

# **Week 11 - DFS app., MST**

Algorithm Analysis

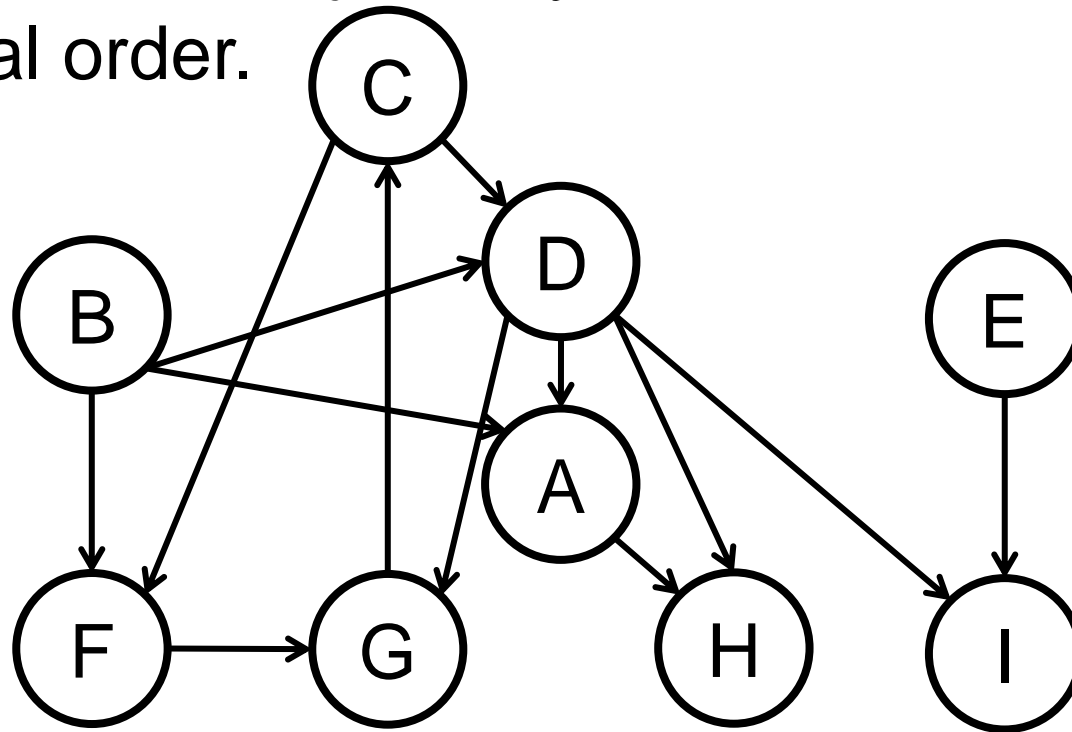
School of CSEE

- Due by 5/24(Wed.)
- Computes transpose of graph using adjacency array/list to save time.

# DFS application

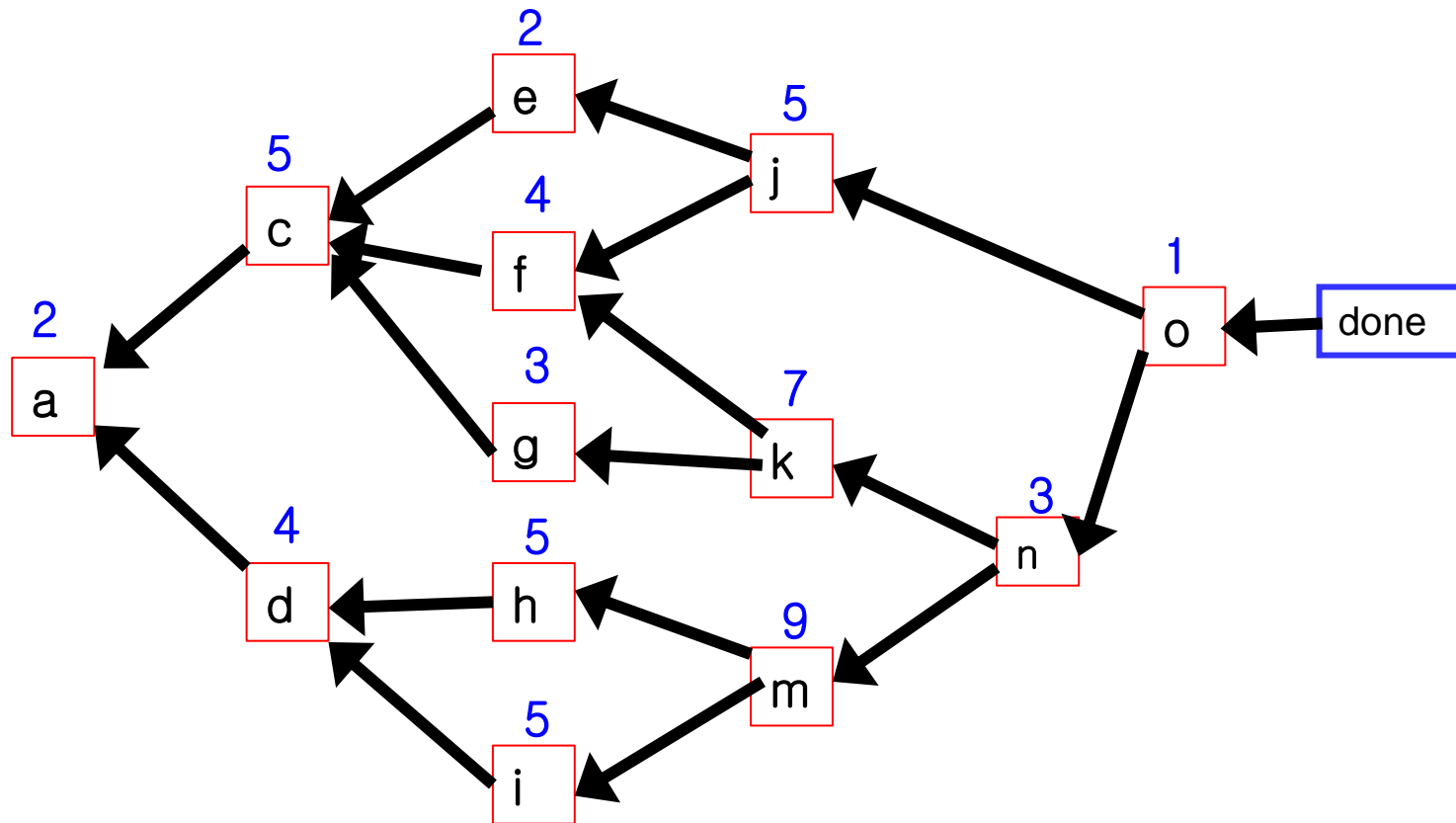
# Exercise 2

From the following directed graph, determine the strongly connected component. Assume that vertices and each adjacency list are in alphabetical order.

