# **Computer Organization**

**Lecture 2 - Performance** 

Reading: 1.4



Photo credit: Jurvetson via flickr

## Roadmap for the term: major topics

- Computer Systems Overview
- Performance
- Assembly Language
- Instruction sets (and Software)
- Logic & Arithmetic
- Processor Implementation
- Memory Systems

#### **Performance Outline**

- ► Motivation 3
- Defining Performance
- Common Performance Metrics
- Benchmarks
- Amdahl's Law

#### **Motivation**

Goal: Learn to "measure, summarize and report" performance of a computer system

- Why study performance?
  - To make intelligent decisions when choosing a system
  - ▶ To make intelligent decisions when designing a system
  - Understand impact of implementation decisions
- Challenges
  - ▶ How do we measure performance accurately?
  - How do we compare performance fairly?

# What's a good measure of performance?

- Response Time
  - How long does it take to complete a single task?
- **▶** Throughput
  - ▶ How many tasks are completed per unit time
- ▶ The measure we use depends on the application

# **Execution Time vs. Throughput**

- Analogy: passenger airplanes (book Figure 1.13)
  - Concorde fastest "response time" for an individual user
  - Boeing 747 highest passenger throughput

Airplane	Passenger Capacity	Cruising Range (miles)	Cruising Speed (mph)	Passenger Throughput (passengers x mph)
Boeing 777	375	4630	610	228,750
Boeing 747	470	4150	610	286,700
Concorde	132	4000	1350	178,200
DC-8-50	146	8720	544	79,424

### **Performance Outline**

- Motivation
- Defining Performance

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- Common Performance Metrics
- Benchmarks
- Amdahl's Law

### **Relative Performance**

▶ For a given program on machine X:

Performance 
$$X = \frac{1}{\text{Execution time } X}$$

**▶** Comparing performance of machines:

Performance<sub>X</sub> > Performance<sub>Y</sub> if Execution Time<sub>X</sub> < Execution Time<sub>Y</sub>

# **Comparing Performance**

▶ We say "X is n times faster than Y" if:

$$\frac{\mathbf{Performance}_{\mathbf{X}}}{\mathbf{Performance}_{\mathbf{Y}}} = n$$

$$\frac{\text{Performance}_{X}}{\text{Performance}_{Y}} = \frac{\text{Execution time}_{Y}}{\text{Execution time}_{X}} = n$$

# **Example - Performance**

Machine	<b>Execution Time</b>
A	15 seconds
В	20 seconds



▶ Which machine is faster?



By how much?

#### **Performance Outline**

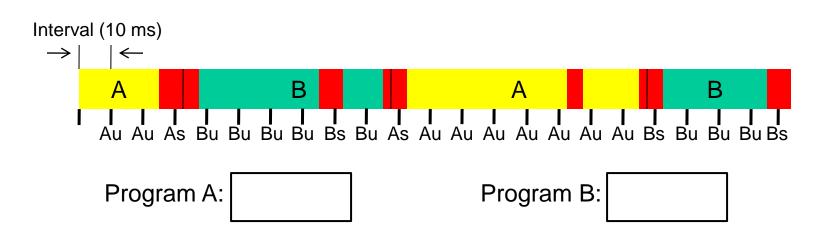
- Motivation
- Defining Performance
- ▶ Common Performance Metrics
- Benchmarks
- Amdahl's Law

# **Elapsed Time / Execution Time**

- ► Elapsed Time ("Wall Clock" Time)
  - ▶ How long does it take the program to complete?
  - System Performance based on elapsed time
- Execution Time (CPU Time)
  - ▶ How much time the CPU spent executing the program
    - User time time CPU spends executing program instructions
    - System time time CPU spends in OS on behalf of program
  - **▶ CPU Performance** based on CPU time

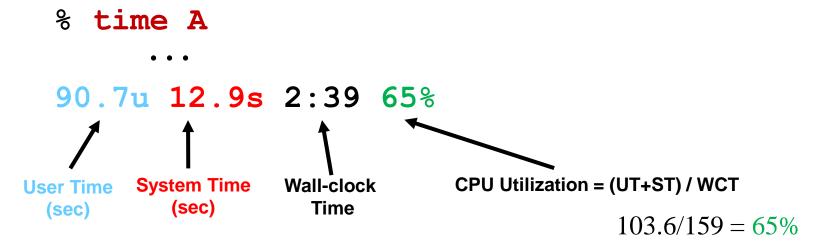
# Measuring CPU Time with the OS

- OS runs multiple programs
  - ▶ Timer interrupts programs at fixed intervals (e.g. 10ms)
  - OS decides whether to switch which program runs
  - Interval counter counts intervals when program is running in user mode (u) and system mode (s)
- Issues
  - Inaccuracy interval counter can miss changes (cancels out for long programs)



