

# SCHUR COMPLEMENTS AND OBSTRUCTIONS FOR LARGE $\varepsilon$ -LIGHT VERTEX SETS IN GRAPHS

ABSTRACT. Let  $G = (V, E)$  be a finite undirected graph (possibly disconnected) with Laplacian  $L$ . For  $S \subseteq V$ , let  $G_S = (V, E(S, S))$  be the graph on the same vertex set but with only edges whose endpoints both lie in  $S$ , and let  $L_S$  be its Laplacian (equivalently, the Laplacian of the induced subgraph  $G[S]$ , padded with isolated vertices in  $V \setminus S$ ). Fix  $\varepsilon \in (0, 1)$ . We call  $S$   $\varepsilon$ -light if  $\varepsilon L - L_S \succeq 0$ .

We study the question of whether there exists a universal constant  $c > 0$  such that every graph contains an  $\varepsilon$ -light set of size at least  $c\varepsilon|V|$  for every  $\varepsilon \in (0, 1)$ . We provide (i) careful kernel/projection bookkeeping for disconnected graphs and equivalent normalized formulations; (ii) sharp extremal examples showing any such  $c$  must satisfy  $c \leq \frac{1}{2}$  and a complete characterization for complete graphs; (iii) explicit counterexamples showing that natural “edgewise effective resistance” and “vertex-star linearization” strategies do not certify  $\varepsilon$ -lightness; and (iv) an exact variational and Schur complement (Kron reduction) characterization that isolates the true obstruction: internal energy in  $G[S]$  must be dominated by the effective coupling of  $S$  to  $V \setminus S$ . Finally, we document why current interlacing-polynomial methods [2] do not directly apply to vertex-induced edge selection, by exhibiting a  $K_3$  example where the expected characteristic polynomial has nonreal roots.

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## 1. INTRODUCTION

Let  $G = (V, E)$  be a finite, undirected, unweighted graph with  $n := |V|$ . Its (combinatorial) Laplacian is the symmetric matrix  $L \in \mathbb{R}^{V \times V}$  defined by

$$L_{uu} = \deg(u), \quad L_{uv} = \begin{cases} -1, & \{u, v\} \in E, \\ 0, & \text{otherwise.} \end{cases}$$

Equivalently,

$$x^\top Lx = \sum_{\{u, v\} \in E} (x_u - x_v)^2 \quad (x \in \mathbb{R}^V).$$

For a subset  $S \subseteq V$ , define  $G_S = (V, E(S, S))$ , the graph obtained by keeping only those edges with both endpoints in  $S$ . Let  $L_S$  be its Laplacian. Note that  $L_S$  is the Laplacian of the induced subgraph  $G[S]$ , padded with isolated vertices outside  $S$ .

**Definition 1.1** ( $\varepsilon$ -light set). Fix  $\varepsilon \in (0, 1)$ . A set  $S \subseteq V$  is  $\varepsilon$ -light if

$$\varepsilon L - L_S \succeq 0.$$

The motivating question is as follows.

**Conjecture 1.2** (Linear-size  $\varepsilon$ -light subsets). *There exists a universal constant  $c > 0$  such that for every finite graph  $G = (V, E)$  and every  $\varepsilon \in (0, 1)$ , there exists an  $\varepsilon$ -light subset  $S \subseteq V$  satisfying*

$$|S| \geq c\varepsilon |V|.$$

At the time of writing we do not know whether Conjecture 1.2 is true. The purpose of this note is to (a) formalize the correct operator-theoretic viewpoint (via Schur complements/Kron reduction) and (b) rigorously record why several seemingly natural approaches fail.