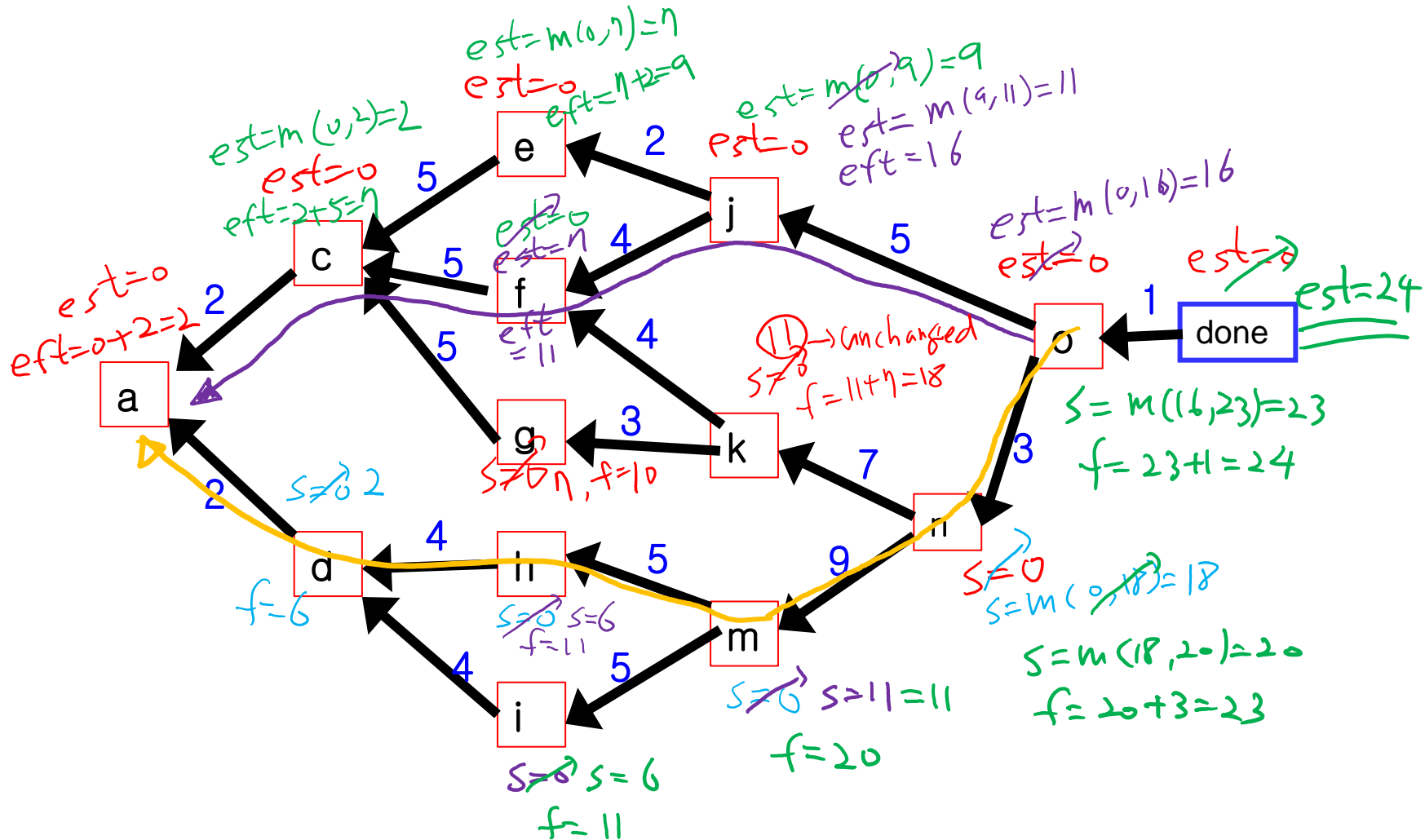


# Exercise 3

Determine 'est' and 'eft' of each node.



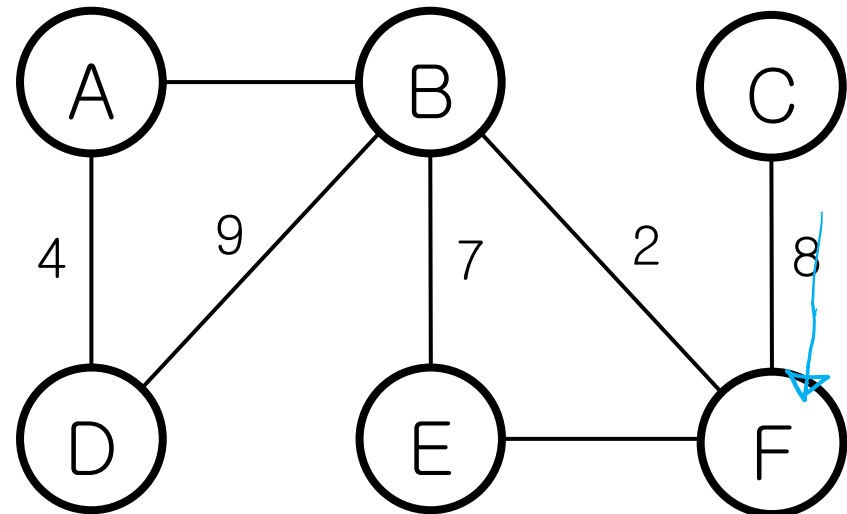
# Exercise 3



# Exercise 4

Determine the value of  $d[E]$  and  $back_F$  when the node B is determined as articulation point. Assume adjacency array and adjacency list is in alphabetical order.

