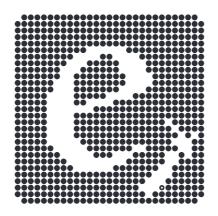
# TECHNICAL REPORTS FROM THE ELECTRONICS GROUP AT THE UNIVERSITY OF OTAGO

### Table of Linear Feedback Shift Registers

by

Roy Ward, Timothy C.A. Molteno

# ELECTRONICS TECHNICAL REPORT No. 2012-1



UNIVERSITY OF OTAGO DUNEDIN, NEW ZEALAND

Online version has

URL: http://www.physics.otago.ac.nz/reports/electronics/ETR2012-1.pdf
The author has homepage: http://www.physics.otago.ac.nz/people/molteno
E-mail: tim@physics.otago.ac.nz
Address: Physics Department, University of Otago, P.O. Box 56, Dunedin, New Zealand

#### **Electronics Group at Otago**

In 1987 Millman and Grabel discarded the historical definition of 'electronics' as the science and technology of the motion of charges, preferring instead the operational definition that the primary concern of people doing electronics is *information processing*. This makes a distinction from *energy processing* practiced in the rest of electrical engineering. The act of information processing is what gets electronics practicioners invloved in the fours 'C's: communication, computation, control, and components. This practical definition seems to describe well the activities within the Electronics Group in the Physics Department at the University of Otago, and the range of topics covered in this technical report series.

In June 2012, research within the Electronics Group include projects on algorithms for sequential inference, lightweight GPS tags for birds, development of radio telescopes, analysis of networks of random resistors, electrical impedance imaging, calibration of numerical models for geothermal fields using Bayesian inference, modelling and sampling of Gaussian processes, and efficient algorithms for Markov chain Monte Carlo applied to inverse problems.

## Table of Linear Feedback Shift Registers

Roy Ward, Timothy C.A. Molteno

#### Abstract

Tables of maximum-cycle Linear Feedback Shift Register (LFSR) taps currently exist in the literature up to n=168 [2]. In this report, we describe a method for generating maximum-cycle Linear Feedback Shift Register designs. It is used to generate n-stage designs, with minimum number of taps, for all  $n \leq 786$  as well as n=1024 and n=2048. These designs are included in this report. This method is computationally efficient, and in addition, can be extended to search for other, non-LFSR, cyclic sequence generators.

## Contents

1	Intr	oducti	on	7
	1.1	Repres	sentation of LFSRs	8
		1.1.1	Cycles	8
		1.1.2	Maximum-cycle LFSRs	9
	1.2	Findin	g Maximum-Cycle LFSRs	9
		1.2.1	Pruning the search tree	9
		1.2.2	Prime Factorisation	10
		1.2.3	The search algorithm	10
2	Tab	le of L	FSR Taps	13
$\mathbf{R}$	efere	nces		19

# List of Figures

1.1	An 8-stage Galois LFSR with cyc	ele size 255.	This LFSR has	s taps at
	positions $8,6,5$ and $4.\ldots$			

# List of Tables

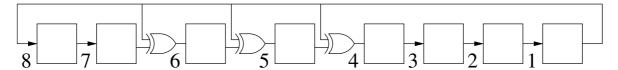
2.1	Shift Registers wit	h Cycle Size	$2^{n}-1$	3
	DIIII I TOO SIDOOLD 1110	1 0,010 0120	9 <b>=</b>	_

## Chapter 1

### Introduction

A Linear Feedback Shift Registers (LFSR) is a shift register where either the outputs of several registers are XORed to provide the input bit to be shifted in (Fibonacci) or where the bit shifted out is XORed to the inputs of several registers (Galois) [6]. The two types are equivalent, so we shall only consider Galois shift registers, as they have a smaller depth (one XOR gate).

We specify a Galois LFSR design by the position of the taps. The taps are the positions (the rightmost position is position 1) of the XOR gates. A tap at position i in an n-stage LFSR would indicate that, at each iteration, the shifted output of the first register would be XORed with the output of the ith register and fed into the input of the next register (at position (i-1)). An n-stage LFSR with a cycle of length  $2^n-1$  is called a maximum-cycle LFSR. Figure 1.1 shows an 8-stage maximum-cycle LFSR with taps at position 8,6,5 and 4.



**Figure 1.1.** An 8-stage Galois LFSR with cycle size 255. This LFSR has taps at positions 8,6,5 and 4.

Alfke [2] presents a table of maximum-cycle n-stage LFSR designs for values of  $n \leq 168$ . This is the largest table in the literature. Clark and Weng [4] for example, show that a shift register can be constructed from its corresponding polynomial. An LFSR will be a maximum-cycle LFSR if and only if the the polynomial represented by the position of the taps is primitive. For instance,  $x^8 + x^6 + x^5 + x^4 + 1$  is a primitive polynomial representing the LFSR 8,6,5,4. Ahmad et. al [1] describe a polynomial-based method to generate maximum-cycle LFSR designs. Their algorithm is designed to find all maximum-cycle LFSR designs and is less efficient than the one presented here for evaluating candidate LFSR designs. It has, to our knowledge, not been applied to values of n greater than 10.

In this paper, we describe a matrix method for generating large n-stage maximum-cycle LFSR designs. This method is efficient, and has been used to generate maximum-cycle LFSR taps for all  $n \le 786$  as well as n = 1024 and n = 2048 [?]. It is feasible

to use this method for all values of n where the prime factors of the corresponding Mersenne number  $2^n - 1$  are known. The matrix method we describe can also be extended to search for other, non-LFSR, cyclic sequence generators.

### 1.1 Representation of LFSRs

The state of an LFSR is a n-vector of 0's and 1's. Motivated by the treatment of Wang et al. [7], an n-stage LFSR can be represented as an  $n \times n$  matrix  $\mathbf{M}$ . Iteration of the LFSR involves multiplication of  $\mathbf{M}$  by the current state vector,  $\mathbf{v}_i$  yielding the next state vector,  $\mathbf{v}_{i+1}$ , i.e.,

$$v_{i+1} = \mathbf{M}v_i$$

The ith iteration from an initial state  $v_0$  can be found by calculating  $\mathbf{M}^i$ , i.e.,

$$oldsymbol{v}_i = \mathbf{M}^i oldsymbol{v}_0$$
 .

