(Deng)

Hi, this is team DD, my name is Deng Kaisheng and this is my teammate Ding Jiayi, (丁哥打招呼)

it’s a great honor to introduce our miniproject to you.

Our project focuses on the paper “External-Memory Exact and Approximate All-Pairs Shortest-Paths in Undirected Graphs”. The paper mainly presented the following 3 algorithms.

1. Present a cache-oblivious exact algorithm for APSP in unweighted undirected graphs

2. Present a cache-aware approximate algorithm for APSP in undirected graphs.

3. Present a cache-aware exact algorithm for APSP in weighted undirected graphs

And for our miniproject, we only discuss the first one, that is…

Before we move a step forward into the algorithm, let’s introduce some background knowledge.

… (background knowledge)

And for today’s video, we will dive deep into the first section of this paper. A cache-oblivious exact algorithm for APSP in unweighted undirected graphs

Before we move a step forward into the new algorithm, we need look back on the MR-BFS algorithm.

(Deng)

Background Knowledge

(Ding)

MR-BFS

BLABLABLA about MR-BFS with graph example.

BLABLABLA with I/O complexity analysis.

After running the MR-BFS, we will get for an unweighted undirected graph G = (V [G],E[G]), the shortest path from a source node s to every node v where s,v ∈ V [G], and we denote this shortest distant by d(u, v).

Before we move further:

We need to clarify 2 observations:

1. d(u, v) = d(v, u) in an undirected graph.
2. For any three vertices u, v and w in G, d(u,w) − d(u, v) ≤ d(v,w) ≤ d(u,w) + d(u, v). Triangle Inequality

(Graph example)

With above Observation, If we sort the adjacency lists we just got in non-decreasing order by d(u, ·), and denote the portion of this sorted list which contains adjacency lists of vertices w with d(u,w) = j by A(j).

Then the adjacency list of any vertex w with d(v,w) = i must reside in some A(j) where i − d(u, v) ≤ j ≤ i + d(u, v).

With this observation, we have the next BFS algorithm.

BLABLABLA about I-BFS with graph example.

BLABLABLA with I/O complexity analysis.

Before we put things together to get our final algorithms, we still have a little step to finish.

Observation 4:

If ET is an Euler Tour of a spanning tree of an unweighted undirected graph G:

1. the number of edges between any two vertices x and y on ET is an upper bound on d(x, y) in G
2. ET has O(V ) edges
3. each vertex of V [G] appears at least once in ET

Ensures that BFS will be performed from each v ∈ V [G].

With this clarified, it’s time to move to the final algorithm.

Firstly, find all vertices of G, and choose a source node to perform MR-BFS. And then for rest vertices, we can use Incremental-BFS to compute the shortest path to all vertices.

BLABLABLA about AP-BFS with graph example.

BLABLABLA with I/O complexity analysis.

Thanks all we are going to share today, thanks for watching!