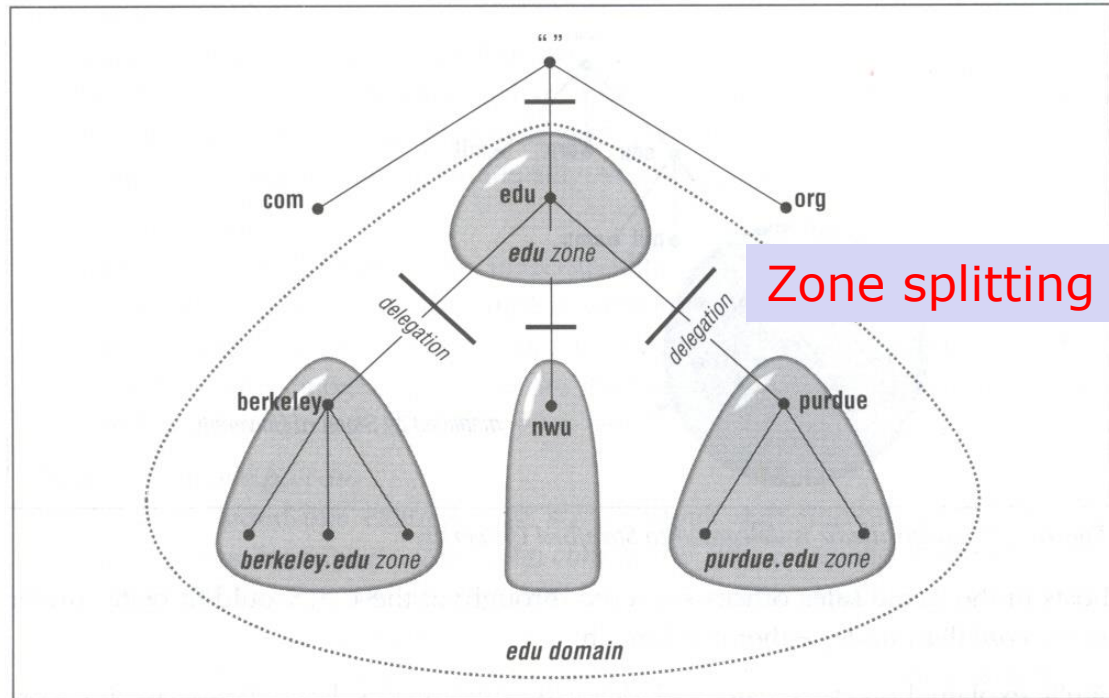


Chap. 8 (added) DNS Security

- DNS Service
- DNS Attacks
- DNSSEC (DNS Security)

DNS name service

- **Domain**: a subtree of a domain name space
- **Zone**: some part of the domain name space
 - **Authoritative name servers** have complete information about a zone
- Distributed management thru delegation



DNS name servers

- Distributed database: no server has all name-to-IP address mappings
- local name servers:
 - each ISP, company has local (default) name server
 - host DNS query first goes to local name server
- authoritative name server:
 - keeps DNS info. for hosts within a specific zone
 - can perform name/address translation for that host's name

DNS name servers

□ root name server:

- contacted by local name server that can not resolve name
- contacts authoritative name server if name mapping not known
- gets mapping
- returns mapping to local name server
- dozen root name servers worldwide

DNS: iterated queries

□ Resolver:

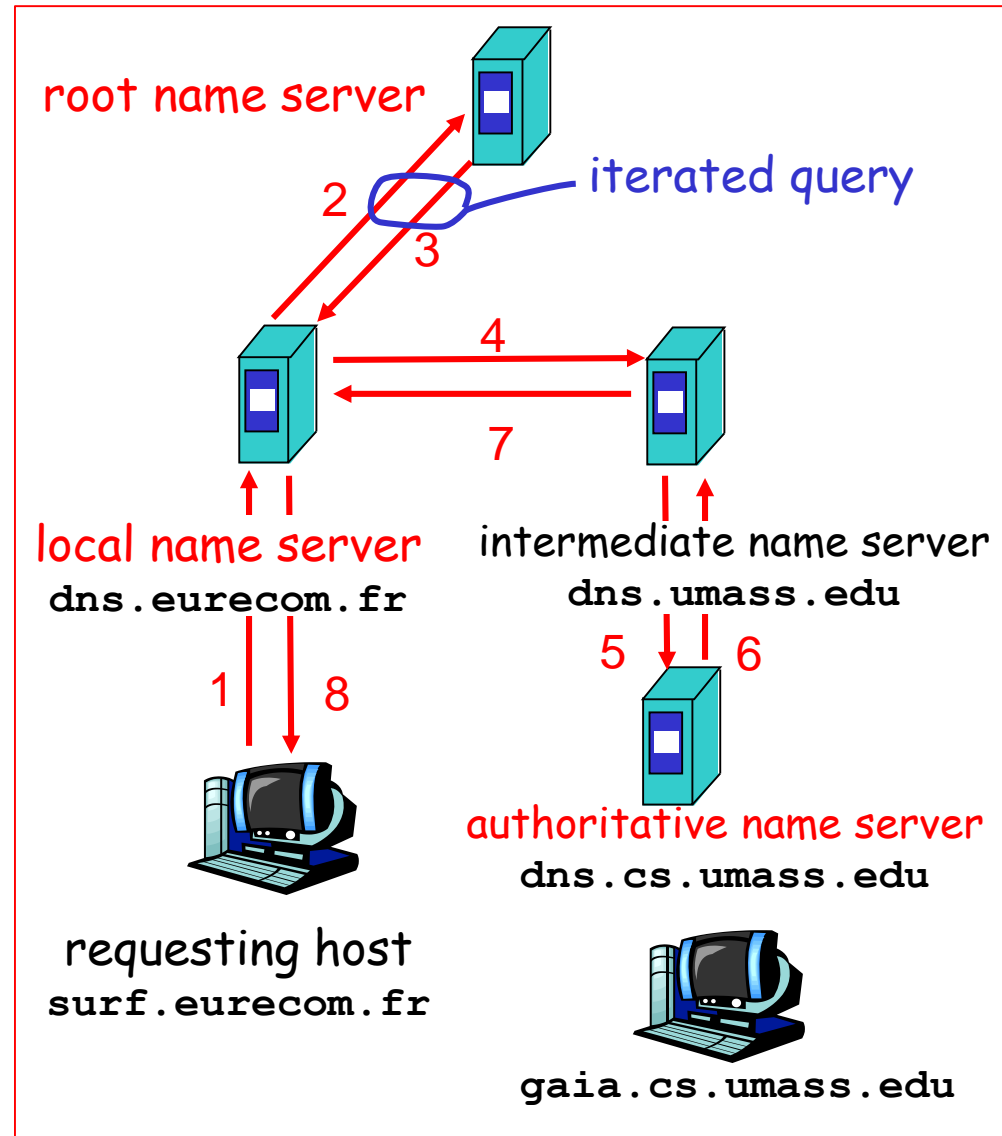
- queries the name server
- interpret the responses
- return result to user

□ recursive query:

- puts burden of name resolution on contacted name server

□ iterated query:

- contacted server replies with name of server to contact



DNS: caching records

□ DNS caching

- Once name server learns mapping, it caches mapping
- cache entries **timeout** (disappear) after some time

DNS Resource Records

- DNS: distributed DB storing resource records (RR)

RR format: (name, value, type, ttl)

- Type=A

- name is hostname
- value is IP address

- Type=NS

- name is domain
- value: IP addr of authoritative name server for this domain

- Type=CNAME

- name is an alias name for some "canonical" name
- value is canonical name

- Type=MX

- value is hostname of mail-server associated with a name (domain name)

DNS protocol, messages

□ DNS protocol : query and reply messages, both with same message format

□ msg header

- identification: 16 bit # to match query and reply
- flags:
 - query or reply
 - recursion desired
 - recursion available
 - reply is authoritative

identification	flags
number of questions	number of answer RRs
number of authority RRs	number of additional RRs
questions (variable number of questions)	
answers (variable number of resource records)	
authority (variable number of resource records)	
additional information (variable number of resource records)	

DNS query

□ Troubleshooting tools for DNS

- **nslookup**
- **dig** (domain information gopher)

```
$ dig cic.ulsan.ac.kr ; looks up A records for
; cic.ulsan.ac.kr
```

```
$ dig ulsan.ac.kr mx ; looks up MX records for ; ulsan.ac.kr
```

```
$ dig @a.root-servers.net . ns
```

DNS query

```
gslacks[pts/3]:~> dig www.cse.ucsc.edu
; <<>> DiG 8.2 <<>> www.cse.ucsc.edu
;; res options: init recurs defnam dnsrch
;; got answer:
;; ->>HEADER<- opcode: RESPONSE, status: NOERROR, id: 4
;; flags: aa rd ra; QUERY: 1, ANSWER: 2, AUTHORITY: 2, ADDITIONAL: 2
;; QUERY SECTION:
;;   www.cse.ucsc.edu, type = A, class = IN
;; ANSWER SECTION:
www.cse.ucsc.edu.      1D IN CNAME   ftp.cse.ucsc.edu.
ftp.cse.ucsc.edu.     1D IN A       128.114.48.173
;; AUTHORITY SECTION:
cse.ucsc.edu.         1D IN NS      services.cse.ucsc.edu.
cse.ucsc.edu.         1D IN NS      fs1.cse.ucsc.edu.
;; ADDITIONAL SECTION:
services.cse.ucsc.edu. 1D IN A       128.114.48.10
fs1.cse.ucsc.edu.     1D IN A       128.114.48.11
.....
```

qr: query packet
aa: authoritative answer
rd: recursion desired
ra: recursion available