### 1.1. A sharp universal upper bound.

**Proposition 1.3** (Any universal constant must satisfy  $c \leq \frac{1}{2}$ ). *If Conjecture 1.2 holds for a universal constant  $c$ , then necessarily  $c \leq \frac{1}{2}$ .*

*Proof.* Let  $G$  be a perfect matching on  $n$  vertices, where  $n$  is even. Then  $L$  is block diagonal with  $n/2$  identical  $2 \times 2$  blocks  $\begin{pmatrix} 1 & -1 \\ -1 & 1 \end{pmatrix}$ . Fix  $\varepsilon \in (0, 1)$  and let  $S \subseteq V$ . If  $S$  contains both endpoints of some matched edge  $e = \{u, v\}$ , then, with  $x := \mathbf{e}_u - \mathbf{e}_v$ , we have  $x^\top L_S x = x^\top L x > 0$  (indeed both equal 4), so  $L_S \preceq \varepsilon L$  fails unless  $\varepsilon \geq 1$ . Therefore, for  $\varepsilon \in (0, 1)$ , every  $\varepsilon$ -light set must contain at most one endpoint from each matched edge, hence  $|S| \leq n/2$ . The requirement  $c\varepsilon n \leq n/2$  for all  $\varepsilon \in (0, 1)$  forces  $c \leq 1/2$ .  $\square$

**1.2. Complete graphs.** Complete graphs show that  $\varepsilon$ -lightness can force  $|S|$  to be *at most* on the order of  $\varepsilon n$ .

**Proposition 1.4** (Complete graphs). *Let  $G = K_n$  and  $\varepsilon \in (0, 1)$ .*

- (i) *If  $|S| \geq 2$ , then  $S$  is  $\varepsilon$ -light if and only if  $|S| \leq \varepsilon n$ .*
- (ii) *If  $|S| \in \{0, 1\}$ , then  $L_S = 0$  and hence  $S$  is  $\varepsilon$ -light for every  $\varepsilon > 0$ .*

*Proof.* If  $|S| \in \{0, 1\}$  then there are no edges with both endpoints in  $S$ , so  $L_S = 0$ .

Assume  $|S| \geq 2$ . The Laplacian of  $K_n$  is  $L = nI - \mathbf{1}\mathbf{1}^\top$ . The padded induced Laplacian  $L_S$  equals the Laplacian of  $K_{|S|}$  on coordinates in  $S$  and is 0 on  $V \setminus S$ . Consider the subspace

$$W := \{x \in \mathbb{R}^V : \text{supp}(x) \subseteq S, \mathbf{1}^\top x = 0\}.$$

On  $W$ , one has  $Lx = nx$  and  $L_S x = |S|x$ . Therefore the maximal generalized eigenvalue of the pencil  $(L_S, L)$  equals  $|S|/n$ , and  $L_S \preceq \varepsilon L$  holds if and only if  $|S|/n \leq \varepsilon$ .  $\square$

## 2. PRELIMINARIES

**2.1. Edge Laplacians and padding.** Fix an arbitrary orientation of edges. For  $e = (u, v)$  write  $b_e := \mathbf{e}_u - \mathbf{e}_v \in \mathbb{R}^V$ . Then the Laplacian decomposes as

$$L = \sum_{e \in E} b_e b_e^\top.$$

If  $F \subseteq E$  is any edge set, we write  $L_F := \sum_{e \in F} b_e b_e^\top$ ; this is the Laplacian of the graph  $(V, F)$ . In particular, for  $S \subseteq V$ ,

$$L_S = L_{E(S, S)} = \sum_{e \in E(S, S)} b_e b_e^\top,$$

and thus  $L_S \preceq L$  for every  $S$ .

**2.2. Kernel, harmonic subspace, and pseudoinverses.** Let  $\ker(L)$  denote the kernel of  $L$ . It is standard that  $\ker(L)$  is the subspace of vectors constant on each connected component of  $G$ . Let  $\mathcal{H} := \ker(L)^\perp$  and let  $\Pi_{\mathcal{H}}$  be the orthogonal projection onto  $\mathcal{H}$ .

