

# Quantum Multi-string Matching

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# Presentation Overview

## 1 Problem Definition

## 2 Classical Algorithms

- Existing Algorithms
- Polynomial Matching
- Polynomial Multiplication

## 3 Quantum Algorithms

- Polynomial Multiplication with QFT

# String Matching

- Text  $S[0 \dots n - 1]$
- Pattern (or key)  $P[0 \dots m - 1]$
- Strings from alphabet  $\Sigma$
- Find set of indices  $i$  such that  $S[i..i + m] = P$

Example:

- $T = \text{"GTAT GATC TC"}$  (ignore spaces)
- $P_1 = \text{"ATCT"}$
- $P_2 = \text{"TGAT"}$
- $P_3 = \text{"ACCC"}$
- $P_1$  matches at 5,  $P_2$  matches at 3,  $P_3$  has no matches
- Multi-string matching: search  $P_1, P_2, \dots$  in  $S$  simultaneously

# Classical Algorithms

- Brute Force  $\rightarrow O(mn)$
- Boyer-Moore  $\rightarrow O(n + m)$ ,  $O(mn)$  worst case
- Knuth-Morris-Pratt  $\rightarrow O(n + m)$  worst case
- Suffix Tree/Array  $\rightarrow O(m)/O(n \log n) + \text{preproc}$
- Karp-Rabin - rolling hash  $\rightarrow O(n + m)$  expected

G	T	A	T	G	A	T	C	T	C
hash-value 84									
	hash-value 194								
		hash-value 6							
			hash-value 18						
				hash-value 95					

- **Polynomial Matching**  $\rightarrow O(n \log n)$

# Polynomial Matching

Calculate a fingerprint for  $P$  and every  $m$  character sequence in  $S$ ;  
**matching fingerprints** suggest pattern match!

$$A \mapsto -3$$

$$C \mapsto 5$$

$$G \mapsto -7$$

$$T \mapsto 11$$

$$\text{"ATCT"} \longrightarrow Tx^3 + Cx^2 + Tx + A$$

$$\longrightarrow 11x^3 + 5x^2 + 11x - 3$$

$$\text{"GTAT GATC TC"} \longrightarrow -7x^{15} + 11x^{14} - 3x^{13} + \dots + 11x^7 + 5x^6$$

Note that  $P$  encoding is reversed (similar to convolution)

# Polynomial Matching

Given polynomials

$$P(x) = \sum_{i=0}^m a_i x^i, S(x) = \sum_{j=0}^n b_j x^j$$

then

$$R(x) = \sum_{k=0}^{m+n} c_k x^k$$

where

$$c_k = \sum_{i=0}^k a_i b_{k-i}, \quad \text{for } 0 \leq k \leq m+n,$$

with the convention that  $a_i = 0$  for  $i > m$  and  $b_j = 0$  for  $j > n$ .

# Polynomial Matching

Coefficients of  $R(x)$  correspond to “dot products” between substrings.

$$P(x) = 11x^3 + 5x^2 + 11x - 3$$

$$S(x) = -7x^{15} + 11x^{14} - 3x^{13} + \dots + 11x^7 + 5x^6$$

$$R(x) = \dots + 71x^8 + 0x^9 + 276x^{10} + 98x^{11} - 4x^{12} + \dots$$

Degrees map to index:

15 yields index 0

14 yields index 1

...

11 yields index 4

10 yields index 5

9 yields index 6

# Polynomial Matching

Coefficients of  $R(x)$  correspond to “dot products” between substrings.

$$P(x) = 11x^3 + 5x^2 + 11x - 3$$

$$S(x) = -7x^{15} + 11x^{14} - 3x^{13} + \dots + 11x^7 + 5x^6$$

$$R(x) = \dots + 71x^8 + 0x^9 + 276x^{10} + 98x^{11} - 4x^{12} + \dots$$

Notice that  $\|P\|^2 = 11^2 + 5^2 + 11^2 + (-3)^2 = 276$

We get exact fingerprint match for single patterns. Let's extend to multiple patterns...



