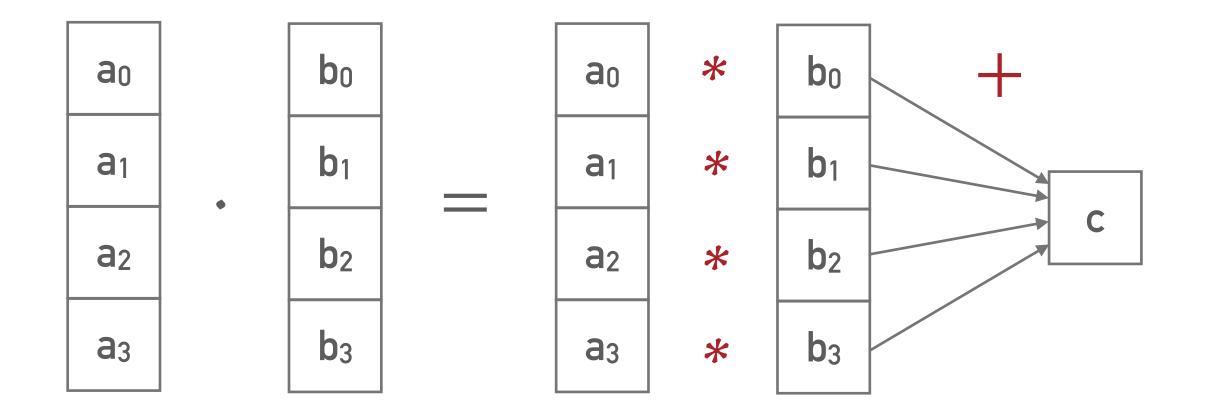


# ME 471/571

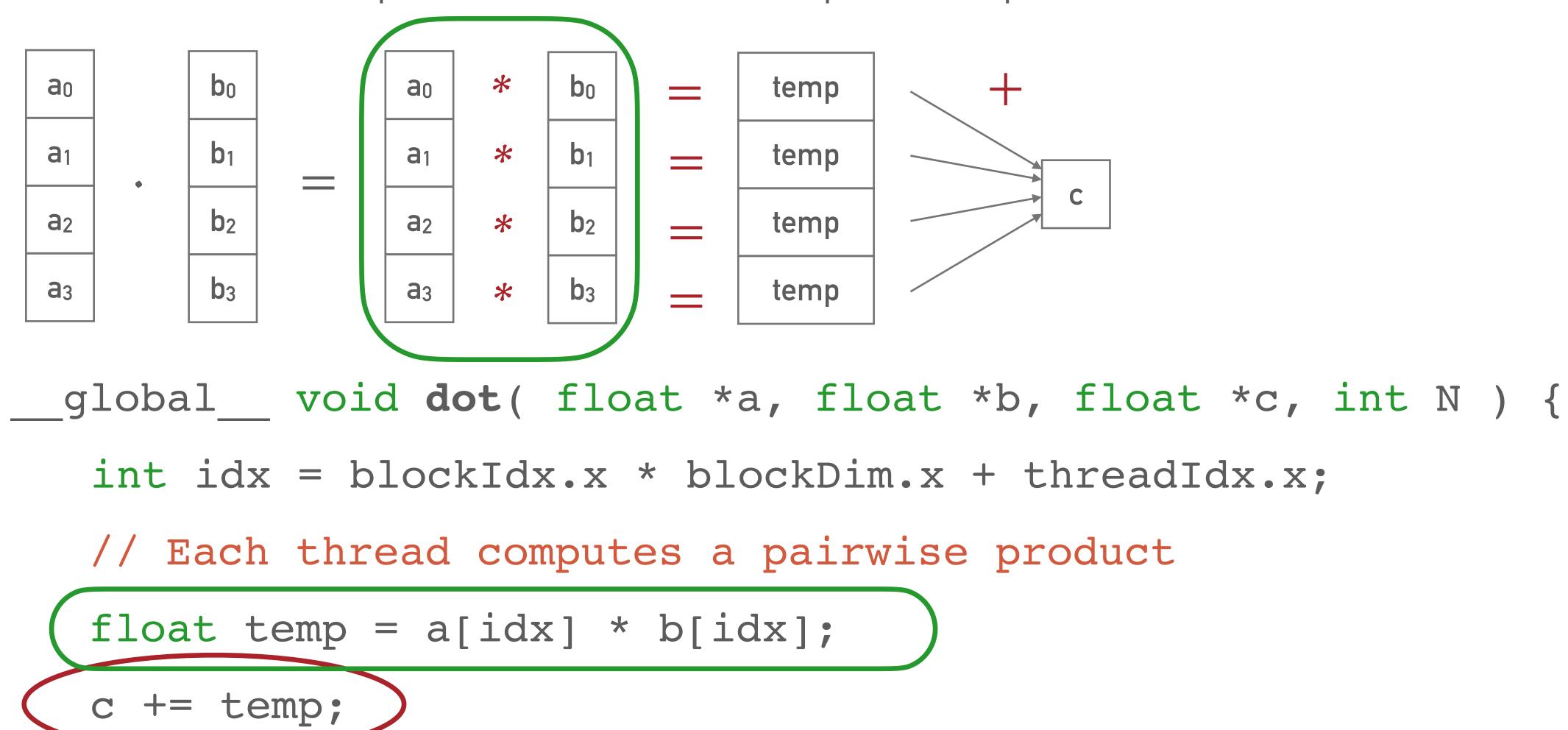
Week 13 - Shared memory and thread synchronization