DNS query

```
gslacks[pts/3]:~> dig www.cse.ucsc.edu
```

```
; <<>> DiG 8.2 <<>> www.cse.ucsc.edu
;; res options: init recurs defnam dnsrch
;; got answer:
;; ->>HEADER<<- opcode: RESPONSE, status: NOERROR, id: 4
;; flags: rd ra; QUERY: 1, ANSWER: 2, AUTHORITY: 2, ADDITIONAL: 2
;; QUERY SECTION:
;;      www.cse.ucsc.edu, type = A, class = IN
;; ANSWER SECTION:
www.cse.ucsc.edu.      23h59m57s IN CNAME  ftp.cse.ucsc.edu.
ftp.cse.ucsc.edu.     23h59m57s IN A    128.114.48.173
;; AUTHORITY SECTION:
cse.ucsc.edu.         15h54m19s IN NS   services.cse.ucsc.edu.
cse.ucsc.edu.         15h54m19s IN NS   fs1.cse.ucsc.edu.
;; ADDITIONAL SECTION:
services.cse.ucsc.edu. 15h54m19s IN A    128.114.48.10
fs1.cse.ucsc.edu.     15h54m19s IN A    128.114.48.11
```

```
.....
```

Zone transfer

□ Name servers for a zone

- **Primary (master) name server:** reads zone data from a file on its host
- **Secondary (slave) name server:** gets zone data from primary name server that is authoritative for the zone

□ Zone transfer:

- When a slave name server starts up, it contacts its master server and pulls the zone data over

```
$ dig @192.168.1.85 dumb.target.net axfr
```

```
; transfer zone dumb.target.net from 192.168.1.85
```

DNS Package: BIND

- DNS package in unix: BIND
 - it has many security problem history
- DNS daemon in unix
 - a process called "named"
 - Configuration file: `/etc/named.conf`

DNS configuration: /etc/named.conf

```
// Config file for name server
```

```
options {  
    directory "/var/named";           zone file directory  
    version "bla bla bla"  
    allow-transfer {192.154.1.30};    secondary name server  
};
```

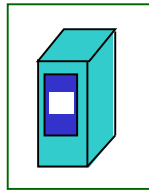
```
zone "." in {  
    type hint;                        contains name and IP address of  
    file "root.hints";               root name server  
};
```

```
zone "foobar.brian.edu" in {          For converting from name to address  
    type master  
    file "pz/db.foobar.brian.edu"
```

```
zone "1.154.192.in-addr.arpa" in {    For converting from address to name  
    type master;  
    file "pz/db.192.154.1";  
};
```

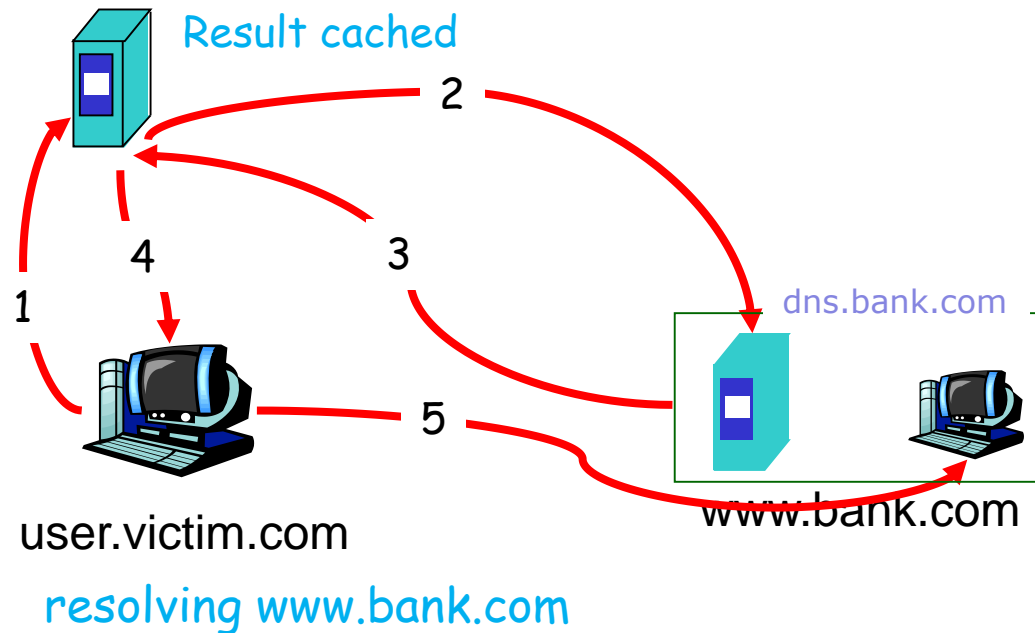
DNS Operation – Normal

dns.hacker.com



- Normally, DNS is resolved with an authoritative server
- (Not shown is lookup to .com root server)