The matrices, M that represent LFSRs, have the form

$$\begin{pmatrix}
0 & 0 & 0 & \dots & 0 & a_n \\
1 & 0 & 0 & \dots & 0 & a_{n-1} \\
0 & 1 & 0 & \dots & 0 & a_{n-2} \\
0 & 0 & 1 & \dots & 0 & a_{n-3} \\
\vdots & \vdots & \vdots & \ddots & \vdots & \vdots \\
0 & 0 & 0 & \dots & 1 & a_1
\end{pmatrix}$$

where  $a_i = 1$  if a tap is present at the *i*th position, and  $a_i = 0$  otherwise. Note that  $a_n = 1$  always.

### 1.1.1 Cycles

An LFSR has a cycle of length l from a state  $\boldsymbol{v}$  if after l iterations, the LFSR returns to the state  $\boldsymbol{v}$ , i.e.,

$$\mathbf{M}^l \mathbf{v} = \mathbf{v} \tag{1.1}$$

This is not equivalent to  $\mathbf{M}^{l-1} = \mathbf{I}$ , where  $\mathbf{I}$  is the  $n \times n$  identity matrix. This is because an LFSR with an l-cycle, may do so starting only from a subset of the possible LFSR states – other starting states might exhibit l'-cyclic behaviour with length  $l' \neq l$  where l and l' are not factors of each other.

An example of this is the 8-stage LFSR with taps at positions 8 and 3. Depending on which starting state is chosen, this design has a cycles of size 217,31,7 and 1 A starting state for the 217-cycle is  $\{1,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,1\}$ , a starting state for the 7-cycle is  $\{1,0,1,1,0,0,1,1\}$  and the starting state for the 1-cycle is  $\{0,0,0,0,0,0,0,0,0,0,0,0,1\}$ . The existence of the 7-cycle means that  $\mathbf{M}^{31-1} \neq \mathbf{I}$ .

### 1.1.2 Maximum-cycle LFSRs

If an *n*-stage LFSR is a maximum-cycle LFSR, then all  $2^n-1$  non-zero states of that LFSR will be visited as the LFSR is iterated. It follows that the condition for an *l*-cycle (Equation 1.1) becomes independent of the initial state  $\boldsymbol{v}$  when  $l=2^n-1$ . Therefore Equation 1.1 becomes

$$\mathbf{M}^{2^{n}-1} = \mathbf{I}.\tag{1.2}$$

Additionally, an LFSR is a maximum-cycle LFSR if there are no smaller cycles, i.e.,

$$\forall k \in \mathbb{Z} : 1 \le k < 2^n - 1, \ \mathbf{M}^k \ne \mathbf{I}$$
 (1.3)

### 1.2 Finding Maximum-Cycle LFSRs

The task of finding a maximum-cycle LFSR can be reduced to the task of finding a LFSR matrix  $\mathbf{M}$ , such that Equation 1.2 and Equation 1.3 hold. At first glance, it seems that to determine whether Equation 1.3 holds, requires a search over all the possible values of k. The computational complexity of such a brute-force search is order  $n^32^n$  and becomes prohibitive for large values of n. We show in the next section how this search can be significantly pruned.

#### 1.2.1 Pruning the search tree

The  $2^n-2$  tests in Equation 1.3, for a  $2^n-1$ -cycle LFSR to have no smaller cycles, can be reduced to only testing factors of  $2^n-1$ .

Consider the set, K, of positive integers that satisfy  $\mathbf{M}^k = \mathbf{I}$ ,

$$K = \{k : 1 \le k \le 2^n - 1, \mathbf{M}^k = \mathbf{I}\}.$$

Assume that there is a cycle with length less than  $2^n-1$ . The set K will have elements less than  $2^n-1$ . Let the smallest such element be  $k_0$ . All multiples if  $k_0$  will also satisfy  $matrix M^k = \mathbf{I}$ , i.e., for all positive integers  $j \in \mathbb{Z}$ ,  $\mathbf{M}^{jk_0} = \mathbf{I}$ .

Assume that there exists  $k_x \in K$  where  $k_x$  is not a multiple of  $k_0$  and where  $\mathbf{M}^{k_x} = \mathbf{I}$  then we can write,

$$k_x = jk_0 + t$$

for some integer j > 0 where  $0 < t < k_0$ . It follows that

$$\mathbf{M}^{k_x} = \mathbf{M}^{jk_0+t} = \mathbf{M}^{jk_0}\mathbf{M}^t = \mathbf{M}^t$$

however by assumption  $\mathbf{M}^{k_x} = \mathbf{I}$ , so  $\mathbf{M}^t = \mathbf{I}$  which violates our assumption that  $k_0$  is the smallest such value. Hence all elements of K must be multiples of  $k_0$ .

In a maximal-cycle LFSR Equation 1.2 holds, and  $\mathbf{M}^{2^{n}-1} = \mathbf{I}$  so we know that  $k_0$  must be a factor of  $2^{n}-1$ . Therefore only values of k that are factors of  $2^{n}-1$  need to be checked in order to establish that Equation 1.3 holds.

#### 1.2.2 Prime Factorisation

A further improvement is still possible by considering the prime factorisation of  $2^{n}-1$ ,

$$2^n - 1 = \prod_{i=1}^m p_i^{k_i},$$

where m is the number of prime factors and the  $p_i$  are the prime factors. Any factor of  $2^n-1$  except  $2^n-1$  itself is a factor of  $\frac{2^n-1}{p_i}$  for some i, so to establish that Equation 1.3 holds, we only need to check that

$$\forall i \in \{1, \dots, m\}, \ \mathbf{M}^{\frac{2^n - 1}{p_i}} \neq \mathbf{I}. \tag{1.4}$$

Thus, if a prime factorisation of  $2^n-1$  is available then we only need to search as many values of k as there are prime factors and Equation 1.2 can be shown to hold with a relatively small amount of computational effort. As there must be fewer than n factors of  $2^n-1$  this search can be done in polynomial time.

