Summary | Real Analysis

Introduction— |

 $| \land | \text{ and } | | \lor | \text{ or } | | \to | \text{ then } | \implies | \text{ implies } | \Leftarrow | \text{ implied by } | | \iff | \text{ if and only if } | \forall | \text{ for all } | \exists | \text{ there exists } | \sim | \text{ not } |$

Let's take $a \rightarrow b$.

- 1. Contrapositive or transposition: $\sim b
 ightarrow \sim a$. This is equivalent to the original.
- 2. Inverse: $\sim a
 ightarrow \sim b$. Does not depend on the original.
- 3. Converse: b
 ightharpoonup a . Does not depend on the original.

$$a \rightarrow b \equiv \sim a \lor b \equiv \sim b \rightarrow \sim a$$

Required proofs

- $oldsymbol{\cdot} \sim orall x \, P(x) \equiv \exists x \sim P(x)$
- $\sim \exists x \, P(x) \equiv \forall x \sim P(x)$
- $\exists x \, \exists y P(x,y) \equiv \exists y \, \exists x P(x,y)$
- $\forall x \, \forall y P(x,y) \equiv \forall y \, \forall x P(x,y)$
- $\cdot \exists x \forall y P(x,y) \implies \forall y \exists x P(x,y)$
- $(A \rightarrow C) \land (B \rightarrow C) \equiv (A \lor B) \rightarrow C$

Methods of proofs

- 1. Just proof what should be proven
- 2. Prove the contrapositive
- 3. Proof by contradiction
- 4. Proof by induction

Proof by contradiction

Suppose $a \implies b$ has to be proven. If $a \land \sim b$ is proven to be false, then, by proof by contradiction, $a \implies b$ can be trivially proven.

Logic behind proof by contradiction

$$egin{aligned} a \wedge \sim b &= F \ &\sim (a \wedge \sim b) = \sim F \ &\sim a ee b = T \ &a &\Longrightarrow b \end{aligned}$$

Set theory

Zermelo-Fraenkel set theory with axiom of Choice(ZFC):9 axioms all together is being used here.

Definitions

- $x \in A^{c} \iff x \notin A$
- $x \in A \cup B \iff x \in A \lor x \in B$
- $x \in A \cap B \iff x \in A \land x \in B$
- $A \subset B = \forall x (x \in A \implies x \in B)$
- $A-B=A\cap B^{\mathrm{c}}$
- $\bullet \ \ A = B \iff ((\forall z \in A \implies z \in B) \land (\forall z \in B \implies z \in A))$

Required proofs

- $(A \cap B)^c = A^c \cup B^c$
- $(A \cup B)^c = A^c \cap B^c$
- $A \cap (B \cup C) = (A \cap B) \cup (A \cap C)$
- $A \cup (B \cap C) = (A \cup B) \cap (A \cup C)$
- $A \subset A \cup B$
- $A \cap B \subset A$

Set of Numbers

Sets of numbers

- Positive integers: $\mathbb{Z}^+ = \{1,2,3,4,\dots\}$.

• Natural integers: $\mathbb{N} = \{0,1,2,3,4,\dots\}$.

- Negative integers: $\mathbb{Z}^- = \{-1, -2, -3, -4, \dots\}$.

• Integers: $\mathbb{Z}=\mathbb{Z}^-\cup\{0\}\cup\mathbb{Z}^+$.

• Rational numbers: $\mathbb{Q}=\left\{rac{p}{q}\Big|q
eq0\land p,q\in\mathbb{Z}
ight\}$.

• Irrational numbers: limits of sequences of rational numbers (which are not rational numbers)

• Real numbers: $\mathbb{R}=\mathbb{Q}^c\cup\mathbb{Q}$.

Complex numbers are not part of the study here.

Continued Fraction Expansion

The process

- · Separate the integer part
- Find the inverse of the remaining part. Result will be greated than 1.
- Repeat the process for the remaining part.

Finite expansion

Take $\frac{420}{69}$ for example.

$$\frac{420}{69} = 6 + \frac{6}{69}$$

$$\frac{420}{69} = 6 + \frac{1}{\frac{69}{6}}$$

$$\frac{420}{69} = 6 + \frac{1}{11 + \frac{3}{6}}$$

$$\frac{420}{69} = 6 + \frac{1}{11 + \frac{1}{2}}$$

As $\frac{420}{69}$ is finite, its continued fraction expansion is also finite. And it can be written as $\frac{420}{69}=[6;11,2]$.