We use  $L^+$  for the Moore–Penrose pseudoinverse, and  $L^{1/2}$ ,  $L^{+1/2}$  for the unique PSD square roots of  $L$  and  $L^+$ . We will use the identities

$$LL^+ = L^+L = \Pi_{\mathcal{H}}, \quad L^{1/2}L^{+1/2} = L^{+1/2}L^{1/2} = \Pi_{\mathcal{H}},$$

and the fact that if  $A \succeq 0$  satisfies  $\ker(L) \subseteq \ker(A)$ , then  $\Pi_{\mathcal{H}}A\Pi_{\mathcal{H}} = A$ .

## 3. EQUIVALENT FORMULATIONS AND KERNEL BOOKKEEPING

A technical pitfall in disconnected graphs is that substituting  $x = L^{1/2}y$  yields  $L^{+1/2}x = \Pi_{\mathcal{H}}y$ , not  $y$ . This section records the correct equivalences.

**Lemma 3.1** (Equivalent normalizations). *Let  $G = (V, E)$  have Laplacian  $L$ , let  $S \subseteq V$ , and let  $L_S$  be the Laplacian of  $G_S = (V, E(S, S))$ . Let  $\mathcal{H} = \ker(L)^\perp$  and  $\Pi_{\mathcal{H}}$  be the orthogonal projection onto  $\mathcal{H}$ . The following are equivalent:*

- (a)  $L_S \preceq \varepsilon L$ .
- (b)  $L^{+1/2}L_S L^{+1/2} \preceq \varepsilon \Pi_{\mathcal{H}}$ .
- (c) For all  $x \in \mathcal{H}$ ,  $x^\top L_S x \leq \varepsilon x^\top L x$ .

*Proof.* (a)  $\Rightarrow$  (b) follows by conjugating with  $L^{+1/2}$  and using  $L^{+1/2}L L^{+1/2} = \Pi_{\mathcal{H}}$ . (b)  $\Rightarrow$  (c) is immediate since  $\Pi_{\mathcal{H}}x = x$  for  $x \in \mathcal{H}$ .

(c)  $\Rightarrow$  (a): let  $y \in \mathbb{R}^V$  be arbitrary and decompose  $y = y_{\mathcal{H}} + y_0$  where  $y_{\mathcal{H}} := \Pi_{\mathcal{H}}y \in \mathcal{H}$  and  $y_0 := y - \Pi_{\mathcal{H}}y \in \ker(L)$ . Then  $y^\top L y = y_{\mathcal{H}}^\top L y_{\mathcal{H}}$  and, crucially,  $L_S y_0 = 0$ . Indeed,  $y_0 \in \ker(L)$  means  $y_0$  is constant on each connected component of  $G$ ; every component of  $G_S$  is contained in a component of  $G$  (since  $G_S$  is a subgraph with the same vertex set), hence  $y_0$  is constant on each component of  $G_S$  and thus  $y_0 \in \ker(L_S)$ . Therefore

$$y^\top (\varepsilon L - L_S) y = y_{\mathcal{H}}^\top (\varepsilon L - L_S) y_{\mathcal{H}} \geq 0$$

by hypothesis (c), proving (a).  $\square$

## 4. TWO CLASSICAL “NAIVE” STRATEGIES AND EXPLICIT COUNTEREXAMPLES

This section records two pitfalls that arise in attempts to prove Conjecture 1.2.

**4.1. Edgewise effective resistance does not control  $\varepsilon$ -lightness.** A tempting heuristic is that if every internal edge  $e \in E(S, S)$  has small effective resistance in  $G$ , then  $L_S$  should be dominated by  $\varepsilon L$ . The following explicit computation shows this is false.

**Example 4.1** (Bounded effective resistances but not  $\varepsilon$ -light). Let  $G$  have vertex set  $\{1, 2, 3, 4, 5, 6\}$ . The induced subgraph on  $\{1, 2, 3, 4\}$  is a clique  $K_4$ . Vertices 5 and 6 are adjacent to each of 1, 2, 3, 4, and there is no edge between 5 and 6. Let  $S := \{1, 2, 3, 4\}$ .