The numbers  $2^n-1$  are known as Mersenne numbers [3]. Implementation of the algorithm described in Equation 1.4 requires a table of the prime factors of Mersenne numbers. A table of all prime factors of the Mersenne numbers M(n) for values of n up to n=786 was generated and is available in machine readable form from Reference [5].

#### 1.2.3 The search algorithm

The search for an n-stage maximum-cycle LFSR is performed by considering all potential designs in order of increasing tap number. Ahmad et. al [1] show in their Theorem 4, that an n-stage LFSR design must have an even number of taps in order to be maximum-cycle. Therefore, our algorithm starts with a search for possible two-tap designs. Each two-tap design with taps at positions i and j is represented by a matrix  $\mathbf{M}_{i,j}$ . The first tap is always at position i=n and represents the feedback from bit 1 to bit n.The two-tap search must check LFSRs with the second tap at positions  $n/2 \leq j < n$  (since, if  $\mathbf{M}_{n,i_1,\dots,i_k}$  is an LFSR, then  $\mathbf{M}_{n,n-i_k,\dots,n-i_1}$  is an LFSR with the same cyclic properties; reflected bit-order and reversed in time). For each pair of tap positions, i, j, Equation 1.2 is checked, and if this test is satisfied, the candidate LFSR design is checked against Equation 1.4 using the prime factors of  $2^n-1$ .

If no two-tap designs are found a four-tap design search is conducted. Once again, the first tap position is n, and a search over the remaining three tap positions is carried out. For each candidate design  $\mathbf{M}_{i,j,k,l}$ , the same tests are done.

The pseudocode for this algorithm, showing the order in which the candidate designs are searched, is shown below.

```
boolean test(M)  \mbox{let } \{p_1...p_m\} \mbox{ = prime factors of } 2^n-1 \mbox{;} \\ \mbox{if } (\mathbf{M}^{2^n-1}=\mathbf{I})
```

```
 \text{ if } (\forall p_i \in \{p_1...p_m\}, \ \mathbf{M}^{\frac{2^n-1}{p_i}} \neq \mathbf{I}) \\ \text{ return true;} \\ \text{ return false;} \\ \\ /\!/ \ \mathit{Two-tap case} \\ \text{ for } j = \{n-1\dots\frac{n}{2}\} \\ \text{ if } (\text{test}(\mathbf{M}_{n,j})) \\ \text{ return } (\texttt{n},\texttt{j}); \\ \\ /\!/ \ \mathit{Four-tap case} \\ \text{ for } 1 = \{n-1\dots\frac{n}{2}\} \\ \text{ for } j = \{n-1\dots l+1\} \\ \text{ for } k = \{l-1\dots j+1\} \\ \text{ if } (\text{test}(\mathbf{M}_{n,j,k,l})) \\ \text{ return } (\texttt{n},\texttt{j},k,l); \\ \end{cases}
```

# Chapter 2

# Table of LFSR Taps

Table 2.1: Shift Registers with Cycle Size  $2^n-1$ 

n	LFSR-2	LFSR_4	n	LFSR-2	LESR_4	n	LFSR-2	LFSR_4
$\frac{n}{2}$	2,1	DI DIC-4	24	DI 510-2	24, 23, 21, 20	46	DI 510-2	46, 40, 39, 38
	$\begin{bmatrix} 2, 1 \\ 3, 2 \end{bmatrix}$		25	25, 22	25, 24, 23, 22		47, 42	47, 46, 43, 42
1	$\begin{bmatrix} 3, 2 \\ 4, 3 \end{bmatrix}$		26	20, 22	26, 25, 24, 20	48	41,42	48, 44, 41, 39
	5, 3	5, 4, 3, 2	27		27, 26, 25, 22		49, 40	49, 45, 44, 43
	· 1	6, 5, 3, 2		20 25	28, 27, 24, 22	50	49,40	50, 48, 47, 46
	6, 5			28, 25				, , ,
7	7,6	7, 6, 5, 4		29,27	29, 28, 27, 25	51	59.40	51, 50, 48, 45
8	0.5	8, 6, 5, 4	30	21 00	30, 29, 26, 24		52,49	52, 51, 49, 46
	· '	9, 8, 6, 5		31,28	31, 30, 29, 28	53		53, 52, 51, 47
	10,7	10, 9, 7, 6	32	00.00	32, 30, 26, 25	54	FF 01	54, 51, 48, 46
	11,9	11, 10, 9, 7		33,20	33, 32, 29, 27	l .	55,31	55, 54, 53, 49
12		12, 11, 8, 6	34		34, 31, 30, 26	56		56, 54, 52, 49
13		13, 12, 10, 9		35, 33	35, 34, 28, 27		57,50	57, 55, 54, 52
14		14, 13, 11, 9		36, 25	36, 35, 29, 28		58,39	58, 57, 53, 52
	15, 14	15, 14, 13, 11	37		37, 36, 33, 31	59		59, 57, 55, 52
16		16, 14, 13, 11	38		38, 37, 33, 32		60, 59	60, 58, 56, 55
	17, 14	17, 16, 15, 14	l	39, 35	39, 38, 35, 32	61		61, 60, 59, 56
	18, 11	18, 17, 16, 13	40		40, 37, 36, 35	62		62, 59, 57, 56
19		19, 18, 17, 14	41	41,38	41, 40, 39, 38	63	63,62	63, 62, 59, 58
20	20,17	20, 19, 16, 14	42		42, 40, 37, 35	64		64, 63, 61, 60
21	21, 19	21, 20, 19, 16	43		43, 42, 38, 37	65	65,47	65, 64, 62, 61
22	22, 21	22, 19, 18, 17	44		44, 42, 39, 38	66		66, 60, 58, 57
23	23, 18	23, 22, 20, 18	45		45, 44, 42, 41	67		67, 66, 65, 62
68	68, 59	68, 67, 63, 61	120		120, 118, 114, 111	172	172, 165	172, 169, 165, 161
69		69, 67, 64, 63	121	121, 103	121, 120, 116, 113	173		173, 171, 168, 165
70		70,69,67,65	122		122, 121, 120, 116	174	174, 161	174, 169, 166, 165
71	71,65	71, 70, 68, 66	123	123, 121	123, 122, 119, 115	175	175, 169	175, 173, 171, 169
72		72, 69, 63, 62	124	124,87	124, 119, 118, 117	176	·	176, 167, 165, 164
73	73,48	73, 71, 70, 69	125	·	125, 120, 119, 118	177	177, 169	177, 175, 174, 172
74	,	74, 71, 70, 67	126		126, 124, 122, 119	178	178,91	178, 176, 171, 170
75		75, 74, 72, 69	127	127, 126	127, 126, 124, 120	179	, i	179, 178, 177, 175
76		76, 74, 72, 71	128	,	128, 127, 126, 121	180		180, 173, 170, 168
77		77, 75, 72, 71	l	129, 124	129, 128, 125, 124	181		181, 180, 175, 174
78		78, 77, 76, 71	l	· · · · · · · · · · · · · · · · · · ·	130, 129, 128, 125	182		182, 181, 176, 174
	79,70	79, 77, 76, 75	131		131, 129, 128, 123		183, 127	183, 179, 176, 175
	1.0,.0	, , ,		L	,,120,120		, - <b> ·</b>	,,

