

Duality

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Introduction

Consider the following problem:

$$\begin{aligned} \min f_0(\underline{x}) \\ \text{s.t.} \\ f_i(\underline{x}) \leq 0, \quad i = 1, 2, \dots, r \\ h_i(\underline{x}) = 0, \quad i = 1, 2, \dots, s \end{aligned}$$

Let p^* be its optimum value. Can we obtain bounds on p^* such that $l < p^* < u$?

As a primitive guess, we could bound p^* by $-\infty / +\infty$, but it would be of no use. We would like to have a tight bound.

An Alternate Problem

The problem considered above is constrained. To state the same optimisation problem without constraints, we could represent it as:

$$\min f_0(\underline{x}) + \sum_{i=1}^r I(f_i(\underline{x})) + \sum_{i=1}^s I_0(h_i(\underline{x}))$$

where,

$$I(y) = \begin{cases} 0 & \text{if } y \leq 0 \\ \infty & \text{otherwise} \end{cases}$$

$$I_0(y) = \begin{cases} 0 & \text{if } y = 0 \\ \infty & \text{otherwise} \end{cases}$$

It is clear from the definitions of I and I_0 , that for any x which is infeasible, i.e not satisfying the constraints of the original problem, the function value is ∞ . Hence, minimisation is implicit within the constraint set.

Due to their nature, the functions I , and I_0 can be called the "displeasure" terms.

Displeasure functions

The functions I and I_0 as defined above are non-differentiable, which is called hard displeasure. This is rather inconvenient as differentiability in optimisation problems is very helpful.

So, we now consider the following definitions:

$$\begin{aligned}
I(y) &= \lambda y \\
I_0(y) &= \nu y \\
I \text{ and } I_0 &\text{ are differentiable}
\end{aligned}$$

Lagrangian

Modifying the proposed alternative representation in the earlier section with soft displeasure, we get:

$$\min f_0(\underline{x}) + \sum_{i=1}^r \lambda_i(f_i(\underline{x})) + \sum_{i=1}^s \nu_i(h_i(\underline{x})), \lambda_i \geq 0$$

This problem is a proxy for the original problem, called the Primal, and is called the Lagrangian.

The Lagrangian w.r.t the Primal(P), is denoted as

$$\mathcal{L}(\underline{x}, \underline{\lambda}, \underline{\nu}) = f_0(\underline{x}) + \sum_{i=1}^r \lambda_i(f_i(\underline{x})) + \sum_{i=1}^s \nu_i(h_i(\underline{x})), \lambda_i \geq 0$$

Properties of the Lagrangian

For any feasible \underline{x} ,

$$f_0(\underline{x}) \geq \mathcal{L}(\underline{x}, \underline{\lambda}, \underline{\nu})$$

Since $p^* = \inf_{\underline{x}} f_0(\underline{x})$,

$$\begin{aligned}
p^* &= \inf_{\underline{x}} f_0(\underline{x}) \geq \inf_{\text{feasible } \underline{x}} \mathcal{L}(\underline{x}, \underline{\lambda}, \underline{\nu}) \\
\implies p^* &\geq \inf_{\underline{x}} \mathcal{L}(\underline{x}, \underline{\lambda}, \underline{\nu}), \quad \forall \lambda_i \geq 0, \nu
\end{aligned}$$

We define $g(\underline{\lambda}, \underline{\nu}) = \inf_{\underline{x}} \mathcal{L}(\underline{x}, \underline{\lambda}, \underline{\nu})$

$$\implies p^* \geq g(\underline{\lambda}, \underline{\nu}), \quad \forall \lambda_i \geq 0, \nu$$

To get the best lower bound, we solve the following optimisation problem:

$$\begin{aligned}
&\max g(\underline{\lambda}, \underline{\nu}) \\
&\text{s.t. } \lambda_i \geq 0
\end{aligned}$$

This problem (D) is called the dual to the problem P.