Then every internal edge  $e \in E(S, S)$  satisfies  $R_{\text{eff}}^G(e) = \frac{1}{3}$ , yet  $S$  is not  $\varepsilon$ -light for  $\varepsilon = \frac{1}{2}$ :

$$\lambda_{\max}(L^{-1/2}L_S L^{-1/2}) \geq \frac{2}{3} > \frac{1}{2}.$$

*Proof.* Consider the internal edge  $e = \{1, 2\}$  and vector  $b := \mathbf{e}_1 - \mathbf{e}_2$ . Vertices 1 and 2 have degree 5 in  $G$ , and a direct calculation gives  $Lb = 6b$ . Hence  $R_{\text{eff}}^G(e) = b^\top L^+ b = \frac{1}{6} \|b\|^2 = \frac{1}{3}$ .

In the induced subgraph  $G[S] = K_4$ , vertices 1 and 2 have degree 3, and similarly  $L_S b = 4b$ . Therefore

$$\frac{b^\top L_S b}{b^\top L b} = \frac{4\|b\|^2}{6\|b\|^2} = \frac{2}{3},$$

so  $L_S \preceq \varepsilon L$  fails for any  $\varepsilon < 2/3$ , in particular for  $\varepsilon = 1/2$ .  $\square$

**4.2. Vertex-star “linearization” overcharges boundary edges.** Another common strategy is to upper bound  $L_S$  by a sum of star Laplacians  $\sum_{u \in S} L_u$ . This necessarily introduces boundary leakage: edges crossing  $(S, V \setminus S)$  appear linearly even though they do not appear in  $L_S$ .

For each  $u \in V$  let  $L_u$  denote the Laplacian of the star graph consisting of all edges incident to  $u$ . Let  $\partial S$  denote the edge boundary, and let  $L_{\partial S}$  denote the Laplacian of the cut graph  $(V, \partial S)$ . The exact identity

$$\sum_{u \in S} L_u = 2L_S + L_{\partial S} \tag{1}$$

implies

$$L_S = \frac{1}{2} \sum_{u \in S} L_u - \frac{1}{2} L_{\partial S} \preceq \frac{1}{2} \sum_{u \in S} L_u,$$

but the inequality  $L_S \preceq \frac{1}{2} \sum_{u \in S} L_u$  can be far from tight.

**Example 4.2** (Boundary leakage in  $K_{m,m}$ ). Let  $G = K_{m,m}$  with bipartition  $V = A \sqcup B$ . Take  $S = A$ . Then  $L_S = 0$  (hence  $S$  is  $\varepsilon$ -light for every  $\varepsilon > 0$ ), but

$$\frac{1}{2} \sum_{u \in S} L_u = \frac{1}{2} L,$$

so any certification of the form  $\frac{1}{2} \sum_{u \in S} L_u \preceq \varepsilon L$  would force  $\varepsilon \geq 1/2$  despite the fact that  $L_S = 0$ .

*Proof.* Since  $A$  is an independent set,  $E(S, S) = \emptyset$  and  $L_S = 0$ .

Every edge of  $K_{m,m}$  has exactly one endpoint in  $A$ , so  $\sum_{u \in A} L_u$  counts each edge Laplacian  $b_e b_e^\top$  exactly once. Hence  $\sum_{u \in A} L_u = \sum_{e \in E} b_e b_e^\top = L$ , giving the claim.  $\square$

## 5. A NORMALIZED OBSTRUCTION AND A CLEAN CONJUGATION

Define the normalized operator

$$M(S) := L^{+1/2} L_S L^{+1/2}.$$

By Lemma 3.1,  $S$  is  $\varepsilon$ -light if and only if

$$M(S) \preceq \varepsilon \Pi_{\mathcal{H}}.$$

The mapping  $S \mapsto M(S)$  is quadratic in vertex indicators and is the source of the major difficulty.