Infinite expansion

For irrational numbers, the expansion will be infinite.

For example π :

$$\pi = 3 + \cfrac{1}{7 + \cfrac{1}{15 + \cfrac{1}{1 + \cfrac{1}{292 + \cdots}}}}$$

Conintued fraction expansion of π is $[3;7,15,1,292,1,1,1,2,1,3,1,14,2,1,1,2,\ldots]$.

Field Axioms

Field Axioms of $\mathbb R$

 $\mathbb{R}
eq \emptyset$ with two binary operations + and \cdot satisfying the following properties

- 1. Closed under addition: $\forall a,b \in \mathbb{R}; a+b \in \mathbb{R}$
- 2. Commutative: $orall a,b\in\mathbb{R}; a+b=b+a$
- 3. Associative: $orall a,b,c\in\mathbb{R}; (a+b)+c=a+(b+c)$
- 4. Additive identity: $\exists 0 \in \mathbb{R} \, orall a \in \mathbb{R}; a+0=0+a=a$
- 5. Additive inverse: $orall a \in \mathbb{R} \, \exists (-a); a+(-a)=(-a)+a=0$
- 6. Closed under multiplication: $orall a,b\in\mathbb{R};a\cdot b\in\mathbb{R}$
- 7. Commutative: $\forall a,b \in \mathbb{R}; a \cdot b = b \cdot a$
- 8. Associative: $orall a,b,c\in\mathbb{R}; (a\cdot b)\cdot c=a\cdot (b\cdot c)$
- 9. Multiplicative identity: $\exists 1 \in \mathbb{R} \, orall a \in \mathbb{R}; a \cdot 1 = 1 \cdot a = a$
- 10. Multiplicative inverse: $orall a \in \mathbb{R} \{0\}\,\exists a^-; a\cdot a^- = a^-\cdot a = 1$

11. Multiplication is distributive over addition: $a \cdot (b+c) = a \cdot b + a \cdot c$

(i) Field

Any set satisfying the above axioms with two binary operations (commonly + and \cdot) is called a **field**. Written as $(\mathbb{R}, +, \cdot)$ is a **Field**. But $(\mathbb{R}, \cdot, +)$ is not a **field**.

Required proofs

The below mentioned propositions can and should be proven using the above-mentioned axioms. $a,b,c\in\mathbb{R}$.

•
$$a \cdot 0 = 0$$

Hint: Start with $a(1+0)$

• Additive identity (
$$\mathbf{0}$$
) is unique

• Multiplicative identity (
$$oldsymbol{1}$$
) is unique

- Additive inverse (
$$-a$$
) is unique for a given $\,a$

- Multiplicative inverse (
$$a^{-1}$$
) is unique for a given $\,a\,$

•
$$a+b=0 \implies b=-a$$

$$\cdot \ a+c=b+c \implies a=b$$

•
$$-(a+b) = (-a) + (-b)$$

•
$$-(-a)=a$$

•
$$ac = bc \implies a = b$$

•
$$ab = 0 \implies a = 0 \lor b = 0$$

$$\cdot \ -(ab) = (-a)b = a(-b)$$

•
$$(-a)(-b) = ab$$

•
$$a \neq 0 \implies (a^{-1})^{-1} = a$$

$$\bullet \ \ a,b\neq 0 \implies ab^{-1}=a^{-1}b^{-1}$$

Field or Not?

	Is field?	Reason (if not)
$(\mathbb{R},+,\cdot)$	True	
$(\mathbb{R},\cdot,+)$	False	Axiom 11 is invalid

	Is field?	Reason (if not)
$(\mathbb{Z},+,\cdot)$	False	Multiplicative inverse doesn't exist
$(\mathbb{Q},+,\cdot)$	True	
$(\mathbb{Q}^c,+,\cdot)$	False	$\sqrt{2}\cdot\sqrt{2} otin\mathbb{Q}^c$
Boolean algebra	False	Additive inverse doesn't exist
$(\{0,1\}, + \bmod 2, \cdot \bmod 2)$	True	
$(\{0,1,2\}, + \bmod 3, \cdot \bmod 3)$	True	
$(\{0,1,2,3\}, + \bmod 4, \cdot \bmod 4)$	False	Multiplicative inverse doesn't exist

Completeness Axiom

Let A be a non empty subset of \mathbb{R} .

