for digital signatures for secure communication and maintain integrity.
End-to-End Encryption	A security method for encrypting data that is only intended for the sender and the receiver to read to ensure privacy and security of data even if it is intercepted.
Private Key	An RSA key generated first; It is used for decrypting data, and the key cannot be openly shared with anyone (hence "private" key).
Public Key	A smaller RSA key derived from the private key pair; It is used for encrypting data, and anyone can use the key to send encrypted messages to the recipient (hence "public" key).
Ciphertext	The scrambled, unreadable text as a result from encryption.
Plaintext	The readable text prior to encryption or as a result from decryption.
Hashing	A one-way algorithm for producing a fixed-length string value with some message as input.
Digital Signature	An algorithm for hashing a message and encrypting it with the private key, which is then sent to and validated by the receiver with the corresponding public key.
Packet	A unit of data, typically broken down from a larger message that is transmitted across a network; It contains the piece of data itself (the payload) and the header containing the route information.
Packet Sniffing	A passive attack where data packets are simply observed while the packets are in transit, and the data is not altered in any way, nor are there additional packets of data from the attacker.
Man-in-the-Middle (MitM)	An attacker that intercepts end-to-end communication by pretending to be both the sender and receiver with the intent to manipulate or steal sensitive data, or insert foreign data.

Table 2: RSA and cybersecurity terminology

to three to four clients and connected to the cloud. Only the second and third edge servers are also directly connected to each other. The point-to-point connections are all assigned a unique IPv4 address before setting up the sockets for data





**Figure 6: Edge Computing Model Simulation in NetAnim**

routing. Each node is then assigned their own socket along with a pre-made message that they will send to the specified destination. In this case, each client will send their message to an edge node, and then the edge node will return a modified message to the corresponding clients.

The end-to-end security version of the edge model is generally the same in terms of the connections, but this time client node 0 and edge node 12 have RSA encryption and digital signatures enabled. To achieve this implementation, an OpenSSL v1.1.1f library built in to the Ubuntu virtual machine was used, and it is an open-source secure sockets layer protocol for ensuring privacy, authentication, and data integrity in any kind of internet communications.

### 3.2 Simulation

When running the code in the command prompt with the `.ns3` command, a `.xml` file is created, which then the NetAnim program is used to open the file and run the edge network simulation. Here is a detailed breakdown of the simulation run (refer to Figure 6 for the simulated model):

1. Edge server connections to client nodes 0 through 9 are assigned IPv4 addresses 10.1.1.1 through 10.1.10.1 respectively, connection of edge nodes 11 and 12 is assigned 10.1.11.1, and cloud connections to edge nodes 10 through 12 are assigned 10.1.12.1 through 10.1.14.1 respectively.

2. Client node 0 will send a unique data packet to edge node 12 while the rest of them will send their unique packets to edge node 11.
3. The data packets reach the edge nodes directly connected to them first.
4. Edge server 11 returns packets to clients 1, 4, and 7 first, while edge server 12 routes packets from clients 2, 5, and 8 to edge server 11, and edge server 10 routes packets from clients 0, 3, 6, and 9 to the cloud.
5. Edge server 11 returns packets to edge server 12 that will then be sent correspondingly to clients 2, 5, and 8, while the cloud routes client packets 3, 6, and 9 to edge server 11 and client packet 0 to edge server 12.
6. Data packets from edge server 11 and 12 are sent to the cloud, where they will then route back to edge server 10 and then distribute to their respective clients.