### 5.1. A projection subtlety for disconnected graphs.

**Proposition 5.1** (Projection bookkeeping for boundary operators). *Let  $S \subseteq V$  and let  $L_{\partial S}$  denote the Laplacian of the cut graph  $(V, \partial S)$ . Let  $y \in \mathbb{R}^V$  be arbitrary and set  $x := L^{1/2} y \in \mathcal{H}$ . Then  $L^{+1/2} x = \Pi_{\mathcal{H}} y$  and*

$$x^\top L^{+1/2} L_{\partial S} L^{+1/2} x = (\Pi_{\mathcal{H}} y)^\top L_{\partial S} (\Pi_{\mathcal{H}} y).$$

Moreover, if  $y_0 := y - \Pi_{\mathcal{H}} y \in \ker(L)$  then  $y_0 \in \ker(L_S) \cap \ker(L_{\partial S})$  and hence

$$L_S(\Pi_{\mathcal{H}} y) = L_S y, \quad L_{\partial S}(\Pi_{\mathcal{H}} y) = L_{\partial S} y.$$

*Proof.* The identity  $L^{+1/2} L^{1/2} = \Pi_{\mathcal{H}}$  gives  $L^{+1/2} x = L^{+1/2} L^{1/2} y = \Pi_{\mathcal{H}} y$ . Then

$$x^\top L^{+1/2} L_{\partial S} L^{+1/2} x = (L^{+1/2} x)^\top L_{\partial S} (L^{+1/2} x) = (\Pi_{\mathcal{H}} y)^\top L_{\partial S} (\Pi_{\mathcal{H}} y).$$

If  $y_0 \in \ker(L)$  then  $y_0$  is constant on each connected component of  $G$  and hence is constant on each connected component of  $G_S$  and of the cut graph  $(V, \partial S)$  (both are subgraphs on  $V$ ). Therefore  $y_0 \in \ker(L_S) \cap \ker(L_{\partial S})$ , so  $L_S y_0 = L_{\partial S} y_0 = 0$  and the final identities follow.  $\square$

**5.2. Star-linearization equivalence.** Define  $A_u := L^{+1/2} L_u L^{+1/2}$ .

**Proposition 5.2** (Star-linearization conjugation). *If  $\sum_{u \in S} A_u \preceq 2\varepsilon \Pi_{\mathcal{H}}$ , then*

$$\frac{1}{2} \sum_{u \in S} L_u \preceq \varepsilon L.$$

*Proof.* Conjugate  $\sum_{u \in S} A_u \preceq 2\varepsilon \Pi_{\mathcal{H}}$  by  $L^{1/2}$ :

$$L^{1/2} \left( \sum_{u \in S} L^{+1/2} L_u L^{+1/2} \right) L^{1/2} \preceq 2\varepsilon L^{1/2} \Pi_{\mathcal{H}} L^{1/2}.$$

Using  $L^{1/2} L^{+1/2} = \Pi_{\mathcal{H}}$  and  $L^{1/2} \Pi_{\mathcal{H}} L^{1/2} = L$ , the left-hand side becomes  $\sum_{u \in S} \Pi_{\mathcal{H}} L_u \Pi_{\mathcal{H}}$ . Since  $\ker(L) \subseteq \ker(L_u)$  for every  $u$  (vectors constant on components of  $G$  are constant on the star at  $u$ ), we have  $\Pi_{\mathcal{H}} L_u \Pi_{\mathcal{H}} = L_u$ . Thus  $\sum_{u \in S} L_u \preceq 2\varepsilon L$ , i.e.,  $\frac{1}{2} \sum_{u \in S} L_u \preceq \varepsilon L$ .  $\square$

## 6. SCHUR COMPLEMENTS AND THE CORRECT VARIATIONAL DUAL

The identity (1) shows that any attempt to replace  $L_S$  by a vertex-linear sum  $\sum_{u \in S} L_u$  must necessarily interact with boundary edges. The correct operator that eliminates the complement is a Schur complement (Kron reduction) [1].

