Hi this is teamDD and we choose this paper: “External-Memory Exact and Approximate All-Pairs Shortest-Paths in Undirected Graphs”. It discusses 3 main contributions.

1. Present a cache-oblivious exact algorithm for APSP in unweighted undirected graphs

2. Present a cache-aware approximate algorithm for APSP in undirected graphs.

3. Present a cache-aware exact algorithm for APSP in weighted undirected graphs

And for today’s video, we will dive deep into the first section of this paper. A cache-oblivious exact algorithm for APSP in unweighted undirected graphs

Before we move a step forward into the new algorithm, we need look back on the MR-BFS algorithm.

BLABLABLA about MR-BFS with graph example.

BLABLABLA with I/O complexity analysis.

After running the MR-BFS, we will get for an unweighted undirected graph G = (V [G],E[G]), the shortest path from a source node s to every node v where s,v ∈ V [G], and we denote this shortest distant by d(u, v).

Before we move further:

We need to clarify 2 observations:

1. d(u, v) = d(v, u) in an undirected graph.
2. For any three vertices u, v and w in G, d(u,w) − d(u, v) ≤ d(v,w) ≤ d(u,w) + d(u, v). Triangle Inequality

(Graph example)

So, If we sort the adjacency lists we just got in non-decreasing order by d(u, ·), and denote the portion of this sorted list which contains adjacency lists of vertices w with d(u,w) = j by A(j).

Then the adjacency list of any vertex w with d(v,w) = I must reside in some A(j) where i − d(u, v) ≤ j ≤ i + d(u, v).

(explain with graph)

AP-BFS Part:

Euler Tour

Increamental-BFS Part: