ECE30030/ITP30010 Database Systems

Advanced SQL

Reading: Chapters 4-5

Charmgil Hong

charmgil@handong.edu

Spring, 2023
Handong Global University



Agenda

- Join
- Views
- Window functions
- Keys

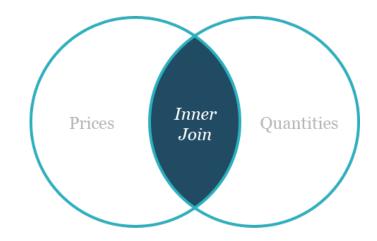
Join Operations

- Join operations take two relations and return another relation
 - A join is a Cartesian product that requires tuples in the two relations match
 - It also specifies the attributes that are present in the result of the join (project)
 - Typically used as subquery expressions in the **FROM** clause
 - Join types
 - INNER JOIN
 - OUTER JOIN
 - Join conditions
 - NATURAL
 - **ON** predicate>
 - USING $(A_1, A_2, ..., A_n)$

- Join: Compare and combine
 - Inner join: Returns matching data from tables
 - Outer join: Returns matching & some dissimilar data from tables

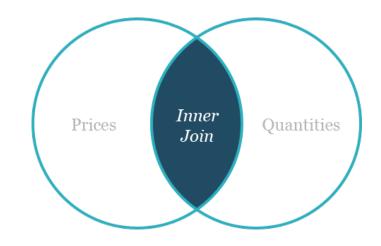
Inner join

Table 1: prices		Table 2: quantities	
product	price	product	quantity
Potatoes	3.00	Potatoes	45
Avocados	4.00	Onions	20
Oranges	5.00	Broccoli	27



• Inner join: matching data

Table 1: prices		Table 2: quantities	
product	price	product	quantity
Potatoes	3.00	Potatoes	45
Avocados	4.00	Onions	20
Oranges	5.00	Broccoli	27

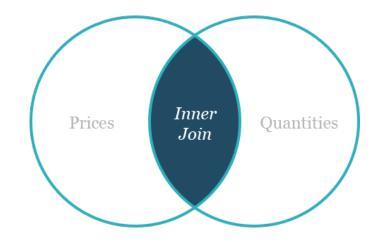


SELECT * **FROM** prices, quantities **WHERE** prices.product = quantities.product;

product	price	product	quantity
Potatoes	3.00	Potatoes	45

• Inner join: matching data

Table 1: prices		Table 2: quantities	
product	price	product	quantity
Potatoes	3.00	Potatoes	45
Avocados	4.00	Onions	20
Oranges	5.00	Broccoli	27



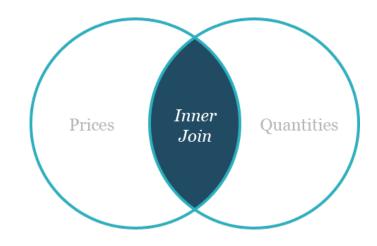
SELECT * FROM prices

JOIN quantities **ON** prices.product=quantities.product;

product	price	product	guantity
Potatoes	3.00	Potatoes	45

• Inner join: matching data

Table 1: prices		Table 2: quantities	
product	price	product	quantity
Potatoes	3.00	Potatoes	45
Avocados	4.00	Onions	20
Oranges	5.00	Broccoli	27

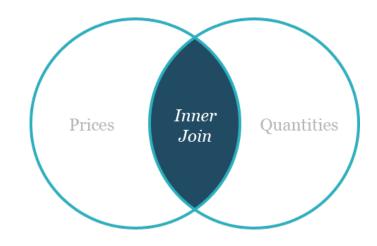


SELECT * **FROM** prices **NATURAL JOIN** quantities;

product	price	quantity
Potatoes	3.00	45

• Inner join: matching data

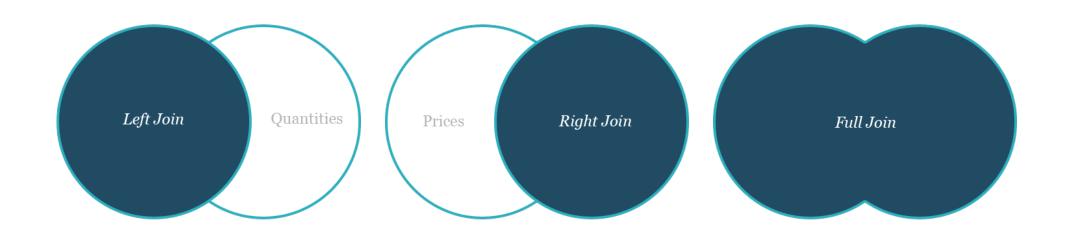
Table 1: prices		Table 2: quantities	
product	price	product	quantity
Potatoes	3.00	Potatoes	45
Avocados	4.00	Onions	20
Oranges	5.00	Broccoli	27



SELECT * **FROM** prices **JOIN** quantities **USING** (product);

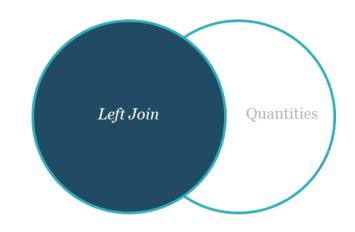
product	price	quantity
Potatoes	3.00	45

- Outer join: matching & some dissimilar data
 - Returns a set of records that include what an inner join would return
 + other rows for which no corresponding match is found in the other table
 - Left (outer) join: "Left" records with no corresponding entry on the "right"
 - Right (outer) join: "Right" records with no corresponding entry on the "left"
 - Full (outer) join: "All" records with no corresponding entry on the other table



• Outer join

Table 1: prices		Table 2: quantities	
product	price	product	quantity
Potatoes	3.00	Potatoes	45
Avocados	4.00	Onions	20
Oranges	5.00	Broccoli	27