# Polynomial Matching

Add patterns together:

$$P(x) = P_1(x) + P_2(x)$$

$$S(x) = -7x^{15} + 11x^{14} - 3x^{13} + \dots + 11x^7 + 5x^6$$

$$R(x) = \dots + 94x^8 + 240x^9 + 272x^{10} + 64x^{11} + 296x^{12} - 60x^{13} \dots$$

No more exact matches, but higher values = likelier index.

# Polynomial Multiplication

Multiple algorithms:

- Naïve polynomial multiplication  $\rightarrow O(n^2)$
- Karatsuba Algorithm  $\rightarrow O(n^{\log_3 2}) = O(n^{1.59})$
- FFT  $\rightarrow O(n \log n)$

## Theorem (Lagrange Interpolation)

*For any  $n$  points  $(x_i, y_i) \in \mathbb{R}^2$  with no two  $x_i$  the same, there exists a unique polynomial  $A(x)$  of degree at most  $n - 1$  that interpolates these points.*

Given polynomials  $A(x)$  and  $B(x)$  of degree  $n$ , we can represent them as

$$A : \{(x_0, A(x_0)), (x_1, A(x_1)), \dots, (x_{n-1}, A(x_{n-1}))\}$$

$$B : \{(x_0, B(x_0)), (x_1, B(x_1)), \dots, (x_{n-1}, B(x_{n-1}))\}$$

Then their product  $C(x) = A(x)B(x)$  is

$$C : \{(x_0, A(x_0)B(x_0)), (x_1, A(x_1)B(x_1)), \dots, (x_n, A(x_{n-1})B(x_{n-1}))\}$$

# Polynomial Multiplication

Algorithm idea:

- ➊ Evaluate  $A(x)$  and  $B(x)$  at  $n$  points
  - $O(n)$  multiplications per point
  - $O(n)$  points
  - $O(n^2)$  overall, no speedup
- ➋ Multiply element-wise
- ➌ Interpolate to find  $C(x)$  coefficients

But this works for any  $n$  inputs. What if we try the  $n^{\text{th}}$  roots of unity?

$$\hat{a}_k = A(\omega_n^k) = \sum_{j=0}^n a_j e^{\frac{2\pi i j k}{n}}$$

for  $0 \leq k \leq n - 1$ . This is **DFT**!

# Polynomial Multiplication with FFT

- 1 Apply FFT to the coefficients of  $A(x)$  and  $B(x)$

With  $\vec{a} = [a_0 \ a_1 \ \dots \ a_{n-1}]^T$  and  $\vec{b} = [b_0 \ b_1 \ \dots \ b_{n-1}]^T$ , then

$$\vec{\alpha} = \text{FFT}(\vec{a}) \quad O(n \log n)$$

$$\vec{\beta} = \text{FFT}(\vec{b}) \quad O(n \log n)$$

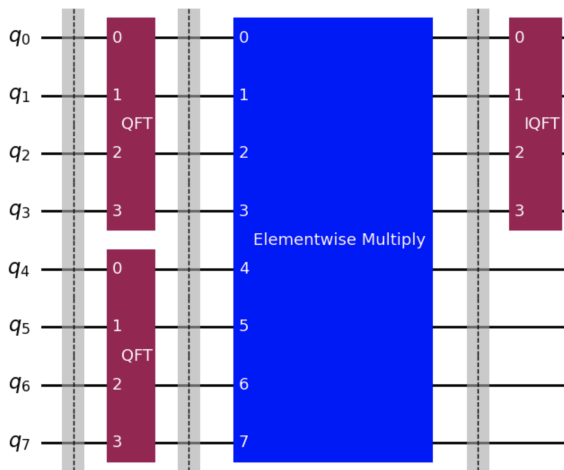
- 2 Perform the Hadamard (element-wise) product of coefficients

$$\begin{aligned} \vec{\gamma} &= \vec{\alpha} \odot \vec{\beta} \\ &= [\alpha_0 \beta_0 \ \alpha_1 \beta_1 \ \dots \ \alpha_{n-1} \beta_{n-1}]^T \end{aligned} \quad O(n)$$