After running the code in the command prompt, it outputs what is happening between the communicating nodes. For example, the NS-3 program will display that client 0 with IP address 10.1.1.1 has sent its message to edge server 12 with IP address 10.1.14.1 and include a timestamp of when that happened, and then once reached, it will output when the message made it and then display the received message:

No.	Time	Source	Destination	Protocol	Length	Info
1	0.000000	10.1.1.1	10.1.14.1	UDP	69	8080 → 8080 Len=39
2	0.000000	10.1.4.1	10.1.13.1	UDP	69	8080 → 8080 Len=39
3	0.000001	10.1.7.1	10.1.13.1	UDP	69	8080 → 8080 Len=39
4	0.000001	10.1.10.1	10.1.13.1	UDP	70	8080 → 8080 Len=40
5	0.004001	10.1.14.1	10.1.1.1	UDP	77	8080 → 8080 Len=47
▶ Frame 1: 69 bytes on wire (552 bits), 69 bytes captured (552 bits)						
0000	00 21 45 00 00 43 00 00	00 00 3f 11 00 00 0a 01	·!E·C· ··?·			
0010	01 01 0a 01 0e 01 1f 90	1f 90 00 2f 00 00 54 68	····· /·Th			
0020	69 73 20 69 73 20 61 20	6e 6f 72 6d 61 6c 20 6d	is is a normal m			
0030	73 67 20 66 72 6f 6d 20	63 6c 69 65 6e 74 20 6e	sg from client n			
0040	6f 64 65 20 31		ode 1			

Figure 7: Message for Normal Edge Model in Wireshark

No.	Time	Source	Destination	Protocol	Length	Info
1	0.000000	10.1.4.1	10.1.13.1	UDP	69	8080 → 8080 Len=39
2	0.000000	10.1.7.1	10.1.13.1	UDP	69	8080 → 8080 Len=39
3	0.000001	10.1.10.1	10.1.13.1	UDP	70	8080 → 8080 Len=40
4	0.000007	10.1.1.1	10.1.14.1	UDP	543	8080 → 8080 Len=513
5	0.004001	10.1.13.1	10.1.4.1	UDP	77	8080 → 8080 Len=47
▶ Frame 4: 543 bytes on wire (4344 bits), 543 bytes captured (4344 bits)						
0000	00 21 45 00 02 1d 00 00	00 00 3f 11 00 00 0a 01	·!E· ··?·			
0010	01 01 0a 01 0e 01 1f 90	1f 90 02 09 00 00 55 45	····· UE			
0020	d8 94 5e e1 87 23 42 2a	39 d1 a4 4c 3f 6a 62 86	·^·#B· 9·L?jb·			
0030	3e 88 0f 8c 50 d8 5c 54	0f 34 cb 5f f5 97 55 92	>·P·\T·4·_·U·			
0040	de 9a 29 09 86 6c b0 88	14 90 5c 6f 35 5f 49 98	·)·l· ·\o5_I·			
0050	59 e9 68 52 9c d6 39 e8	30 f6 2b 93 07 d8 38 c6	Y·hR·9·0·+·8·			
0060	90 9c 0a ec 4c 7a 82 ac	8f 66 dc ec c6 f3 b1 fb	·····Lz· ·f·			
0070	d7 5c 4d 64 14 76 bb 01	dd 8d 63 ef f8 5b df df	·\Md·v· ·c·[·			
0080	61 1f e0 6e fe 43 08 80	66 3c b0 05 5b 74 84 c5	a·n·C· f<·[t·			
0090	a2 dc d6 3b f7 b2 09 28	bc cf c5 d0 ad 75 de 8b	·;·;(· ···u·			
00a0	2d c6 85 38 ae 1f 85 9b	2a fe 79 2f 14 87 9b b0	-·8· ···*·y/·			
00b0	c0 4d f7 af 4f f1 e5 9d	15 18 46 37 c5 3b c2 d8	·M·@· ···F7·;·			
00c0	e1 87 9f 9e e2 73 47 bf	7b 05 c2 16 91 7f 14 11	·····sG· {·			
00d0	97 c5 8d 22 86 e4 a2 bb	8f 91 1d 3a 97 60 94 6d	··"· ···:·m			
00e0	e3 23 d0 78 10 39 04 21	d6 1c 7a 2c 3e 0f 4a a4	·#·x·9·!· ·z,>·J·			
00f0	78 f7 d0 37 15 9a 37 39	79 2c 7b 6b 9b 22 f4 89	x·7·79 y,{k·"			
0100	b9 d8 8b 96 55 2f f8 a5	6d ba 46 b2 28 5e 89 fc	··U/· m·F·(^·			
0110	42 92 c9 16 a0 b8 74 ed	e4 bf 4b 4f 3d 80 3a 4c	B· ···t· ·K0=·:L			
0120	77 b7 6a ed cd 19 2c 3c	0d 59 be e3 4e 57 b5 9e	w·j· ···<·Y·NW·			
0130	bf b9 f6 57 f7 39 b0 5c	81 33 4e b8 7a 84 26 83	··W·9·\·3N·z·&·			
0140	69 b4 40 dd 7b 63 46 cb	1e ca 55 39 4b ed c8 17	i·@·{cF· ·U9K·			
0150	8e 0e 62 65 7d 07 ce 58	64 bb 48 45 6d a4 2e ee	·be}·X d·HEm·			
0160	18 45 54 0b 1e 77 1c 13	6c 69 3f 4e b6 be 1e b3	·ET·w· li?N·			
0170	c4 13 7e c2 23 c4 c9 2b	c2 bc d1 12 0b 3a ed 3d	··~·#·+· ···:·=			
0180	06 6b ac f1 75 c9 e4 db	15 d1 d2 b1 ed aa eb 9f	·k·u· ···			
0190	da a3 fc eb bb aa 23 0b	22 67 77 be 71 6e 46 62	·····#· "gw·qnFb			
01a0	22 6f 85 f6 b0 2e 64 9c	39 26 a8 b1 ae 2f c9 42	"o· ···d· 9&· ·/·B			
01b0	b0 8f 03 81 64 e7 a9 b4	98 ee b7 1a c2 77 1b eb	·····d· ···w·			
01c0	41 0d 81 22 32 48 9c b3	35 54 17 77 c7 4f de d5	A· "2H· ·5T·w·0·			
01d0	fa 60 ca 30 95 6f ce 42	26 a6 ca 1a 89 b1 60 f0	·`·0·o·B &· ···`·			
01e0	18 75 93 bb 9c 1c 43 bf	60 7c 83 68 9a d6 07 b5	·u· ···C· ` ·h·			
01f0	36 9c 8f 6b 03 a6 5b f7	33 ae 58 14 9a 04 de 47	6·k· [· 3·X· ·G			
0200	59 73 19 cc 99 11 04 45	1e e9 3e bd b7 7f 2e bc	Ys· ···E· ·>· ·			
0210	1e 08 5e c2 5f 0c 27 12	f6 92 0f 67 f8 2b e8	··^·_·'· ···g·+			