- u is the upper bound of A if: $\forall a \in A; a \leq u$
- $oldsymbol{A}$ is bounded above if $oldsymbol{A}$ has an upper bound
- Maximum element of A : $\max A = u$ if $u \in A$ and u is an upper bound of A
- Supremum of $A \, \sup A$, is the smallest upper bound of A
- Maximum is a supremum. Supremum is not necessarily a maximum.
- l is the lower bound of A if: $orall a \in \mathit{A}; a \geq \mathit{l}$
- $oldsymbol{A}$ is bounded below if $oldsymbol{A}$ has a lower bound
- Minimum element of A : $\min A = l$ if $l \in A$ and l is a lower bound of A
- Infimum of A $\inf A$, is the largest lower bound of A
- Minimum is a infimum. Infimum is not necessarily a minimum.

Theorems

Let A be a non empty subset of $\mathbb R$.

- Say u is an upper bound of A . Then $u=\sup A$ iff: $orall \epsilon>0$ $\exists a\in A;\ a+\epsilon>u$
- Say l is a lower bound of A . Then $l=\inf A$ iff: $orall \epsilon>0$ $\exists a\in A;\ a-\epsilon< l$

Required proofs

- sup(a,b) = b
- inf(a,b)=a

Completeness axioms of real numbers

- Every non empty subset of ${\mathbb R}$ which is bounded above has a supremum in ${\mathbb R}$
- Every non empty subset of ${\mathbb R}$ which is bounded below has a infimum in ${\mathbb R}$
- (i) Note
- ${\mathbb Q}$ doesn't have the completeness property.

Completeness axioms of integers

- Every non empty subset of ${\mathbb Z}$ which is bounded above has a maximum
- Every non empty subset of ${\mathbb Z}$ which is bounded below has a minimum

Two important theorems

- $\exists a \ \forall \epsilon > 0, a < \epsilon \implies a \leq 0$
- $\forall \epsilon > 0 \; \exists a, a < \epsilon \implies a \leq 0$

Order Axioms

- Trichotomy: $orall a, b \in \mathbb{R}$ exactly one of these holds: a > b , a = b , a < b
- Transitivity: $orall a, b, c \in \mathbb{R}; a < b \wedge b < c \implies a < c$
- Operation with addition: $\forall a,b \in \mathbb{R}; a < b \implies a+c < b+c$
- Operation with mutliplication: $orall a, b, c \in \mathbb{R}; a < b \land 0 < c \implies ac < bc$

Definitions

•
$$a < b \equiv b > a$$

•
$$a \le b \equiv a < b \lor a = b$$

$$\cdot \ a \neq b \equiv a < b \lor a > b$$

$$oldsymbol{\cdot} \; |x| = egin{cases} x & ext{if } x \geq 0, \ -x & ext{if } x < 0 \end{cases}$$

Triangular inequalities

$$|a| - |b| \le |a + b| \le |a| + |b|$$

$$\Big||a|-|b|\Big|\leq |a+b|$$

(i) Proof Hint

For first:

• Use
$$-|a| \leq a \leq |a|$$

For second:

• Use the below substitutions in first conclusion

$$\circ$$
 $a = a - b \land b = b$

$$a = b - a \land b = a$$

Required proofs

$$m{\cdot} \ \ orall a,b,c \in \mathbb{R}; a < b \wedge c < 0 \implies ac > bc$$

•
$$1 > 0$$

•
$$-|a| \le a \le |a|$$

• Triangular inequalities

Theorems

- $\exists a \ \forall \epsilon > 0, \ a < \epsilon \implies a \leq 0$
- $\exists a \ \forall \epsilon > 0, \ 0 \leq a < \epsilon \implies a = 0$

(!) Caution

$$\forall \epsilon > 0 \; \exists a, \, a < \epsilon \implies a \leq 0 \; \text{is not valid.}$$

Let A be a non-empty subset of $\mathbb R$ which is bounded above and has an upper bound u.

$$u = \sup A \iff orall \epsilon > 0 \, \exists a \in A, \, a > u - \epsilon$$

Let A be a non-empty subset of $\mathbb R$ which is bounded below and has an lower bound m.

$$m = \inf A \iff orall \epsilon > 0 \, \exists a \in A, \, a < m + \epsilon$$

Relations

Definitions

- Cartesian Product of sets A,B $A imes B = \{(a,b) | a \in A, b \in B\}$
- Ordered pair $(a,b)=\{\{a\},\{a,b\}\}$

Relation

Let $A,B
eq \emptyset$. A relation R:A o B is a non-empty subset of A imes B.