- 3 Apply Inverse FFT

$$\vec{c} = \text{IFFT}(\vec{\gamma}) \quad O(n \log n)$$

# Polynomial Multiplication with QFT



Polynomial multiplication for degree up to  $n = 2^4 - 1 = 15$

# Elementwise Multiply

Is there such a unitary operator? Define

$$|\alpha\rangle = \alpha_0|00\rangle + \alpha_1|01\rangle + \alpha_2|10\rangle + \alpha_3|11\rangle$$

$$|\beta\rangle = \beta_0|00\rangle + \beta_1|01\rangle + \beta_2|10\rangle + \beta_3|11\rangle$$

$$|\alpha \odot \beta\rangle = C(\alpha_0\beta_0|00\rangle + \alpha_1\beta_1|01\rangle + \alpha_2\beta_2|10\rangle + \alpha_3\beta_3|11\rangle)$$

for some normalization constant  $C$

$$U|x\rangle|y\rangle = U|x\rangle|x \odot y\rangle$$

If  $|y\rangle = \frac{1}{2}(|00\rangle + |01\rangle + |10\rangle + |11\rangle) \longrightarrow$  violates **no-cloning theorem**

# Elementwise Multiply

Allow probabilistic operator:

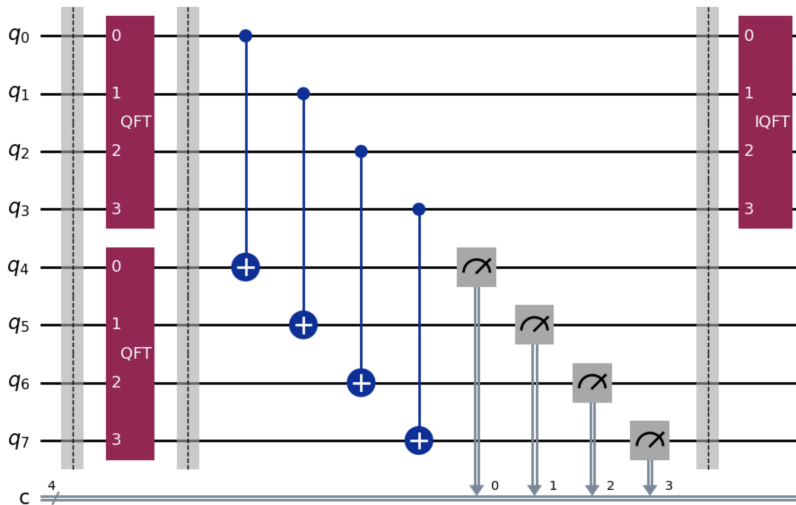
$$\begin{aligned} |\alpha\rangle|\beta\rangle &= \alpha_0\beta_0|0000\rangle + \alpha_0\beta_1|0001\rangle + \dots \\ &\quad \alpha_1\beta_0|0100\rangle + \alpha_1\beta_1|0101\rangle + \dots \\ &\quad \dots + \alpha_2\beta_2|1010\rangle + \dots \\ &\quad \dots + \alpha_3\beta_3|1111\rangle + \dots \end{aligned}$$

Want when **first two bits** = **last two bits**.

Use CNOT to test for equality.



# Elementwise Multiply



# Elementwise Multiply

$$\begin{aligned} |\alpha\rangle|\beta\rangle &= \alpha_0\beta_0|0000\rangle + \alpha_0\beta_1|0001\rangle + \dots \\ &\quad \alpha_1\beta_0|0100\rangle + \alpha_1\beta_1|0101\rangle + \dots \\ &\quad \dots + \alpha_2\beta_2|1010\rangle + \dots \\ &\quad \dots + \alpha_3\beta_3|1111\rangle + \dots \end{aligned}$$

$$\begin{aligned} \text{CNOT}_{0,2}\text{CNOT}_{1,3} |\alpha\rangle|\beta\rangle &= \alpha_0\beta_0|0000\rangle + \alpha_0\beta_1|0001\rangle + \dots \\ &\quad \alpha_1\beta_0|0101\rangle + \alpha_1\beta_1|0100\rangle + \dots \\ &\quad \dots + \alpha_2\beta_2|1000\rangle + \dots \\ &\quad \dots + \alpha_3\beta_3|1100\rangle + \dots \end{aligned}$$

Measure the last two qubits, success when  $|00\rangle \rightarrow$  elementwise multiply in first two qubits!