$\underline{\lambda}$ and $\underline{\nu}$ are called dual variables or Lagrange multipliers.

Properties of the Dual Objective

$$\begin{aligned} g(\underline{\lambda}, \underline{\nu}) &= \inf_{\underline{x}} \mathcal{L}(\underline{x}, \underline{\lambda}, \underline{\nu}) \\ &= \inf_{\underline{x}} f_0(\underline{x}) + \sum_{i=1}^r \lambda_i(f_i(\underline{x})) + \sum_{i=1}^s \nu_i(h_i(\underline{x})) \end{aligned}$$

The dual objective g is the infimum of an affine function of λ and ν , i.e. the minimum of an affine (which is also concave).

$\implies g(\underline{\lambda}, \underline{\nu})$ is always a concave function

Since the dual is the maximization of a concave function, it is always a convex optimisation problem.

Duality

Let d^* denote the optimal value of the dual problem.

Weak Duality

$$p^* \geq d^*$$

We began with trying to find bounds on the optimum value of the primal. We defined the Lagrangian, and showed that it's infimum lower bounds p^* .

Since the dual problem is maximization of this lower bound (to make it tighter), $p^* \geq d^*$ always.

Strong Duality

This is the case when,

$$p^* = d^*$$

Illustration with Linear Programs

Consider the following linear optimization problem

$$\begin{aligned} \min \quad & \underline{c}^T \underline{x} \\ \text{s.t.} \quad & \\ & A\underline{x} = \underline{b} \\ & \underline{x} \geq 0 \end{aligned}$$

Equivalently, converting constraints to the notation used above

$$\begin{aligned} \min \quad & \underline{c}^T \underline{x} \\ \text{s.t.} \quad & \\ & A\underline{x} = \underline{b} \\ & -\underline{x} \leq 0 \end{aligned}$$

The Lagrangian ($\mathcal{L}(\underline{\lambda}, \underline{\nu})$) is given by

$$\begin{aligned} & \underline{c}^T \underline{x} + \sum_{i=1}^n \lambda_i (-x_i) + \sum_{i=1}^m \nu_i (a_i^T \underline{x} - b_i), \quad \underline{a}_i \text{ is the } i\text{th row vector of } A \\ &= \underline{c}^T \underline{x} - \underline{\lambda}^T \underline{x} + \underline{\nu}^T (A\underline{x} - \underline{b}) \\ &= (\underline{c}^T - \underline{\lambda}^T + \underline{\nu}^T A) \underline{x} - \underline{\nu}^T \underline{b} \\ &= (\underline{c} - \underline{\lambda} + A^T \underline{\nu})^T \underline{x} - \underline{\nu}^T \underline{b} \end{aligned}$$

The dual objective $g(\underline{\lambda}, \underline{\nu})$ is

$$\inf_{\underline{x}} \mathcal{L}(\underline{x}, \underline{\lambda}, \underline{\nu}) = \begin{cases} -\infty, & \text{if } \underline{c} - \underline{\lambda} + A^T \underline{\nu} \neq 0 \\ -\underline{\nu}^T \underline{b}, & \text{if } \underline{c} - \underline{\lambda} + A^T \underline{\nu} = 0 \end{cases}$$

Therefore, the Dual problem is

$$\begin{aligned} \max_{\underline{\lambda}, \underline{\nu}} \quad & g(\underline{\lambda}, \underline{\nu}) \\ \text{s.t.} \quad & \\ & \lambda_i \geq 0 \\ & \underline{c} - \underline{\lambda} + A^T \underline{\nu} = 0 \end{aligned}$$

Equivalently,

$$\begin{aligned} \max_{\underline{\lambda}, \underline{\nu}} \quad & g(\underline{\lambda}, \underline{\nu}) \\ \text{s.t.} \quad & \underline{c} + A^T \underline{\nu} \geq 0 \end{aligned}$$