Chapter 1: roadmap

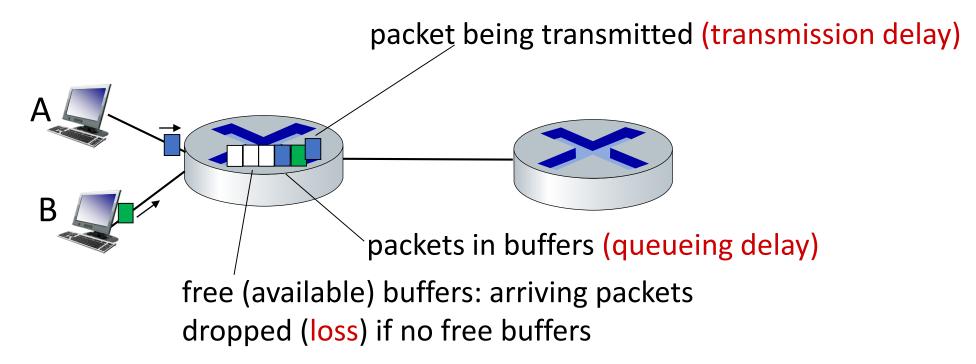
- What is the Internet?
- What is a protocol?
- Network edge: hosts, access network, physical media
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- Performance: loss, delay, throughput
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- Protocol layers, service models
- History



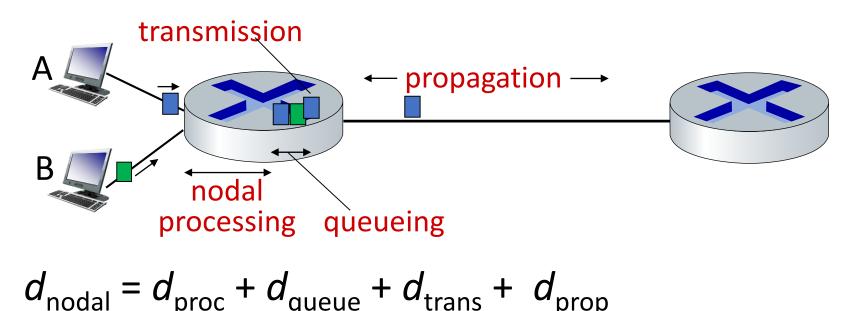
How do packet loss and delay occur?

packets queue in router buffers

- packets queue, wait for turn
- arrival rate to link (temporarily) exceeds output link capacity: packet loss



Packet delay: four sources



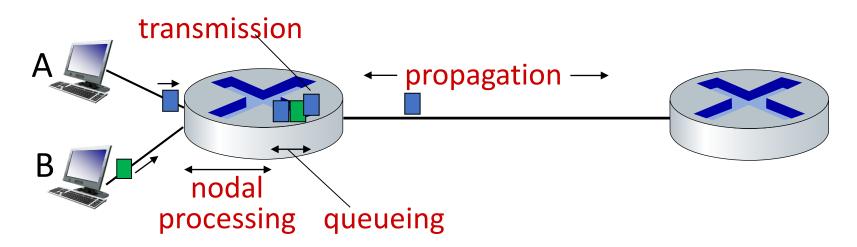
$$d_{\text{proc}}$$
: nodal processing

- check bit errors
- determine output link
- typically < msec</p>

d_{queue}: queueing delay

- time waiting at output link for transmission
- depends on congestion level of router

Packet delay: four sources



$$d_{\text{nodal}} = d_{\text{proc}} + d_{\text{queue}} + d_{\text{trans}} + d_{\text{prop}}$$

d_{trans} : transmission delay:

- L: packet length (bits)
- R: link transmission rate (bps)

$$d_{trans} = L/R$$

$$d_{trans} \text{ and } d_{prop}$$

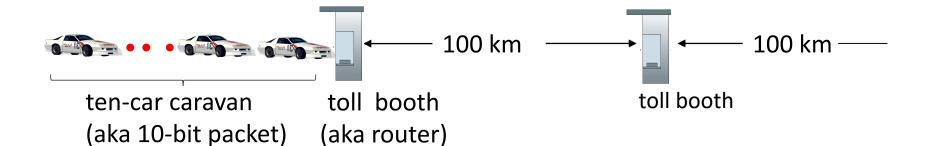
$$very \text{ different}$$

d_{prop} : propagation delay:

- d: length of physical link
- s: propagation speed (~2x10⁸ m/sec)

^{*} Check out the online interactive exercises: http://gaia.cs.umass.edu/kurose_ross

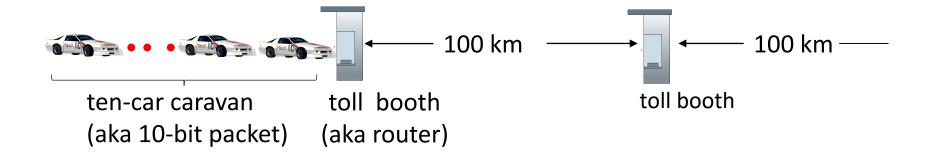
Caravan analogy



- cars "propagate" at 100 km/hr
- toll booth takes 12 sec to service car (bit transmission time)
- car ~ bit; caravan ~ packet
- Q: How long until caravan is lined up before 2nd toll booth?

- time to "push" entire caravan through toll booth onto highway = 12*10 = 120 sec
- time for last car to propagate from 1st to 2nd toll both: 100km/(100km/hr) = 1 hr
- *A:* 62 minutes

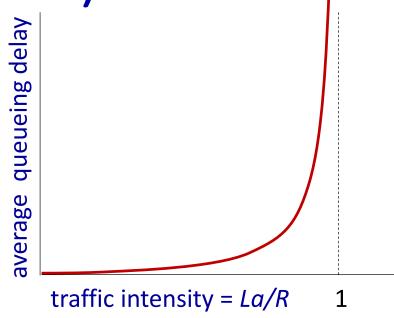
Caravan analogy



- suppose cars now "propagate" at 1000 km/hr
- and suppose toll booth now takes one minute to service a car
- Q: Will cars arrive to 2nd booth before all cars serviced at first booth?
 A: Yes! after 7 min, first car arrives at second booth; three cars still at first booth

Packet queueing delay (revisited)

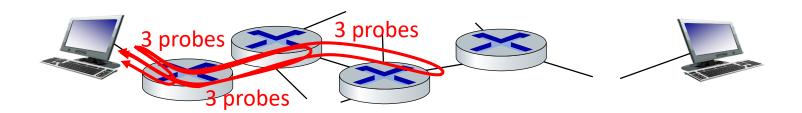
- R: link bandwidth (bps)
- L: packet length (bits)
- a: average packet arrival rate
- La/R ~ 0: avg. queueing delay small
- La/R -> 1: avg. queueing delay large
- La/R > 1: more "work" arriving is more than can be serviced - average delay infinite!





"Real" Internet delays and routes

- what do "real" Internet delay & loss look like?
- traceroute program: provides delay measurement from source to router along end-end Internet path towards destination. For all i:
 - sends three packets that will reach router i on path towards destination (with time-to-live field value of i)
 - router *i* will return packets to sender
 - sender measures time interval between transmission and reply