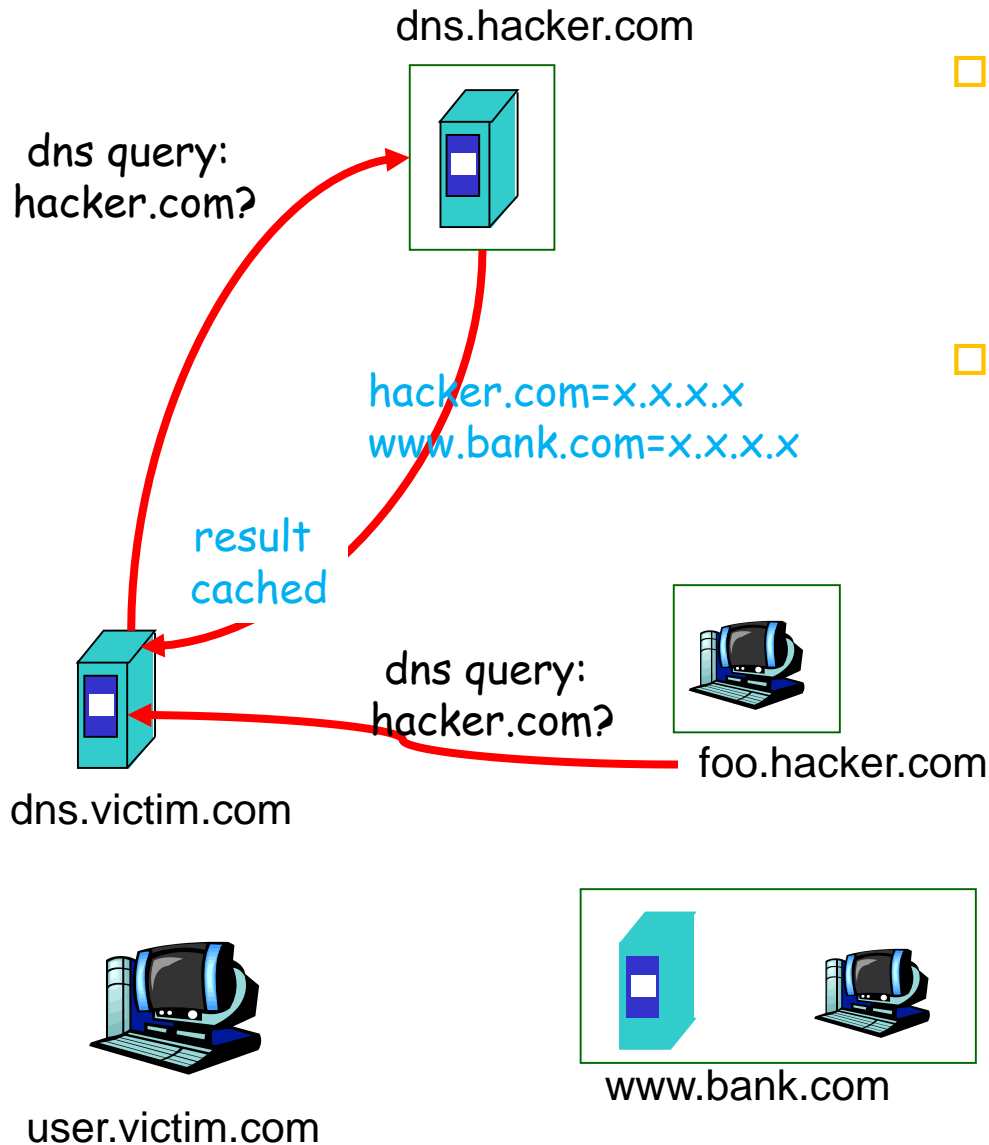
dns.victim.com



DNS Cache Poisoning

- URL spoofing attack:
 - register a similar name of someone you are attacking
 - e.g., `www.ibn.com`, ...
- Cache poisoning is a more sophisticated version of the same idea
- The two key ideas for the attack are:
 - The query number (and reply id) are often predictable if you can learn earlier IDs
 - DNS caches previous results

Old Version of the Attack

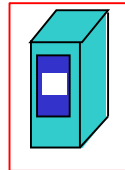


- On older versions of BIND, replies that came in could include additional lookup information ([additional info section](#))
- This additional lookup information would be stored for future use

Old Version of the Attack

- The next time a user from victim.com looks for the bank, is redirected to the hacker's site
- All the hacker has to do is create a mock-up of the banks site to get passwords, etc.
- The old attack no longer work since BIND was patched to ignore additional information

dns.hacker.com



result
cached

hacker.com=x.x.x.x
www.bank.com=x.x.x.x

x.x.x.x



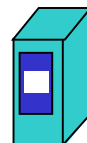
www.hacker.com

dns.victim.com

dns query:
www.bank.com?