n	LFSR-2	LFSR-4	n	LFSR-2	LFSR-4	n	LFSR-2	LFSR-4
80		80, 78, 76, 71	132	132, 103	132, 130, 127, 123	184		184, 177, 176, 175
	81,77	81, 79, 78, 75	133	102, 100	133, 131, 125, 124		185, 161	185, 184, 182, 177
82		82, 78, 76, 73		134,77	134, 133, 129, 127	186		186, 180, 178, 177
83		83, 81, 79, 76	135	135, 124	135, 132, 131, 129	187		187, 182, 181, 180
	84,71	84, 83, 77, 75	136		136, 134, 133, 128	188		188, 186, 183, 182
85	.,	85, 84, 83, 77	137	137, 116	137, 136, 133, 126	189		189, 187, 184, 183
86		86, 84, 81, 80	138		138, 137, 131, 130	190		190, 188, 184, 177
	87,74	87, 86, 82, 80	139		139, 136, 134, 131	191	191, 182	191, 187, 185, 184
88	,	88, 80, 79, 77	140	140, 111	140, 139, 136, 132	192	,	192, 190, 178, 177
	89, 51	89, 86, 84, 83	141	,	141, 140, 135, 128		193, 178	193, 189, 186, 184
90	,	90, 88, 87, 85	142	142, 121	142, 141, 139, 132	l .	194, 107	194, 192, 191, 190
91		91, 90, 86, 83	143	,	143, 141, 140, 138	195	,	195, 193, 192, 187
92		92, 90, 87, 86	144		144, 142, 140, 137	196		196, 194, 187, 185
93	93, 91	93, 91, 90, 87	145	145,93	145, 144, 140, 139	197		197, 195, 193, 188
94	94, 73	94, 93, 89, 88	146	,	146, 144, 143, 141	198	198, 133	198, 193, 190, 183
95	95,84	95, 94, 90, 88	147		147, 145, 143, 136	199	199, 165	199, 198, 195, 190
96	,	96, 90, 87, 86	148	148, 121	148, 145, 143, 141	200	,	200, 198, 197, 195
97	97,91	97, 95, 93, 91	149	,	149, 142, 140, 139	201	201, 187	201, 199, 198, 195
98	98,87	98, 97, 91, 90	150	150,97	150, 148, 147, 142	202	202, 147	202, 198, 196, 195
99		99, 95, 94, 92	151	151, 148	151, 150, 149, 148	203		203, 202, 196, 195
100	100,63	100, 98, 93, 92	152		152, 150, 149, 146	204		204, 201, 200, 194
101		101, 100, 95, 94	153	153, 152	153, 149, 148, 145	205		205, 203, 200, 196
102		102, 99, 97, 96	154		154, 153, 149, 145	206		206, 201, 197, 196
103	103,94	103, 102, 99, 94	155		155, 151, 150, 148	207	207, 164	207, 206, 201, 198
104		104, 103, 94, 93	156		156, 153, 151, 147	208		208, 207, 205, 199
105	105,89	105, 104, 99, 98	157		157, 155, 152, 151	209	209,203	209, 207, 206, 204
106	106, 91	106, 105, 101, 100	158		158, 153, 152, 150	210		210, 207, 206, 198
107		107, 105, 99, 98	159	159, 128	159, 156, 153, 148	211		211, 203, 201, 200
1	108,77	108, 103, 97, 96	160		160, 158, 157, 155	212	212, 107	212, 209, 208, 205
109		109, 107, 105, 104	161	161, 143	161, 159, 158, 155	213		213, 211, 208, 207
110		110, 109, 106, 104			162, 158, 155, 154	214		214, 213, 211, 209
111	111, 101	111, 109, 107, 104			163, 160, 157, 156	215	215, 192	215, 212, 210, 209
112		112, 108, 106, 101			164, 159, 158, 152	216		216, 215, 213, 209
	113, 104	113, 111, 110, 108			165, 162, 157, 156			217, 213, 212, 211
114		114, 113, 112, 103			166, 164, 163, 156	l .	218,207	218, 217, 211, 210
115		115, 110, 108, 107		167, 161	167, 165, 163, 161	219		219, 218, 215, 211
116		116, 114, 111, 110			168, 162, 159, 152	220		220, 211, 210, 208
117		117, 116, 115, 112		· · · · · · · · · · · · · · · · · · ·		221		221, 219, 215, 213
	118,85	118, 116, 113, 112		170, 147	170, 169, 166, 161	222		222, 220, 217, 214
	119, 111	119, 116, 111, 110			171, 169, 166, 165	l .	223,190	223, 221, 219, 218
224		224, 222, 217, 212			276, 275, 273, 270	328		328, 323, 321, 319
	225, 193	225, 224, 220, 215			277, 274, 271, 265	l .	329,279	329, 326, 323, 321
226		226, 223, 219, 216		· · · · · · · · · · · · · · · · · · ·	278, 277, 274, 273	330		330, 328, 323, 322
227				[279, 274]	279, 278, 275, 274	331		331, 329, 325, 321
228		228, 226, 217, 216	ll .		280, 278, 275, 271	l .	· · · · · · · · · · · · · · · · · · ·	332, 325, 321, 320
229		229, 228, 225, 219					333,331	333, 331, 329, 325
230		230, 224, 223, 222	282	282,247	282, 278, 277, 272	334		334, 333, 330, 327