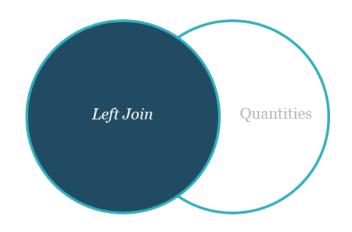
SELECT * FROM prices
LEFT OUTER JOIN quantities
ON prices.product=quantities.product;

product	price	product	quantity
Potatoes	3.00	Potatoes	45
Avocados	4.00	NULL	NULL
Oranges	5.00	NULL	NULL



• Outer join

Table 1: prices		Table 2: quantities	
product	price	product	quantity
Potatoes	3.00	Potatoes	45
Avocados	4.00	Onions	20
Oranges	5.00	Broccoli	27



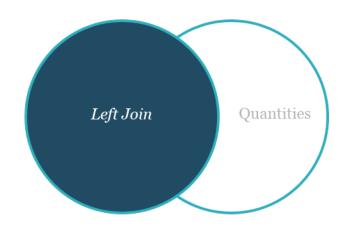
SELECT * FROM prices
LEFT OUTER JOIN quantities
ON prices.product=quantities.product;

product	price	product	quantity
Potatoes	3.00	Potatoes	45
Avocados	4.00	NULL	NULL
Oranges	5.00	NULL	NULL



• Outer join

Table 1: prices		Table 2: quantities	
product	price	product	quantity
Potatoes	3.00	Potatoes	45
Avocados	4.00	Onions	20
Oranges	5.00	Broccoli	27



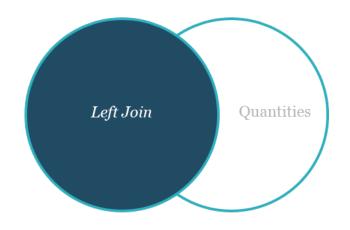
SELECT * FROM prices **NATURAL LEFT OUTER JOIN** quantities;

product	price	quantity
Potatoes	3.00	45
Avocados	4.00	NULL
Oranges	5.00	NULL



• Outer join

Table 1: prices		Table 2: quantities	
product	price	product	quantity
Potatoes	3.00	Potatoes	45
Avocados	4.00	Onions	20
Oranges	5.00	Broccoli	27



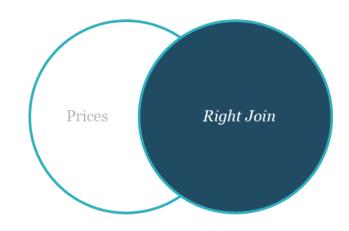
SELECT * **FROM** prices **LEFT OUTER JOIN** quantities **USING** (product);

product	price	quantity
Potatoes	3.00	45
Avocados	4.00	NULL
Oranges	5.00	NULL



• Outer join

Table 1: prices		Table 2: quantities	
product	price	product	quantity
Potatoes	3.00	Potatoes	45
Avocados	4.00	Onions	20
Oranges	5.00	Broccoli	27



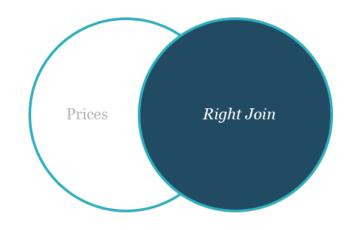
SELECT * FROM prices
RIGHT OUTER JOIN quantities
ON prices.product=quantities.product;

product	price	product	quantity
Potatoes	3.00	Potatoes	45
NULL	NULL	Onions	20
NULL	NULL	Broccoli	27



• Outer join

Table 1: prices		Table 2: quantities	
product	price	product	quantity
Potatoes	3.00	Potatoes	45
Avocados	4.00	Onions	20
Oranges	5.00	Broccoli	27



SELECT * FROM prices **NATURAL RIGHT OUTER JOIN** quantities;

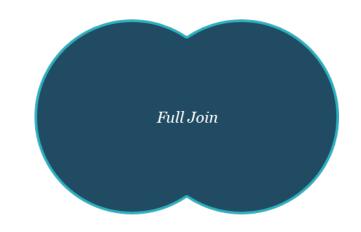
⇔ SELECT * FROM prices
RIGHT OUTER JOIN quantities
USING (product);

product	quantity	price
Potatoes	45	3.00
Onions	20	NULL
Broccoli	27	NULL



• Outer join

Table 1: prices		Table 2: quantities	
product	price	product	quantity
Potatoes	3.00	Potatoes	45
Avocados	4.00	Onions	20
Oranges	5.00	Broccoli	27



SELECT * FROM prices

LEFT OUTER JOIN quantities

ON prices.product=quantities.product

UNION

SELECT * FROM prices

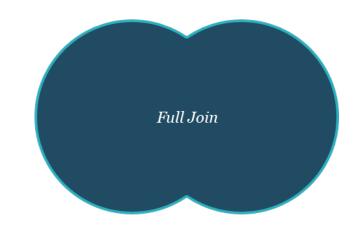
RIGHT OUTER JOIN quantities

ON prices.product=quantities.product;

product	price	product	quantity
Potatoes	3.00	Potatoes	45
Avocados	4.00	NULL	NULL
Oranges	5.00	NULL	NULL
NULL	NULL	Onions	20
NULL	NULL	Broccoli	27

• Outer join

Table 1: prices		Table 2: quantities	
product	price	product	quantity
Potatoes	3.00	Potatoes	45
Avocados	4.00	Onions	20
Oranges	5.00	Broccoli	27



SELECT * FROM prices
FULL OUTER JOIN quantities
ON prices.product=quantities.product;

product	price	product	quantity
Potatoes	3.00	Potatoes	45
Avocados	4.00	NULL	NULL
Oranges	5.00	NULL	NULL
NULL	NULL	Onions	20
NULL	NULL	Broccoli	27

Running Example

• Relations: student, takes

Į₽ ID	‡	"≣ name	📭 dept_name	\$ i≣ tot_cred :
00128		Zhang	Comp. Sci.	102
12345		Shankar	Comp. Sci.	32
19991		Brandt	History	80
23121		Chavez	Finance	110
44553		Peltier	Physics	56
45678		Levy	Physics	46
54321		Williams	Comp. Sci.	54
55739		Sanchez	Music	38
70557		Snow	Physics	0
76543		Brown	Comp. Sci.	58
76653		Aoi	Elec. Eng.	60
98765		Bourikas	Elec. Eng.	98
98988		Tanaka	Biology	120