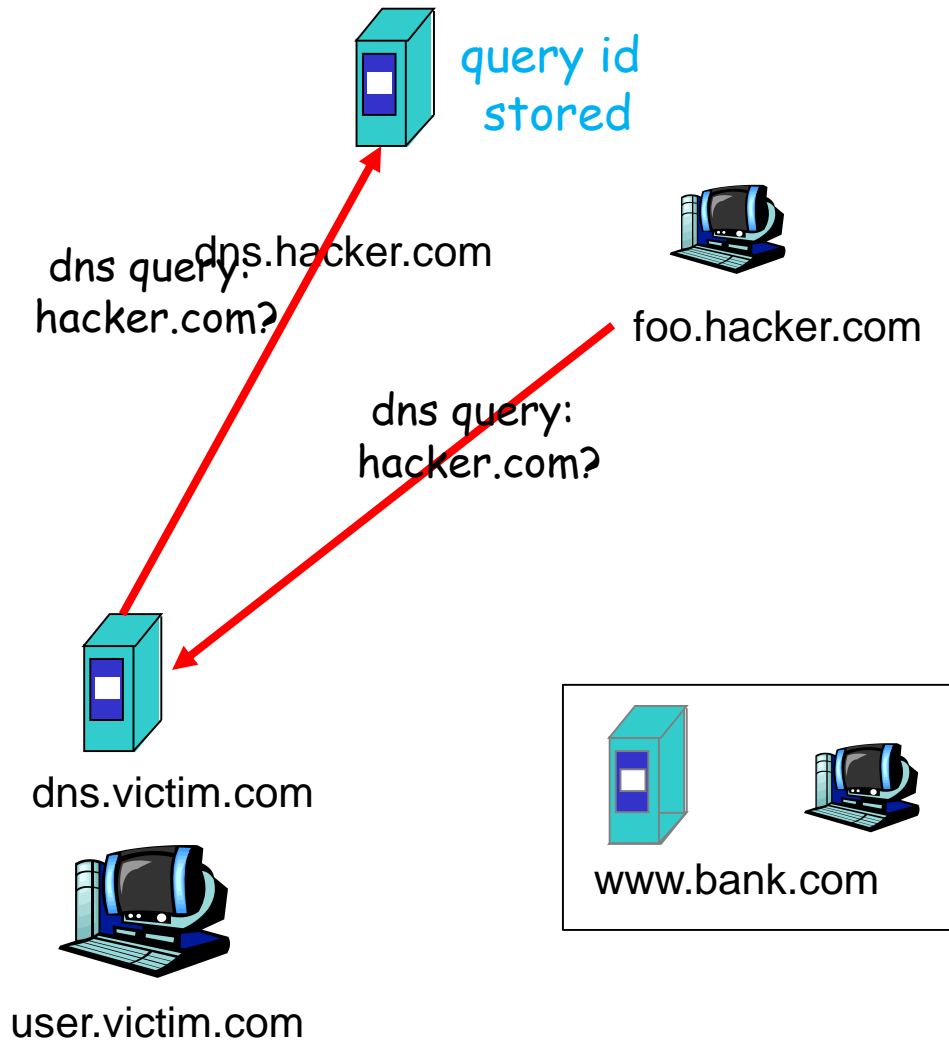
user.victim.com



www.bank.com

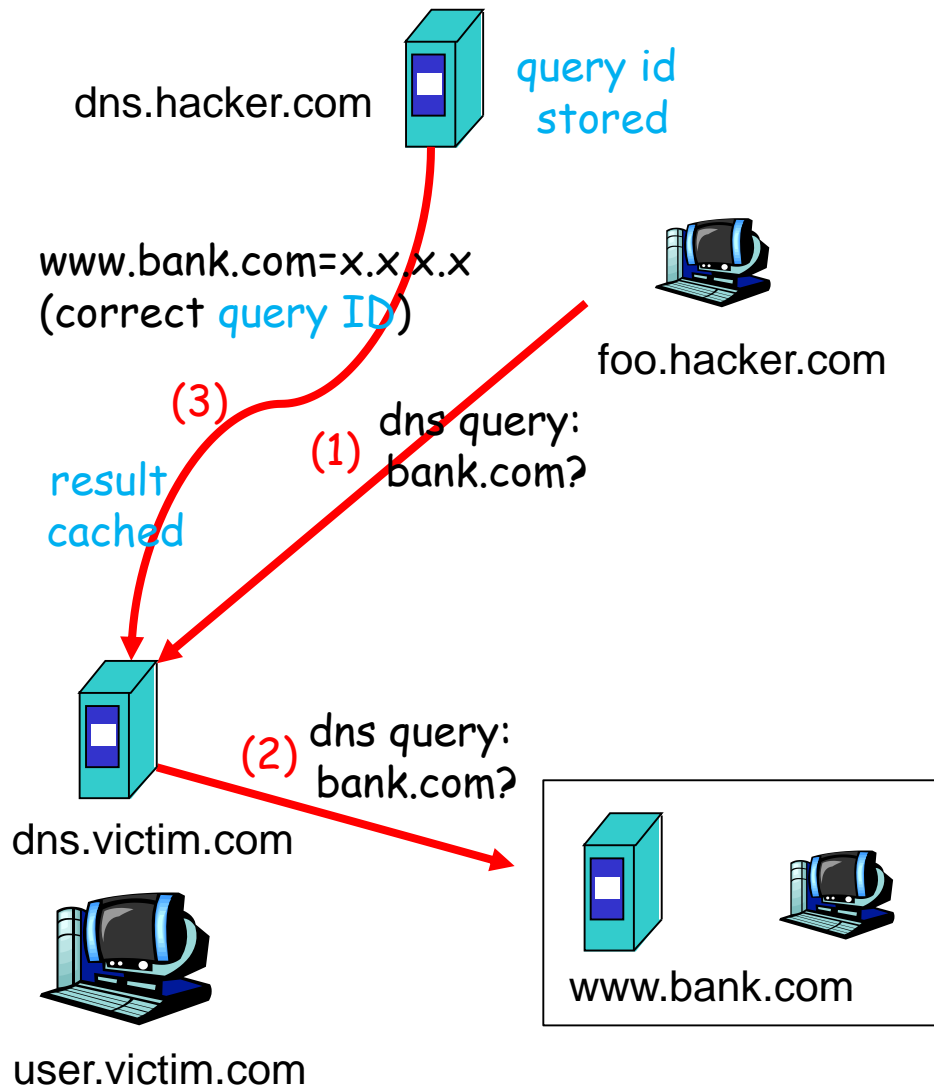


The Attack - more



- The first step of the attack is to learn victim.com's current **query id** number
- The simplest way to do that is to get the victim to query the attacker's DNS machine
- This querying can be repeated many times to know how the query ids change over time