n	LFSR-2	LFSR-4	n	LFSR-2	LFSR-4	n	LFSR-2	LFSR-4
231		231, 229, 227, 224	283	21 010 2	283, 278, 276, 271	335	21 010 2	335, 333, 328, 325
232	201, 200	232, 228, 223, 221		284 165	284, 279, 278, 276	336		336, 335, 332, 329
	233, 159	233, 232, 229, 224	285	201,100	285, 280, 278, 275	l .	337, 282	337, 336, 331, 327
234	· /	234, 232, 225, 223		286, 217	286, 285, 276, 271	338	001,202	338, 336, 335, 332
235	254, 205	235, 234, 229, 226		· ·	287, 285, 282, 281	339		339, 332, 329, 323
	236, 231	236, 229, 228, 226	288	201,210	288, 287, 278, 277	340		340, 337, 336, 329
237	250, 251	237, 236, 233, 230		289, 268	289, 286, 285, 277	341		341, 336, 330, 327
238		238, 237, 236, 233	290	209, 200	290, 288, 287, 285		342, 217	342, 341, 340, 331
	239, 203	239, 238, 232, 227	291		291, 286, 280, 279		· · · · · · · · · · · · · · · · · · ·	343, 338, 335, 333
240	255, 205	240, 237, 235, 232		292, 195	292, 291, 289, 285	344	343, 200	344, 338, 334, 333
241	241, 171	241, 237, 233, 232	293	252, 150	293, 292, 287, 282		345, 323	345, 343, 341, 337
242	241, 111	242, 241, 236, 231		294, 233	294, 292, 291, 285	346	040,020	346, 344, 339, 335
243		243, 242, 238, 235		,	295, 293, 291, 290	347		347, 344, 337, 336
244		244, 243, 240, 235	296	250, 241	296, 292, 287, 285	348		348, 344, 341, 340
245		245, 244, 241, 239		297, 292	297, 296, 293, 292	349		349, 347, 344, 343
246		246, 245, 244, 235	298	201,202	298, 294, 290, 287	l .	350, 297	350, 340, 337, 336
	247, 165	247, 245, 243, 238	299		299, 295, 293, 288			351, 348, 345, 343
248	211,100	248, 238, 234, 233		300, 293	300, 290, 288, 287	352	001,011	352, 346, 341, 339
	249 163	249, 248, 245, 242	301	500,200	301, 299, 296, 292		353 284	353, 349, 346, 344
250	· · · · · · · · · · · · · · · · · · ·	250, 247, 245, 240		302, 261	302, 297, 293, 290	354	000, 204	354, 349, 341, 340
251	200, 111	251, 249, 247, 244	303	002,201	303, 297, 291, 290	355		355, 354, 350, 349
	252, 185	252, 251, 247, 241	304		304, 303, 302, 293	356		356, 349, 347, 346
253	202, 100	253, 252, 247, 246		305, 203	305, 303, 299, 298	357		357, 355, 347, 346
254		254, 253, 252, 247	306	300,200	306, 305, 303, 299	358		358, 351, 350, 344
	255, 203	255, 253, 252, 250	307		307, 305, 303, 299		359, 291	359, 358, 352, 350
256	200,200	256, 254, 251, 246	308		308, 306, 299, 293	360	000,201	360, 359, 335, 334
	257, 245	257, 255, 251, 250	309		309, 307, 302, 299	361		361, 360, 357, 354
258	· · · · · · · · · · · · · · · · · · ·	258, 254, 252, 249	310		310, 309, 305, 302		362, 299	362, 360, 351, 344
259		259, 257, 253, 249	311		311, 308, 306, 304	363	, , , , , ,	363, 362, 356, 355
260		260, 253, 252, 250	312		312, 307, 302, 301		364, 297	364, 363, 359, 352
261		, , , , , , , , , , , , , , , , , , ,		313, 234	313, 312, 310, 306	365	,	365, 360, 359, 356
262		262, 258, 254, 253					366, 337	366, 362, 359, 352
	263, 170		315	,	315, 314, 306, 305		· '	367, 365, 363, 358
264	,	, , , , , , , , , , , , , , , , , , ,		316, 181	316, 309, 305, 304	368	,	368, 361, 359, 351
265	265, 223	265, 263, 262, 260	317	,	317, 315, 313, 310		369, 278	369, 367, 359, 358
266	· /	266, 265, 260, 259	318		318, 313, 312, 310		· '	370, 368, 367, 365
267	,	267, 264, 261, 259		319, 283	319, 318, 317, 308	371	,	371, 369, 368, 363
	268, 243	268, 267, 264, 258	320		320, 319, 317, 316	372		372, 369, 365, 357
269	,	269, 268, 263, 262		321, 290	321, 319, 316, 314	373		373, 371, 366, 365
	270, 217	270, 267, 263, 260			322, 321, 320, 305	374		374, 369, 368, 366
271	· · · · · · · · · · · · · · · · · · ·	271, 265, 264, 260	323	,	323, 322, 320, 313		375, 359	375, 374, 368, 367
272	,	272, 270, 266, 263	324		324, 321, 320, 318	376	, i	376, 371, 369, 368
273	273,250	273, 272, 271, 266	325		325, 323, 320, 315	377	377,336	377, 376, 374, 369
274	· · · · · · · · · · · · · · · · · · ·	274, 272, 267, 265	326		326, 325, 323, 316		· · · · · · · · · · · · · · · · · · ·	378, 374, 365, 363
275		275, 266, 265, 264	327	327,293	327, 325, 322, 319	379		379, 375, 370, 369
380	380, 333	380, 377, 374, 366	432		432, 429, 428, 419	484	484,379	484, 483, 482, 470
381		381, 380, 379, 376	433	433,400	433, 430, 428, 422	485		485, 479, 469, 468