ID :	📭 course_id		_id 🗼 📭 semester	‡	📭 year ᠄	≣ grade	‡
00128	CS-101	1	Fall		2017	Α	
00128	CS-347	1	Fall		2017	A-	
12345	CS-101	1	Fall		2017	С	
12345	CS-190	2	Spring		2017	Α	
12345	CS-315	1	Spring		2018	Α	
12345	CS-347	1	Fall		2017	Α	
19991	HIS-351	1	Spring		2018	В	
23121	FIN-201	1	Spring		2018	C+	
44553	PHY-101	1	Fall		2017	B-	
45678	CS-101	1	Fall		2017	F	
45678	CS-101	1	Spring		2018	B+	
45678	CS-319	1	Spring		2018	В	
54321	CS-101	1	Fall		2017	A-	
54321	CS-190	2	Spring		2017	B+	
55739	MU-199	1	Spring		2018	A-	
76543	CS-101	1	Fall		2017	Α	
76543	CS-319	2	Spring		2018	Α	
76653	EE-181	1	Spring		2017	С	
98765	CS-101	1	Fall		2017	C -	
98765	CS-315	1	Spring		2018	В	
98988	BIO-101	1	Summer		2017	Α	
98988	BIO-301	1	Summer		2018	<null></null>	

Running Example

• Relations: course, instructor

course_id :	i title :	indept_name ;	⊞ credits :
BIO-101	Intro. to Biology	Biology	4
BIO-301	Genetics	Biology	4
BIO-399	Computational Biology	Biology	3
CS-101	Intro. to Computer Science	Comp. Sci.	4
CS-190	Game Design	Comp. Sci.	4
CS-315	Robotics	Comp. Sci.	3
CS-319	Image Processing	Comp. Sci.	3
CS-347	Database System Concepts	Comp. Sci.	3
EE-181	Intro. to Digital Systems	Elec. Eng.	3
FIN-201	Investment Banking	Finance	3
HIS-351	World History	History	3
MU-199	Music Video Production	Music	3
PHY-101	Physical Principles	Physics	4

₽ ID	‡	p name	\$ 📭 dept_name	‡	⊯ salary ‡
10101		Srinivasan	Comp. Sci.		65000.00
12121		Wu	Finance		90000.00
15151		Mozart	Music		40000.00
22222		Einstein	Physics		95000.00
32343		El Said	History		60000.00
33456		Gold	Physics		87000.00
45565		Katz	Comp. Sci.		75000.00
58583		Califieri	History		62000.00
76543		Singh	Finance		80000.00
76766		Crick	Biology		72000.00
83821		Brandt	Comp. Sci.		92000.00
98345		Kim	Elec. Eng.		80000.00



Natural Join

 Natural join matches tuples with the same values for all common attributes, and retains only one copy of each common column

• E.g., List the names of students along with the course ID of the courses that

they took

SELECT name, course_id
 FROM student, takes
 WHERE student.ID = takes.ID;

Same query in SQL with natural join:

SELECT name, course_id
 FROM student NATURAL JOIN takes;

mame ÷	course_id ÷
Zhang	CS-101
Zhang	CS-347
Shankar	CS-101
Shankar	CS-190
Shankar	CS-315
Shankar	CS-347
Brandt	HIS-351
Chavez	FIN-201
Peltier	PHY-101
Levy	CS-101
Levy	CS-101
Levy	CS-319
Williams	CS-101
Williams	CS-190
Sanchez	MU-199
Brown	CS-101
Brown	CS-319
Aoi	EE-181
Bourikas	CS-101
Bourikas	CS-315
Tanaka	BIO-101
Tanaka	BIO-301

Natural Join

- The **FROM** clause can have multiple relations combined using natural join:
 - SELECT A_1 , A_2 , ... A_n FROM r_1 NATURAL JOIN r_2 NATURAL JOIN ... NATURAL JOIN r_n WHERE P;

Caveat

• E.g., (Incorrect)

SELECT dept_name, course_id, name, title, credits **FROM** student **NATURAL JOIN** takes **NATURAL JOIN** course;

dept_name ÷	i course_id ÷	mame ;	i title ÷	⊞ credits :
Biology	BIO-101	Tanaka	Intro. to Biology	4
Biology	BIO-301	Tanaka	Genetics	4
Comp. Sci.	CS-101	Zhang	Intro. to Computer Science	4
Comp. Sci.	CS-101	Shankar	Intro. to Computer Science	4
Comp. Sci.	CS-101	Williams	Intro. to Computer Science	4
Comp. Sci.	CS-101	Brown	Intro. to Computer Science	4
Comp. Sci.	CS-190	Shankar	Game Design	4
Comp. Sci.	CS-190	Williams	Game Design	4
Comp. Sci.	CS-315	Shankar	Robotics	3
Comp. Sci.	CS-319	Brown	Image Processing	3
Comp. Sci.	CS-347	Zhang	Database System Concepts	3
Comp. Sci.	CS-347	Shankar	Database System Concepts	3
Elec. Eng.	EE-181	Aoi	Intro. to Digital Systems	3
Finance	FIN-201	Chavez	Investment Banking	3
History	HIS-351	Brandt	World History	3
Music	MU-199	Sanchez	Music Video Production	3
Physics	PHY-101	Peltier	Physical Principles	4

Caveat

- Beware of unrelated attributes with same name getting equated incorrectly
 - E.g., List the names of students along with the titles of courses that they have taken
 - Correct

SELECT name, title **FROM** student **NATURAL JOIN** takes, course **WHERE** takes.course_id = course.course_id;

Incorrect

SELECT *name*, *title* **FROM** *student* **NATURAL JOIN** *takes* **NATURAL JOIN** *course*;

 This query omits all (student name, course title) pairs where the student takes a course in a department other than the student's own department

Natural Join with USING Clause

- To avoid the danger of equating attributes erroneously, use the USING construct
 - USING: allows us to specify exactly which columns should be equated
 - E.g., SELECT name, title
 FROM (student NATURAL JOIN takes) JOIN course USING (course_id)