The Attack - more



- Early versions of DNS servers deterministically incremented the ID field; it was patched by **random query IDs**
- **Birthday attack**
 - 16-bit query ID (only 65,536 options)
 - the resolver sends many queries, with different IDs, at the same time
 - Send hundreds of reply with random transaction IDs at the same time
 - Increase the probability of making a correct guess

Defenses

- New versions of BIND have harder to predict query IDs
- **DNSSEC** : RFC 2535
 - a secure version of DNS with RSA signed DNS records
- **Authentication of DNS responses**
 - Each DNS response has a signature of the requested RR
 - Resolver authenticates response using the public key of the authoritative name server

Defenses

□ Split-split defense:

- One DNS server for resolving names for users inside your domain; this server doesn't respond to outside queries
- Another separate DNS server is setup for responding to queries from outsiders
- The two never exchange information
- Your users are not subject to poisoned information

Defenses in /etc/named.conf

```
// name server config file
options {
    directory "/var/named";
    version "version bla bla"
    allow-transfer {192.154.1.30};
    allow-query { any; };
    allow-recursion { 192.154.1.0/24; };
};

zone "foobar.brian.edu" in {
    type master
    file "pz/db.foobar.brian.edu"
    allow-query { 192.154.1.0/24; }
};

zone "1.154.192.in-addr.arpa" in {
    type master;
    file "pz/db.192.154.1";
    allow-query { 192.154.1.0/24; }
};
```

Global access control list
(ACL): Only this subnet
can query us

Zone specific ACL: take
precedence over a global ACL

Homework #3

□ BIND package 기능 조사:

- BIND package 기능
- /etc/named.conf 설정파일 기능
- 호스트에서 name resolution 하는 과정: OS별(windows, linux) 조사
- 파일 제출: "hw3-학번-이름.hwp"

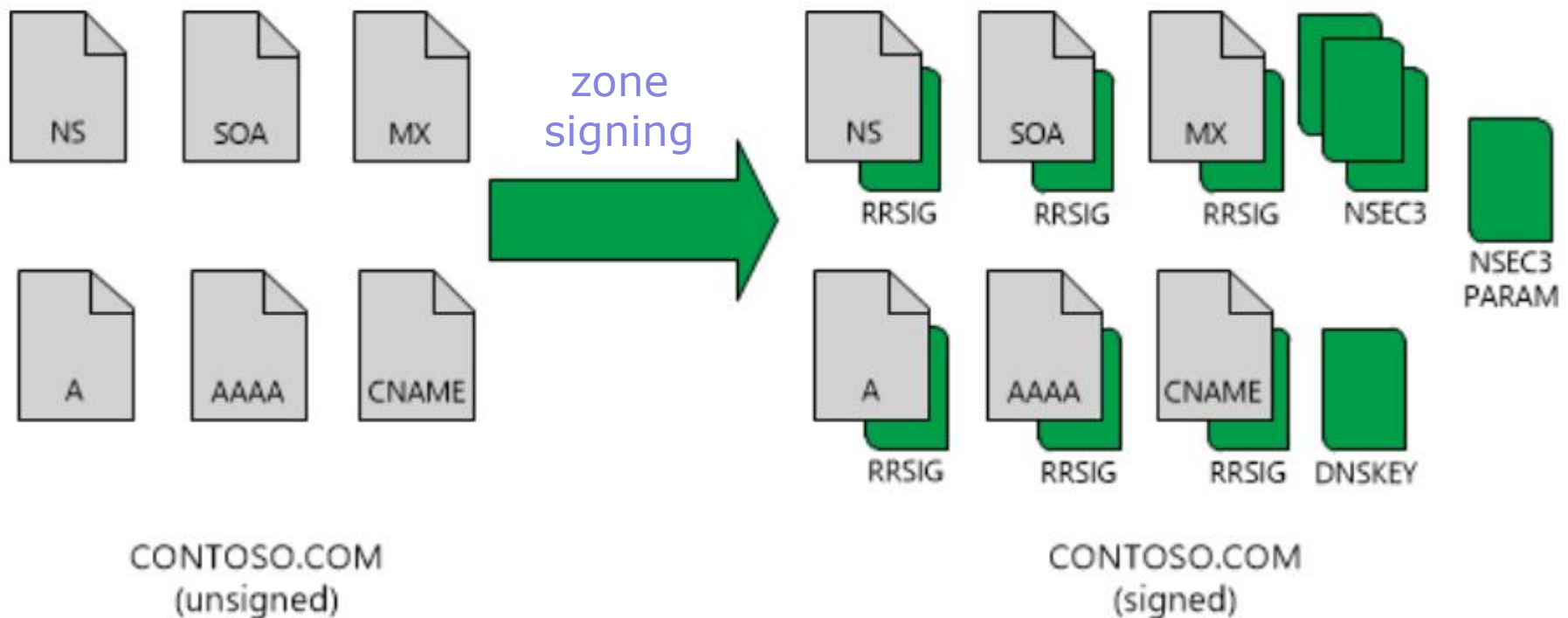
DNSSEC (DNS Security)