Consider a computation of dot product:



```
void dotProduct(float *A, float *B, float *C, const int N)
for (int idx = 0; idx < N; idx++)
   *C += A[idx] * B[idx];</pre>
```

Parallel threads have no problem with the multiplication part:



# DATA RACE CONDITION

#### thread 0

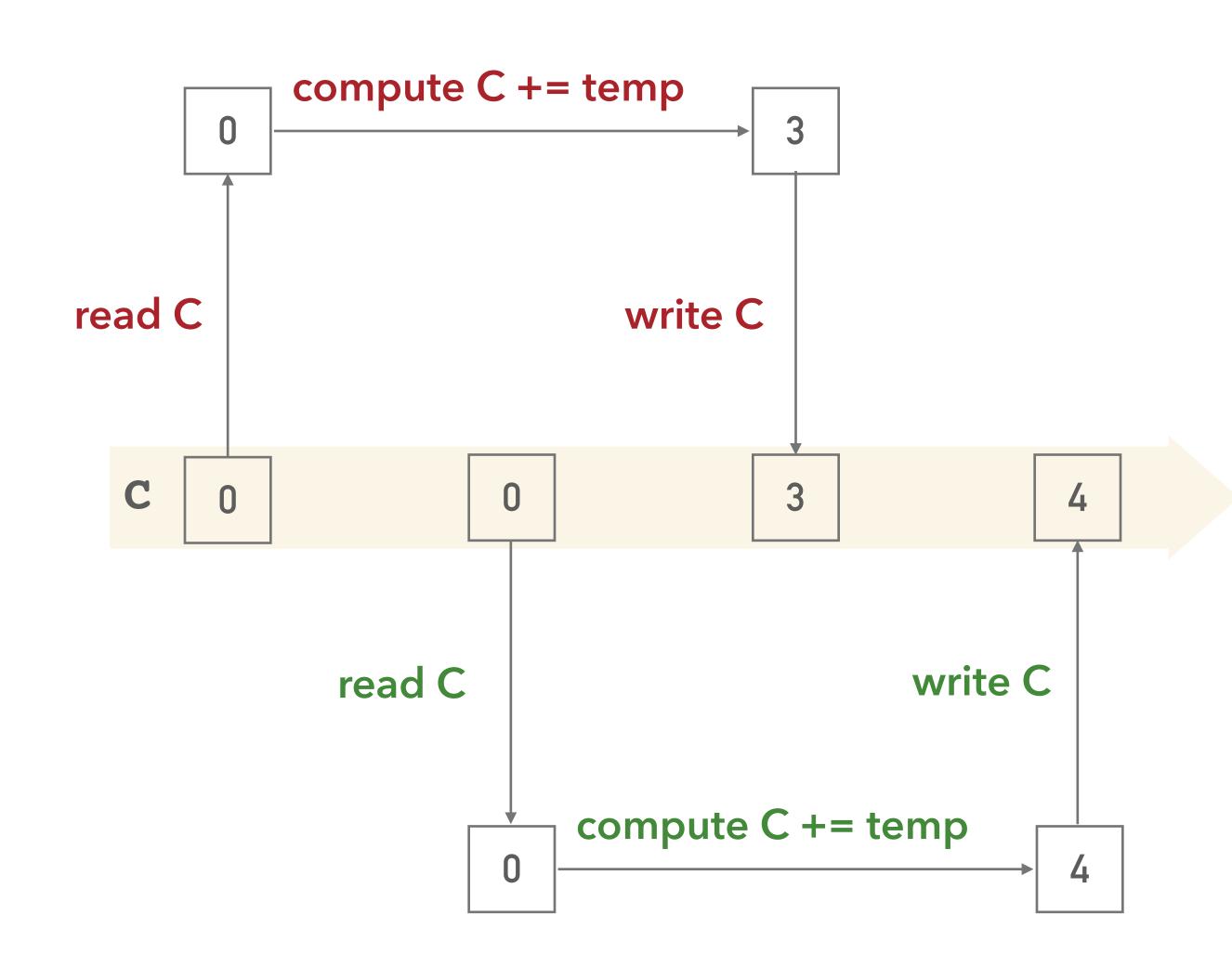
temp = 3

### global memory

C = 0

#### thread 1

temp = 4



We cannot guarantee when each thread will read the memory, which leads to the data race condition.