# **Measuring CPU Time with Unix/Linux**

- Interval-based timing in the "real world"
- Linux/Unix time command



### **Clocks and Performance**

▶ Clock signal - controls sequential circuit operation

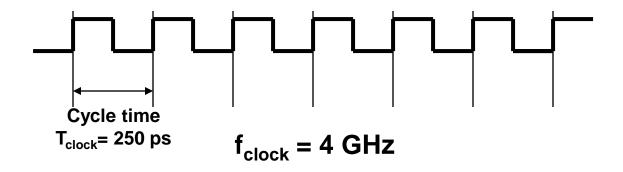


Clock frequency f<sub>clock</sub> = 1/T<sub>clock</sub>

초당 클락이 몇 번 바뀌는가?

#### **Clocks and Performance - CPU Time**

How do we relate clock to CPU Time?



CPU time = 
$$\frac{\text{clock cycles}}{\text{program}} \times \frac{\text{seconds}}{\text{clock cycle}}$$

- $1GHz = 10^9Hz$
- $1ps = 10^{-12}s$
- $1 \text{ns} = 10^{-9} \text{s}$

## **Measuring Elapsed Time with Counters**

#### Modern microprocessors include cycle counters

**Ex) Intel Pentium & Later Processors** 

- 64-bit Time Stamp Counter (TSC)
- Counts clock cycles since reset
- rdtsc instruction moves TSC to {edx,eax} registers
- Measure elapsed time by
  - ✓ Reading TSC at start
  - ✓ Reading TSC at end
  - ✓ Subtract & multiply by clock period

Get the number of clocks during program execution

#### Issues

- Context changes
- Impact of caching

# **Example - Clock Cycles**

Machine	<b>Execution Time</b>	Clock Freq.
А	15 seconds	3 GHz
В	20 seconds	2.5 GHz

WOF	many	CIOCK	cycles	does	A exec	ute?	

▶ How many clock cycles does B execute?

# **Clock Cycles per Instruction (CPI)**

- Consider the 68HC11 ...
  - ▶ ADDA 3 cycles (IMM) -> 5 cycles (IND, Y)
  - MUL 10 cycles
  - ▶ IDIV 41 cycles
- More complex processors have other issues...
  - Pipelining parallel execution (CPI=1!), but stalls occur
  - Memory system: cache misses, page faults, etc.
- How can we combine these into an overall metric?



# **Clock Cycles per Instruction (CPI)**

- Average number of clock cycles per instruction
- Measured for an entire program

$$CPU Clock Cycles = \frac{Instructions}{Program} \times \frac{Average Clock Cycles}{Instruction}$$

$$CPI = \frac{Average \ Clock \ Cycles}{Instruction} = \frac{CPU \ Clock \ Cycles/Program}{Instructions/Program}$$

# **Example - CPI**

Machine	<b>Clock Cycles</b>	Instructions
А	1.35 X 10 <sup>10</sup>	2.6 X 10 <sup>9</sup>
В	1.2 X 10 <sup>10</sup>	3.0 X 10 <sup>9</sup>

▶ What is the CPI of A?

▶ What is the CPI of B?

# **CPI Example**

- Computer A: Cycle Time = 250ps, CPI = 2.0
- Computer B: Cycle Time = 500ps, CPI = 1.2 (Same ISA)
- Which is faster, and by how much?

$$\begin{aligned} \text{CPU Time}_{A} &= \underbrace{\text{Instruction Count}} \times \text{CPI}_{A} \times \text{Cycle Time}_{A} \\ &= \text{IC} \times 2.0 \times 250 \text{ps} = \text{IC} \times 500 \text{ps} & \text{A is faster...} \end{aligned}$$

$$\begin{aligned} \text{CPU Time}_{B} &= \underbrace{\text{Instruction Count}} \times \text{CPI}_{B} \times \text{Cycle Time}_{B} \\ &= \text{IC} \times 1.2 \times 500 \text{ps} = \text{IC} \times 600 \text{ps} \end{aligned}$$

$$\begin{aligned} &\stackrel{\text{CPU Time}_{B}}{\text{CPU Time}_{A}} &= \underbrace{\frac{\text{IC} \times 600 \text{ps}}{\text{IC} \times 500 \text{ps}}} = 1.2 & \text{...by this much} \end{aligned}$$