Figure 8: Message for Encrypted Edge Model in Wireshark

"This is a normal msg from client node 1." In return, the edge server will send a message back to the client that says "This is a return msg from 10.1.14.1 to 10.1.1.1." The same will happen for all the client and edge nodes in parallel. The messages can also be viewed using Wireshark if the data capture for the connections is enabled in the code. You can refer to Figures 7 and 8 shown in section 4.

## 4. RESULTS

As mentioned earlier, the command prompt displays the timestamps for each time a node sends or receives data from another node, and this is important when you want to find out how much latency is added when including RSA encryption and digital signatures. According to Figures 9 and 10, the encryption and digital signatures add a delay of about 0.00003

```

10.1.9.1 received msg from 10.1.13.1 at +8.008e+09ns
Received msg: This is a return msg from 10.1.13.1 to 10.1.9.1

10.1.1.1 received msg from 10.1.14.1 at +8.012e+09ns
Received msg: This is a return msg from 10.1.14.1 to 10.1.1.1

10.1.4.1 received msg from 10.1.13.1 at +8.012e+09ns
Received msg: This is a return msg from 10.1.13.1 to 10.1.4.1

10.1.7.1 received msg from 10.1.13.1 at +8.012e+09ns
Received msg: This is a return msg from 10.1.13.1 to 10.1.7.1

10.1.10.1 received msg from 10.1.13.1 at +8.01201e+09ns
Received msg: This is a return msg from 10.1.13.1 to 10.1.10.1

kbjackson@kbjackson-VirtualBox:~/ns-allinone-3.42/ns-3.42$