I≣ name ‡	I≣ title	‡
Tanaka	Intro. to Biology	
Tanaka	Genetics	
Zhang	Intro. to Computer Science	
Shankar	Intro. to Computer Science	
Levy	Intro. to Computer Science	
Williams	Intro. to Computer Science	
Brown	Intro. to Computer Science	
Bourikas	Intro. to Computer Science	
Levy	Intro. to Computer Science	
Shankar	Game Design	
Williams	Game Design	
Shankar	Robotics	
Bourikas	Robotics	
Levy	Image Processing	
Brown	Image Processing	
Zhang	Database System Concepts	
Shankar	Database System Concepts	
Aoi	Intro. to Digital Systems	
Chavez	Investment Banking	
Brandt	World History	
Sanchez	Music Video Production	
Peltier	Physical Principles	



JOIN ... ON

- The ON condition allows a general predicate over the relations being joined
 - Written like a WHERE clause predicate
 - E.g., SELECT *
 FROM student JOIN takes ON student.ID = takes.ID
 - The **ON** condition specifies that a tuple from *student* matches a tuple from *takes* if their *ID* values are equal
 - Equivalent to:
 SELECT name, course_id
 FROM student, takes
 WHERE student.ID = takes.ID;

Inner Join

- Inner join: Does not preserve nonmatched tuples
 - Tables are joined based on common columns mentioned in the ON or USING clause
 - One can specify the condition with an ON or USING construct
- C.f., Natural join: assumes the join condition to be where samenamed columns in both tables match
 - One cannot use ON or USING
 - In the result of a natural join, repeated columns are avoided

Natural Join

 Natural join: Some tuples in either or both relations being joined may be lost

• SELECT *
FROM course NATURAL JOIN prereq;

page course_id	‡ ⊞ title	: I≣ dept_name ÷	⊞ credits :	prereq_id ÷
BIO-301	Genetics	Biology	4	BIO-101
CS-190	Game Design	Comp. Sci.	4	CS-101

Examples

• Tables

ROLL_NO	NAME
1	HARSH
2	PRATIK
3	RIYANKA
4	DEEP
5	SAPTARHI
6	DHANRAJ
7	ROHIT
8	NIRAJ

COURSE_ID	ROLL_NO
1	1
2	2
2	3
3	4
1	5
4	9
5	10
4	11

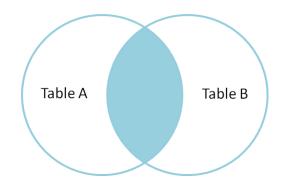
Student

StudentCourse

Examples

- Inner join
 - SELECT StudentCourse.COURSE_ID, Student.NAME
 FROM Student
 INNER JOIN StudentCourse
 ON Student.ROLL_NO = StudentCourse.ROLL_NO;

COURSE_ID	NAME
1	HARSH
2	PRATIK
2	RIYANKA
3	DEEP
1	SAPTARHI

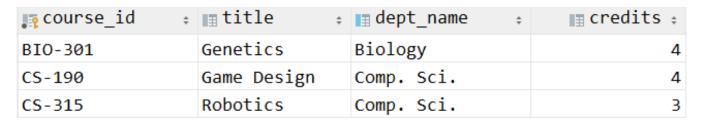


Outer Join

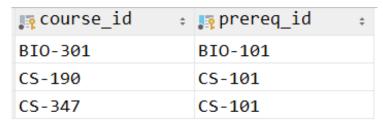
- An extension of the join operation that avoids loss of information
 - Outer join preserves those tuples that would be lost in a join by creating tuples in the result containing null values
 - Computes the join and then adds tuples form one relation that does not match tuples in the other relation to the result of the join
- Three forms of outer join:
 - LEFT OUTER JOIN
 - RIGHT OUTER JOIN
 - FULL OUTER JOIN

Running Example

• Relation course



Relation prereq



- course is missing CS-347
- prereq is missing CS-315

Inner Join with NATURAL

 Natural join: Some tuples in either or both relations being joined may be lost

SELECT *
 FROM course NATURAL JOIN prereq;

i course_id ÷	i title	dept_name ;	i credits :	prereq_id ÷
BIO-301	Genetics	Biology	4	BIO-101
CS-190	Game Design	Comp. Sci.	4	CS-101

Left Outer Join with NATURAL

- Left outer join: Preserves tuples only in the relation named before (to the left of) the operation
- SELECT *
 FROM course NATURAL LEFT OUTER JOIN prereq;

i course_id :	: I≣ title ÷	dept_name ÷	i credits :	prereq_id ÷
BIO-301	Genetics	Biology	4	BIO-101
CS-190	Game Design	Comp. Sci.	4	CS-101
CS-315	Robotics	Comp. Sci.	3	<null></null>

Right Outer Join with NATURAL

- Right outer join: Preserves tuples only in the relation named after (to the right of) the operation
- SELECT *
 FROM course NATURAL RIGHT OUTER JOIN prereq;

⊞ course_id	‡	prereq_id	‡	i title	‡	dept_name	‡	⊞ credits :
BIO-301		BIO-101		Genetics		Biology		4
CS-190		CS-101		Game Design		Comp. Sci.		4
CS-347		CS-101		<null></null>		<null></null>		<null></null>

Full Outer Join with NATURAL

SELECT *
 FROM course NATURAL FULL OUTER JOIN prereq;

page course_id	‡ title	dept_name	≣ credits	<pre> prereq_id</pre>
BIO-301	Genetics	Biology		4 BIO-101
CS-190	Game Design	Comp. Sci.		4 CS-101
CS-315	Robotics	Comp. Sci.		3 <null></null>
CS-347	<null></null>	<null></null>	<nul< td=""><td>L> CS-101</td></nul<>	L> CS-101

- MySQL does NOT support FULL join
 - Alternative: use the UNION of left and right joins
 SELECT course_id, title, dept_name, credits, prereq_id
 FROM course NATURAL LEFT OUTER JOIN prereq
 UNION
 SELECT course_id, title, dept_name, credits, prereq_id
 FROM course NATURAL RIGHT OUTER JOIN prereq;
 - In order to perform UNION properly, the attributes of both join queries must be aligned

