# cmath.sty: An Infrastructure for building Inline Content Math in STEX\*

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#### Abstract

The cmath package is a central part of the STEX collection, a version of TEX/LATEX that allows to markup TEX/LATEX documents semantically without leaving the document format, essentially turning TEX/LATEX into a document format for mathematical knowledge management (MKM).

This package supplies an infrastructure that allows to build content math expresions (strict content MathML or OpenMath objects) in the text. This is needed whenever the head symbols of expressions are variables and can thus not be treated via the \symdef mechanism in STFX.

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### 1 Introduction

STEX allows to build content math expressions via the \symdef mechanism [KGA14] if their heads are constants. For instance, if we have defined \symdef{lt}[2] {#1<#2} in the module relation1, then an invocation of \lt3a will be transformed to

```
<OMA>
  <OMS cd="relation1" name="lt"/>
  <OMI>3</OMI>
  <OMV name="a"/>
  </OMA>
```

If the head of the expression (i.e. the function symbol in this case) is a variable, then we cannot resort to a  $\symdef$ , since that would define the functional equivalent of a logical constant. Sometimes, LATEXML can figure out that when we write f(a,b) that f is a function (especially, if we declare them to be via the functions= key in the dominating statement environment [Koh14]). But sometimes, we want to be explicit, especially for n-ary functions and in the presence of elided elements in argument sequences. A related problem is markup for complex variable names, such as  $x_{\text{left}}$  or  $ST^*$ .

The cmath package supplies the LATEX bindings that allow us to achieve this.

### 2 The User Interface

### 2.1 Variable Names

In mathematics we often use complex variable names like x',  $g_n$ ,  $f^1$ ,  $\widetilde{\phi}_i^j$  or even foo; for presentation-oriented LATEX, this is not a problem, but if we want to generate content markup, we must show explicitly that those are complex identifiers (otherwise the variable name foo might be mistaken for the product  $f \cdot o \cdot o$ ). In careful mathematical typesetting, sin is distinguished from sin, but we cannot rely on this effect for variable names.

\vname

\vname

\vname identifies a token sequence as a name, and allows the user to provide an ASCII (XML-compatible) identifier for it. The optional argument is the identifier, and the second one the LaTeX representation. The identifier can also be used with \vnref for referencing. So, if we have used \vnname[xi]{x\_i}, then we can later use \vnref{xi} as a short name for \vname{x\_i}. Note that in output formats that are capable of generating structure sharing, \vnref{xi} would be represented as a cross-reference.

\livar

EdN:1

Since indexed variable names make a significant special case of complex identifiers, we provides the macros \livar that allows to mark up variables with lower indices. If \livar is given an optional first argument, this is taken as a name. Thus \livar[foo]{x}1 is "short" for \vname[foo]{x\_1}. The macros \livar,

<sup>\</sup>livar

 $<sup>^{1}\</sup>mathrm{EDNote}$ : DG: Do we know whether using the same name in two vname invocations, would refer to two instances of the same variable? Presumably so, since the names are the same? We should make this explicit in the text. A different variable would e.g. have a name "xi2", but the same body

| \nappa{f}{a_1,a_2,a_3}                               | $f(a_1, a_2, a_3)$                            |
|--|---|
| $\label{landan} $$ \operatorname{f}_{a_1}_{a_n}$$    | $f(a_1,\ldots,a_n)$                           |
| \symdef{eph}[1]{e_{#1}^{\varphi(#1)}}\nappf{g}\eph14 | $g(e_1^{\varphi(1)},\ldots,e_4^{\varphi(4)})$ |
| \nappli{f}a1n  | $f(a_1,\ldots,a_n)$                           |

Figure 1: Application Macros

\ulivar \primvar \pprimvar serve the analogous purpose for variables with upper indices, and \ulivar for upper and lower indices. Finally, \primvar and \pprimvar do the same for variables with primes and double primes (triple primes are bad style).