...by this much

# **The Performance Equation**

#### ▶ The "Iron Law" of Performance

$$\frac{\text{CPU time}}{\text{program}} = \frac{\text{\# instructions}}{\text{program}} \times \frac{\text{clock cycles}}{\text{instruction}} \times \frac{\text{seconds}}{\text{clock cycle}}$$

$$= \frac{\text{Instruction Count} \times (\text{CPI})}{\text{Clock Rate}}$$

#### **CPI in More Detail**

Class 1: CPI = 3

Class 2: CPI = 5

10 Ins

30 Ins

▶ If *n* different instruction classes take different numbers of cycles per instruction (CPI)

Clock Cycles = 
$$\sum_{i=1}^{n} (CPI_i \times Instruction Count_i)$$
  
 $CC = 10 \times 3 + 30 \times 5 = 180$   
 $CPI = 180/40 = 4.5$ 

Weighted average CPI

$$CPI = \frac{Clock \ Cycles}{Instruction \ Count} = \sum_{i=1}^{n} \left( CPI_i \times \frac{Instruction \ Count}{Instruction \ Count} \right)$$
Relative frequency

 $CPI = 3 \times 10/40 + 5 \times 30/40 = 180/40 = 4.5$ 

# **Another CPI Example**

 Consider two compiled code sequences using instructions in classes A, B, C

Class	Α	В	С
CPI for class	1	2	3
IC in sequence 1	20	10	20
IC in sequence 2	40	10	10

- Which sequence executes more instructions?
- Which sequence is faster?
- What is the CPI for each sequence?

## **CPI Example**

Solution: calculate total cycles for each sequence, then CPI

Clock Cycles = 
$$\sum_{i=1}^{n} (CPI_i \times Instrunction Count_i)$$

Class	Α	В	С
CPI for class	1	2	3
IC in sequence 1	20	10	20
IC in sequence 2	40	10	10

- **▶** Sequence 1: IC = 50
  - Clock Cycles =  $20 \times 1 + 10 \times 2 + 20 \times 3 = 100$
  - Avg. CPI = 100/50 = 2.0

- **▶** Sequence 2: IC = 60
  - Clock Cycles =  $40 \times 1 + 10 \times 2 + 10 \times 3 = 90$
  - Avg. CPI = 90/60 = 1.5

# **Performance Summary**

#### The "BIG PICTURE"

CPU time = 
$$\frac{\text{\#instructions}}{\text{program}} \times \frac{\text{clock cycles}}{\text{instruction}} \times \frac{\text{seconds}}{\text{clock cycle}}$$

#### Performance depends on

- Algorithm: affects Instruction Count (IC), possibly CPI
- Programming language: affects IC, CPI
- Compiler: affects IC, CPI
- ► Instruction set architecture: affects IC, CPI, T<sub>clock</sub>
- ► Implementation Technology and Design: affects CPI, T<sub>clock</sub>

# **Clock Cycles and Performance - Example**

- Program runs on Computer A:
  - ▶ CPU Time: 10 seconds
  - ▶ Clock (rate): 400MHz = 400x10<sup>6</sup> cycles/sec
- Computer B can run clock faster
  - ▶ But, requires 1.2 X clock cycles to perform same task
  - Desired CPU Time: 6 Seconds
- What should the clock frequency of Computer B be to reach this target?
- Key to approach: Performance equation

CPU time = 
$$\frac{\text{clock cycles}}{\text{program}} \times \frac{\text{sec onds}}{\text{clock cycle}} = \frac{\text{clock cycles}}{\text{clock frequency}}$$

## Clock Cycles and Performance - Example (cont'd)

•	First step: find clock cycles executed by Computer A
,	Second step: find clock cycles executed by Computer

### Clock Cycles and Performance - Example (cont'd)

 Third step: given clock cycles and CPU time, solve for clock rate of Computer B

# **Summary**

- Response time vs. Throughput
- Elapsed time vs. execution time (CPU time)
- How does OS measure CPU time?
- What is the problem if interrupt interval is too short in measuring CPU time?
- Why is it difficult to measure the execution time of a program?
- Why do we use the number of clock cycles to measure the elapsed time?
- ▶ Why do we use CPI to measure CPU time?