**6.1. Block notation and a key decomposition.** Fix  $S \subseteq V$  and let  $T := V \setminus S$ . Write the Laplacian in block form

$$L = \begin{pmatrix} L_{SS} & L_{ST} \\ L_{TS} & L_{TT} \end{pmatrix}.$$

Let  $L_{G[S]}$  denote the *unpadded* Laplacian of the induced subgraph  $G[S]$  (an  $|S| \times |S|$  matrix). Then

$$L_{SS} = L_{G[S]} + D_{\partial S}, \tag{2}$$

where  $D_{\partial S}$  is the diagonal matrix of boundary degrees on  $S$  (i.e.  $(D_{\partial S})_{uu} = |\{v \in T : \{u, v\} \in E\}|$ ).

### 6.2. Variational characterization of the Schur complement.

**Lemma 6.1** (Schur complement via energy minimization). *Let  $L \succeq 0$  be a graph Laplacian, partitioned into blocks as above. Fix  $x_S \in \mathbb{R}^S$  and consider*

$$\min_{x_T \in \mathbb{R}^T} \begin{pmatrix} x_S \\ x_T \end{pmatrix}^\top \begin{pmatrix} L_{SS} & L_{ST} \\ L_{TS} & L_{TT} \end{pmatrix} \begin{pmatrix} x_S \\ x_T \end{pmatrix}. \quad (3)$$

*Then the minimum is attained, equals*

$$x_S^\top (L_{SS} - L_{ST} L_{TT}^+ L_{TS}) x_S,$$

*and a canonical minimizer is  $x_T^* = -L_{TT}^+ L_{TS} x_S$ .*

*Proof.* Expand the objective:

$$Q(x_T) = x_S^\top L_{SS} x_S + 2x_T^\top L_{TS} x_S + x_T^\top L_{TT} x_T.$$

We first show  $\text{im}(L_{TS}) \subseteq \text{im}(L_{TT})$ , which guarantees solvability of  $L_{TT} x_T = -L_{TS} x_S$ . Let  $z \in \ker(L_{TT})$ . For any  $t \in \mathbb{R}$  consider  $v = (x_S^\top, tz^\top)^\top$ . Since  $L \succeq 0$ ,

$$0 \leq v^\top L v = x_S^\top L_{SS} x_S + 2t z^\top L_{TS} x_S + t^2 z^\top L_{TT} z = x_S^\top L_{SS} x_S + 2t z^\top L_{TS} x_S$$

for all  $t \in \mathbb{R}$ . Hence the linear term vanishes:  $z^\top L_{TS} x_S = 0$  for all  $x_S$ , so  $z^\top L_{TS} = 0$ . Thus  $\ker(L_{TT}) \subseteq \ker(L_{TS})$ , and taking orthogonal complements yields  $\text{im}(L_{TS}) \subseteq \text{im}(L_{TT})$ .

Therefore  $L_{TT} x_T = -L_{TS} x_S$  is solvable. Any such solution is a stationary point of  $Q$ , and since  $L_{TT} \succeq 0$  the function  $Q$  is convex in  $x_T$ , so any stationary point is a minimizer. The pseudoinverse choice  $x_T^* = -L_{TT}^+ L_{TS} x_S$  is a solution because  $L_{TT} L_{TT}^+$  is the orthogonal projection onto  $\text{im}(L_{TT}) \supseteq \text{im}(L_{TS})$ . Substituting  $x_T^*$  into  $Q$  yields the asserted minimum value.  $\square$

### 6.3. Kron reduction and $\varepsilon$ -lightness.

**Definition 6.2** (Kron reduction). The *Kron-reduced Laplacian* (Schur complement) of  $L$  onto  $S$  is

$$L_{\text{Kron}}(S) := L_{SS} - L_{ST} L_{TT}^+ L_{TS}.$$

**Theorem 6.3** ( $\varepsilon$ -lightness via Kron reduction). *Let  $G = (V, E)$  have Laplacian  $L$ . Let  $S \subseteq V$  with complement  $T$  and induced Laplacian  $L_{G[S]}$ . Then  $S$  is  $\varepsilon$ -light if and only if*

$$L_{G[S]} \preceq \varepsilon L_{\text{Kron}}(S). \quad (4)$$

*Equivalently, using (2),*

$$(1 - \varepsilon) L_{G[S]} \preceq \varepsilon (D_{\partial S} - L_{ST} L_{TT}^+ L_{TS}). \quad (5)$$

