End Term Examination.

I VI PARIS

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tor an XYZ company instead of stoxing their employees to ples in a heap file, with a choose to store it with an index on the empname are a choose to store it with an index on the emphase and index is possible only when clustured index is possible only when clustured index is possible only when there exert employee will have a unique name exert employee will have a unique name. The this is ensured, the toples will be organted this is ensured, the toples will be organted the toples will be organted.

Instead, they can use empid ous a considering everyone already has a unique considering everyone which is empidical according to them which is empidical according to their empidical according to the empidical according to the

But using both Rields Cempnoume and empty)
as indexes may not possible. This operation
as indexes may not possible to have one
may fail. But it is possible to have one
clustered index of cound, one mon clustered
index. Thus XY2 company can store their
employees data.

K. V. L. Amyu)ha, (abcsoy8

2) ppi is important in representing information
in pams because it is used to describe
external and logical schemas.

pmi is used to access and update data;
it is not important for representing data.

True or interleave the actions of different transactions instead of executing transactions one after the other. To improve the execution time of users queries, transactions from these users are interleaved. A orms is typically shared among many users. By interleaving queries, users do not have to wait for other user's transactions to complete fully before their own transaction begins.

find appear a your more continued.

THE WAY

without, intexteaving, if A begins to transaction that will take to seconds to complete, and user B wants to begin ou transaction, user B have to wait an additional coseconds for user A's transaction to complete before the dutabase would begin processing user B's request.

To avoid this waiting time and to imposs execution time, transactions are interteaved.

How the since to the data base, a user must guarantee that with a user must grantee that his or her must guarantee that withdraw 1 transaction accurately models the amount or person removes prom this or her account on accurately models the account or person removes prom this or her account

b) (onsistency in doutabase systems refers to the requirement that any igiven dutabase transaction must change affected data only in allowed ways. A DBMS must gravantee that transactions are executed fully and independently of other transactions: An essential property of a pams is that a transaction should execute atomically, or as if it is the only transaction running. Also, transactions will either complete fully, or will be aborted and the database returned to its initial starte. This ensures that the dartabase remains consistant.

S) Yes, we can determine the key of relation with the help of instance, For relation with on a one to many relation example, in a one to many relation we can consider the column attribute we can consider the column attribute with unique values as a primary key.

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7) Let the two suppliers be R1, R2 000 12000
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Relabional algebra query:

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P(Rz, carbolog)

MRI. RidGRI. Pid = Rz. Pid nRI. Sid! = R, Sidlen

Sal query:

SELECT C. Sto

FROM carbalog C

WHERE Exists (SELECT CI. SID

FROM carbalog C,

WHERE CI. PID AND CI. SID != C. SID
```

Because, this relational algebra statement does not return anything because of the sequence of projection operators, once the sid is projected, it is the only field in the set. Therefore, projecting on same will not return anything.

```
K. U. 2. Amount (ab ceous)

9) The following view query on Emp Country by updated automatically by updating Emp;

CREATE VIEW senior Emp (cid maume, age, Salaby)

AS SELECT F. eid, F. Ename, E. age, E. Salaby)

FROM Emp E where E. age > 50.
```