# DATA RACE CONDITION

### compute C += temp

#### thread 0

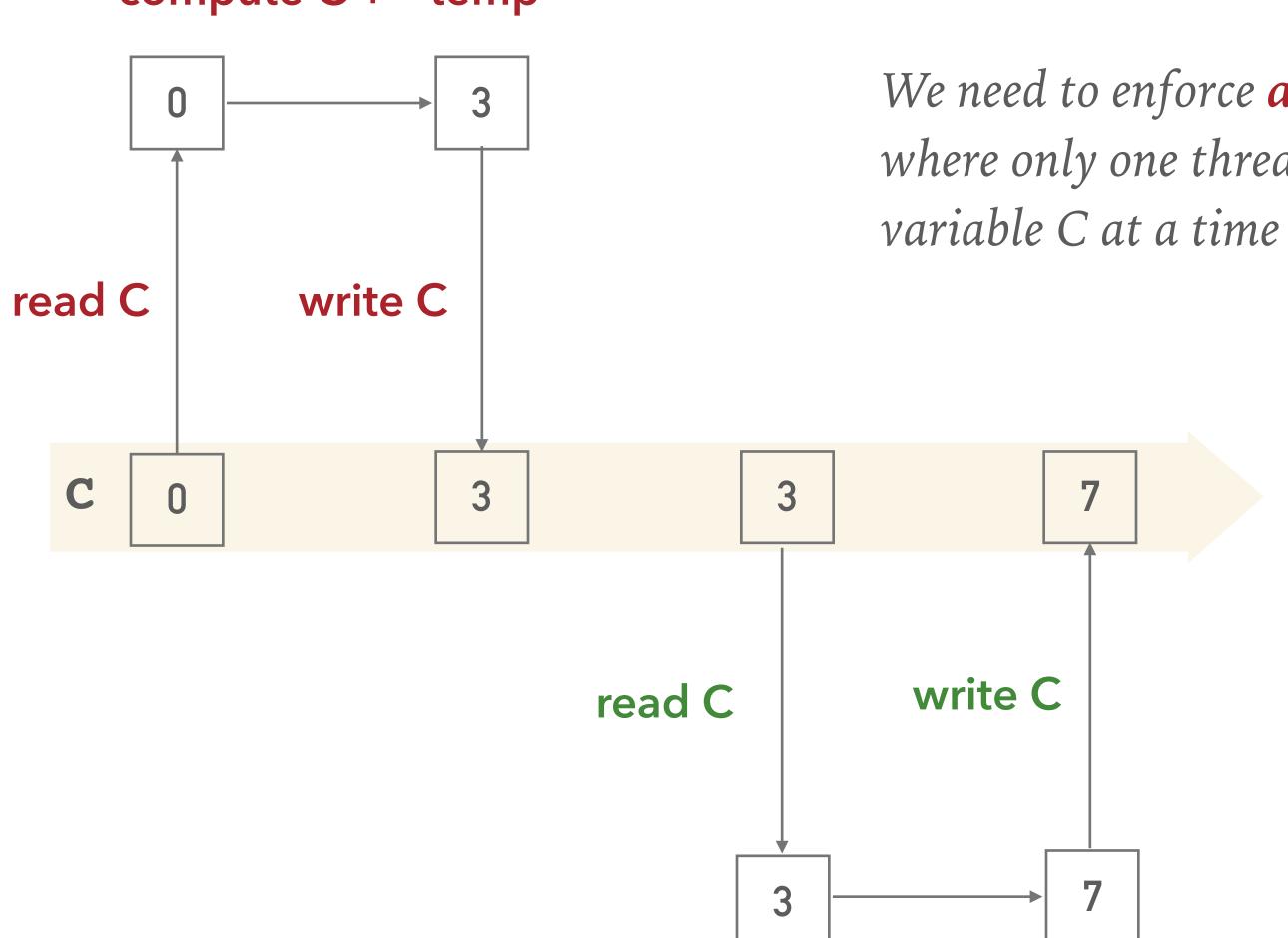
temp = 3

### global memory

C = 0

#### thread 1

temp = 4



We need to enforce atomic operation, where only one thread get's access to

compute C += temp

# ATOMIC OPERATIONS

Atomic operations ensure that only one thread can access and modify a memory location. Other threads need to wait until atomic operation is completed.