*Proof.* By Lemma 3.1(c),  $S$  is  $\varepsilon$ -light if and only if for all  $x \in \mathbb{R}^V$ ,

$$x^\top L_S x \leq \varepsilon x^\top L x.$$

Since  $L_S$  only contains edges with both endpoints in  $S$ , we have  $x^\top L_S x = x_S^\top L_{G[S]} x_S$ . Fix  $x_S$  and minimize the right-hand side over  $x_T$  using Lemma 6.1. We obtain

$$x_S^\top L_{G[S]} x_S \leq \varepsilon \cdot \min_{x_T} \begin{pmatrix} x_S \\ x_T \end{pmatrix}^\top L \begin{pmatrix} x_S \\ x_T \end{pmatrix} = \varepsilon x_S^\top L_{\text{Kron}}(S) x_S$$

for all  $x_S \in \mathbb{R}^S$ , which is exactly (4). Rearranging using  $L_{\text{Kron}}(S) = L_{SS} - L_{ST} L_{TT}^+ L_{TS}$  and (2) yields (5).  $\square$

**Remark 6.4** (Interpretation). Equation (5) isolates the true analytic content of  $\varepsilon$ -lightness: internal Laplacian energy in  $G[S]$  must be dominated by the *effective coupling of  $S$  to its complement* captured by  $D_{\partial S} - L_{ST} L_{TT}^+ L_{TS}$ , rather than by naive linear surrogates that overcount boundary edges.

## 7. OBSTRUCTIONS TO CURRENT POLYNOMIAL/INTERLACING APPROACHES

A natural hope is to partition  $V$  into  $r$  parts and show that some part  $S_i$  satisfies  $L_{G[S_i]} \preceq O(1/r) L$ , then take  $r \approx 1/\varepsilon$  and pick the largest part. This resembles matrix paving and suggests using interlacing-polynomial methods [2]. However, two independent obstacles arise: (i) arithmetic leakage in rounding  $r$  from  $\varepsilon$ ; and (ii) the induced-subgraph Laplacian depends quadratically on vertex indicators, violating the independence required for the mixed characteristic polynomial method.

### 7.1. A rounding leak in dyadic reductions.

**Lemma 7.1** (A ceiling-function leak). *Assume hypothetically that for every integer  $r \geq 2$  and every graph  $H$  on  $n$  vertices, there exists  $S \subseteq V(H)$  with*

$$|S| \geq \frac{n}{r} \quad \text{and} \quad L_{H[S]} \preceq \frac{2}{r} L_H.$$

*Then Conjecture 1.2 would hold with the universal constant  $c = \frac{1}{3}$ . Moreover, this dyadic rounding mechanism alone cannot yield  $c = \frac{1}{2}$ .*

*Proof.* Fix  $\varepsilon \in (0, 1)$  and set  $r := \lceil 2/\varepsilon \rceil$ . Then  $2/r \leq \varepsilon$ , so the assumed  $S$  is  $\varepsilon$ -light. Its size satisfies  $|S| \geq n/r$ . Since  $r < \frac{2}{\varepsilon} + 1 = \frac{2+\varepsilon}{\varepsilon}$ , we have

$$\frac{1}{r} > \frac{\varepsilon}{2+\varepsilon} \geq \frac{\varepsilon}{3}.$$

Thus  $|S| \geq (\varepsilon/3)n$ , i.e.  $c = 1/3$ . As  $\varepsilon \rightarrow 1$ , the bound  $\varepsilon/(2+\varepsilon) \rightarrow 1/3$ , showing that this mechanism cannot recover  $c = 1/2$ .  $\square$

**7.2. Quadratic dependence breaks interlacing families.** In the MSS framework [2], one controls the spectrum of a random sum of independent rank-one PSD matrices by analyzing a mixed characteristic polynomial and exploiting real-rootedness and interlacing. For vertex-induced subgraphs, edge indicators are *quadratic* in vertex indicators and highly dependent.