Examples

• Tables

ROLL_NO	NAME
1	HARSH
2	PRATIK
3	RIYANKA
4	DEEP
5	SAPTARHI
6	DHANRAJ
7	ROHIT
8	NIRAJ

COURSE_ID	ROLL_NO
1	1
2	2
2	3
3	4
1	5
4	9
5	10
4	11

Student

StudentCourse

Examples

- Left join
 - SELECT Student.NAME, StudentCourse.COURSE_ID
 FROM Student
 LEFT JOIN StudentCourse
 ON StudentCourse.ROLL_NO = Student.ROLL_NO;

NAME	COURSE_ID
HARSH	1
PRATIK	2
RIYANKA	2
DEEP	3
SAPTARHI	1
DHANRAJ	NULL
ROHIT	NULL
NIRAJ	NULL

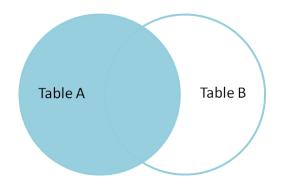
- Right join
 - SELECT Student.NAME, StudentCourse.COURSE_ID
 FROM Student
 RIGHT JOIN StudentCourse
 ON StudentCourse.ROLL_NO = Student.ROLL_NO;

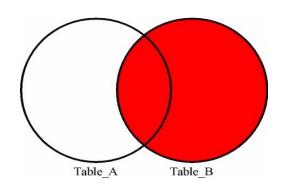
NAME	COURSE_ID
HARSH	1
PRATIK	2
RIYANKA	2
DEEP	3
SAPTARHI	1
NULL	4
NULL	5
NULL	4

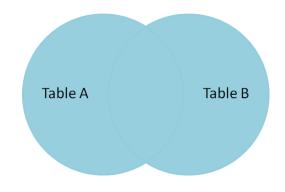
Examples

- Full join
 - SELECT Student.NAME, StudentCourse.COURSE_ID
 FROM Student
 FULL JOIN StudentCourse
 ON StudentCourse.ROLL_NO = Student.ROLL_NO;

NAME	COURSE_ID
HARSH	1
PRATIK	2
RIYANKA	2
DEEP	3
SAPTARHI	1
DHANRAJ	NULL
ROHIT	NULL
NIRAJ	NULL
NULL	9
NULL	10
NULL	11



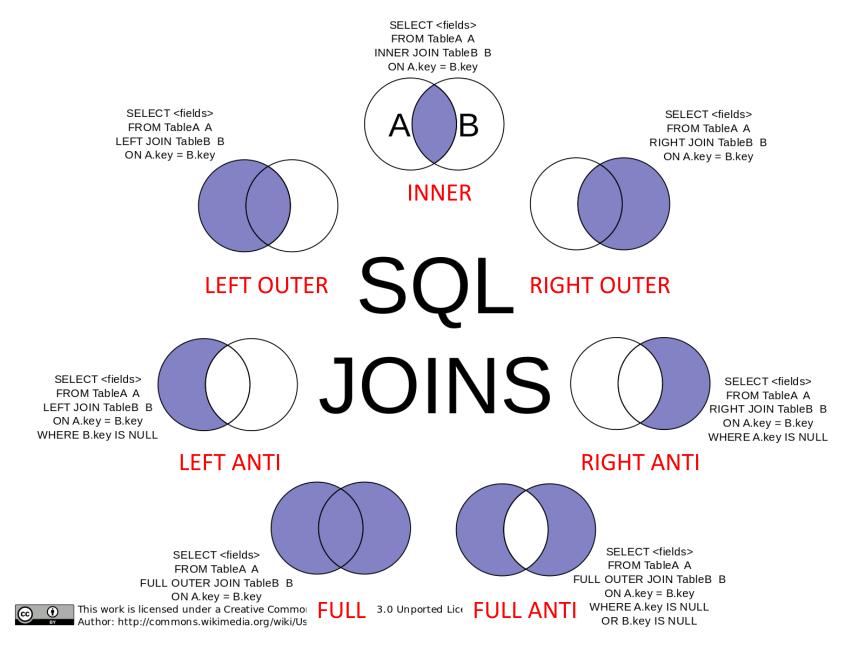




Join Types and Conditions

- Join type: Defines how tuples in each relation that do not match any tuples in the other relation are treated
 - INNER JOIN
 - LEFT OUTER JOIN
 - RIGHT OUTER JOIN
 - FULL OUTER JOIN
- Join condition: Defines which tuples in the two relations match
 - NATURAL
 - **ON** predicate>
 - USING $(A_1, A_2, ..., A_n)$

Join Types



Join Condition

- Join condition
 - NATURAL: Joins two tables based on same attribute name and datatypes
 - SELECT * FROM course NATURAL JOIN prereq;
 - ON predicate>: Joins two tables based on the column(s) explicitly specified in the ON clause
 - SELECT * FROM course JOIN prereq ON course.course_id = prereq.prereq_id;
 - USING (A₁, A₂, ..., A_n): Joins two tables based on common attribute name(s) listed next to USING
 - SELECT * FROM course
 JOIN prereq USING (course_id)

Inner Join vs. Natural Join

INNER JOIN	NATURAL JOIN
Joins two tables on the basis of the column which is explicitly specified in the ON clause	Joins two tables based on same attribute name and datatypes
The resulting table will contain all the attribute of both the tables (including duplicate columns)	The resulting table will contain all the attribute of both the tables but keep only one copy of each common column
Only those records will return which exists in both tables	If there is no indication of LEFT, RIGHT, or FULL, it returns the rows based on the common column

^{*} Source: https://www.geeksforgeeks.org/difference-between-natural-join-and-inner-join-in-sql/



Inner Join vs. Natural Join

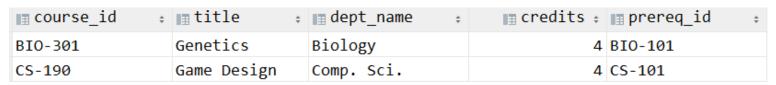
- Inner join
 - SELECT * FROM course
 INNER JOIN prereq ON course.course_id = prereq.prereq_id;
- Natural join
 - SELECT *
 FROM course NATURAL JOIN prereq
 ON course.course_id = prereq.prereq_id; ← NOT VALID!