atom	nicAdd	addition
alun	$\Pi$ C $\Pi$ UU	MMILLOIL

➤ atomicSub subtraction

➤ atomicMin minimum

➤ atomicMax maximum

> atomicInc increment

➤ atomicDec decrement

➤ atomicExch exchange

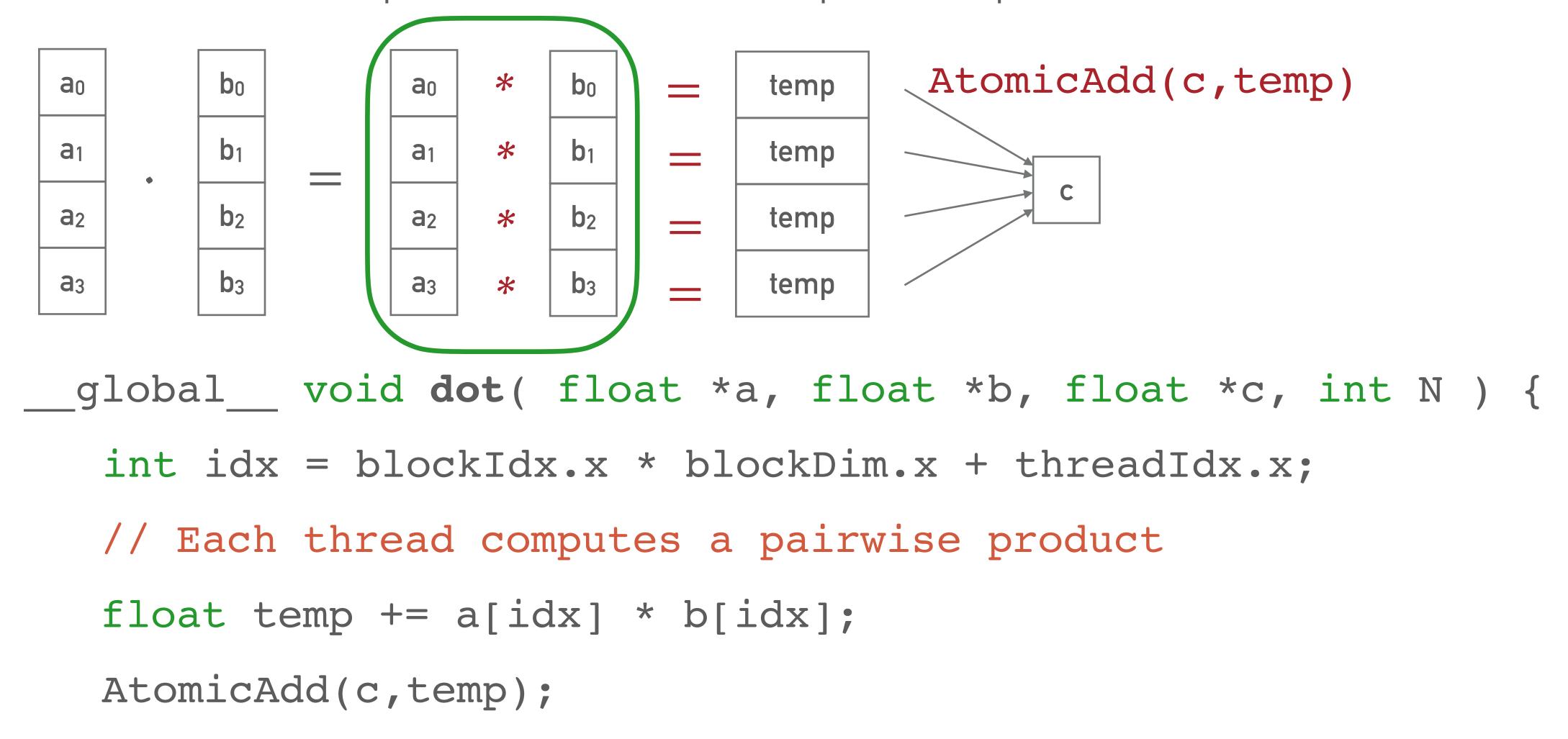