```

Figure 9: Output section without RSA

```

10.1.9.1 received msg from 10.1.13.1 at +8.008e+09ns
Received msg: This is a return msg from 10.1.13.1 to 10.1.9.1

10.1.4.1 received msg from 10.1.13.1 at +8.012e+09ns
Received msg: This is a return msg from 10.1.13.1 to 10.1.4.1

10.1.7.1 received msg from 10.1.13.1 at +8.012e+09ns
Received msg: This is a return msg from 10.1.13.1 to 10.1.7.1

10.1.10.1 received msg from 10.1.13.1 at +8.012e+09ns
Received msg: This is a return msg from 10.1.13.1 to 10.1.10.1

10.1.1.1 received encrypted msg from 10.1.14.1 at +8.01203e+09ns
Received hash: ba38ff8803022b77ff5b7c117a21c6411330e39a7f61f54ec95744a789b77adbdone
Check hash: ba38ff8803022b77ff5b7c117a21c6411330e39a7f61f54ec95744a789b77adbdone
Hashes match!
Decrypted msg: This is a SECURE msg from edge node 3

kbjackson@kbjackson-VirtualBox:~/ns-allinone-3.42/ns-3.42$

```

Figure 10: Output section with RSA

nanoseconds compared to the edge model without RSA. In Wireshark, when put side by side, it also shows how large the packet size of the encrypted and signed message is in Figure 7 compared to the packet size of a normal message in Figure 8, and you can see the scrambled text that would be impossible to decipher without the correct keys. With this comparison, it can be assumed that the increased packet size from using the RSA algorithms is a major factor in why there is that small delay difference.

## 5. CHALLENGES

### 5.1 Errors and Struggles

```

10.1.9.1 received msg from 10.1.13.1 at +8.008e+09ns
Received msg: This is a return msg from 10.1.13.1 to 10.1.9.1

10.1.4.1 received msg from 10.1.13.1 at +8.012e+09ns
Received msg: This is a return msg from 10.1.13.1 to 10.1.4.1

10.1.7.1 received msg from 10.1.13.1 at +8.012e+09ns
Received msg: This is a return msg from 10.1.13.1 to 10.1.7.1

10.1.10.1 received msg from 10.1.13.1 at +8.012e+09ns
Received msg: This is a return msg from 10.1.13.1 to 10.1.10.1

10.1.1.1 received encrypted msg from 10.1.14.1 at +8.01203e+09ns
Decryption error: error:04067072:rsa routines:rsa_ossl_public_decrypt:padding check failed
Check hash: done
Check hash: 77b18e8d9375e30112e6a236d07f820b6bd267504f1a57883178e5b73f0cfcfdone
Invalid msg from mismatched hash! Integrity at risk!

kbjackson@kbjackson-VirtualBox:~/ns-allinone-3.42/ns-3.42$

```

Figure 11: Output section of failed RSA

One major problem that was encountered is performing the proper key-exchange operation. It was attempted by including two separate functions that only send and receive the public keys respectively, and those would be the first functions to execute for the edge and client nodes using RSA. However, when sent, the packet containing the key failed to make it

to its destination, so when sending "encrypted" messages to each other for the rest of the run, there is an error when the nodes attempt to decrypt them. This can be seen in the Wireshark screenshot in Figure 12, where we see the top part of the public key data belonging to edge server 12 with the 10.1.14.1 IP address, and the green message with a black background stating that the destination was "unreachable". The work-around for this was to simply swap the public keys when passing them into the node setup function for their sockets, though this would not be most accurate to a real key-exchange.

Another major problem was the handling of digital signatures. Despite Figure 10 showing the successful hash matching, there are also times where the hash decryption and the matching fails. This is because the digital signature was appended to the encrypted message with a colon character in-between them as the separator before it was sent to the receiver, and in the digital signature, similar to the encrypted text, it also contained colon characters, so when the receiving side separates the message and the signature, it does not separate at the correct colon, and both the digital signature and encrypted message would therefore fail at decryption and hash matching. A possible solution could be to put the encrypted message and digital signature into a structure before inserting it to a packet, but this was not attempted in time.