n	LFSR-2	LFSR-4	n	LFSR-2	LFSR-4	n	LFSR-2	LFSR-4
			434	21 210 2	434, 429, 423, 422	486	21 010 2	486, 481, 478, 472
	383, 293	383, 382, 378, 374			435, 430, 426, 423		487, 393	487, 485, 483, 478
384	000,200	384, 378, 369, 368	ll .	436 271	436, 432, 431, 430	488	101,000	488, 487, 485, 484
1	385, 379		437	100,211	437, 436, 435, 431		489, 406	489, 484, 483, 480
386	· ′	, , ,		438, 373	438, 436, 432, 421	490	· · ·	490, 485, 483, 481
387	300,303		439	439, 390	439, 437, 436, 431	491	150,211	491, 488, 485, 480
388		388, 387, 385, 374		100,000	440, 439, 437, 436	492		492, 491, 485, 484
389		389, 384, 380, 379		441, 410	441, 440, 433, 430	493		493, 490, 488, 483
	390, 301	, , , , , , , , , , , , , , , , , , ,	442	111, 110	442, 440, 437, 435		494, 357	494, 493, 489, 481
1	· ′	, , ,	443		443, 442, 437, 433		495, 419	495, 494, 486, 480
392	001,000	, , , , , , , , , , , , , , , , , , ,	444		444, 435, 432, 431	496	150, 115	496, 494, 491, 480
	303 386		445		445, 441, 439, 438		497, 419	497, 493, 488, 486
	,	394, 392, 387, 386		446 341	446, 442, 439, 431	498	131, 113	498, 495, 489, 487
395	094, 209	395, 390, 389, 384		440,341 $447,374$	447, 446, 441, 438	499		499, 494, 493, 488
	396, 371	396, 392, 390, 389		111,011	448, 444, 442, 437	500		500, 499, 494, 490
397	550,571	397, 392, 387, 385		449, 315	449, 446, 440, 438	501		501, 499, 497, 496
398		398, 393, 392, 384	ll .	450,371	450, 443, 438, 434	502		502, 498, 497, 494
	399, 313	, , , ,	451	400,071	451, 450, 441, 435		503, 500	503, 502, 501, 500
400	555,515	400, 398, 397, 395			452, 448, 447, 446	504	303,300	504, 502, 490, 483
	401, 249	, , ,	453		453, 449, 447, 438		505, 349	505, 500, 497, 493
402	401, 249	402, 399, 398, 393			454, 449, 445, 444		506, 411	506, 501, 494, 491
403		403, 398, 395, 394		455,417	455, 453, 449, 444	507	500,411	507, 504, 501, 494
	404, 215	404, 400, 398, 397		400,411	456, 454, 445, 433		508, 399	508, 505, 500, 495
405	404, 210			457, 441	457, 454, 449, 446	509	300, 333	509, 506, 502, 501
	406, 249	406, 402, 397, 393		457,441 $458,255$	458, 453, 448, 445	510		510, 501, 500, 498
407	· · ·		459	400, 200	459, 457, 454, 447		511,501	511, 509, 503, 501
408	407,550		460	460, 399	460, 459, 455, 451	512	311,501	512, 510, 507, 504
409	400 322	409, 406, 404, 402	461	100,000	461, 460, 455, 454		513, 428	513, 505, 503, 500
410	103, 322	, , , ,	462	462, 389	462, 457, 451, 450	514	010, 420	514, 511, 509, 507
411		411, 408, 401, 399		462,369 $463,370$	463, 456, 455, 452	515		515, 511, 508, 501
	412 265	412, 409, 404, 401		100,010	464, 460, 455, 441	516		516, 514, 511, 509
413	112,200	413, 407, 406, 403		465 406		l .		517, 515, 507, 505
414		414, 405, 401, 398	ll .	100, 100	466, 460, 455, 452		518 485	518, 516, 515, 507
	415 313	415, 413, 411, 406			467, 466, 461, 456		· · ·	519, 517, 511, 507
416	110,010	416, 414, 411, 407			468, 464, 459, 453	520	010, 110	520, 509, 507, 503
	417.310	417, 416, 414, 407			469, 467, 464, 460		521, 489	521, 519, 514, 512
418	111,010	418, 417, 415, 403	ll .	470.321	470, 468, 462, 461	522	021, 100	522, 518, 509, 507
419		419, 415, 414, 404		· '	471, 469, 468, 465	523		523, 521, 517, 510
420		420, 412, 410, 407			472, 470, 469, 461		524, 357	524, 523, 519, 515
421		421, 419, 417, 416			473, 470, 467, 465	525	21,001	525, 524, 521, 519
	422, 273	422, 421, 416, 412		474, 283	· / / /	526		526, 525, 521, 517
1	,	423, 420, 418, 414		,	475, 471, 467, 466		527, 480	527, 526, 520, 518
424		424, 422, 417, 415		476, 461		528	], 100	528, 526, 522, 517
	425, 413	425, 422, 421, 418		,	477, 470, 462, 461		529, 487	529, 528, 525, 522
426	, 113	426, 415, 414, 412		478, 357	478, 477, 474, 472	530	, , 10,	530, 527, 523, 520
427		427, 422, 421, 416		· '	479, 475, 472, 470	531		531, 529, 525, 519
	428, 323	428, 426, 425, 417		,	480, 473, 467, 464		532, 531	532, 529, 528, 522
	1-0,020	,,,,,	1-00				552,551	,,,