### 2.2 Applications

\nappa

To construct a content math application of the form  $f(a_1, ..., a_n)$  with concrete arrugments  $a_i$  (i.e. without elisions), then we can use the \nappa maro. If we have elisions in the arguments, then we have to interpret the arguments as a sequence of argument constructors applied to the respective positional indexes. We can mark up this situation with the \nappf macro: \nappf{\langle fun\rangle}{\langle const\rangle}{\langle (first\rangle)}{\langle (last\rangle)} \text{ where \$\langle const\rangle}\$ is a macro for the constructor is presented as \$\langle fun\rangle (\langle const\rangle \langle first\rangle, ..., \langle const\rangle \langle last\rangle)\$; see Figure 1 for a concrete example, and Figure 1.<sup>2</sup>

 $\n$ 

\nappe

EdN:2

For a simple elision in the argument list, we can use nappe macro:  $\texttt{nappe}\{\langle fun\rangle\}\{\langle firstarg\rangle\}\{\langle lastarg\rangle\}$  will be formatted as  $\langle fun\rangle(\langle firstarg\rangle, \ldots, \langle lastarg\rangle)$ . Note that this is quite unsemantic (we have to guess the sequence), so the use of nappe is discouraged.

### 2.3 Binders

### 3 Limitations

In this section we document known limitations. If you want to help alleviate them, please feel free to contact the package author. Some of them are currently discussed in the STEX TRAC [sTeX].

1. none reported yet

### 4 The Implementation

The cmath package generates to files: the LATEXML bindings (all the code between <code>\\*package</code>) and <code>\/package</code>) and the LATEXML bindings (between <code>\\*Itxml</code>) and <code>\/Itxml</code>). We keep the corresponding code fragments together, since the documentation applies to both of them and to prevent them from getting out of sync.

For LATEXML, we initialize the package inclusions.

1 (\*ltxml)

 $<sup>^2{\</sup>rm EDNoTE}$ : MK@MK: we need a meta-cd cmath with the respective notation definition here. It is very frustrating that we cannot even really write down the axiomatization of

```
EdN:3
```

Example 1: Application Macros

```
2 # -*- CPERL -*-
3 package LaTeXML::Package::Pool;
4 use strict;
5 use LaTeXML::Package;
6 (/ltxml)
```

### 4.1 Package Options

We declare some switches which will modify the behavior according to the package options. Generally, an option xxx will just set the appropriate switches to true (otherwise they stay false).<sup>3</sup>

```
7 \langle *package \rangle
8 \ProcessOptions
9 \langle /package \rangle
```

### 4.2 Variable Names

Avname a name macro; the first optional argument is an identifier  $\langle id \rangle$ , this is standard for IATEX, but for LATEXML, we want to generate attributes xml:id="cvar. $\langle id \rangle$ " and name=" $\langle id \rangle$ ". However, if no id was given in we default them to xml:id="cvar. $\langle count \rangle$ " and name="name.cvar. $\langle count \rangle$ ".

10  $\langle *package \rangle$ 11 \newcommand\vname[2][]{#2%

<sup>12 \</sup>def\@opt{#1}%
13 \ifx\@opt\@empty\else\expandafter\gdef\csname MOD@name@#1\endcsname{#2}\fi}

 $<sup>^3\</sup>mathrm{EdNote}$ : we have no options at the moment

```
14 (/package)
        15 (*ltxml)
        16 \; \text{\# return:} \; \text{unique ID for variable}
        17 sub cvar_id {
            my ($id) = @_;
            $id = ToString($id);
            if (!$id) {
               $id=LookupValue('cvar_id') || 0;
        21
               AssignValue('cvar_id', $id + 1, 'global'); }
        22
             "cvar.$id"; }#$
        23
        24 DefConstructor('\vname[]{}',
            "<ltx:XMWrap role='ID' xml:id='&cvar_id(#1)'>#2</ltx:XMWrap>",
            requireMath=>1);
        27 (/ltxml)
\vnref
        28 (*package)
        29 \def\vnref#1{\csname MOD@name@#1\endcsname}
        30 (/package)
        31 \langle *ltxml \rangle
        32 # \vnref{<reference>}
        33 DefMacro('\vnref{}','\@XMRef{cvar.#1}');
        35 (/ltxml)
        4
\uivar constructors for variables.
        37 \newcommand\primvar[2][]{\vname[#1]{#2^\prime}}
        38 \mbox{ $$\mbox{$1$} {\#2^{\rm prime}}} \label{eq:prime} $$
        39 \newcommand\uivar[3][]{\vname[#1]{{#2}^{#3}}}
        40 \newcommand\livar[3][]{\vname[#1]{{#2}_{#3}}}
        41 \newcommand\ulivar[4][]{\vname[#1]{{#2}^{#3}_{#4}}}
        42 (/package)
        43 (*ltxml)
        44 # variants for declaring variables
        45 DefMacro('\uivar[]{}{}',
                                         '\vname[#1]{{#2}^{#3}}');
        46 DefMacro('\livar[]{}{}',
                                         '\ne [#1] {{#2}_{#3}}');
        47 DefMacro('\ulivar[]{}{}\', '\vname[#1]{{#2}^{#3}_{#4}}');
        48 DefMacro('\primvar[]{}',
                                         '\vname[#1]{#2^\prime}');
        49 DefMacro('\pprimvar[]{}',
                                         '\vname[#1]{#2^{\prime\prime}}');
        51 (/ltxml)
```

