

Graph Theory Important Questions

Spanning trees in complete graph is equal to $n^{(n-2)}$ (where n is no of sides or regularity in complete graph). So, spanning trees in complete graph K_4 will be $4^{(4-2)}$. i.e. $4^2 = 16$. Spanning trees in a bipartite graph $K_{m,n}$ is equal to $m^{(n-1)} * n^{(m-1)}$. So, spanning trees in $K_{2,2}$ will be $2^{(2-1)} * 2^{(2-1)}$. i.e. $2^1 * 2^1 = 4$. So, option (C) is correct.

The 2^n vertices of a graph G corresponds to all subsets of a set of size n , for $n \geq 6$. Two vertices of G are adjacent if and only if the corresponding sets intersect in exactly two elements.

The maximum degree of a vertex in G is:

(A) $\binom{n}{2} 2^{n/2}$

(B) 2^{n-2}

(C) $3 \cdot 2^{n-3}$

(D) 2^{n-1}

Answer: (C)

Explanation:

$$\max_k (k \binom{n}{2} 2^{n-k}) = 3 \binom{n}{2} 2^{n-3} = 3 \cdot 2^{n-3}.$$

Ans: C

Let $S = \{1, 2, 3, \dots, n\}$ and G be a simple graph where every subset of S is a vertex in G . There will be an edge between two sub-set of S when they intersect in exactly 3 elements. What will be the degree of a vertex containing 4 elements?

$$C(4,3) \times 2^{(n-4)}$$

Let G be a complete undirected graph on 6 vertices. If vertices of G are labeled, then the number of distinct cycles of length 4 in G is equal to

- A** 15
- B 30
- C** 45
- D 360

Ans: C

$$C(6,4) \times ((4-1)!)/2.$$

As vertices are labeled and they are circular, we need to consider arrangements too.

G is a graph on n vertices and $2n - 2$ edges. The edges of G can be partitioned into two edge-disjoint spanning trees. Which of the following is NOT true for G ?

A

For every subset of k vertices, the induced subgraph has at most $2k - 2$ edges

B

The minimum cut in G has at least two edges

C

There are two edge-disjoint paths between every pair of vertices



There are two vertex-disjoint paths between every pair of vertices



Which of the following graphs has an Eulerian circuit?



Any k -regular graph where k is an even number.



A complete graph on 90 vertices



The complement of a cycle on 25 vertices



None of the above

Not every k -regular graph is connected even if degree is even.

A graph $G = (V, E)$ satisfies $|E| \leq 3|V| - 6$.
The min-degree of G is defined as

$$\min_{v \in V} \{\text{degree}(v)\}$$

. Therefore, min-degree of G
cannot be

A 3

B 4

C 5

 6

$$e \leq 3n-6 \quad \text{--- (1)}$$

~~8(a)~~ wkt

$$\delta(a) \leq \frac{2e}{n}$$

$$e \geq \frac{n \cdot \delta(a)}{2} \quad \text{--- (2)}$$

from (1) & (2)

$$\frac{n \delta(a)}{2} \leq e \leq 3n-6$$

a) $\delta(a) = 3$

$$\frac{3n}{2} \leq e \leq 3n-6$$

b) $\delta(a) = 4$

$$2n \leq e \leq 3n-6$$

c) $\delta(a) = 5$

$$\frac{5n}{2} \leq e \leq 3n-6$$

d) $\delta(a) = 6$

$$3n \leq e \leq 3n-6$$

$$3n \leq 3n-6$$

$$0 \leq -6$$

This is not possible

$\therefore \text{opt (d)}$

In an undirected graph G with n vertices, vertex 1 has degree 1, while each vertex $2, \dots, n-1$ has degree 10 and the degree of vertex n is unknown, Which of the following statement must be TRUE on the graph G ?

- a. There is a path from vertex 1 to vertex n .
- b. There is a path from vertex 1 to each vertex $2, \dots, n-1$.
- c. Vertex n has degree 1.
- d. The diameter of the graph is at most $\frac{n}{10}$
- e. All of the above choices must be TRUE

Ans: a

**There can't exist a single vertex of odd degree.
So vertex 'n' must have odd degree so it must be connect to some vertices from 2 to 'n-1'.**

When a graph has exactly 2 vertices of odd degree, there exists a path b/w those two vertices.

Let $G = (V, E)$ be a simple undirected graph on n vertices. A colouring of G is an assignment of colours to each vertex such that endpoints of every edge are given different colours. Let $\chi(G)$ denote the chromatic number of G , i.e. the minimum numbers of colours needed for a valid colouring of G . A set $B \subseteq V$ is an independent set if no pair of vertices in B is connected by an edge. Let $\alpha(G)$ be the number of vertices in a largest possible independent set in G . In the absence of any further information about G we can conclude.

- A. $\chi(G) \geq \alpha(G)$
- B. $\chi(G) \leq \alpha(G)$
- C. $\alpha(G) \geq \frac{n}{\chi(G)}$
- D. $\alpha(G) \leq \frac{n}{\chi(G)}$
- E. None of the above.