### 5.2 Potential Issues Moving Forward

Looking forward to the evolution of network security, there is a possibility that RSA key sizes of 2048 bits will no longer be strong enough minimum requirement for security, so we would have to bump up the bits requirement. Depending on how large the key size would have to be by then along with how evolved technology would be overall, this could raise a latency issue and potentially affect task efficiency of many devices using the edge. Another alternative is to use a different kind of asymmetric cryptography algorithm like elliptic curve cryptography (ECC). ECC uses smaller key sizes that offer about the same level of security as RSA, and the algorithm has been gaining popularity, but the main issue is that it is more complicated to implement compared to RSA, and if an error is made, it could hurt the security of the network using the algorithm [1].

## 6. CONCLUSION

For this project, we saw that using RSA cryptography could benefit the security and integrity of data in transit through an edge network if implemented correctly. With that said, there are a few things for this project that could still be further improved in the future in addition to the issues mentioned earlier in section 5.1, such as:

- programming a key rotation requirement that enforces key pair renewal after a certain amount of time to improve key management practices used here.
- including different kinds of attacks in the simulation to further test out the network security between the targeted communicating parties.
- including WiFi connections instead of only point-to-point connections to better represent today's real-life networks.

No.	Time	Source	Destination	Protocol	Length	Info
1	0.000000	10.1.1.1	10.1.14.1	UDP	456	8080 → 8080 Len=426
2	0.004004	10.1.14.1	10.1.1.1	ICMP	58	Destination unreachable (Port unreachable)
3	1.999794	10.1.4.1	10.1.13.1	UDP	69	8080 → 8080 Len=39
4	1.999794	10.1.7.1	10.1.13.1	UDP	69	8080 → 8080 Len=39
Frame 1: 456 bytes on wire (3648 bits), 456 bytes captured (3648 bits)						
0000	00 21 45 00 01 c6 00 00	00 00 3f 11 00 00 0a 01	!E.....?.....			
0010	01 01 0a 01 0e 01 1f 90	1f 90 01 b2 00 00 2d 2d	-----			
0020	2d 2d 2d 42 45 47 49 4e	20 52 53 41 20 50 55 42	---BEGIN RSA PUB			
0030	4c 49 43 20 4b 45 59 2d	2d 2d 2d 2d 0a 4d 49 49	LIC KEY- ----MII			
0040	42 43 67 4b 43 41 51 45	41 6e 39 78 36 32 62 4c	BCgKCAQE An9x62bL			
0050	44 46 75 34 47 73 63 47	4f 6c 4b 39 4d 50 71 35	DFu4GscG 0LK9MPq5			
0060	56 6d 31 6a 4b 54 37 63	67 33 2f 63 59 66 6d 77	Vm1jKT7c g3/cYfmw			
0070	54 51 61 68 78 49 47 71	4c 74 36 69 57 0a 79 39	TQahxIGq Lt6iW-y9			
0080	46 73 4d 62 35 4e 57 77	66 30 59 4e 56 61 32 4f	FsMb5NwW f0YNVa20			
0090	4c 77 6d 49 6a 31 4f 78	66 45 61 75 50 34 44 47	LwmIj10x fEauP4DG			

Figure 12: Failed public key exchange attempt

Computers and technology have come a long way since the very first electronic computer from 1946, and as technology evolved, resource demand becomes greater. With edge computing implemented, devices connected to the edge have greatly improve in performance and efficiency even today, but the security and integrity of potentially sensitive data that is transmitted throughout the edge network will always be a critical issue for the world to address as the edge is vulnerable to different kinds of cyberattacks without the use of some kind of cryptographic algorithm like RSA to secure communicating parties. As technology continues to evolve, so will security risks, and we must keep up to maintain strong network security.

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