- $aRb \equiv (a,b) \in R$
- Domain of $R \colon dom(R) = A$
- Codomain of R : codom(R) = B
- Range of R : $ran(R) = \{y | (x,y) \in R\}$
- $ran(R) \subseteq B$
- Pre-range of R : $preran(R) = \{x \, | \, (x,y) \in R\}$
- $preran(R) \subseteq A$
- $R(a) = \{b \mid (a,b) \in R\}$

Everywhere defined

 $oldsymbol{R}$ is everywhere defined

$$\iff A = dom(R) = preran(R)$$

$$\iff \forall a \in A, \ \exists b \in B; \ (a,b) \in R.$$

Onto

 $oldsymbol{R}$ is onto

$$\iff B = codom(R) = ran(R)$$

$$\iff \forall b \in B \,\exists a \in A \, (a,b) \in R$$

Aka. surjection.

Inverse

Inverse of R: $R^{-1} = \{(b,a) \mid (a,b) \in R\}$

Types of relation

one-many

$$\iff \exists a \in A, \ \exists b_1, b_2 \in B \ ((a,b_1),(a,b_2) \in R \ \land \ b_1 \neq b_2)$$

Not one-many

$$\iff orall a \in A, \, orall b_1, b_2 \in B \; ((a,b_1),(a,b_2) \in R \implies b_1 = b_2)$$

many-one

$$\iff \exists a_1,a_2 \in A, \, \exists b \in B \ ((a_1,b),(a_2,b) \in R \, \wedge \, a_1
eq a_2)$$

Not many-one

$$\iff orall a_1, a_2 \in A, \, orall b \in B \; ((a_1,b),(a_2,b) \in R \implies a_1 = a_2)$$

many-many

iff R is **one-many** and **many-one**.

one-one

iff R is not one-many and not many-one. Aka. injection.

Bijection

When a relation is **onto** and **one-one**.

Functions

A function $f:A \to B$ is a relation $f:A \to B$ which is <u>everywhere defined</u> and <u>not one-many</u>.

• dom(f) = A = preran(f)

Inverse

For a function $f\colon A o B$ to have its inverse relation $f^{-1}\colon B o A$ be also a function, we need:

- f is onto
- $m{f}$ is <u>not many-one</u> (in other words, $m{f}$ must be <u>one-one</u>)

The above statement is true for all unrestricted function $m{f}$ that has an inverse $m{f}^{-1}$:

$$f(f^{-1}(x)) = x = f^{-1}(f(x)) = x$$

Composition

Composition of relations

Let $R:A \to B$ and $S:B \to C$ are 2 relations. Composition can be defined when $\operatorname{ran}(R) = \operatorname{preran}(S)$.

Say ran(R) = preran(S) = D. Composition of the 2 relations is written as:

$$S \circ R = \{(a,c) \, | \, (a,b) \in R, \, (b,c) \in S, \, b \in D\}$$

Composition of functions

Let f:A o B and g:B o C be 2 functions where f is onto.

$$g \circ f = \{(x,z) \, | \, (x,y) \in f, \, (y,z) \in g, \, y \in B\} = g(f(x))$$

Countability

A set A is countable **iff** $\exists f: A \to Z^+$, where f is a one-one function.

Examples

- Countable: Any finite set, \mathbb{Z}, \mathbb{Q}
- Uncountable: $\mathbb R$, Any open/closed intervals in $\mathbb R$.

Transitive property

Say $B\subset A$.

 $A ext{ is countable } \implies B ext{ is countable }$

B is not countable $\implies A$ is not countable

Limits

$$\lim_{x o a}f(x)=L$$
 iff:

$$orall \epsilon > 0 \; \exists \delta > 0 \; orall x \; (0 < |x-a| < \delta \implies |f(x)-L| < \epsilon)$$

Defining δ in terms of a given ϵ is enough to prove a limit.