### 4.3 Applications

 $\ne$ 

EdN:4

 $<sup>^4\</sup>mathrm{EdNote}$ : the following macros are just ideas, they need to be implemented and documented

```
52 (*package)
53 \mbox{newcommand} \mbox{nappa[2]{#1(#2)}}
54 \newcommand \nappe [3] {\nappa {\#1} {\#2, \ldots, \#3}} \\
55 \newcommand\nappf[4]{\nappe{#1}{#2{#3}}{#2{#4}}}
56 \newcommand\nappli[4]{\nappe{#1}{#2_{#3}}{#2_{#4}}}
57 \newcommand\nappui [4] {\nappe{#1}{#2^{#3}}{#2^{#4}}}
58 (/package)
59 (*ltxml)
60 # \nappa{<function>}{<(const)(,\1)*>}
61 # @#1(#2)
62 DefConstructor('\nappa{}{}',
    "<ltx:XMApp>"
      ."<ltx:XMTok meaning='#1' />"
64
      ."<ltx:XMArg>#2</ltx:XMArg>"
65
    ."</ltx:XMApp>");
66
67
68 # \@napp@seq{<function>}{start <const>}{end <const>}
69 # @#1(@sequence(#2,sequencefromto,#3))
70 DefConstructor('\@napp@seq{}{}{}',
71
   "<ltx:XMApp>"
   ."<ltx:XMTok meaning='#1' />"
72
   ."<ltx:XMArg>"
73
      ."<ltx:XMApp>"
74
        ."<ltx:XMTok meaning='sequence' />"
75
        ."<ltx:XMArg>#2</ltx:XMArg>"
76
        ."<ltx:XMArg><ltx:XMTok meaning='sequencefromto' /></ltx:XMArg>"
77
        ."<ltx:XMArg>#3</ltx:XMArg>"
78
      ."</ltx:XMApp>"
79
    ."</ltx:XMArg>"
80
    ."</ltx:XMApp>");
81
82
83 DefMacro('\nappe{}{}\',
                              '\@napp@seq{#1}{#2}{#3}');
84 DefMacro('\nappf{}{}{}',
                              '\@napp@seq{#1}{#2{#3}}{#2{#4}}');
86 \ \tt DefMacro('\nappui{}{}{}', '\napp@seq{#1}{#2^{#3}}{#2^{#4}}');
88 (/ltxml)
```

### 4.4 Binders

### 4.5 Finale

Finally, we need to terminate the file with a success mark for perl. 89 (ltxml)1;

## Index

Numbers written in italic refer to the page where the corresponding entry is described; numbers underlined refer to the code line of the definition; numbers in roman refer to the code lines where the entry is used.

LATEXML, 3-5 XML, 3

## References

- [KGA14] Michael Kohlhase, Deyan Ginev, and Rares Ambrus. modules.sty: Semantic Macros and Module Scoping in sTeX. Tech. rep. Comprehensive TEX Archive Network (CTAN), 2014. URL: http://www.ctan.org/get/macros/latex/contrib/stex/modules/modules.pdf.
- [Koh14] Michael Kohlhase. omtext: Semantic Markup for Mathematical Text Fragments in LATEX. Tech. rep. Comprehensive TEX Archive Network (CTAN), 2014. URL: http://www.ctan.org/tex-archive/macros/latex/contrib/stex/omtext/omtext.pdf.
- [sTeX] Semantic Markup for LATEX. Project Homepage. URL: http://trac.kwarc.info/sTeX/ (visited on 02/22/2011).