(Discovery time of vertex A is 1.)



# Articulation Point

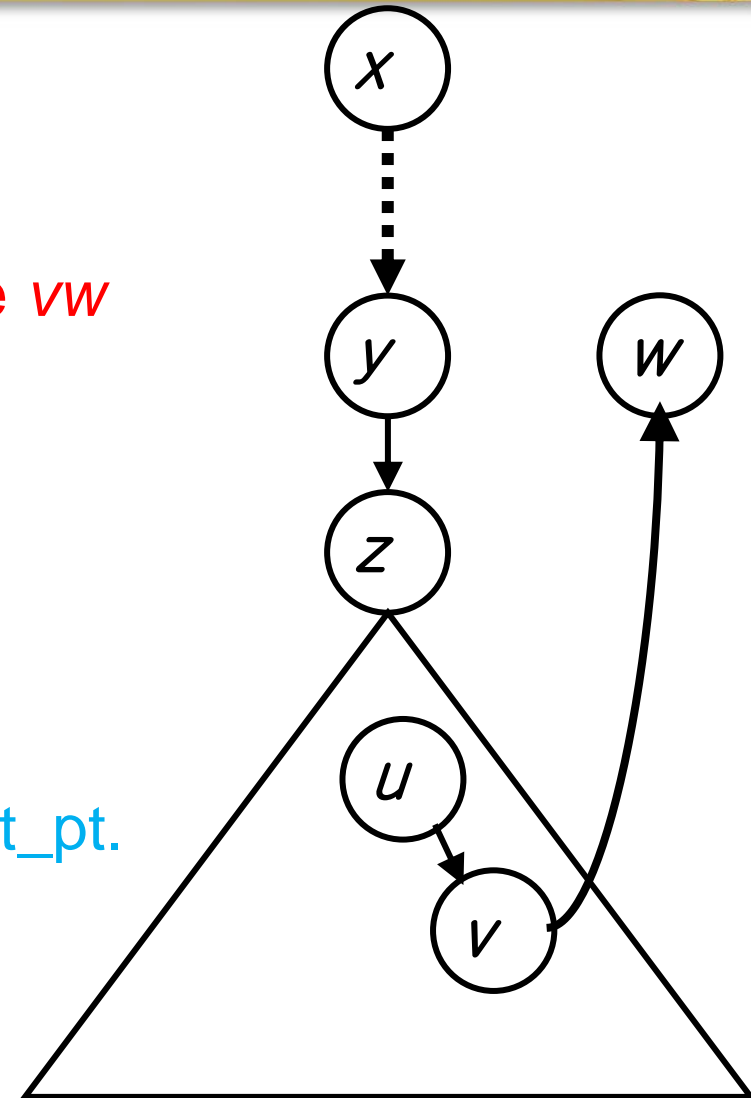
For all  $v < y$ , check back edge  $vw$ .

- a. *back* is initialized to  $d[v]$ .
- b. When encountered with back edge  $vw$   
 $back_v = \min(back_v, d(w))$
- c. When backtracking from  $v$  to  $u$ ,  
 $back_u = \min(back_u, back_v)$

When backing up from  $z$  to  $y$ ,

if all vertices  $\leq z$ ,  $back_z \geq d(y)$ . :  $y$  is art\_pt.

else  $back_z < d(y)$ . :  $y$  isn't art\_pt.