One sided limits

 $\lim_{x o a^+}f(x)=L$ iff:

$$orall \epsilon > 0 \; \exists \delta > 0 \; orall x \; (0 < x - a < \delta \implies |f(x) - L| < \epsilon)$$

 $\lim_{x o a^-}f(x)=L$ iff:

$$orall \epsilon > 0 \; \exists \delta > 0 \; orall x \; (-\delta < x - a < 0 \implies |f(x) - L| < \epsilon)$$

 $\lim_{x o a}f(x)=L^+$ iff:

$$orall \epsilon > 0 \; \exists \delta > 0 \; orall x \; (0 < |x - a| < \delta \implies 0 \le f(x) - L < \epsilon)$$

 $\lim_{x o a}f(x)=L^-$ iff:

$$orall \epsilon > 0 \; \exists \delta > 0 \; orall x \; (0 < |x-a| < \delta \implies -\epsilon < f(x) - L \le 0)$$

Limits including infinite

 $\lim_{x o\infty}f(x)=L$ iff:

$$orall \epsilon > 0 \; \exists N > 0 \; orall x \; (x > N \implies |f(x) - L| < \epsilon)$$

 $\lim_{x o -\infty}f(x)=L$ iff:

$$orall \epsilon > 0 \; \exists N > 0 \; orall x \; (x < -N \implies |f(x) - L| < \epsilon)$$

$$\lim_{x o a}f(x)=\infty$$
 iff:

$$orall M>0 \; \exists \delta>0 \; orall x \; (0<|x-a|<\delta \implies f(x)>M)$$

$$\lim_{x o a}f(x)=-\infty$$
 iff:

$$orall M>0 \; \exists \delta>0 \; orall x \; (0<|x-a|<\delta \implies f(x)<-M)$$

Indeterminate forms

- $\begin{array}{cc} \bullet & \frac{0}{0} \\ \bullet & \frac{\infty}{\infty} \end{array}$
- $\cdot \infty \cdot 0$
- $\infty \infty$
- $\cdot \infty^0$
- $\cdot 0^{0}$
- 1∞

Continuity

A function f is continuous at a iff:

$$\lim_{x o a}f(x)=f(a)$$

$$orall \epsilon > 0 \; \exists \delta > 0 \; orall x \; (|x-a| < \delta \implies |f(x) - f(a)| < \epsilon)$$

One-side continuous

A function f is continuous from right at a iff:

$$\lim_{x o a^+}f(x)=f(a)$$

A function $m{f}$ is continuous from left at $m{a}$ iff:

$$\lim_{x o a^-}f(x)=f(a)$$

Continuous on an open interval

A function f is continuous in (a,b) **iff** f is continuous on every $c \in (a,b)$.

Continuous on a closed interval

A function f is continuous in [a,b] iff f is:

- continuous on every $c \in (a,b)$
- right-continuous at $oldsymbol{a}$
- left-continuous at $oldsymbol{b}$

Uniformly continuous

Suppose a function f is continuous on (a,b). f is uniformly continuous on (a,b) iff:

$$orall \epsilon > 0 \; \exists \delta > 0 \; ext{s.t.} \; |x-y| < \delta \implies |f(x)-f(y)| < \epsilon$$

If a function f is continuous on [a, b], f is uniformly continuous on [a, b].

∴ Todo

Is this section correct? I am not 100% sure.

Continuity Theorems

Extreme Value Theorem

If f is continuous on [a,b], f has a maximum and a minimum in [a,b].

(i) Proof Hint

Proof is quite hard.

Intermediate Value Theorem

Let f is continuous on [a,b]. If $\exists u$ such that f(a)>u>f(b) or f(a)< u< f(b): $\exists c\in (a,b)$ such that f(c)=u.

(i) Proof Hint

Proof the case when u=0. Otherwise define a new function g(x) such that middle part of the above inequality has a 0 in the place of u.

Sandwich (or Squeeze) Theorem

Let:

- For some $\delta > 0$: $orall x(0 < |x-a| < \delta \implies f(x) \le g(x) \le h(x))$
- $oldsymbol{\cdot} \lim_{x o a}f(x)=\lim_{x o a}h(x)=L\in\mathbb{R}$

Then $\lim_{x o a} g(x) = L$.

(i) Note

Works for any kind of x limits.

"No sudden changes"

Positive

Let f be continuous on a and f(a)>0

$$\implies \exists \delta > 0; \forall x (|x-a| < \delta \implies f(x) > 0)$$

(i) Proof Hint

Take
$$\epsilon = rac{f(a)}{2}$$

Negative

Let f be continuous on a and f(a) < 0

$$\implies \exists \delta > 0; \forall x (|x - a| < \delta \implies f(x) < 0)$$

(i) Proof Hint

Take
$$\epsilon = -rac{f(a)}{2}$$

Differentiability

A function f is differentiable at a iff:

$$\lim_{x o a}rac{f(x)-f(a)}{x-a}=L\in\mathbb{R}=f'(a)$$

f'(a) is called the derivative of f at a.