n	LFSR-2	LFSR-4	n	LFSR-2	LFSR-4	n	LFSR-2	LFSR-4
429		429, 422, 421, 419	481	481, 343	481, 480, 472, 471	533		533, 531, 530, 529
430		430, 419, 417, 415	482		482, 477, 476, 473	534		534, 533, 529, 527
431	431, 311	431, 430, 428, 426	483		483, 479, 477, 474	535		535, 533, 529, 527
536		536, 533, 531, 529		588, 437	588, 577, 572, 571	640		640, 638, 637, 626
	537, 443	537, 536, 535, 527	589		<b>589</b> , <b>586</b> , <b>585</b> , <b>579</b>		641,630	641, 640, 636, 622
538		538, 537, 536, 533		590, 497	590, 588, 587, 578		<i>'</i>	642, 636, 633, 632
539		539, 535, 534, 529	591		591, 587, 585, 582	643	, , , ,	643, 641, 640, 632
	540, 361	540, 537, 534, 529	592		592, 591, 573, 568	644		644, 634, 633, 632
541	,	541, 537, 531, 528		593, 507	593, 588, 585, 584	645		645, 641, 637, 634
542		542, 540, 539, 533		594, 575	594, 586, 584, 583	646	646, 397	646, 635, 634, 633
543	543, 527	543, 538, 536, 532	595	,	595, 594, 593, 586	1	· '	647, 646, 643, 642
544	,	544, 538, 535, 531	596		596, 592, 591, 590	648	,	648, 647, 626, 625
545	545, 423	545, 539, 537, 532	597		597, 588, 585, 583	649	649,612	649, 648, 644, 638
546		546, 545, 544, 538	598		598, 597, 592, 591	650	650,647	650, 644, 635, 632
547		547, 543, 540, 534	599	599, 569	599, 593, 591, 590	651		651, 646, 638, 637
548		548, 545, 543, 538	600		600, 599, 590, 589	652	652,559	652, 647, 643, 641
549		549, 546, 545, 533	601	601,400	601, 600, 597, 589	653		653, 646, 645, 643
550	550,357	550, 546, 533, 529	602		602, 596, 594, 591	654		654, 649, 643, 640
551	551,416	551, 550, 547, 542	603		603,600,599,597	655	655, 567	655, 653, 639, 638
552		552, 550, 547, 532	604		604, 600, 598, 589	656		656, 646, 638, 637
553	553,514	553, 550, 549, 542	605		605,600,598,595	657	657,619	657, 656, 650, 649
554		554, 551, 546, 543	606		606, 602, 599, 591	658	658,603	658, 651, 648, 646
555		555, 551, 546, 545	607	607,502	607,600,598,595	659		659, 657, 655, 644
556	556,403	556, 549, 546, 540	608		608,606,602,585	660		660, 657, 656, 648
557		557, 552, 551, 550	609	609,578	609, 601, 600, 597	661		661, 657, 650, 649
558		558, 553, 549, 544	610	610,483	610,602,600,599	662	662,365	662, 659, 656, 650
	559,525	559, 557, 552, 550	611		611,609,607,601	1	663,406	663, 655, 652, 649
560		560, 554, 551, 549	612		612,607,602,598	664		664, 662, 660, 649
	561,490	561, 558, 552, 550	613		613,609,603,594	1	665,632	665,661,659,654
562		, , , ,	614		614, 613, 612, 607	666		666,664,659,656
563				615,404	615, 614, 609, 608	667		667,664,660,649
	564,401	, , , ,	616		616, 614, 602, 597	668		668, 658, 656, 651
565		, , , , , , , , , , , , , , , , , , ,		617,417	617,612,608,607	669		669, 667, 665, 664
	· · · · · · · · · · · · · · · · · · ·	, , , ,	618		618, 615, 604, 598		<i>'</i>	670, 669, 665, 664
567	567,424	567, 563, 557, 556	619		619, 614, 611, 610		671,656	671, 669, 665, 662
568	<b>X</b> 00 400	, , , ,	620		620, 619, 618, 611	672		672, 667, 666, 661
	· · · · · · · · · · · · · · · · · · ·	569, 568, 559, 557	621		621, 616, 615, 609		673,645	673, 666, 664, 663
570	570,503			· ·	622, 612, 610, 605	674		674, 671, 665, 660
571		, , , , , , , , , , , , , , , , , , , ,		623,555	623, 614, 613, 612	675	050 105	675, 674, 672, 669
572		, , , ,	624	COF 100	624, 617, 615, 612		676,435	676, 675, 671, 664
573	FF4 F03	573, 569, 567, 563		625,492	625, 620, 617, 613	677		677, 674, 673, 669
	· · · · · · · · · · · · · · · · · · ·	574, 569, 565, 560	626		626, 623, 621, 613	678	070 010	678, 675, 673, 663
575	575, 429	575, 572, 570, 569	627	COO 407	627, 622, 617, 613		679,613	679, 676, 667, 661
576	F77 FF0	576, 573, 572, 563		628,405	628, 626, 617, 616	680		680, 679, 650, 645
	577, 552	577, 575, 574, 569	629		629, 627, 624, 623	681		681, 678, 672, 670
578		578, 562, 556, 555	630	691 994	630, 628, 626, 623	682		682, 681, 679, 675
579		579, 572, 570, 567	631	031,324	631,625,623,617	683		683, 682, 677, 672

