## The 4th Project Report

## Samira Vaez Barenji

Student ID Number: 9921384

In this project, you can find a solution to the maximum number of edge disjoint paths problem in a graph. Two paths are said to be edge disjoint if they don't share any edge.

This problem can be solved by reducing it to maximum flow problem, following the below steps:

- 1) Consider the given source and destination as source and sink in flow network. Assign unit capacity to each edge.
- 2) Find the maximum flow from source to sink.
- 3) The maximum flow is equal to the maximum number of edge-disjoint paths.

When we run Ford-Fulkerson, we reduce the capacity by a unit. Therefore, the edge cannot be used again. So, the maximum flow is equal to the maximum number of edge-disjoint paths. We use Edmonds-Karp implementation of Ford-Fulkerson algorithm to find the maximum flow from source to sink.

So, we can use a maximum flow algorithm to find k edge-disjoint, s-t paths in a graph. Embedded within any flow of value k on a unit-capacity graph there are k edge-disjoint paths. In other words, the value of the flow or the capacity of the minimum cut gives us the number of edge disjoint paths.

The idea of Edmonds-Karp is to use BFS in Ford Fulkerson implementation as BFS always picks a path with minimum number of edges. When BFS is used, the worst-case time complexity can be reduced to O(VE<sup>2</sup>). This implementation uses adjacency matrix representation though where BFS takes O(V<sup>2</sup>) time, the time complexity of the above implementation is O(EV<sup>3</sup>), where E is the number of edges and V is the number of vertices.

You can find the resulting output image below for a sample unweighted graph.