One-side differentiable

Left differentiable

A function f is left-differentiable at a iff:

$$\lim_{x o a^-}rac{f(x)-f(a)}{x-a}=L\in\mathbb{R}=f'_-(a)$$

Right differentiable

A function $m{f}$ is right-differentiable at $m{a}$ iff:

$$\lim_{x o a^+}rac{f(x)-f(a)}{x-a}=L\in\mathbb{R}=f'_+(a)$$

Differentiability implies continuity

f is differentiable at $a \implies f$ is continuous at a

(i) Proof Hint

Use $\delta = min(\delta_1, rac{\epsilon}{1+|f'(a)|})$.

(i) Note

Suppose f is differentiable at a. Define g:

$$g(x) = \left\{ egin{array}{ll} rac{f(x) - f(a)}{x - a}, & x
eq a \ f'(a), & x = a \end{array}
ight.$$

g is continuous at a.

Extreme Values

Suppose $f:[a,b] o \mathbb{R}$, and $F=f([a,b])=\Big\{f(x)\mid x\in [a,b]\Big\}$. Minimum and maximum values of f are called the extreme values.

Maximum

Maximum of the function f is f(c) where $c \in [a,b]$ iff:

$$\forall x \in [a,b], \ f(c) \geq f(x)$$

aka. Global Maximum. Maximum doesn't exist always.

Local Maximum

A Local maximum of the function f is f(c) where $c \in [a,b]$ iff:

$$\exists \delta \ \ orall x \, (0 < |x-c| < \delta \implies f(c) \geq f(x))$$

Global maximum is obviously a local maximum.

The above statement can be simplified when c=a or c=b.

When c = a:

$$\exists \delta \ \ \forall x \, (0 < x - c < \delta \implies f(c) \geq f(x))$$

When c = b:

$$\exists \delta \ \ orall x \left(-\delta < x - c < 0 \ \Longrightarrow \ f(c) \geq f(x)
ight)$$

Minimum

Minimum of the function f is f(c) where $c \in [a,b]$ iff:

$$\forall x \in [a,b], \ f(c) \leq f(x)$$

aka. Global Minimum. Minimum doesn't exist always.

Local Minimum

$$\exists \delta \ \ orall x \, (0 < |x-c| < \delta \implies f(c) \leq f(x))$$

Global minimum is obviously a local maximum.

The above statement can be simplified when c=a or c=b.

When c = a:

$$\exists \delta \ \ \forall x \, (0 < x - c < \delta \implies f(c) \leq f(x))$$

When c = b:

$$\exists \delta \ \ orall x \left(-\delta < x - c < 0 \ \Longrightarrow \ f(c) \leq f(x)
ight)$$

Special cases

f is continuous

Then by Extreme Value Theorem, we know f has a minimum and maximum in [a,b].

f is differentiable

- If f(a) is a local maximum: $f'_+(a) \leq 0$
- If f(b) is a local maximum: $f_-'(b) \geq 0$
- $oldsymbol{\cdot}$ $c \in (a,b)$ and If f(c) is a local maximum: f'(c)=0
- If f(a) is a local minimum: $f'_+(a) \geq 0$
- If f(b) is a local minimum: $f'(b) \leq 0$
- $c \in (a,b)$ and If f(c) is a local minimum: f'(c)=0

Critical point

 $c \in [a,b]$ is called a critical point iff:

$$f'(c) = 0 \quad \lor \quad f'(c) \text{ is undefined}$$

Other Theorems

Rolle's Theorem

Let f be continuous on [a,b] and differentiable on (a,b). And f(a)=f(b). Then:

$$\exists c \in (a,b) \text{ s.t. } f'(c) = 0$$

(i) Proof Hint

By Extreme Value Theorem, maximum and minimum exists for f.

Consider 2 cases:

- 1. Both minimum and maximum exist at $\,a\,$ and $\,b\,$.
- 2. One of minimum or maximum occurs in (a,b) .