Real Internet delays and routes

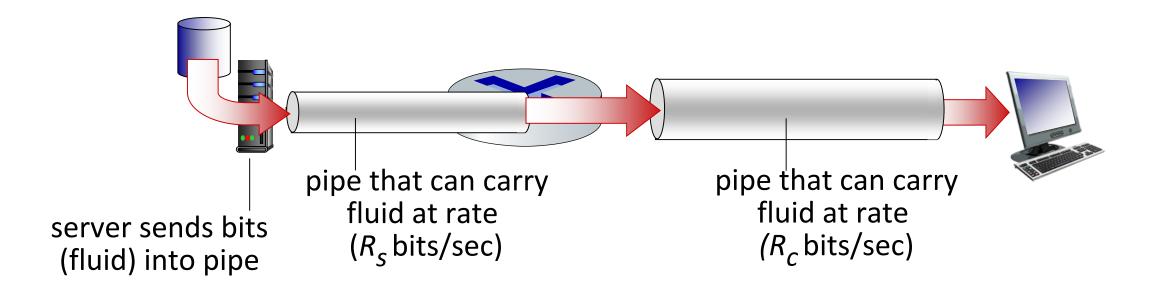
traceroute: gaia.cs.umass.edu to www.eurecom.fr

```
3 delay measurements from
                                       gaia.cs.umass.edu to cs-gw.cs.umass.edu
1 cs-gw (128.119.240.254) 1 ms 1 ms 2 ms
2 border1-rt-fa5-1-0.gw.umass.edu (128.119.3.145) 1 ms 1 ms 2 ms 4 delay measurements
                                                                    to border1-rt-fa5-1-0.gw.umass.edu
3 cht-vbns.gw.umass.edu (128.119.3.130) 6 ms 5 ms 5 ms
4 jn1-at1-0-0-19.wor.vbns.net (204.147.132.129) 16 ms 11 ms 13 ms
5 jn1-so7-0-0.wae.vbns.net (204.147.136.136) 21 ms 18 ms 18 ms
6 abilene-vbns.abilene.ucaid.edu (198.32.11.9) 22 ms 18 ms 22 ms
7 nycm-wash.abilene.ucaid.edu (198.32.8.46) 22 ms 22 ms trans-oceanic link
8 62.40.103.253 (62.40.103.253) 104 ms 109 ms 106 ms
9 de2-1.de1.de.geant.net (62.40.96.129) 109 ms 102 ms 104 ms
10 de.fr1.fr.geant.net (62.40.96.50) 113 ms 121 ms 114 ms
                                                                          looks like delays
11 renater-gw.fr1.fr.geant.net (62.40.103.54) 112 ms 114 ms 112 ms
                                                                          decrease! Why?
12 nio-n2.cssi.renater.fr (193.51.206.13) 111 ms 114 ms 116 ms
13 nice.cssi.renater.fr (195.220.98.102) 123 ms 125 ms 124 ms 14 r3t2-nice.cssi.renater.fr (195.220.98.110) 126 ms 126 ms 124 ms
15 eurecom-valbonne.r3t2.ft.net (193.48.50.54) 135 ms 128 ms 133 ms
16 194.214.211.25 (194.214.211.25) 126 ms 128 ms 126 ms
                   * means no response (probe lost, router not replying)
19 fantasia.eurecom.fr (193.55.113.142) 132 ms 128 ms 136 ms
```

^{*} Do some traceroutes from exotic countries at www.traceroute.org

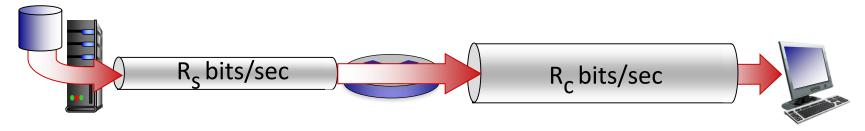
Throughput

- throughput: rate (bits/time unit) at which bits are being sent from sender to receiver
 - instantaneous: rate at given point in time
 - average: rate over longer period of time

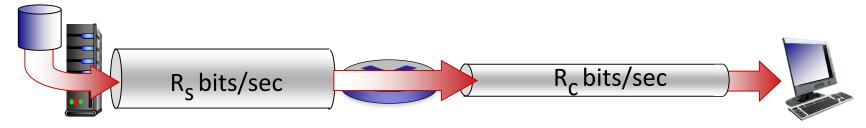


Throughput

 $R_s < R_c$ What is average end-end throughput?