□ DNSSEC:

- Security extension of DNS service
- Authoritative name servers secure their zones by performing [zone signing](#)
- provides [end-to-end authentication](#) using digital signatures b/w a resolver and an authoritative server
- defines a set of [new resource record types](#) and modifications to the existing DNS protocol
- RFC 4033, 4034, 4035

DNSSEC (DNS Security)

□ DNSSEC:



DNSSEC (DNS Security)

□ DNSSEC:

- each DNSSEC-enabled authoritative NS can have two public keys
- ZSK (Zone Signing Key):
 - used to sign the RRset of the Zone
- KSK (Key Signing Key):
 - used to sign DNSKEY RRs
 - can be used as a trust anchor

DNSSEC (DNS Security)

□ DNSSEC Resource Records:

- DNSKEY resource record: defines the public key

domain name	TTL	Class	Type	Flags	Proto	Alg	Public key
corpxyz.com	1296000	IN	DNSKEY	257	3	5	20181231235959hiZsq1gPtq4vSymSxBsqzueQW4jrjCBsCZBvwQMgE07dxaOeTpwpagI7XhOjlarzM8nTf1PJ+4av1Kr0EfwS0tEAwD7Isvt2vW24cE

Flags: DNSKEY-flag(bit7), KEY-type(bit15: ZSK or KSK);
256 (ZSK), 257 (KSK)

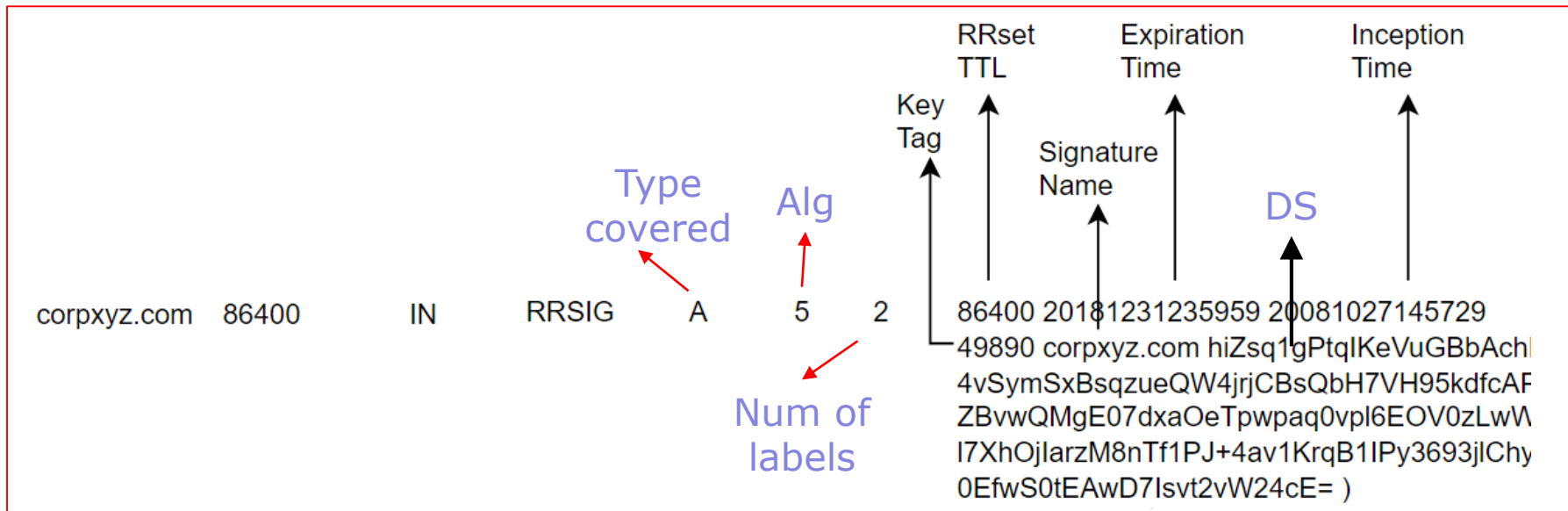
Proto: DNSSEC protocol type (3)

Alg: type of the public key algorithm
(RSA/MD5, DSA, RSA/SHA1, etc.)

DNSSEC (DNS Security)

□ DNSSEC Resource Records:

- RRSIG resource record: defines the signature of an RRset



Type covered: RR type covered by the RRSIG record

Alg: cryptographic algorithm used to create the DS (RSA/MD5, RSA/SHA-1, ...)