We record an explicit  $K_3$  computation showing that even the *expected* characteristic polynomial can have complex roots.

**Example 7.2** (Complex roots in an expected induced-subgraph characteristic polynomial). Let  $G = K_3$  and assign each vertex independently to one of  $r = 2$  parts. Let  $S$  be the first part and define  $M(S) := L^{+1/2} L_S L^{+1/2}$  acting on the harmonic subspace  $\mathcal{H} = \ker(L)^\perp$  (which has dimension 2 here). Then

$$\mathbb{E}[\det(xI_{\mathcal{H}} - M(S))] = x^2 - \frac{1}{2}x + \frac{1}{8},$$

whose roots are  $\frac{1}{4} \pm i\frac{1}{4}$  and hence are nonreal.

*Proof.* For  $K_3$ , the nonzero Laplacian eigenvalues equal 3, so on  $\mathcal{H}$  one has  $L^{+1/2} = \frac{1}{\sqrt{3}}I_{\mathcal{H}}$ .

There are three cases for the random set  $S$ :

- If  $|S| \leq 1$ , then  $L_S = 0$  and hence  $M(S) = 0$  on  $\mathcal{H}$ , so  $\det(xI_{\mathcal{H}} - M(S)) = x^2$ .
- If  $|S| = 2$ , then  $G[S]$  is a single edge. The unpadding edge Laplacian has nonzero eigenvalue 2 on its (one-dimensional) orthogonal complement of constants; after padding and restricting to  $\mathcal{H}$ , the resulting operator has eigenvalues  $\{2/3, 0\}$  on  $\mathcal{H}$ . Thus  $\det(xI_{\mathcal{H}} - M(S)) = x(x - 2/3) = x^2 - \frac{2}{3}x$ .
- If  $|S| = 3$ , then  $L_S = L$  and  $M(S) = L^{+1/2} L L^{+1/2} = \Pi_{\mathcal{H}}$ , which is  $I_{\mathcal{H}}$  on  $\mathcal{H}$ . Thus  $\det(xI_{\mathcal{H}} - M(S)) = (x - 1)^2 = x^2 - 2x + 1$ .

Under independent 2-coloring,  $\mathbb{P}(|S| = 3) = 1/8$ ,  $\mathbb{P}(|S| = 2) = 3/8$ , and  $\mathbb{P}(|S| \leq 1) = 4/8$ . Therefore

$$\mathbb{E}[\det(xI_{\mathcal{H}} - M(S))] = \frac{1}{8}(x^2 - 2x + 1) + \frac{3}{8}\left(x^2 - \frac{2}{3}x\right) + \frac{4}{8}x^2 = x^2 - \frac{1}{2}x + \frac{1}{8}.$$

Its discriminant is  $(-1/2)^2 - 4 \cdot (1/8) = 1/4 - 1/2 = -1/4 < 0$ , so the roots are nonreal.  $\square$

**Remark 7.3** (Consequences for interlacing). If a family of real-rooted polynomials has a common interlacing, then every convex combination is real-rooted. Therefore Example 7.2 certifies that the characteristic polynomials arising from vertex-induced parts in this simplest instance do not admit the interlacing structure required by the MSS method. This does not preclude other polynomial techniques, but it rules out a direct import of [2].

## 8. DISCUSSION AND OPEN DIRECTIONS

The Schur complement characterization in Theorem 6.3 suggests that Conjecture 1.2 is fundamentally a statement about finding a large terminal set  $S$  for which the internal Laplacian is dominated by the Kron-reduced Laplacian of the full graph onto  $S$ . Any successful proof must simultaneously:

- control the quadratic dependence of internal edges on vertex selection, and
- exploit the effective boundary operator  $D_{\partial S} - L_{ST} L_{TT}^+ L_{TS}$  rather than a lossy linear surrogate.

We hope the framework and counterexamples here serve as a clean starting point.

## REFERENCES

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