n	LFSR-2	LFSR-4	n	LFSR-2	LFSR-4	n	LFSR-2	LFSR-4
580		580, 579, 576, 574	632		632, 629, 619, 613	684		684, 681, 671, 666
581		581, 575, 574, 568	633	633,532	633, 632, 631, 626	685		685, 684, 682, 681
582	582,497	582, 579, 576, 571	634	634,319	634, 631, 629, 627	686	686,489	686, 684, 674, 673
583	583,453	583, 581, 577, 575	635		635, 631, 625, 621	687	687,674	687, 682, 675, 673
584		584, 581, 571, 570	636		636, 632, 628, 623	688		688, 682, 674, 669
585	585,464	585, 583, 582, 577	637		637, 636, 628, 623	689	689,675	689, 686, 683, 681
586		586, 584, 581, 579	638		638, 637, 633, 632	690		690, 687, 683, 680
587		587, 586, 581, 576	639	639,623	639, 636, 635, 629	691		691,689,685,678
692	692,393	692, 687, 686, 678	725		725, 720, 719, 716	758		758, 757, 746, 741
693		693, 691, 685, 678	726	726,721	726, 725, 722, 721	759	759,661	759, 757, 756, 750
694		694, 691, 681, 677	727	727,547	727, 721, 719, 716	760		760, 757, 747, 734
695	695,483	695, 694, 691, 686	728		728, 726, 725, 724	761	761,758	761, 760, 759, 758
696		696, 694, 686, 673	729	729,671	729, 726, 724, 718	762	762,679	762, 761, 755, 745
697	697,430	697, 689, 685, 681	730	730,583	730, 726, 715, 711	763		763, 754, 749, 747
698	698,483	698, 690, 689, 688	731		731, 729, 725, 723	764		764, 761, 759, 758
699		699, 698, 689, 684	732		732, 729, 728, 725	765		765, 760, 755, 754
700		700,698,695,694	733		733, 731, 726, 725	766		766, 757, 747, 744
701		701, 699, 697, 685	734		734, 724, 721, 720	767	767,599	767, 763, 760, 759
702	702,665	702, 701, 699, 695	735	735,691	735, 733, 728, 727	768		768, 764, 751, 749
703		, , , ,	736		736, 730, 728, 723	769	769,649	769, 763, 762, 760
704		704, 701, 699, 692	737	737,732	737, 736, 733, 732	770		770, 768, 765, 756
705	705,686	705, 704, 698, 697	738	738,391	738, 730, 729, 727	771		771, 765, 756, 754
706		706, 697, 695, 692	739		739, 731, 723, 721	772	772,765	772, 767, 766, 764
707		707, 702, 699, 692		740,587		773		773, 767, 765, 763
708	708,421	, , , , , , , , , , , , , , , , , , ,	741		741, 738, 733, 732	774	774,589	774,767,760,758
709		, , , , , , , , , , , , , , , , , , ,	742		742, 741, 738, 730	775	775,408	775, 771, 769, 768
710			743	743,653	743, 742, 731, 730	776		776, 773, 764, 759
711	711,619	, , , , , , , , , , , , , , , , , , ,	744		744, 743, 733, 731	777	777,748	
712				· · · · · · · · · · · · · · · · · · ·	745, 740, 738, 737	778	778,403	
	713,672	713, 706, 703, 696		746,395		779		779, 776, 771, 769
714	714,691	714,709,707,701			747, 743, 741, 737	780		780, 775, 772, 764
715		715, 714, 711, 708			748, 744, 743, 733	781		781,779,765,764
	716,533	716, 706, 705, 704			749, 748, 743, 742			782, 780, 779, 773
717		717,716,710,701			750, 746, 741, 734		783,715	783, 782, 776, 773
718		718, 717, 716, 713		751,733	, , , ,	784		784, 778, 775, 771
719	719,569	719, 711, 710, 707			752, 749, 732, 731	1	785,693	785, 780, 776, 775
720		720, 718, 712, 709		· /	, , , ,	786		786, 782, 780, 771
721	1			754,735	754, 742, 740, 735	1024		1024, 1015, 1002, 1001
722	722,491	722, 721, 718, 707			755, 754, 745, 743			2048, 2035, 2034, 2029
723		, , , , , , , , , , , , , , , , , , ,		756,407	, , , ,	4096		4096, 4095, 4081, 4069
724		724, 719, 716, 711	757		757, 756, 751, 750			

## References

- [1] A. Ahmad and A.M. Elabdalla. An efficient method to determine linear feedback connections in shift registers that generate maximal length pseudo-random up and down binary sequences. *Computers and Electrical Engineering*, 23(1):33–39, 1997.
- [2] Peter Alfke. Application Note: Efficient Shift Registers, LFSR Counters, and Long Pseudo- Random Sequence Generators. Technical report, Xilinx Inc., San Jose, CA, 1996. App. note XApp052.
- [3] J. Brillhart, D.H. Lehmer, J.L. Selfridge, B. Tuckerman, and S.S. Wagstaff Jr. Factorizations of  $b^n \pm 1$ , b= 2, 3, 5, 6, 7, 10, 11, 12 up to High Powers. *Contemporary Mathematics*, 22, 2002.
- [4] D.W. Clark and L.J. Weng. Maximal and near-maximal shift register sequences: efficient event-counters and easy discrete logarithms. *Computers, IEEE Transactions on*, 43(5):560–568, 1994.
- [5] T.C.A. Molteno and R. W. Ward. Table of prime factors of Mersenne numbers. Technical report, University of Otago, Dunedin, New Zealand, 2007. http://www.physics.otago.ac.nz/px/research/electronics/papers/technical-reports/mersenne\_factor\_table.pdf.
- [6] M.J.B. Robshaw. Stream Ciphers. RSA Labratories, 25, 1995.
- [7] L.T. Wang and E.J. McCluskey. Hybrid designs generating maximum-length sequences. Computer-Aided Design of Integrated Circuits and Systems, IEEE Transactions on, 7(1):91–99, 1988.



Electronics Group
Department of Physics
University of Otago
elec.otago.ac.nz