➤ atomicCAS compare and swap

atomicAdd\_block atomicAdd\_system

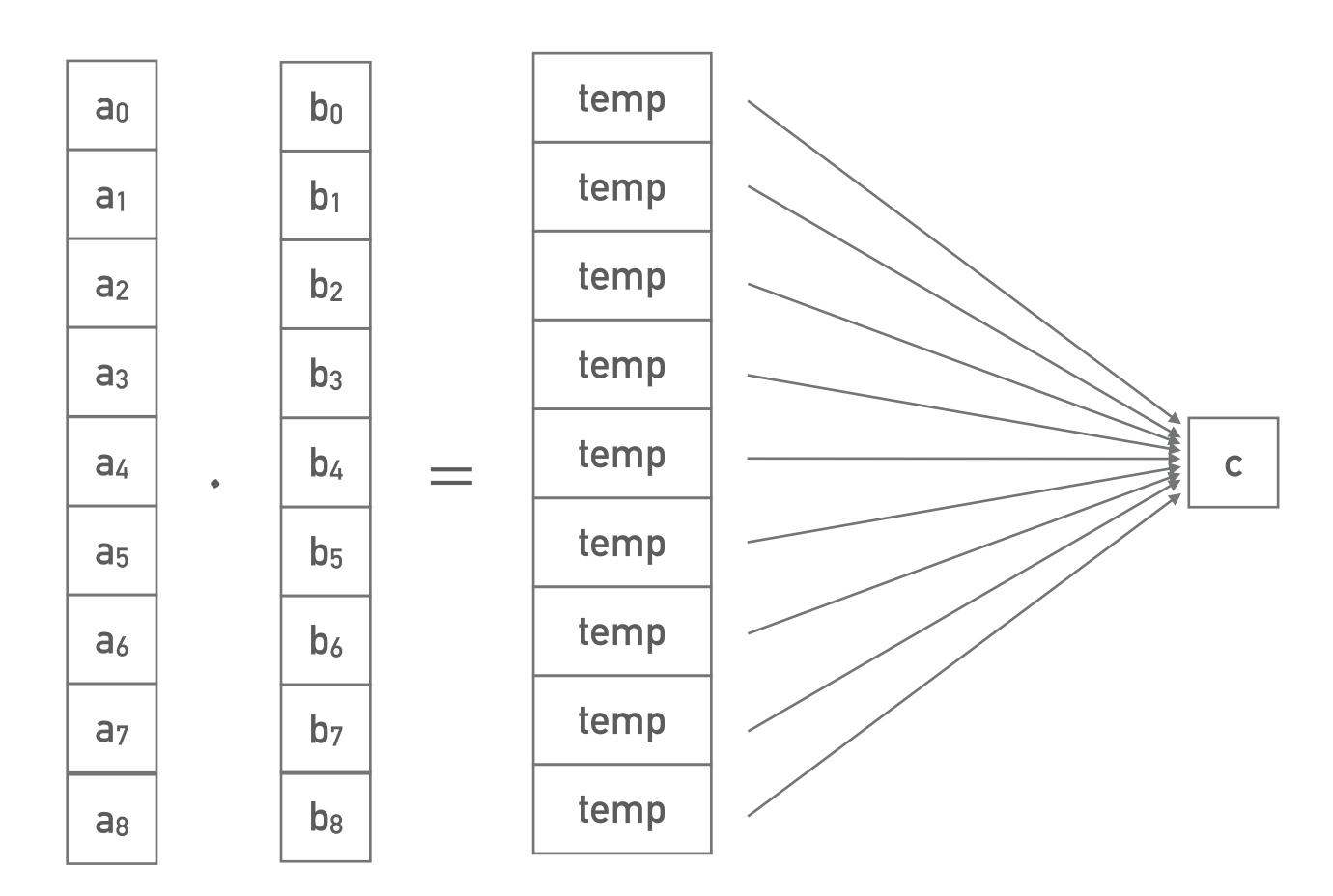
atomic with respect to block only

atomic with respect to GPU and CPU (for unified memory)