Inner Join vs. Natural Join

- Inner join
 - SELECT * FROM course
 INNER JOIN prereq ON course.course_id = prereq.course_id;
 - Equivalent to:
 SELECT * FROM course
 JOIN prereq ON course_id = prereq.course_id;

<pre> course.course_id</pre>	i title ;	dept_name ;	i credits :	<pre> prereq.course_id ‡ </pre>	prereq_id ÷
BIO-301	Genetics	Biology	4	BIO-301	BIO-101
CS-190	Game Design	Comp. Sci.	4	CS-190	CS-101

- Natural join
 - SELECT *
 FROM course NATURAL JOIN prereq;





Outer Join vs. Natural Join

- Right outer join
 - SELECT *
 FROM course NATURAL RIGHT OUTER JOIN prereq;
 - Figure 1 Equivalent to:

 SELECT *

 FROM course RIGHT OUTER JOIN prereq

 USING (course_id);

⊞ course_id	prereq_id ÷	i title ;	dept_name ÷	⊞ credits :
BIO-301	BIO-101	Genetics	Biology	4
CS-190	CS-101	Game Design	Comp. Sci.	4
CS-347	CS-101	<null></null>	<null></null>	<null></null>

 SELECT *
 FROM course RIGHT OUTER JOIN prereq
 ON course.course_id = prereq.course_id;

<pre> course.course_id</pre>	i title ÷	dept_name :	p credits :	<pre> prereq.course_id ‡</pre>	prereq_id ÷
BIO-301	Genetics	Biology	4	BIO-301	BIO-101
CS-190	Game Design	Comp. Sci.	4	CS-190	CS-101
<null></null>	<null></null>	<null></null>	<null></null>	CS-347	CS-101



Outer Join vs. Natural Join

- Left outer join
 - SELECT *
 FROM course NATURAL LEFT OUTER JOIN prereq;
 - Equivalent to:
 SELECT *
 FROM course LEFT OUTER JOIN prereq

USING (course_id);

i course_id	i title ;	dept_name	i credits :	prereq_id ÷
BIO-301	Genetics	Biology	4	BIO-101
CS-190	Game Design	Comp. Sci.	4	CS-101
CS-315	Robotics	Comp. Sci.	3	<null></null>

 SELECT *
 FROM course LEFT OUTER JOIN prereq
 ON course.course_id = prereq.course_id;

<pre> course.course_id</pre>	i title ‡	dept_name ;	i credits :	prereq.course_id :	prereq_id ÷
BIO-301	Genetics	Biology	4	BIO-301	BIO-101
CS-190	Game Design	Comp. Sci.	4	CS-190	CS-101
CS-315	Robotics	Comp. Sci.	3	<null></null>	<null></null>



Natural Joins Are Often Avoided

- Natural joins are often avoided in practice, because:
 - Natural joins are not particularly readable (by most SQL coders) and possibly not supported by various tools/libraries
 - Natural joins are not informative; you cannot tell what columns are being joined on without referring to the schema
 - Your join conditions are invisibly vulnerable to schema changes
 - Even if there are multiple natural join columns and one such column is removed from a table, the query will still execute
 - But the result may not be correct and this change in behavior will be silent
 - Hardly worth the effort; you are only saving about 10 seconds by not typing specific conditions

ECE30030/ITP30010 Database Systems

Term Project

Charmgil Hong

charmgil@handong.edu

Spring, 2023
Handong Global University



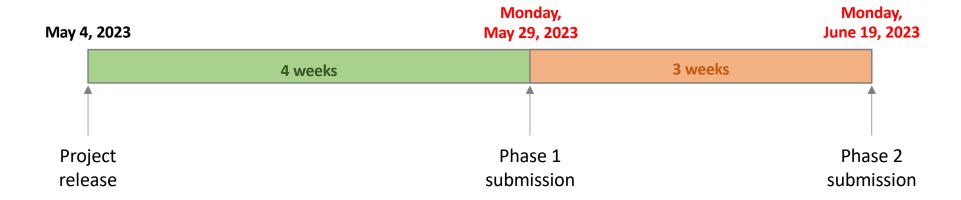
Term Project

Goals

- To practice the concepts and underlying mechanisms of database management system with an actual database instance
- To represent database designs in modeling languages and analyze the designs with respect to given constraints
- To articulate the relational database language (structured query language)
- To exercise the optimization and evaluation of the database performance
- In this project, each team will be given a large chunk of data that is completely unnormalized
 - Your objective is to design a "good" database schema that can accommodate the provided data without any loss of information
 - "Good" in that...
 - Efficient in terms of space and time complexity

Term Project Overview

Planned timeline



- Phase 1 "space" submission: Monday, May 29, 2023
- Phase 2 "time" submission: Monday, June 19, 2023