# Algorithmic Analysis

Let  $n$  be the degree of the polynomial produced by multiplying  $A(x)$  and  $B(x)$ .

Initial QFT:  $O(n^2)$  gates.

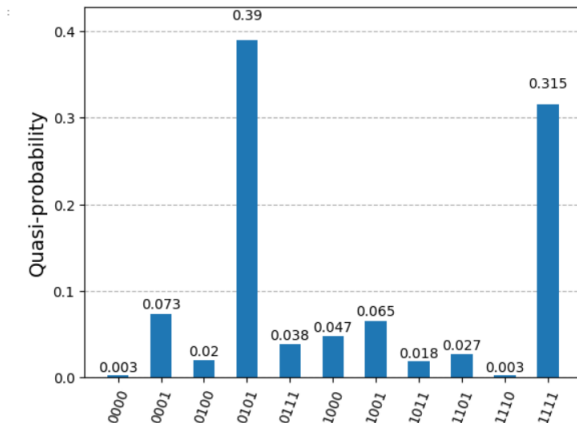
Element-wise multiplication:  $O(n)$  CNOT gates, but requires an expected  $O(2^n)$  measurements for a result where the last  $n$  qubits are all zero.

Final IQFT:  $O(n^2)$  gates.

So the total cost is  $O(n^2)$  gates, but expected  $O(2^n)$  shots.

# Qiskit Results

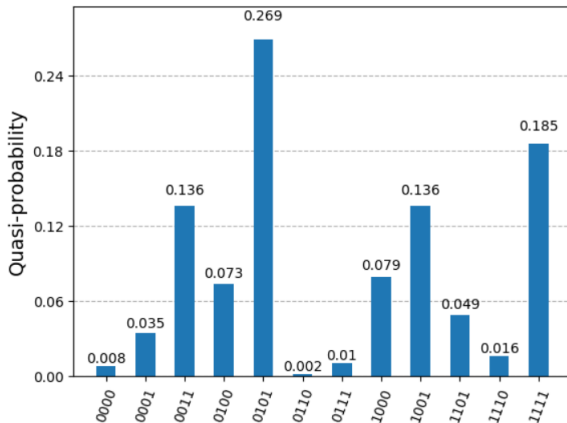
Text: "GTAT GATC TC" Key: "ATCT"



Notice that the highest quasi-probability is for state  $0101_2 = 5$ !

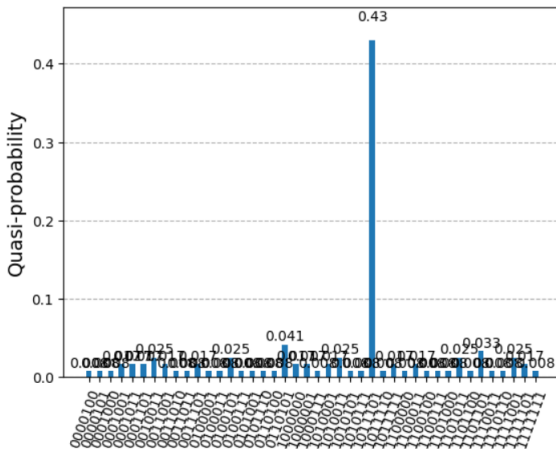
# Multi-String Matching

- Text: "GTAT GATC TC"
- Key 1: "ATCT"
- Key 2: "TGAT"
- Key 3: "ACCC"



# More Qiskit Results

- Text: 90 randomly generated characters  $\in \{A, T, C, G\}$
- Key: 52 characters



Add some slides describing number of shots, probability distribution, runtime.

Current code uses statevector, not sure how many shots would be good. Need min  $2^n$  just to get elementwise. TODO - try actual simulator

Try some larger texts

# References



John Smith (2022)

Publication title

*Journal Name* 12(3), 45 – 678.



Annabelle Kennedy (2023)

Publication title

*Journal Name* 12(3), 45 – 678.