Parallel threads have no problem with the multiplication part:



# **BLOCKS AND THREADS**

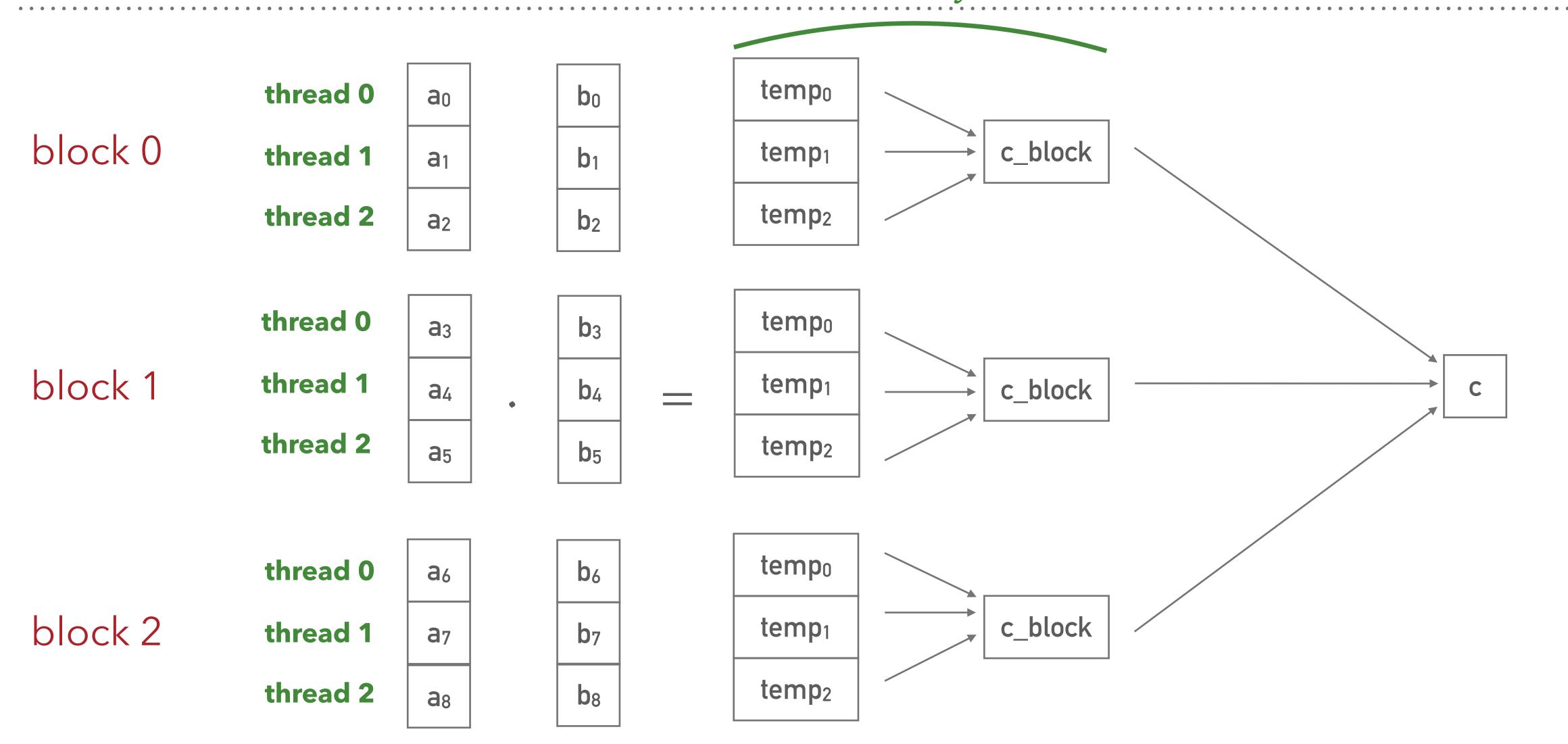


# **BLOCKS AND THREADS**

	thread 0	<b>a</b> <sub>0</sub>		b <sub>0</sub>		temp	
block 0	thread 1	<b>a</b> <sub>1</sub>		b <sub>1</sub>		temp	
	thread 2	<b>a</b> <sub>2</sub>		b <sub>2</sub>		temp	
			ı		l		
	thread 0	<b>a</b> <sub>3</sub>		<b>b</b> <sub>3</sub>		temp	
block 1	1 thread 1	<b>a</b> <sub>4</sub>	•	b <sub>4</sub>		temp	C
	thread 2	<b>a</b> <sub>5</sub>		<b>b</b> <sub>5</sub>		temp	
			1				
	thread 0	<b>a</b> <sub>6</sub>		b <sub>6</sub>		temp	
block 2	k 2 thread 1	<b>a</b> <sub>7</sub>		<b>b</b> <sub>7</sub>		temp	
	thread 2	<b>a</b> <sub>8</sub>		b <sub>8</sub>		temp	
			I	$\overline{}$	I		

# **BLOCKS AND THREADS**

shared memory within a block