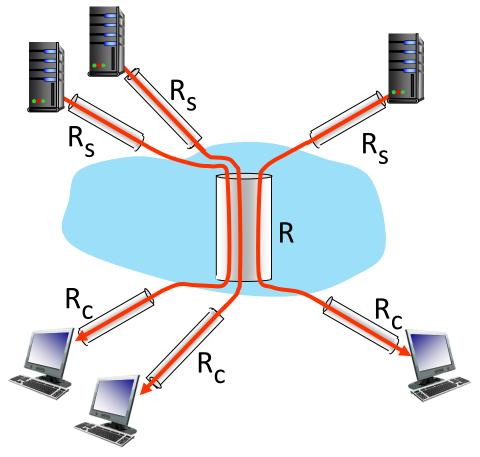
 $R_s > R_c$ What is average end-end throughput?



bottleneck link

link on end-end path that constrains end-end throughput

Throughput: network scenario



10 connections (fairly) share backbone bottleneck link *R* bits/sec

- per-connection endend throughput: min(R_c, R_s, R/10)
- in practice: R_c or R_s is often bottleneck

^{*} Check out the online interactive exercises for more examples: http://gaia.cs.umass.edu/kurose_ross/

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Network security

- field of network security:
 - how bad guys can attack computer networks
 - how we can defend networks against attacks
 - how to design architectures that are immune to attacks
- Internet not originally designed with (much) security in mind
 - original vision: "a group of mutually trusting users attached to a transparent network" ©
 - Internet protocol designers playing "catch-up"
 - security considerations in all layers!

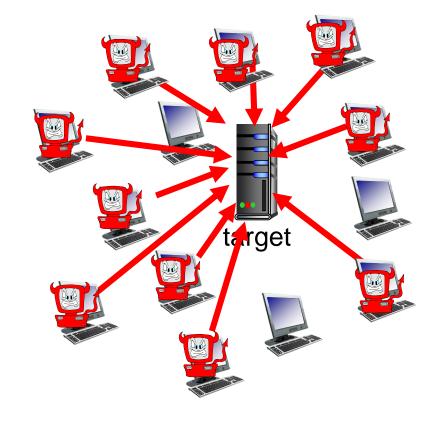
Bad guys: malware

- malware can get in host from:
 - *virus:* self-replicating infection by receiving/executing object (e.g., e-mail attachment)
 - worm: self-replicating infection by passively receiving object that gets itself executed
- spyware malware can record keystrokes, web sites visited, upload info to collection site
- infected host can be enrolled in botnet, used for spam or distributed denial of service (DDoS) attacks

Bad guys: denial of service

Denial of Service (DoS): attackers make resources (server, bandwidth) unavailable to legitimate traffic by overwhelming resource with bogus traffic

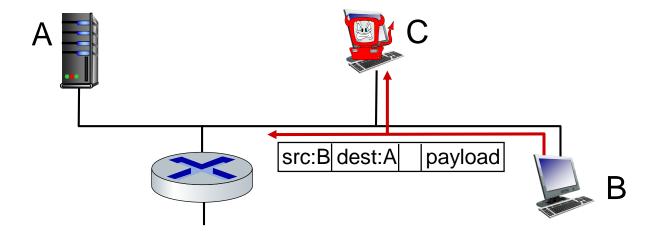
- 1. select target
- 2. break into hosts around the network (see botnet)
- 3. send packets to target from compromised hosts



Bad guys: packet interception

packet "sniffing":

- broadcast media (shared Ethernet, wireless)
- promiscuous network interface reads/records all packets (e.g., including passwords!) passing by

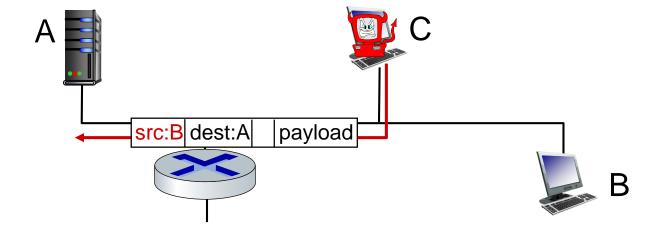




Wireshark software used for our end-of-chapter labs is a (free) packet-sniffer

Bad guys: fake identity

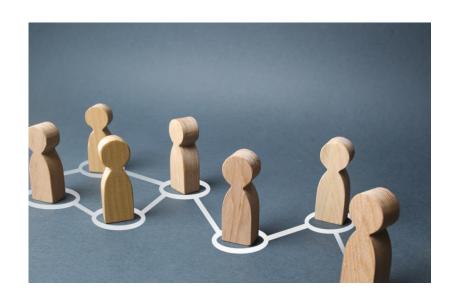
IP spoofing: send packet with false source address



... lots more on security (throughout, Chapter 8)

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Protocol "layers" and reference models

Networks are complex, with many "pieces":

- hosts
- routers
- links of various media
- applications
- protocols
- hardware, software

Question:

is there any hope of organizing structure of network?

