## Turing machines: examples

CS 580

 $A = \{0^{2^n} | n \ge 0\}$ ; i.e. the language consisting of all strings of 0s whose length is a power of 2.

## **Solution:** The machine is:

M = "On input string w:

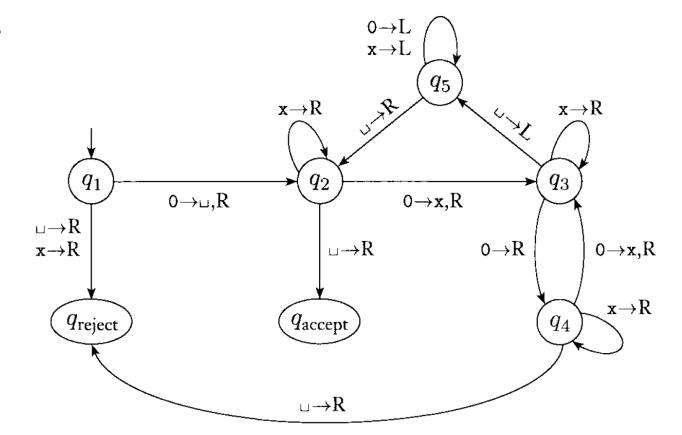
- 1. Sweep left to right across the tape, crossing off every other 0.
- 2. If in step 1 the tape contained a single 0, accept.
- 3. If in step 1 the tape contained more than a single 0 and the number of 0s was odd, reject.
- 4. Return the head to the left-hand end of the tape.
- 5. Go to step 1."

**Explanation:** Each iteration of stage 1 cuts the number of 0s in half.

## Formal description: $M = (Q, \Sigma, \Gamma, \delta, q_1, q_{reject})$ : $Q = \{q_1, q_2, q_3, q_4, q_5, q_{accept}, q_{reject}\}$ $\Sigma = \{0\}$ $\Gamma = \{0, x, \sqcup\}$

 $\delta$ : described via state diagram. The start state is  $q_1$ , the accept state is  $q_{accept}$ , and the

reject state is



$$B = \{w \# w \mid w \in \{0,1\}^*\}$$

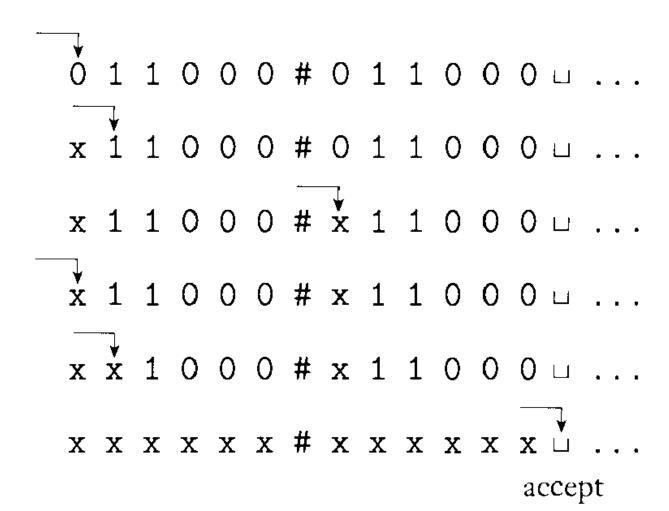
$$B = \{w \# w \mid w \in \{0,1\}^*\}$$

M = "On input string w:

- 1. Zig-zag across the tape to corresponding positions on either side of the # symbol to check whether these positions contain the same symbol. If they do not, or if no # is found, reject. Cross off symbols as they are checked to keep track of which symbols correspond.
- 2. When all symbols to the left of the # have been crossed off, check for any remaining symbols to the right of the #. If any symbols remain, reject; otherwise, accept."

$$B = \{w \# w \mid w \in \{0,1\}^*\}$$

Snapshots of the machine on input 011000#011000.



$$B = \{w \# w \mid w \in \{0,1\}^*\}$$

State diagram.