Mean Value Theorem

Let f be continuous on [a,b] and differentiable on (a,b). Then:

$$\exists c \in (a,b) ext{ s.t. } f'(c) = rac{f(b) - f(a)}{b - a}$$

(i) Proof Hint

- Define $g(x) = f(x) \left(rac{f(a) f(b)}{a b}
 ight) x$
- g(a) will be equal to g(b)
- ullet Use Rolle's Theorem for $oldsymbol{g}$

Cauchy's Mean Value Theorem

Let f and g be continuous on [a,b] and differentiable on (a,b), and $orall x \in (a,b) \ g'(x)
eq 0$ Then:

$$\exists c \in (a,b) ext{ s.t. } rac{f'(c)}{g'(c)} = rac{f(b) - f(a)}{g(b) - g(a)}$$

(i) Proof Hint

- Define $h(x) = f(x) \Big(rac{f(a) f(b)}{g(a) g(b)}\Big)g(x)$
- h(a) will be equal to h(b)
- ullet Use Rolle's Theorem for h

Mean value theorem can be obtained from this when g(x)=x.

Generalized MVT for Riemann Integrals

Let f,g be continuous on [a,b] ($\Longrightarrow f,g$ are integrable), and g does not change sign on (a,b). Then $\exists \zeta \in (a,b)$ such that:

$$\int_a^b f(x)g(x)\mathrm{d}x = f(\zeta)\int_a^b g(x)\mathrm{d}x$$

(i) Proof Hint

- Use Extreme value theorem for $m{f}$
- Multiply by g(x) . Then integrate. Then divide by $\int_a^b g(x)$.
- Use intermediate value theorem to find $f(\zeta)$

L'Hopital's Rule

(i) Note

Be careful with the pronunciation.

- It's not "Hospital's Rule", there are no "s"
- It's not "Hopital's Rule" either, there is a "L'"

L'Hopital's Rule can be used when all of these conditions are met. (here δ is some positive number). Select the appropriate x ranges.

- 1. Either of these conditions must be satisfied
 - $\circ \ f(a) = g(a) = 0$
 - $\circ \ \lim f(x) = \lim g(x) = 0$
 - $\circ \lim f(x) = \lim g(x) = \infty$
- 2. f,g are continuous on $x\in [a,a+\delta]$
- 3. f,g are differentiable on $x\in(a,a+\delta)$
- 4. $g'(x) \neq 0$ on $x \in (a, a + \delta)$
- 5. $\lim_{x o a^+} rac{f'(x)}{g'(x)} = L \in \mathbb{R}$

Then:
$$\lim_{x o a^+} rac{f(x)}{g(x)} = L$$

(i) Note

L'Hopital's rule can be proven using Cauchy's Mean Value Theorem.

It is valid for all types of "x limits".

Higher Order Derivatives

Suppose f is a function defined on (a,b). f is n times differentiable or n-th differentiable **iff**:

$$\lim_{x o a}rac{f^{(n-1)}(x)-f^{(n-1)}(a)}{x-a}=L\in\mathbb{R}=f^{(n)}(a)$$

Here $f^{(n)}$ denotes n-th derivative of f. And $f^{(0)}$ means the function itself. $f^{(n)}(a)$ is the n-th derivative of f at a.

(i) Note

f is n-th differentiable at $a \implies f^{(n-1)}$ is continuous at a

Taylor's Theorem

Let f is n+1 differentiable on (a,b). Let $c,x\in(a,b)$. Then $\exists \zeta \text{ s.t.}$:

$$f(x) = f(c) + \sum_{k=1}^n rac{f^{(k)}(c)}{k!} (x-c)^k + rac{f^{(n+1)}(\zeta)}{(n+1)!} (x-c)^{n+1}$$

Mean value theorem can be derived from taylor's theorem when n=0.

(i) Proof Hint

$$F(t) = f(t) + \sum_{k=1}^{n} rac{f^{(k)}(t)}{k!} (x-t)^{k}$$

$$G(t) = (x-t)^{n+1}$$

- Define F,G as mentioned above
- Consider the interval $\left[c,x
 ight]$
- Use <u>Cauchy's mean value theorem</u> for F,G after making sure the conditions are met.