KUBiC: Korean Unification Bigdata Center

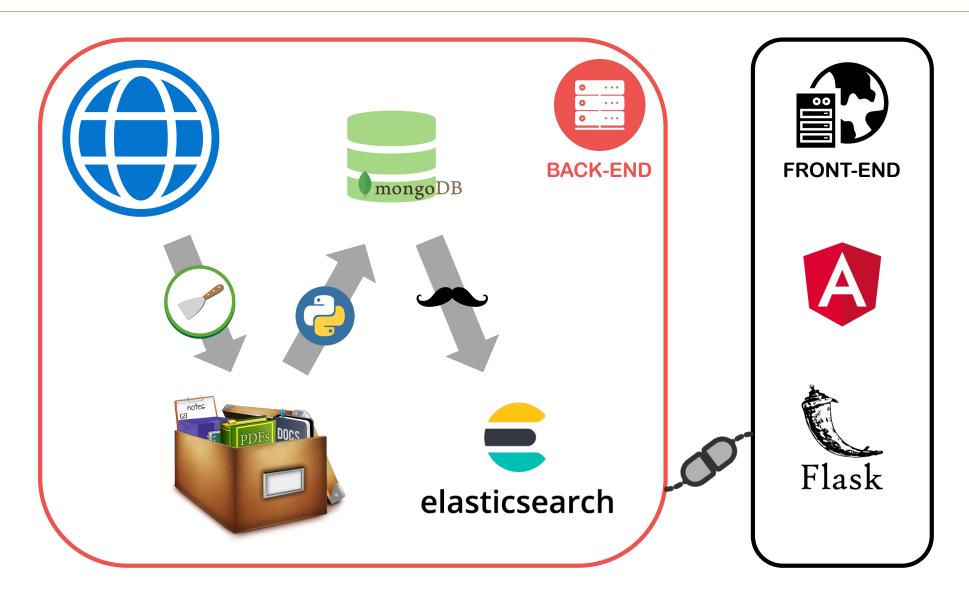
- Term-Project data is provided by the KUBiC project team
- A government-funded project on a data-center development focusing on the Korean unification
 - URL: https://kubic.handong.edu/
 - Data archive + search engine + web-based analysis tools, specialized on the Korean unification and North Korea research
 - Contains a lot of academic papers and government reports on the relevant topics



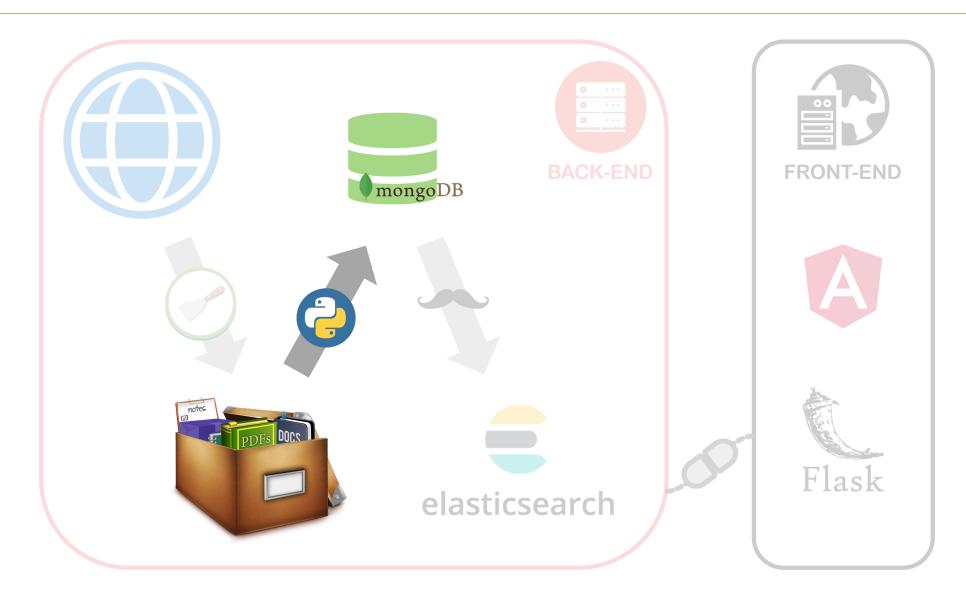




KUBiC: Korean Unification Bigdata Center



KUBiC: Korean Unification Bigdata Center



Term Project

Background

- You will be given large chunks of data snapshot from the KUBIC database, that consist of one SQL dump file and two csv files
 - core.sql
 - 116,320 records, 42 columns (approx. 2.45 GB)
 - Completely unnormalized
 - tfidf.csv
 - TF-IDF analysis of the service documents
 - 877,490 records, 4 columns (approx. 170.6 MB)
 - rcmd.csv
 - Cosince similarity analysis of the service documents
 - 1,000,000 records, 3 columns (approx. 126.8 MB)
- SQL dump file: Ordinary text file, written in the SQL syntax
 - Contains a record of the table structure and/or the data from a database
 - Often used for backing up a database so that its contents can be restored in the event of data loss

Provided Data

• Core

- Collection of core meta-data about the web-documents that KUBIC contains
- Also contains the bulletin boards, user information, saved documents of each user
- 116,320 records, 42 columns (2.45GB)
- Completely unnormalized

⊞ core	
I ≣_id	
I ∄ isAdmin	
II isApiUser	
II≣ name	
I ∄ email	
I ∄ inst	
I≣ status	
I≣ userId	
I ≣ registeredDate	
I modifiedDate	
I isActive	
II≣ type	
II≣ title	
II content	
I writerName	
I writerEmail	
II regDate	
I modDate	
III docID	
II isMainAnnounce	
II category	
II userEmail	
II≣ keyword	
II savedDate	
I savedDocHashKeys	
■ post_title	
■ post_writer	
II post_date	
■ post_body	
■ published_institution	
■ published_institution_url	
■ top_category	
■ original_url	
III file_download_url	
II file_name	
II file_id_in_fsfiles	
☐ file_extracted_content ☐	
III timestamp	char(255)
II hash_key	char(255)
II topic	
■ docTitle	
II hashKey	

Provided Data

- tfidf
 - TF-IDF analysis of the service documents
 - For Phase 1, this table is supposed to be empty

```
frequency

docID char(255)

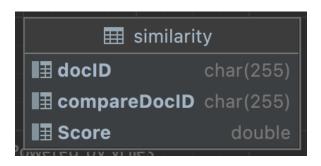
docTitle char(255)

fill tfidfWord char(255)

Score double
```

Provided Data

- simscores
 - Cosince similarity analysis of the service documents
 - For Phase 1, this table is supposed to be empty



Term Project

- Phase 1 requirements
 - Design and implment a database that can effectively accommodate the entire data without any loss
 - You and your team will need to draw E-R diagrams and conduct a number of normalization processes
 - Import the data; there should be no missing portion
 - You will be asked to create and submit views
 - Make the database size as small as possible!
- Phase 2 requirements
 - Optimize the database using
 - Denormalization
 - Indexing