Independence number : Size of largest maximum independent set. $\alpha(G)$ (it covers all adjacent vertices)

Chromatic Number : Minimum No. of color required to properly color the graph. $\chi(G)$

The vertices of G can be partitioned into $\chi(G)$ monochromatic classes. Each class is an independent set, and hence cannot have size larger than $\alpha(G)$.

$\alpha(G) \chi(G) \geq n$. (its a theorem),

option C.

2.11.1 Line Graph: GATE2013-26 [top](#)<https://gateoverflow.in/1537>

The line graph $L(G)$ of a simple graph G is defined as follows:

There is exactly one vertex $v(e)$ in $L(G)$ for each edge e in G .

For any two edges e and e' in G , $L(G)$ has an edge between $v(e)$ and $v(e')$, if and only if e and e' are incident with the same vertex in G .

Which of the following statements is/are TRUE?

- (P) The line graph of a cycle is a cycle.
- (Q) The line graph of a clique is a clique.
- (R) The line graph of a planar graph is planar.
- (S) The line graph of a tree is a tree.

- A. P only
- B. P and R only
- C. R only
- D. P, Q and S only

gate2013 graph-theory normal line-graph

Answer

2.11.2 Line Graph: TIFR2017-B-13 [top](#)<https://gateoverflow.in/95821>

For an undirected graph $G = (V, E)$, the line graph $G' = (V', E')$ is obtained by replacing each edge in E by a vertex, and adding an edge between two vertices in V' if the corresponding edges in G are incident on the same vertex. Which of the following is TRUE of line graphs?

- A. the line graph for a complete graph is complete
- B. the line graph for a connected graph is connected

- C. the line graph for a bipartite graph is bipartite
- D. the maximum degree of any vertex in the line graph is at most the maximum degree in the original graph
- E. each vertex in the line graph has degree one or two

asked 2017 graph-theory line-graph

Answer

Answers: Line Graph

2.11.1 Line Graph: GATE2013-26 tag

<https://gateoverflow.in/1537>



P) **True**. Because every edge in cycle graph will become a vertex in new graph $L(G)$ and every vertex of cycle graph will become an edge in new graph.

R) False. We can give counter example. Let G has 5 vertices and 9 edges which is a planar graph. Assume degree of one vertex is 2 and of all others are 4. Now, $L(G)$ has 9 vertices (because G has 9 edges) and 25 edges. (See below). But for a graph to be planar $|E| \leq 3|V| - 6$.

For 9 vertices $|E| \leq 3 * 9 - 6$

$\Rightarrow |E| \leq 27 - 6$

$\Rightarrow |E| \leq 21$. But $L(G)$ has 25 edges and so is not planar.

As R) is False option (B), (C) are eliminated.

<http://www.personal.kent.edu/~rmuhamma/GraphTheory/MyGraphTheory/planarity.htm>

S) False. By counter example. Try drawing a simple tree which has a Root node, Root node has one child A, node A has two child B and C. Draw its Line graph acc. to given rules in question you will get a cycle graph of 3 vertices.

So (D) also not correct.

\therefore option (A) is correct.

For a graph G with n vertices and m edges, the number of vertices of the line graph $L(G)$ is m , and the number of edges of $L(G)$ is half the sum of the squares of the **degrees** of the vertices in G , minus m .

34 votes

-- prashant singh (115 points)

2.11.2 Line Graph: TIFR2017-B-13 tag

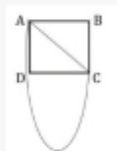
<https://gateoverflow.in/95871>

option B is the right answer

We can solve this question by eliminating options

Option (A) False

Let us take a complete graph of 4 vertices:



In a line graph, no. of edges is

$$\sum_{i=0}^n d_i C_2, \quad d_i \text{ is degree of each vertex}$$

$$= \frac{3 \times 2}{2} + \frac{3 \times 2}{2} + \frac{3 \times 2}{2} + \frac{3 \times 2}{2} = 3 \times 4 = 12$$

No. of vertices in line graph = No. of edges in original graph

No. of vertices in line graph = 6

So, no. of edges to make complete graph with 6 vertices = $\frac{6 \times 5}{2} = 3 \times 5 = 15$

But for given line graph from complete graph of 4 vertices we have only 12 edges.
Contradiction.

Option (B) True

1. Smallest line graph for original graph one edge

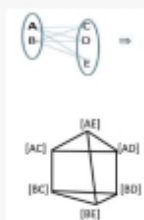


which is also connected graph

If a graph is connected with more than one edge, it will never be disconnected



Option (C) False

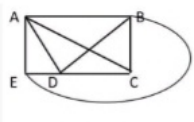


This cannot be 2-colorable and hence is not bipartite.

Option (D) False

Because line graph degree of vertex depends on the attribute

e.g., $[A, B]$ as point in line graph, then degree of this vertex depends on degree of A and degree of B in the original graph.



I'm drawing degree for a point $[AB]$ in line graph.



So, this is wrong.

Option (E) False (wrong as proved in above option (D))

Last modified: May 21, 2020