The above equation can be written like:

$$f(x) = T_n(x,c) + R_n(x,c)$$

Taylor Polynomial

This part of the above equation is called the Taylor polynomial. Denoted by $T_n(x,c)$.

$$T_n(x,c) = f(c) + \sum_{k=1}^n rac{f^{(k)}(c)}{k!} (x-c)^k$$

Remainder

Denoted by $R_n(x,c)$.

$$R_n(x,c) = rac{f^{(n+1)}(\zeta)}{(n+1)!} (x-c)^{n+1}$$

Integral form of the remainder

$$R_n(x,c)=rac{1}{n!}\int_c^x f^{(n+1)}(t)(x-t)^n\mathrm{d}t$$

(i) Proof Hint

- Method 1: Use integration by parts and mathematical induction.
- Method 2: Use Generalized MVT for Riemann Integrals where:

$$\circ$$
 $F=f^{(n+1)}$

$$G = (x-t)^n$$

(i) Note

When n=1:

$$f(x) = f(c) + f'(c)(x-c) + rac{f''(\zeta)}{2!}(x-c)^2$$

$$f(x) - ext{Tangent line} = rac{f''(\zeta)}{2!} (x-c)^2$$

From this: $f''(c) > 0 \implies$ a local minimum is at c. Converse is **not** true.

Sequences

A sequence on a set A is a function $u:\mathbb{Z}^+ o A$.

Image of the n is written as u_n . A sequence is indicated by one of these ways:

$$\left\{u_n\right\}_{n=1}^{\infty}$$
 or $\left\{u_n\right\}$ or $\left(u_n\right)_{n=1}^{\infty}$

Convergence

Converging

A sequence $ig(u_nig)_{n=1}^\infty$ is converging (to $L\in\mathbb{R}$) iff: $\lim_{n o\infty}u_n=L$

$$orall \epsilon > 0 \; \exists N \in \mathbb{Z}^+ \; orall n \; (n > N \implies |u_n - L| < \epsilon)$$

(i) Note

$$orall x \in \mathbb{R} ~~ \lim_{n o \infty} rac{x^n}{n!} = 0$$

Diverging

A sequence is diverging **iff** it is not converging.

$$\lim_{n o\infty}u_n=egin{cases}\infty\-\infty\ ext{undefined,} & ext{when }u_n ext{ is osciallating} \end{cases}$$

Increasing or Decreasing

A sequence $\left(u_{n}
ight)$ is

- Increasing **iff** $u_n \geq u_m$ for n>m
- Decreasing **iff** $u_n \leq u_m$ for n>m
- Monotone iff either increasing or decreasing
- Strictly increasing ${ ext{iff}}\ u_n>u_m$ for n>m
- Strictly decreasing ${\sf iff}\ u_n < u_m$ for n>m

Convergence test

Increasing and bounded above

Let $ig(u_nig)$ be increasing and bounded above. Then $ig(u_nig)$ is converging (to $\sup\ \{u_n\}$).

(i) Proof Hint

- $\{u_n\}$ has a $\sup u_n (=s)$
- Prove: $\lim_{n o \infty} u_n = s^-$

Decreasing and bounded below

Let $ig(u_nig)$ be decreasing and bounded below. Then $ig(u_nig)$ is converging (to $\inf\ \{u_n\}$).

(i) Proof Hint

- $\{u_n\}$ has a $\inf u_n (=l)$
- Prove: $\lim_{n o \infty} u_n = l^+$

Newton's method of finding roots

Suppose f is a function. To find its roots:

- Select a point x_0
- Draw a tangent at $oldsymbol{x_0}$
- Choose x_1 which is where the tangent meets y=0
- · Continue this process repeatedly

$$x_{n+1}=x_n-rac{f(x_n)}{f'(x_n)}$$

Series

Let $\left(u_{n}
ight)$ be a sequence, and a series (a new sequence) can be defined from it such that:

$$s_n = \sum_{k=1}^n u_k$$

Convergence

If (s_n) is converging:

$$\lim_{n o\infty}s_n=\lim_{n o\infty}\sum_{k=1}^nu_k=\sum_{k=1}^\infty u_k=S\in\mathbb{R}$$

Direct Comparison Test

Let $0 < u_k < v_k$ and $\sum_{k=1}^\infty v_k$ is converging. Then $s_n = \sum_{k=1}^\infty u_k$ is converging.

(i) Proof Hint

- $\sum_{k=1}^n u_k$ and $\sum_{k=1}^n v_k$ are increasing $\sum_{k=1}^\infty v_k$ converges to its supremum

(i) Example

Example usage of this is proving the convergence of $\sum_{k=1}^{\infty} rac{1}{k!}$, by using $k! \geq 2^{k-1}$ $\text{ for all } k \geq 0.$

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