Data Files

- Core
 - https://drive.google.com/file/d/1BUTHZv0AgZPUEaOna3loxUklSXO8VfZ5/view?usp=sharing
- Tfidf
 - https://drive.google.com/file/d/1MUNteBF58NZHNLOf31ZN90BkE0MSUS8 H/view?usp=sharing
- Rcmds
 - https://drive.google.com/file/d/14QpCNHPQEucieDK6iWBYjY_Xflz2DWKW/view?usp=sharing

Technical Resources

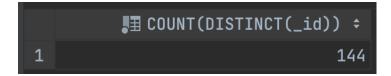
- Upon completion, submit your result to LMS. Each submission should have the following items:
 - Dump of the database (in .sql)
 - How to create a SQL dump?
 - https://dev.mysql.com/doc/refman/8.0/en/mysqldump.html
 - https://dev.mysql.com/doc/refman/8.0/en/mysqldump-sql-format.html
 - Report Documents (in .pdf)
 - How to attack this problem?
 - DDL query and result for View instruction
 - ER Diagram of your database
 - Submission should be one .zip file

TA's are up for help

- Chanju Lee (이찬주), Seohwee (박서휘): Data-specific questions
- Jihyeon Song (송지현), Harim Kim (김하림): SQL and DBMS functionalities-related questions

- Goal: Design and implement a database instance that is efficient in space
 - You are expected to conduct a database design using ERD and apply the normalization theory
 - We will check the correctness and completeness of your data by examining the output of the views suggested in next slides
 - The database size on the physical storage will be estimated; the smallest 10% teams will earn bonus points (maximum +7%)
 - Before the submission, each team is expected to run several iterations of design, implement, data import, and internal evaluation

- Views to create (and submit)
 - 1. View: userCount
 - Count the number of users in the database
 - SELECT * FROM userCount



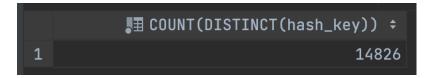
• The column name may vary

- Views to create (and submit)
 - 2. View: boardCount
 - Count the number of bulletins on the board
 - **SELECT** * **FROM** boardCount

```
$ COUNT(DISTINCT(title)) ÷
1 37
```

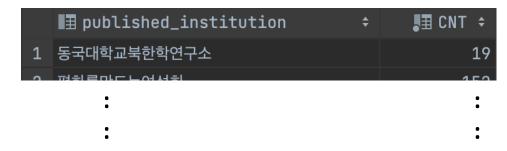
• The column name may vary

- Views to create (and submit)
 - 3. View: docCount
 - Count the number of documents that are stored
 - **SELECT** * **FROM** docCount



• The column name may vary

- Views to create (and submit)
 - View: instPubInfo
 - List the names of publisher institutes and their numbers of publications (sort the results in ascending order of the number of publications)
 - **SELECT** * **FROM** instPubInfo



- Views to create (and submit)
 - 5. View: docInfo
 - List the posting title, post author name and affiliation, posted date, and top category tag
 - SELECT * FROM docInfo



:

- Views to create (and submit)
 - 6. View: bulletinSummary
 - List all bulletin titles, author names (writer names), and posted dates
 - SELECT * FROM bulletinSummary



- Views to create (and submit)
 - 7. View: docSummary
 - Count the number of documents per each of top category values; show the results in descending order of the counts and put their ranks
 - SELECT * FROM docSummary



- Views to create (and submit)
 - 8. View: fileSummary
 - Show the attached file information by summarizing their timestamp, file ID, filename, and download url
 - SELECT * FROM fileSummary

```
Y- WHERE
                                 ≡ • ORDER BY
                  III timestamp

‡ II file_name
                                           김정은 정권의 대남정책 및 통일담론 : 텍스트마이닝을 이용한 분석 http://unibook.unikorea.go.kr
                    608591d4f879c5b21a2fa295
                    60859191f879c5b21a2fa16d International Journal of Korean Unification S...
                                                                                     http://unibook.unikorea.go.kr/
 2021-04-26 12:55:59
                    6085910ef879c5b21a2f9ee1
                                           평화의 심리학 : 한국인의 평화인식
                                                                                      http://unibook.unikorea.go.kr/
                                           북한인권 책임규명 방안과 과제 : 로마규정 관할범죄에 대한 형사소: http://unibook.unikorea.go.kr/
  2021-04-26 12:52:39 60859046f879c5b21a2f99b6
                                           통일 이후 통합방안 : 민족주의와 편익을 넘어선 통일담론의 모색 「http://unibook.unikorea.go.kr,
```

- A query to check the size of your database instance
 - SELECT table_schema AS 'DatabaseName',
 ROUND(SUM(data_length+index_length)/1024, 1) AS 'Size(KB)'
 FROM information_schema.tables
 WHERE table_schema = 'YOUR DATABASE NAME'
 GROUP BY table_schema;
- A query to check each table size from your database
 - SELECT TABLE_SCHEMA, TABLE_NAME,
 ROUND(DATA_LENGTH/(1024), 1) AS 'data(KB)',
 ROUND(INDEX_LENGTH/(1024), 1) AS 'idx(KB)'
 FROM information_schema.tables
 WHERE TABLE_TYPE = 'BASE TABLE'
 AND TABLE_SCHEMA = 'YOUR DATABASE SIZE';

- What to submit
 - A report including
 - ER diagram of the implemented database
 - List of all tables and their attributes with precise notions of data types and integrity constraints
 - Description of the requested views
 - Size of the resulting table (in counts)
 - The screenshots of the table header and first five records
 - Summary of the database size and table sizes (in Kilobytes)
 - A zipped MySQL dump file containing all the database implementations including the database schema, records, views, etc.

- Resources
 - How to create a dump file
 - MySQL Workbench https://dev.mysql.com/doc/workbench/en/wb-admin-export-import-management.html
 - DataGrip https://www.jetbrains.com/help/datagrip/export-data-in-ide.html
 - HeidiSQL https://www.heidisql.com/screenshots.php?which=export_sql
 - SequelAce https://sequelpro.com/docs/ref/working-with-data