### Shared memory:

- > is shared among the threads, but private to each block
- > is extremely fast (think cache, but shared among threads)

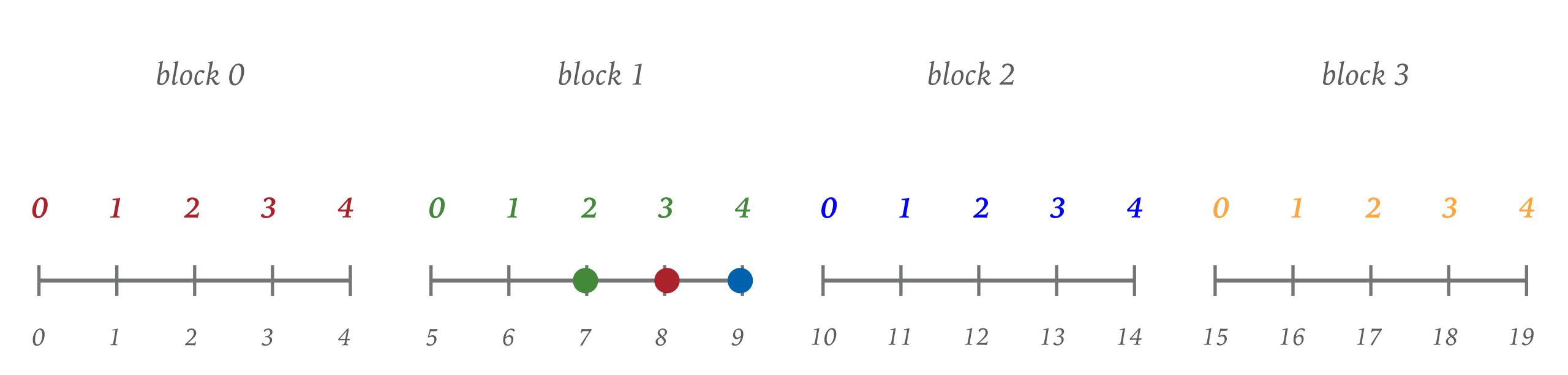
```
__global__ void dot( int *a, int *b, int *c, int N ) {
   int idx = blockIdx.x * blockDim.x + threadIdx.x;

// Each thread computes a pairwise product
   __shared__ int temp[THREADS_PER_BLOCK];
   temp[threadIdx.x] = a[idx] * b[idx];
}
```

```
global void dot( int *a, int *b, int *c, int N ) {
   int idx = blockIdx.x * blockDim.x + threadIdx.x;
// Each thread computes a pairwise product
   shared int temp[TREADS PER BLOCK];
   temp[threadIdx.x] = a[idx] * b[idx];
   syncthreads();
   // Thread 0 sums up the pairwise products
   if(threadIdx.x == 0) {
      float block c = 0;
      for (int i = 0; i < N; i++)
         block c += temp[i];
      atomicAdd(c,block c);
```