.... or at least our discussion of networks?

Example: organization of air travel

ticket (purchase)

baggage (check)

gates (load)

runway takeoff

airplane routing

ticket (complain)

baggage (claim)

gates (unload)

runway landing

airplane routing

airplane routing

airline travel: a series of steps, involving many services

Example: organization of air travel

ticket (purchase)	ticketing service	ticket (complain)	
baggage (check)	baggage service	baggage (claim)	
gates (load)	gate service	gates (unload)	
runway takeoff	runway service	runway landing	
airplane routing	routing service	airplane routing	

layers: each layer implements a service

- via its own internal-layer actions
- relying on services provided by layer below

Q: describe in words the service provided in each layer above

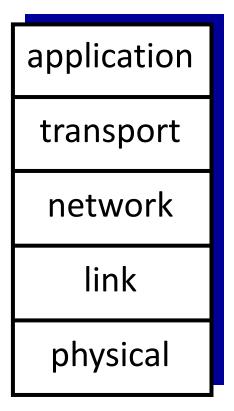
Why layering?

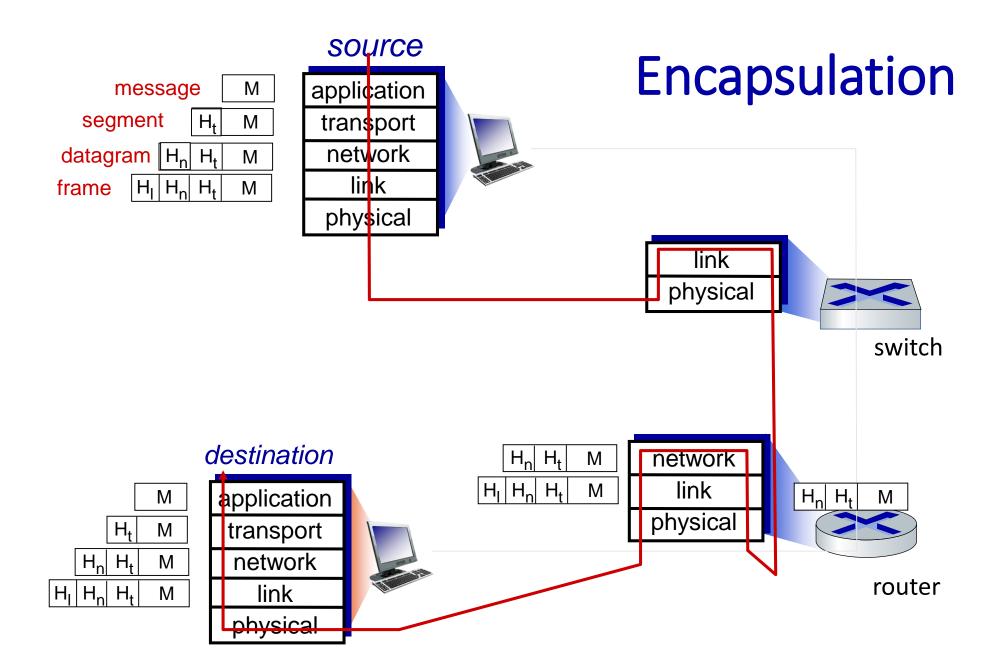
dealing with complex systems:

- explicit structure allows identification, relationship of complex system's pieces
 - layered *reference model* for discussion
- modularization eases maintenance, updating of system
 - change in layer's service implementation: transparent to rest of system
 - e.g., change in gate procedure doesn't affect rest of system
- layering considered harmful?
- layering in other complex systems?