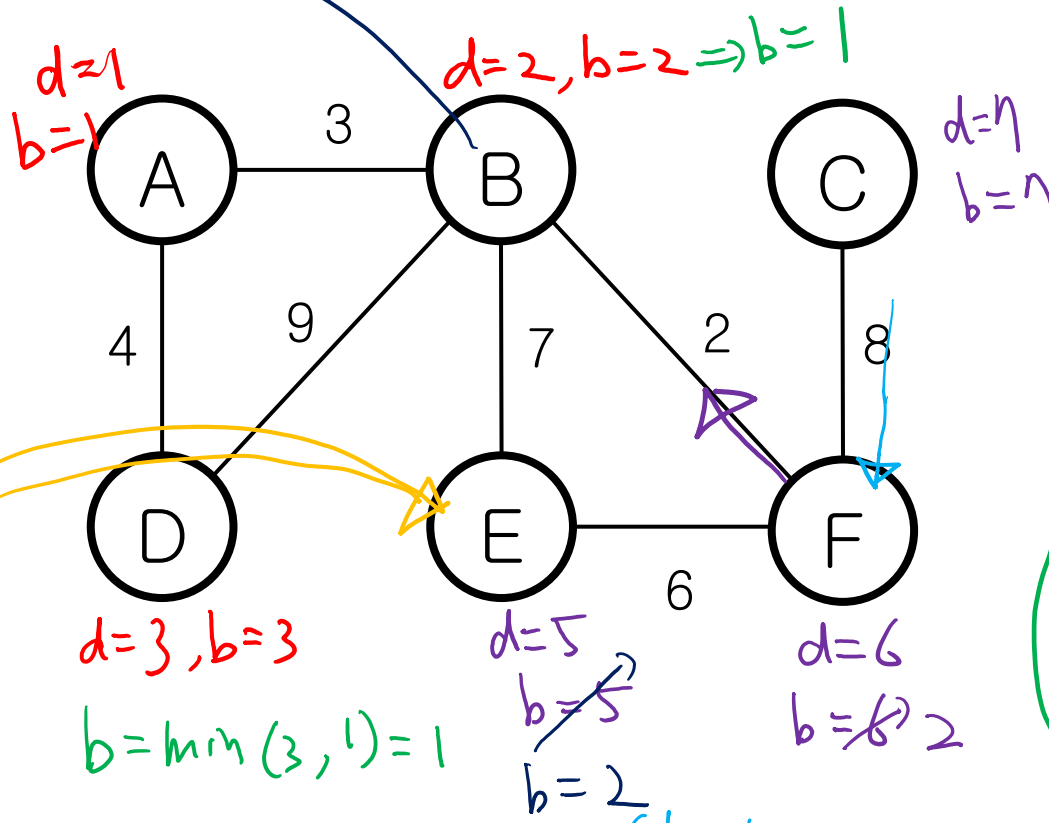
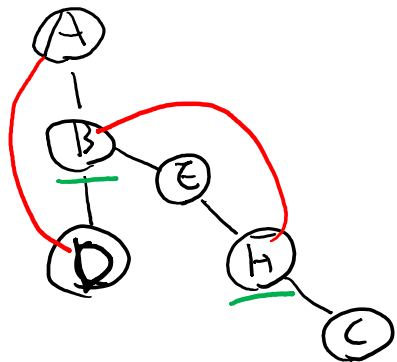




When backing up from  $z$  to  $y$ ,  
 if all vertices  $\leq z$ ,  $back_z \geq d(y)$ . :  $y$  is art\_pt.  
 else  $back_z < d(y)$ . :  $y$  isn't art\_pt.

$(back_E = 2) \geq d_B = 2$  : B is A.P.

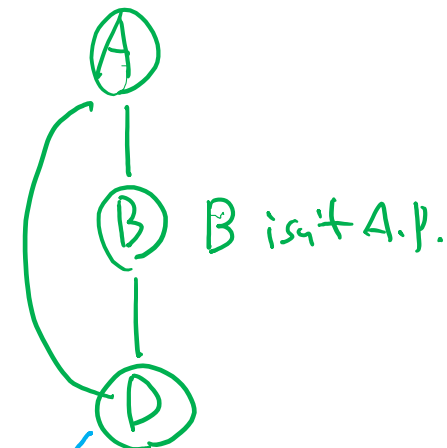
$(back_D = 1) < (d_B = 2)$  : B isn't A.P.



$(back_H = 2)$   
 $< (d_E = 5)$   
 : E isn't A.P.

$b = \min(3, 1) = 1$

$(back_C = 7) \geq (d_F = 6)$  : F is A.P.



MST

# Exercise 1

What is data structure used for simple implementation of Prim's algorithm?