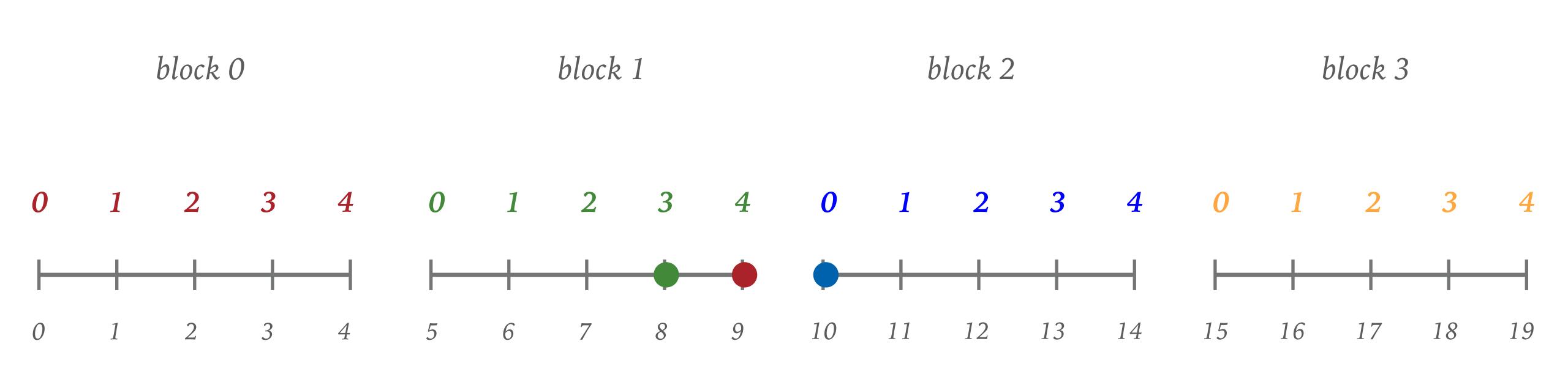
# FINITE DIFFERENCE ON GPUS

$$\frac{\partial u}{\partial x} \approx \frac{u_{i+1} - u_{i-1}}{2\Delta x}$$



# FINITE DIFFERENCE ON GPUS

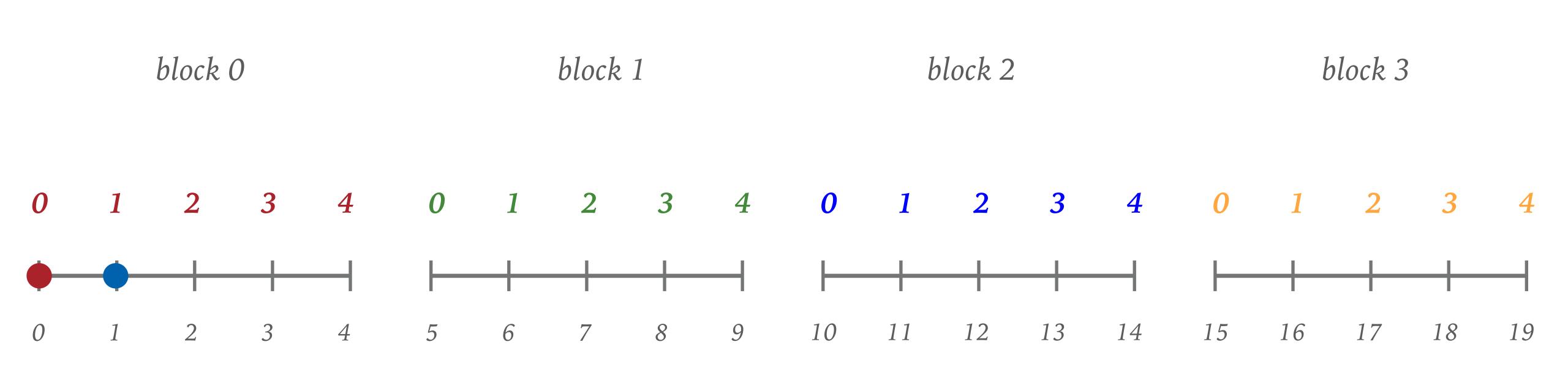
$$\frac{\partial u}{\partial x} \approx \frac{u_{i+1} - u_{i-1}}{2\Delta x}$$



this is fine, since we have global memory

# FINITE DIFFERENCE ON GPUS

$$\frac{\partial u}{\partial x} \approx \frac{u_{i+1} - u_{i-1}}{2\Delta x}$$



we need to do boundary points using one-sided difference, as usual