Internet protocol stack

- application: supporting network applications
 - IMAP, SMTP, HTTP
- transport: process-process data transfer
 - TCP, UDP
- network: routing of datagrams from source to destination
 - IP, routing protocols
- link: data transfer between neighboring network elements
 - Ethernet, 802.11 (WiFi), PPP
- physical: bits "on the wire"





Chapter 1: roadmap

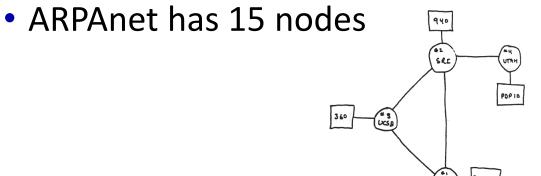
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1961-1972: Early packet-switching principles

- 1961: Kleinrock queueing theory shows effectiveness of packet-switching
- 1964: Baran packet-switching in military nets
- 1967: ARPAnet conceived by Advanced Research Projects Agency
- 1969: first ARPAnet node operational

- **1972**:
 - ARPAnet public demo
 - NCP (Network Control Protocol) first host-host protocol
 - first e-mail program



THE ARPA NETWORK

1972-1980: Internetworking, new and proprietary nets

- 1970: ALOHAnet satellite network in Hawaii
- 1974: Cerf and Kahn architecture for interconnecting networks
- 1976: Ethernet at Xerox PARC
- late70's: proprietary architectures: DECnet, SNA, XNA
- late 70's: switching fixed length packets (ATM precursor)
- 1979: ARPAnet has 200 nodes

Cerf and Kahn's internetworking principles:

- minimalism, autonomy no internal changes required to interconnect networks
- best-effort service model
- stateless routing
- decentralized control

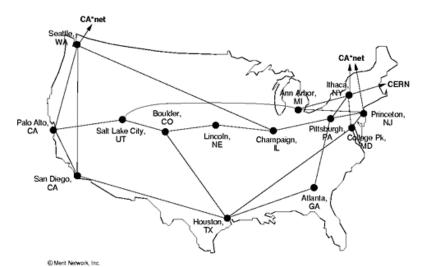
define today's Internet architecture

1980-1990: new protocols, a proliferation of networks

- 1983: deployment of TCP/IP
- 1982: smtp e-mail protocol defined
- 1983: DNS defined for nameto-IP-address translation
- 1985: ftp protocol defined
- 1988: TCP congestion control

- new national networks: CSnet, BITnet, NSFnet, Minitel
- 100,000 hosts connected to confederation of networks

NSFNET T1 Network 1991



Introduction: 1-29

1990, 2000s: commercialization, the Web, new applications

- early 1990s: ARPAnet decommissioned
- 1991: NSF lifts restrictions on commercial use of NSFnet (decommissioned, 1995)
- early 1990s: Web
 - hypertext [Bush 1945, Nelson 1960's]
 - HTML, HTTP: Berners-Lee
 - 1994: Mosaic, later Netscape
 - late 1990s: commercialization of the Web

late 1990s – 2000s:

- more killer apps: instant messaging, P2P file sharing
- network security to forefront
- est. 50 million host, 100 million+ users
- backbone links running at Gbps

2005-present: more new applications, Internet is "everywhere"

- ~18B devices attached to Internet (2017)
 - rise of smartphones (iPhone: 2007)
- aggressive deployment of broadband access
- increasing ubiquity of high-speed wireless access: 4G/5G, WiFi
- emergence of online social networks:
- Facebook: ~ 2.5 billion users
- service providers (Google, FB, Microsoft) create their own networks
 - bypass commercial Internet to connect "close" to end user, providing "instantaneous" access to search, video content, ...
- enterprises run their services in "cloud" (e.g., Amazon Web Services, Microsoft Azure)

Chapter 1: summary

We've covered a "ton" of material!

- Internet overview
- what's a protocol?
- network edge, access network, core
 - packet-switching versus circuitswitching
 - Internet structure
- performance: loss, delay, throughput
- layering, service models
- security
- history

You now have:

- context, overview, vocabulary, "feel" of networking
- more depth, detail, and fun to follow!

Additional Chapter 1 slides

Wireshark