Num of labels: the number of labels in the owner name of the signed records

Key tag: key tag value of the DNSKEY RR that validates the signature

RRset TTL: TTL value of the RRset covered by the RRSIG record

DS: Base64 encoding of the digital signature

DNSSEC (DNS Security)

□ DNSSEC Resource Records:

- DS (Delegation Signer) resource record:
 - contains a hash of a child zone's KSK and can be used as a trust anchor
 - creates a secure delegation point for a signed subzone

corpxyz.com	86400	IN	DS	25924	5	1	49D2801B! B622B1F8'
				Key tag	Alg	Digest type	Digest value

Key tag: key tag value of the DNSKEY RR to which this DS RR refers

Alg: the algorithm of the DNSKEY RR to which this DS RR refers

Digest type: algorithm used to create the digest; 1 (SHA-1), 2 (SHA-256)

Digest : digest value of the DNSKEY RR to which this DS RR refers

DNSSEC (DNS Security)

□ DNSSEC Resource Records:

- a chain of trust thru DS (Delegation Signer) resource record

```
A      server1.corp100.com
A      ftp.corp100.com
A      sales.corp100.com
RRSIG  A  5  2  86400....
DS    25924 5  1 ←
```

corp100.com

```
A      server2.sales.corp100.com
A      ftp1.sales.corp100.com
RRSIG  A  5  2  86400....
DS    25854 5  1 ←
DNSKEY 256
DNSKEY 257
```

sales.corp100.com

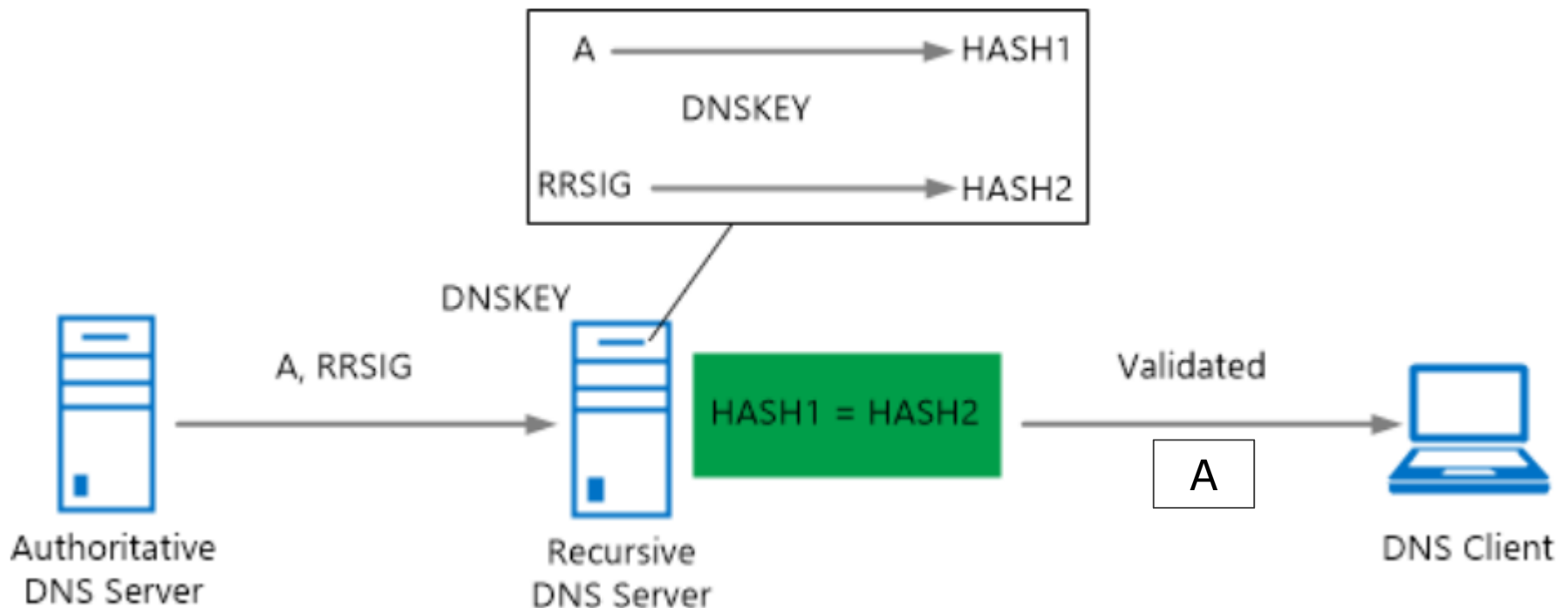
```
A      server3.nw.sales.corp100.com
A      ftp1.nw.sales.corp100.com
RRSIG  A  5  2  86400....
DNSKEY 256
DNSKEY 257
```

nw.sales.corp100.com

DNSSEC (DNS Security)

Validation process:

- A DNS server receives RRset and DNSKEY, RRSIG RRs
- uses the DNSKEY RR to validate responses from the authoritative DNS server by decrypting digital signatures



DNSSEC (DNS Security)

□ Zone signing process:

